

Note on Miller's Recurrence Algorithm

Hisayoshi SHINTANI

(Received March 10, 1965)

1. Introduction

In this paper, we are concerned with a recurrence algorithm, originated by J. C. P. Miller [1]¹⁾, for computing a solution f_n of a second-order difference equation

$$(1.1) \quad y_{n-1} = a_n y_n + b_n y_{n+1} \quad (b_n \neq 0; n = 1, 2, \dots),$$

in the case where (1.1) has a second solution g_n which ultimately grows much faster than f_n [6]. This algorithm is used for computing Bessel functions [1, 2, 4, 9], Legendre functions [8], repeated integrals of the error function [3], and so on.

Let $P_n(k)$ be defined by the formula

$$(1.2) \quad P_n(k-1) = a_k P_n(k) + b_k P_n(k+1) \quad (k = n+1, n, \dots, 1),$$

where

$$(1.3) \quad P_n(n) = 1, \quad P_n(n+1) = 0, \quad P_n(n+2) = 1/b_{n+1}.$$

Then Miller's algorithm is applied in the following two ways:

1°. when the normalizing condition

$$(1.4) \quad m_0 f_0 + m_1 f_1 + \dots = c \quad (c \neq 0)$$

is known, put

$$(1.5) \quad S_n(k) = \frac{c P_n(k)}{R_n} \quad (k = 0, 1, \dots, n),$$

where

$$(1.6) \quad R_n = \sum_{j=0}^n m_j P_n(j).$$

2°. when f_0 is known and $f_0 \neq 0$, put

1) Numbers in square brackets refer to the references listed at the end of this paper.

$$(1.7) \quad T_n(k) = \frac{f_0 P_n(k)}{P_n(0)} \quad (k = 0, 1, \dots, n).$$

Under suitable conditions, which are reported to have been obtained by Gautschi [3, 6], it is valid that

$$(1.8) \quad S_n(k), T_n(k) \rightarrow f_k \quad \text{as } n \rightarrow \infty.$$

In the sequel, we consider the case where (1.8) holds.

There arises the question how large n should be in order to obtain the approximate values of f_k ($k = 0, 1, \dots, N+1$) to the desired accuracy. Such a value n will depend on the value N , the desired accuracy, the coefficients a 's and b 's, and so on. Until now theoretical bounds for the starting value n have been obtained for spherical Bessel functions [2], and for the repeated integrals of the error function [3], and empirical bounds have been obtained for $J_k(x)$ [5]. In the case where such a bound is not known, usually Miller's algorithm is applied repeatedly for different values of n ; the results obtained are compared in accordance with a preassigned tolerance and the process is repeated with n increased by a fixed amount until the criteria for acceptance are satisfied [9].

In the first part of this paper, recurrence formulas are derived for generating $P_n(k)$ and R_n for increasing n with k fixed. By generating $S_n(N)$ and $S_n(N+1)$ for increasing n through these formulas, the approximate values of f_N and f_{N+1} can be obtained to the desired accuracy and then the approximate values of f_k ($k = N-1, \dots, 1, 0$) can be generated through (1.2). This process seems to be more efficient than the above iterative process.

In the second part of this paper, we consider the case where $a_r > 0$ and $b_r > 0$ ($r = 1, 2, \dots$), and show the methods for generating the approximate values of f_k ($k = 0, 1, \dots, N+1$) to the desired relative accuracy.

2. Recurrence formulas

We shall first show the following

THEOREM 1. $P_n(k)$ ($n = k-1, k, \dots$) satisfy the recurrence formula

$$(2.1) \quad P_{n+1}(k) = a_{n+1}P_n(k) + b_nP_{n-1}(k) \quad (n+1 \geq k \geq 0),$$

where

$$(2.2) \quad P_k(k) = 1, \quad P_{k-1}(k) = 0, \quad P_{k-2}(k) = 1/b_{k-1}.$$

Proof. Since by definition

$$(2.3) \quad P_{n-1}(n+1) = 1/b_n, \quad P_n(n+1) = 0, \quad P_{n+1}(n+1) = 1,$$

