

## EXISTENCE OF TRICONNECTED GRAPHS WITH PRESCRIBED DEGREES

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**Necessary and sufficient conditions for the existence of a  $p$ -connected (linear undirected) graph with prescribed degrees  $d_1, d_2, \dots, d_n$  are known for  $p = 1, 2$ . In this paper we solve this problem for  $p = 3$ .**

Let  $d_1, d_2, \dots, d_n$  be positive integers and let  $d_1 \leq d_2 \leq \dots \leq d_n$ .

LEMMA. *If a triconnected graph  $G$  exists with degrees  $d_1, d_2, \dots, d_n$ , then*

- (1)  $d_i \geq 3$ .
- (2)  $d_1, d_2, \dots, d_n$  is graphical, i.e., there exists a graph with these degrees.
- (3)  $d_n + d_{n-1} \leq m - n + 4$  where  $2m = \sum_{i=1}^n d_i$ .
- (4) If  $d_n + d_{n-1} = m - n + 4$ , then  $m \geq 2n - 2$ .

*Proof.* (1) and (2) are evident. To prove (3), let  $x_n, x_{n-1}$  be the vertices of  $G$  with degrees  $d_n$  and  $d_{n-1}$  respectively. Then the number of edges in  $G - \{x_n, x_{n-1}\}$  is  $m - (d_n + d_{n-1} - 1)$  or  $m - (d_n + d_{n-1})$  according as  $x_n, x_{n-1}$  are adjacent or not adjacent in  $G$ . Also  $G - \{x_n, x_{n-1}\}$  is connected, so (3) follows. If now  $d_n + d_{n-1} = m - n + 4$ , then

$$2m \geq d_n + d_{n-1} + 3(n - 2) = m + 2n - 2.$$

This completes the proof of the lemma.

THEOREM. *Conditions (1) to (4) of the lemma are necessary and sufficient for the existence of a triconnected graph with degrees  $d_1, d_2, \dots, d_n$ .*

*Proof.* Necessity was proved in the lemma.

To prove sufficiency, first let conditions (1), (3) be satisfied and let  $d_n + d_{n-1} = m - n + 4 = n + \lambda$  where  $2 \leq \lambda \leq n - 2$ . Let  $k$  be the number of  $d_i$  such that  $1 \leq i \leq n - 2$  and  $d_i = 3$ . Then define

$$e_i = d_i - 2 \text{ for } i = k + 1, \dots, n - 2.$$

Then we have

$$\sum_{i=1}^{n-2} d_i = 2m - d_n - d_{n-1} = 3n + \lambda - 8,$$

$$\sum_{i=k+1}^{n-2} e_i = 3n + \lambda - 8 - 3k - 2(n - 2 - k) = n + \lambda - k - 4.$$

Define now  $\eta = n - 2 - \lambda$  and  $\varepsilon = k - \eta$ . Then  $\eta \geq 0$ , and  $\varepsilon \geq 2$  since

$$\begin{aligned} 2m &\geq m - n + 4 + 3k + 4(n - 2 - k) \\ &= m + 3n - k - 4 \end{aligned}$$

and so

$$\lambda = m - 2n + 4 \geq n - k.$$

Write now

$$e_i = \begin{cases} 1 & \text{for } i = 1, 2, \dots, \varepsilon, \\ 2 & \text{for } i = \varepsilon + 1, \dots, k, \\ d_i - 2 & \text{for } i = k + 1, \dots, n - 2. \end{cases}$$

Then  $\sum_{i=1}^{n-2} e_i = 2(n-3)$  and so there exists a tree  $T$  with degrees  $e_1, \dots, e_{n-2}$ , attained by the vertices  $x_1, \dots, x_{n-2}$ , say, in that order [2]. Take two more vertices  $x_{n-1}$  and  $x_n$  and join them. Also join each of  $x_{n-1}, x_n$  to  $x_i$  for  $i = 1, \dots, \varepsilon, k + 1, \dots, n - 2$ . Of the  $\eta$  vertices  $x_{\varepsilon+1}, \dots, x_k$ , join  $d_{n-1} - 1 - \varepsilon - n + 2 + k$  to  $x_{n-1}$  and the rest ( $d_n - 1 - \varepsilon - n + 2 + k$  in number) to  $x_n$ . Note that

$$d_{n-1} - 1 - \varepsilon - n + 2 + k = d_{n-1} - \lambda - 1 \geq 0.$$

The graph we thus obtain has degrees  $d_1, \dots, d_n$  and is triconnected since any vertex of  $T$  with degree in  $T$  less than 3 is joined to either  $x_{n-1}$  or  $x_n$ .