and

$$(2.4) \quad P_{n-1}(n) = 0, \quad P_n(n) = 1, \quad P_{n+1}(n) = a_{n+1},$$

(2.1) is valid for $k = n + 1, n, \dots, q$ ($q > 0$). Hence suppose that (2.1) holds for $k = n + 1, n, \dots, q$ ($q > 0$). Then we have

$$(2.5) \quad P_{n+1}(q) = a_{n+1}P_n(q) + b_nP_{n-1}(q),$$

and

$$(2.6) \quad P_{n+1}(q+1) = a_{n+1}P_n(q+1) + b_nP_{n-1}(q+1).$$

On the other hand, from (1.2) it follows that

$$(2.7) \quad P_{n+1}(q-1) = a_qP_{n+1}(q) + b_qP_{n+1}(q+1).$$

Substituting (2.5) and (2.6) into (2.7), we obtain

$$(2.8) \quad \begin{aligned} P_{n+1}(q-1) &= a_{n+1}[a_qP_n(q) + b_qP_n(q+1)] + \\ &\quad + b_n[a_qP_{n-1}(q) + b_qP_{n-1}(q+1)] \\ &= a_{n+1}P_n(q-1) + b_nP_{n-1}(q-1). \end{aligned}$$

This proves the theorem.

Next put

$$(2.9) \quad U_n(k) = \sum_{j=k}^n m_j P_n(j) \quad (n \geq k).$$

Then we have the following

THEOREM 2. $U_n(k)$ ($n = k, k + 1, \dots$) satisfy the recurrence formula

$$(2.10) \quad U_{n+1}(k) = a_{n+1}U_n(k) + b_nU_{n-1}(k) + m_{n+1} \quad (n \geq k),$$

where

$$(2.11) \quad U_{k-1}(k) = 0, \quad U_k(k) = m_k.$$

Proof. Since

$$(2.12) \quad U_k(k) = m_k P_k(k) = m_k$$

and

$$(2.13) \quad U_{k+1}(k) = m_k P_{k+1}(k) + m_{k+1} P_{k+1}(k+1) = m_k a_{k+1} + m_{k+1},$$

(2.10) is valid for $n = k$.

For $n > k$, we have

$$(2.14) \quad \begin{aligned} U_{n+1}(k) &= \sum_{j=k}^n m_j P_{n+1}(j) + m_{n+1} P_{n+1}(n+1) \\ &= \sum_{j=k}^n m_j [a_{n+1} P_n(j) + b_n P_{n-1}(j)] + m_{n+1} \\ &= a_{n+1} \sum_{j=k}^n m_j P_n(j) + b_n \sum_{j=k}^{n-1} m_j P_{n-1}(j) + m_{n+1} \\ &= a_{n+1} U_n(k) + b_n U_{n-1}(k) + m_{n+1}, \end{aligned}$$

because

$$(2.15) \quad P_{n+1}(n+1) = 1, \quad P_{n-1}(n) = 0.$$

Thus the theorem has been proved.

Since $R_n = U_n(0)$, from this theorem we obtain the following

COROLLARY. $R_n (n = 0, 1, 2, \dots)$ satisfy the recurrence formula

$$(2.16) \quad R_{n+1} = a_{n+1} R_n + b_n R_{n-1} + m_{n+1} \quad (n \geq 0),$$

where

$$(2.17) \quad R_{-1} = 0, \quad R_0 = m_0.$$

Making use of (2.1) and (2.16), we can generate $P_n(k)$, $S_n(k)$, and $T_n(k)$ for increasing n with k fixed. Hence, for a specified k , we can obtain the approximate value of f_k to the desired accuracy by increasing n .

When $P_n(N)$ and $P_n(N+1)$ are obtained by means of (2.1), we can use (1.2) to generate $P_n(j)$ ($j = N-1, \dots, 1, 0$). In that case, if $U_n(N)$ is computed, R_n can be obtained by the formula

$$(2.18) \quad R_n = \sum_{j=0}^{N-1} m_j P_n(j) + U_n(N).$$

Since by (1.5), (1.7) and (1.8)

$$(2.19) \quad \frac{P_n(k+1)}{P_n(k)} \rightarrow \frac{f_{k+1}}{f_k} \quad \text{as } n \rightarrow \infty \quad (f_k \neq 0),$$

we have the following

THEOREM 3. *The ratio f_{k+1}/f_k ($f_k \neq 0$) can be expanded into the continued fraction as follows:*

$$(2.20) \quad \frac{f_{k+1}}{f_k} = \frac{1}{a_{k+1} + \frac{b_{k+1}}{a_{k+2} + \frac{b_{k+2}}{a_{k+3} + \dots}}} \quad (k \geq 0).$$

Proof. By (2.19) it suffices to show that

$$(2.21) \quad \frac{P_{m+1}(k+1)}{P_{m+1}(k)} = \frac{1}{a_{k+1} + \frac{b_{k+1}}{a_{k+2} + \dots + \frac{b_{m-1}}{a_m + \frac{b_m}{a_{m+1}}}}} \quad (m = k + 1, k + 2, \dots).$$

By (2.1), (2.3) and (2.4) we have

$$(2.22) \quad \frac{P_{k+2}(k+1)}{P_{k+2}(k)} = \frac{a_{k+2}}{a_{k+2}a_{k+1} + b_{k+1}} = \frac{1}{a_{k+1} + \frac{b_{k+1}}{a_{k+2}}},$$

so that (2.21) is valid for $m = k + 1$. Hence suppose that (2.21) holds for $m = k + 1, k + 2, \dots, n$. Then it is valid that

$$(2.23) \quad \frac{a_{n+1}P_n(k+1) + b_nP_{n-1}(k+1)}{a_{n+1}P_n(k) + b_nP_{n-1}(k)} = \frac{1}{a_{k+1} + \frac{b_{k+1}}{a_{k+2} + \dots + \frac{b_{n-1}}{a_n + \frac{b_n}{a_{n+1}}}}}.$$

Replacing a_{n+1} and b_n in (2.23) with $a_{n+2}a_{n+1} + b_{n+1}$ and $a_{n+2}b_n$ respectively, we have

$$(2.24) \quad \frac{(a_{n+2}a_{n+1} + b_{n+1})P_n(k+1) + a_{n+2}b_nP_{n-1}(k+1)}{(a_{n+2}a_{n+1} + b_{n+1})P_n(k) + a_{n+2}b_nP_{n-1}(k)} \\ = \frac{1}{a_{k+1} + \frac{b_{k+1}}{a_{k+2} + \dots + \frac{b_{n-1}}{a_n + \frac{a_{n+2}b_n}{a_{n+2}a_{n+1} + b_{n+1}}}}}.$$

Since by (2.1)

$$(2.25) \quad (a_{n+2}a_{n+1} + b_{n+1})P_n(r) + a_{n+2}b_nP_{n-1}(r) \quad (r = k, k + 1) \\ = a_{n+2}[a_{n+1}P_n(r) + b_nP_{n-1}(r)] + b_{n+1}P_{n-1}(r) \\ = a_{n+2}P_{n+1}(r) + b_{n+1}P_{n+1}(r) = P_{n+2}(r),$$

and

$$(2.26) \quad \frac{a_{n+2}b_n}{a_{n+2}a_{n+1} + b_{n+1}} = \frac{b_n}{a_{n+1} + \frac{b_{n+1}}{a_{n+2}}},$$

(2.21) is valid also for $m = n + 1$.

Now we shall show the examples to which the above results can be applied.