Next let conditions (1), (2) be satisfied and let

$$d_n + d_{n-1} \leq m - n + 3.$$

Then  $d_n < m - n + 2$ , so there exists a biconnected graph  $G$  with degrees  $d_1, d_2, \dots, d_n$  [2]. If  $G$  is not triconnected, let  $x_i, x_j$  be two vertices such that  $G - \{x_i, x_j\}$  is disconnected. Let  $C_1, C_2, \dots$  be the components of  $G - \{x_i, x_j\}$ . By (1),  $|C_g| \geq 2$  for  $g = 1, 2, \dots$ . Also by hypothesis,

$$m - d_i - d_j \geq n - 3,$$

so it follows that one of the components, say  $C_1$ , contains a cycle.

We first prove that there exists an edge  $(x, y)$  in  $C_1$  and two chains  $\mu_1, \mu'_1$  of  $G$  connecting  $x$  and  $y$  such that  $(x, y), \mu_1, \mu'_1$  are disjoint except for  $x$  and  $y$ , and  $\mu_1$  is contained in  $C_1$ . Since  $G$  is biconnected, there exists a chain connecting  $x_i$  and  $x_j$  with all intermediate vertices in  $C_2$ .

If now two vertices  $x, y$  with degree two in  $C_1$  are adjacent and belong to a cycle of  $C_1$ , the required edge is  $(x, y)$ . So we may take

that no two vertices of degree two in  $C_1$  can belong to a block (on more than two vertices) and be adjacent. Let  $B$  be any block of  $C_1$  which is not an edge. If some cycle of  $B$  has a chord  $(x, y)$ , then  $(x, y)$  is the required edge. Otherwise, by the results of [1], two vertices  $y, z$  of degree two in  $B$  will be adjacent to a vertex  $x$  of degree three in  $B$ . If  $w$  is another vertex of  $B$  adjacent to  $x$ , then there is a chain connecting  $w$  to  $y$  in  $B - \{x\}$ . This chain together with  $(x, w)$  may be taken as  $\mu_1$ . To get  $\mu'_1$ , go from  $x$  to  $z$  along  $(x, z)$ , from  $z$  to  $x_i$  or  $x_j$  (through another block of  $C_1$  at  $z$  if necessary), then to  $y$ . Thus  $(x, y)$  is the required edge.

Let now  $(x, y)$  be an edge of  $C_1$  chosen as explained above. If  $C_2$  is a tree, take any edge  $(u, v)$  of  $C_2$ . Then  $(u, v)$  is a chord of a cycle of  $G$ . If  $C_2$  is not a tree, choose an edge  $(u, v)$  of  $C_2$  such that there are chains  $\mu_2, \mu'_2$  of  $G$  connecting  $u$  and  $v$ ,  $(u, v), \mu_2, \mu'_2$  are disjoint except for  $u, v$ , and  $\mu_2$  is contained in  $C_2$ .

We define  $f_G(s, t)$  to be the number of components of  $G - \{s, t\}$ . Now we will make a modification on  $G$  so that the degrees of the vertices are unaltered,  $f(x_i, x_j)$  decreases and  $f(s, t)$  does not increase for any two vertices  $s$  and  $t$ .

First we associate with  $x$ , a subset  $A(x)$  of  $\{x_i, x_j\}$  by the following rule.  $x_i \in A(x)$  if and only if there is a chain  $\nu$  connecting  $x$  to  $x_i$  with all intermediate vertices in  $C_1$  such that  $\nu$  is disjoint with  $(x, y)$  and  $\mu_1$  except for  $x$ . Similarly  $A(y)$  is defined. If  $C_2$  is a tree, put  $A(u) = A(v) = \{x_i, x_j\}$ . Otherwise  $A(u), A(v)$  are defined in a manner similar to that of  $A(x)$  and  $A(y)$ . Now  $A(x), A(y)$  are made nonempty by a proper choice of  $\mu_1$ , and  $A(u), A(v)$  are made nonempty by a proper choice of  $\mu_2$  (in case  $C_2$  is not a tree).

Now suppress the edges  $(x, y), (u, v)$  and join  $x$  to one of  $u, v$  and  $y$  to the other as follows. Join  $x$  to  $u$  if  $A(x) \neq A(u)$  and  $A(y) \neq A(v)$  whenever such a choice is possible. Let the new graph thus obtained be  $H$ . To be specific we take that  $x$  is joined to  $u$  in  $H$ .