EXAMPLE 1. Bessel functions of the first kind $J_k(x)$ ($k = 0, 1, \dots$) satisfy the recurrence formula [9]

$$(2.27) \quad J_{n-1}(x) = \frac{2n}{x} J_n(x) - J_{n+1}(x)$$

with the normalizing condition

$$(2.28) \quad J_0(x) + 2 \sum_{k=1}^{\infty} J_{2k}(x) = 1.$$

Hence we can use (2.1) and (2.16) to obtain the approximate values of $J_0(x)$ and $J_1(x)$ to the desired accuracy without knowing previously the starting value n . They can be used also to determine the empirical bound for the starting value n for $J_0(x)$ and $J_1(x)$. Once such a bound is obtained, we can use (1.2) to generate the approximate values of $J_0(x)$, $J_1(x)$ and so on efficiently.

EXAMPLE 2. Let

$$(2.29) \quad i^n \operatorname{erfc} x = \int_x^{\infty} i^{n-1} \operatorname{erfc} t \, dt \quad (n = 0, 1, \dots),$$

where

$$(2.30) \quad i^{-1} \operatorname{erfc} x = \frac{2}{\sqrt{\pi}} e^{-x^2},$$

and put

$$(2.31) \quad y_n = i^{n-1} \operatorname{erfc} x.$$

Then y_n ($n = 0, 1, \dots$) satisfy the recurrence formula [3]

$$(2.32) \quad y_{n-1} = 2xy_n + 2ny_{n+1}.$$

Since $y_1 = \operatorname{erfc} x$, it is valid that

$$(2.33) \quad T_n(1) \rightarrow \operatorname{erfc} x \quad \text{as } n \rightarrow \infty,$$

where

$$(2.34) \quad T_n(1) = \frac{2}{\sqrt{\pi}} e^{-x^2} \frac{P_n(1)}{P_n(0)}.$$

Hence we can use (2.1) to obtain the approximate value of $\operatorname{erfc} x$.

On the other hand, from (6.20) it follows that

$$(2.35) \quad \operatorname{erfc} x = \frac{2}{\sqrt{\pi}} e^{-x^2} \left[\frac{1}{2x+} \frac{2}{2x+} \frac{4}{2x+} \frac{6}{2x+} \dots \right].$$

J. Patry and J. Keller [7] obtained the expansion

$$(2.36) \quad \operatorname{erfc} x = e^{-x^2} \left[\frac{1}{c_0x+} \frac{1}{c_1x+} \frac{1}{c_2x+} \dots \right]$$

where

$$(2.37) \quad c_0 = \sqrt{\pi}, \quad c_1 = \frac{2}{\sqrt{\pi}}, \quad c_{n+1} = \frac{2}{c_n+} \frac{2}{c_{n-1}}.$$

As is easily seen, this is equivalent to (2.35), but (2.35) is simpler than (2.36).

3. Case of positive coefficients

In this paragraph, we are concerned with the case where

$$(3.1) \quad a_n > 0, \quad b_n > 0 \quad (n = 1, 2, \dots).$$

This condition is satisfied, for instance, by the recurrence formulas for $I_n(x)$, $i_n(x)$ and $i^n \operatorname{erfc} x$. Our problem is how to generate f_k ($k = 0, 1, \dots, N+1$) to the desired relative accuracy. To that end we need the following

LEMMA. *Put*

$$(3.2) \quad \frac{f_{N+1}}{f_N} = r$$

and

$$(3.3) \quad \frac{P_n(N+1)}{P_n(N)} = r_n = r(1 + e_n) \quad (n \geq N+1).$$

Then it is valid that

$$(3.4) \quad \frac{P_n(k)}{f_k} = (1 + d_k e_n) \frac{P_n(N)}{f_N} \quad (k = 0, 1, \dots, N+1),$$

$$(3.5) \quad 0 = d_N < d_{N-2} < \dots < d_0 < \dots < d_{N-1} < d_{N+1} = 1,$$

and

$$(3.6) \quad -1 = e_N < e_{N+2} < \dots < 0 < \dots < e_{N+3} < e_{N+1},$$

where

$$(3.7) \quad d_k = \frac{rb_N P_{N-1}(k)}{P_N(k) + rb_N P_{N-1}(k)}.$$

Proof. It is easy to show by induction that

$$(3.8) \quad P_n(k) = P_n(N)P_N(k) + b_N P_n(N+1)P_{N-1}(k)$$

and

$$(3.9) \quad f_k = f_N P_N(k) + b_N f_{N+1} P_{N-1}(k) \quad (k = 0, 1, \dots, N+1).$$