First we show that  $H$  is biconnected. Obviously  $G_1 = G - (x, y)$  is biconnected. Now we show that  $(u, v)$  is a chord of a cycle of  $G_1$ . If  $C_2$  is a tree, then the cycle is

$$(u, x) + \mu_1[x, y] + (y, v) + [v, \dots, p_1] + (p_1, x_i) + (x_i, p_2) + [p_2, \dots, u]$$

where  $p_1, p_2$  are suitable pendant vertices of  $C_2$ . Otherwise the cycle is

$$\mu_2[u, v] + \mu'_2[v, u]$$

where if  $\mu'_2$  contains the edge  $(x, y)$ , then  $(x, y)$  is replaced by  $\mu_1[x, y]$  and the resulting cycle is made elementary.

Trivially now  $f_G(x_i, x_j) = f_H(x_i, x_j) + 1$ . Next we will show that

$$(5) \quad f_G(s, t) \geq f_H(s, t)$$

for any two vertices  $s$  and  $t$ . For this it is enough to show that  $x, y$  are connected and  $u, v$  are connected in  $H - \{s, t\}$ .

First let  $s = x_i$ . Now  $x, y, u, v$  belong to a cycle in  $H - \{x_i\}$ , so (5) follows. So we may take  $\{s, t\} \cap \{x_i, x_j\} = \emptyset$ .

Now let  $s = x$ . Then to prove (5) it is enough to show that  $u, v$  are connected in  $H - \{x, t\}$  when  $t \neq u$  and  $t \neq v$ . This is evident if  $C_2$  is a tree or  $t \notin \mu_2$ . So let  $t \in \mu_2$  and  $C_2$  be not a tree. If  $A(u) \cap A(v) \neq \emptyset$ , there is a chain connecting  $u, v$  in  $H - \{x, t\}$ . So we take without loss of generality  $A(u) = x_j$  and  $A(v) = x_i$ . If now  $x_j \in A(y)$ , then  $u, v$  are connected through  $x_j$  and  $y$  in  $H - \{x, t\}$ . So we take  $A(y) = x_i$ . If  $x_j \in A(x)$ , then  $y$  would not have been joined to  $v$ , so  $A(x) = x_i$ . Now in  $G$ ,  $x_j$  is connected to some vertex  $z$  of  $\mu_1$  by a chain with all intermediate vertices belonging to  $C_1$  but not to  $\mu_1$ . Now we obtain a chain connecting  $u, v$  in  $H - \{x, t\}$  by going from  $u$  to  $x_j$ ,  $x_j$  to  $z$ ,  $z$  to  $y$  along  $\mu_1$ ,  $y$  to  $x_i$ , and  $x_i$  to  $v$ . Thus we may take  $\{s, t\} \cap \{x_i, x_j, x, y\} = \emptyset$ .

Next let  $s = u$ . If  $t \notin \mu_1$ , then (5) is trivial, so let  $t \in \mu_1$ . Suppose first that  $C_2$  is a tree. Then we obtain a chain connecting  $x, y$  in  $H - \{u, t\}$  by going from  $x$  to  $x_i$  or  $x_j$ , then to  $v$  through a suitable pendant vertex of  $C_2$  and then to  $y$ . If  $C_2$  is not a tree, the situation is similar to that of the preceding paragraph. Thus we take  $\{s, t\} \cap \{x_i, x_j, x, y, u, v\} = \emptyset$ .

If none of  $s, t$  belongs to  $\mu_1$ , then (5) is trivial. So let  $s \in \mu_1$ .

Suppose now that  $C_2$  is a tree. Then for any fixed vertex  $t$ , there are chains in  $H - \{s, t\}$  from one of  $u, v$  to both  $x_i$  and  $x_j$ , and a chain from the other (of the vertices  $u, v$ ) to  $x_i$  or  $x_j$ . Hence  $u, v$  are connected and (5) follows.

Suppose next that  $C_2$  is not a tree. Obviously we may take  $s \in \mu_1$  and  $t \in \mu_2$ . If now  $A(x) \cap A(y) \neq \emptyset$  or  $A(u) \cap A(v) \neq \emptyset$ , then again (5) follows. So we may take  $A(x) = x_i$ ,  $A(y) = x_j$ ,  $A(u) = x_j$ ,  $A(v) = x_i$ . Now we obtain a chain connecting  $x, y$  in  $H - \{s, t\}$  by going from  $x$  to  $u$ ,  $u$  to  $x_j$ ,  $x_j$  to  $y$ . This proves (5) completely.

Now by a repeated application of the above procedure we reduce the graph until finally  $f(s, t) = 1$  for any two vertices. The final graph has degrees  $d_1, d_2, \dots, d_n$  and is triconnected and this completes the proof of the theorem.

Perhaps necessary and sufficient conditions, similar to the conditions (1) to (4) above, for the existence of a  $p$ -connected graph with prescribed degrees  $d_1, d_2, \dots, d_n$  can be obtained for all  $p \geq 3$ , but the authors have not yet succeeded in this.

## REFERENCES

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