From these, (3.2) and (3.3) we have

$$(3.10) \quad \frac{P_n(k)}{f_k} = \frac{P_n(N)}{f_N} \cdot \frac{P_N(k) + r(1 + e_n)b_N P_{N-1}(k)}{P_N(k) + rb_N P_{N-1}(k)} \\ = \frac{P_n(N)}{f_N} \left[1 + \frac{rb_N P_{N-1}(k)}{P_N(k) + rb_N P_{N-1}(k)} \cdot e_n \right].$$

This proves (3.4).

Next, substituting (3.3) into

$$(3.11) \quad P_{n+1}(N+1) = a_{n+1}P_n(N+1) + b_n P_{n-1}(N+1),$$

we have

$$(3.12) \quad r_{n+1}P_{n+1}(N) = r_n a_{n+1}P_n(N) + r_{n-1}b_n P_{n-1}(N).$$

From this and

$$(3.13) \quad P_{n+1}(N) = a_{n+1}P_n(N) + b_n P_{n-1}(N),$$

it follows that

$$(3.14) \quad r_{n+1} - r_n = (r_{n-1} - r_n) \frac{b_n P_{n-1}(N)}{P_{n+1}(N)}$$

and

$$(3.15) \quad r_{n+1} - r_{n-1} = (r_n - r_{n-1}) \frac{a_{n+1}P_n(N)}{P_{n+1}(N)}.$$

Since

$$(3.16) \quad r_N = 0, \quad r_{N+1} = \frac{1}{a_{N+1}}, \quad P_m(N) > 0 \quad (m \geq N),$$

from (3.14) and (3.15) we have

$$(3.17) \quad 0 = r_N < r_{N+2} < \dots < r < \dots < r_{N+3} < r_{N+1}.$$

This proves (3.6).

Lastly, from (3.7) we can easily deduce the relations

$$(3.18) \quad d_k - d_{k-1} = (d_k - d_{k+1})b_k \frac{P_N(k+1) + r b_N P_{N-1}(k+1)}{P_n(k-1) + r b_N P_{N-1}(k-1)}$$

and

$$(3.19) \quad d_{k+1} - d_{k-1} = (d_{k-1} - d_k)a_k \frac{P_N(k) + r b_N P_{N-1}(k)}{P_n(k+1) + r b_N P_{N-1}(k+1)}.$$

Since $d_N = 0$ and $d_{N+1} = 1$, from (3.18) and (3.19) follows (3.5). This completes the proof of the lemma.

Now we shall show the following

THEOREM 4. *Let*

$$(3.20) \quad f_0^* = f_0(1+c) \quad (|c| \leq c_0 < 1)$$

and

$$(3.21) \quad T_n^*(k) = \frac{f_0^* P_n(k)}{P_n(0)} = (1 + e_{n,k})f_k \quad (k = 0, 1, \dots, N+1),$$

where f_0^* is an approximate value of f_0 . Then, for n such that

$$(3.22) \quad n = N + 1 + 2q \quad (q \geq 1),$$

it is valid that

$$(3.23) \quad (1+c)(1-e_n) < 1 + e_{n,k} < (1+c)(1+e_n) \quad (k = 0, 1, \dots, N+1)$$

and

$$(3.24) \quad \frac{r_n - r_{n+2}}{r_{n+2}} < e_n < \frac{r_n - r_{n-1}}{r_{n-1}}.$$

Proof. From (3.20), (3.21) and (3.4), it follows that

$$(3.25) \quad 1 + e_{n,k} = (1+c) \frac{1+d_k e_n}{1+d_0 e_n} = (1+c) \left[1 + \frac{(d_k - d_0)}{1+d_0 e_n} \cdot e_n \right].$$

We consider the case where n satisfies (3.22). Then it is valid that

$$(3.26) \quad 1 + d_0 e_n > 1,$$

because $d_0 > 0$ and $e_n > 0$ by (3.5) and (3.6). Further, from (3.5), it follows that

$$(3.27) \quad |d_k - d_0| < 1.$$

Hence we have the inequality

$$(3.28) \quad \left| \frac{d_k - d_0}{1 + d_0 e_n} \right| < 1,$$

and (3.23) is proved.

On the other hand, since by (3.6)

$$(3.29) \quad -1 < e_{n-1} < 0, \quad 0 < e_{n+2} < e_n,$$

we have

$$(3.30) \quad \frac{r_n - r_{n-1}}{r_{n-1}} = \frac{e_n - e_{n-1}}{1 + e_{n-1}} > e_n,$$

and

$$(3.31) \quad \frac{r_n - r_{n+2}}{r_{n+2}} = \frac{e_n - e_{n+2}}{e_{n+2}} < e_n.$$

Thus the theorem has been proved.

Now, by (3.24), it holds that

$$(3.32) \quad e_{n,k} < e_n(1 + c) + c \leq e_n(1 + c_0) + c_0$$

and

$$(3.33) \quad e_{n,k} > -e_n(1 + c) + c \geq -e_n(1 + c_0) - c_0.$$

Hence we have the following

COROLLARY. *Under the condition (3.22), if for a positive number μ ($\mu > c_0$)*

$$(3.34) \quad e_n \leq \frac{\mu - c_0}{1 + c_0},$$

then the inequality

$$(3.35) \quad |e_{n,k}| < \mu$$

is valid for $k=0, 1, \dots, N+1$.

Next we shall show the following

THEOREM 5. *Let*

$$(3.36) \quad S_n(k) = f_k(1 + s_{n,k})$$

and suppose that, for a positive number μ ,

$$(3.37) \quad |s_{n,N}| \leq \mu, \quad |s_{n,N+1}| \leq \mu.$$

Then the inequality

$$(3.38) \quad |s_{n,k}| \leq \mu$$

is valid for $k=0, 1, \dots, N+1$.

Proof. From (3.4) and (1.5) it follows that

$$(3.39) \quad s_{n,k} = s_{n,N} + d_k e_n(1 + s_{n,N}) \quad (k = 0, 1, \dots, N+1)$$

and

$$(3.40) \quad e_n(1 + s_{n,N}) = s_{n,N+1} - s_{n,N},$$

because $d_{N+1}=1$. Substituting (3.40) into (3.39), we obtain

$$(3.41) \quad s_{n,k} = d_k s_{n,N+1} + (1 - d_k) s_{n,N}.$$

Since by (3.5)

$$(3.42) \quad 0 \leq d_k \leq 1 \quad (k = 0, 1, \dots, N+1),$$

from (3.41) and (3.37) follows (3.38).

Now we are in a position to apply theorems 4 and 5 for generating the approximate values of f_k ($k=0, 1, \dots, N+1$) such that

$$(3.43) \quad |s_{n,k}| \leq \mu \quad (k = 0, 1, \dots, N+1)$$

for a preassigned positive number μ . For this purpose, the following three methods can be considered.

Method 1. Generate $P_n(0)$, R_n , $P_n(N)$ and $P_n(N+1)$ for increasing n until

the inequalities

$$(3.44) \quad |s_{n,0}| \leq \frac{\mu}{2}$$

and

$$(3.45) \quad |e_n| \leq \frac{\mu}{2+\mu}$$

are valid for n satisfying (3.22), and then compute $S_n(k)$ ($k=0, 1, \dots, N+1$) by (1.2).

Method 2. Generate R_n , $P_n(N)$ and $P_n(N+1)$ until the condition (3.37) is satisfied, and then compute $S_n(k)$ ($k=0, 1, \dots, N+1$) by (1.2).

When there is known a bound $M(\mu)$ such that the inequality

$$(3.46) \quad n \geq M(\mu)$$

implies (3.44), the following method becomes possible.

Method 3. Generate $P_n(N)$, $P_n(N+1)$ and $U_n(N)$ until (3.45) and (3.46) are valid for n satisfying (3.22), and then compute $S_n(k)$ ($k=0, 1, \dots, N+1$) by (1.2) and (2.18).

Among the three methods, the last one seems to be the most efficient, and the methods 1 and 2 can be applied for determining the empirical bound $M(\mu)$ with $N=0$.

EXAMPLE 3. Let $I_n(x)$ ($n=0, 1, \dots$) be the modified Bessel functions of the first kind and put

$$(3.47) \quad y_n = e^{-x} I_n(x)$$

for a fixed value of x . Then they satisfy the recurrence formula [9]

$$(3.48) \quad y_{n-1} = \frac{2n}{x} y_n + y_{n+1}$$

with the normalizing condition

$$(3.49) \quad y_0 + 2 \sum_{j=1}^{\infty} y_j = 1.$$

Generating $S_n(0)$ and $S_n(1)$ for increasing n until they were in the state of numerical convergence [10] for $x=0.01, 0.05(0.05)1.0, 1.5(0.5)10, 15(5)100$, and $110(10)500$, we obtained the following empirical bounds for a digital computer with 39 bits mantissa:

$$(3.50) \quad M(x, 10^{-6}) = \begin{cases} x + 9 - \frac{62}{39x + 10} & (0 < x \leq 10) \\ 0.1x + 74 - \frac{6270}{x + 105} & (10 < x \leq 500), \end{cases}$$

$$(3.51) \quad M(x, 10^{-8}) = \begin{cases} x + 12 - \frac{83}{27x + 10} & (0 < x \leq 10) \\ 0.1x + 99 - \frac{10800}{x + 130} & (10 < x \leq 500), \end{cases}$$

$$(3.52) \quad M(x, 10^{-10}) = \begin{cases} x + 16 - \frac{44}{5x + 4} & (0 < x \leq 10) \\ 0.15x + 83 - \frac{4514}{x + 65} & (10 < x \leq 500). \end{cases}$$

These bounds mean that the inequalities

$$(3.53) \quad |s_{n,0}^*| \leq \frac{\mu}{2}, \quad |s_{n,1}^*| \leq \frac{\mu}{2}$$

are valid approximately provided $n \geq M(x, \mu)$, where $s_{n,0}^*$ and $s_{n,1}^*$ are the relative errors of $S_n(0)$ and $S_n(1)$ to the computed values of $e^{-x}I_0(x)$ and $e^{-x}I_1(x)$ respectively.

References

- [1] British Association for the Advancement of Science: *Bessel functions, Part II*, Mathematical Tables vol. X, Cambridge (1952).
- [2] Corbató, F. J. and J. L. Uretsky: *Generation of spherical Bessel functions in digital computers*, J. Assoc. Comput. Mach., **6** (1959), 366-375.
- [3] Gautschi, W.: *Recursive computation of the repeated integrals of the error function*, Math. Comput., **15** (1961), 227-232.
- [4] Goldstein, M. and R. M. Thaler: *Recurrence techniques for the calculation of Bessel functions*, Math. Tables Aids Comput., **13** (1959), 102-108.
- [5] Lance, G. N.: *Numerical methods for high speed computers*, London (1960).
- [6] Olver, F. W. J.: *Error analysis of Miller's recurrence algorithm*, Math. Comput., **18** (1964), 65-74.
- [7] Patry, J. and J. Keller: *Zur Berechnung des Fehlerintegrals*, Numer. Math., **6** (1964), 89-97.
- [8] Rotenberg, A.: *The calculation of toroidal harmonics*, Math. Comput., **14** (1960), 274-276.
- [9] Stegun, I. A. and M. Abramowitz: *Generation of Bessel functions on high speed computers*, Math. Tables Aids Comput., **11** (1957), 255-257.
- [10] Urabe, M.: *Convergence of numerical iteration in solution of equations*, J. Sci. Hiroshima Univ., Ser. A, **19** (1956), 479-489.

*Department of Mathematics
Faculty of Science
Hiroshima University*

