

## ON THE REDUCED LENGTH OF A POLYNOMIAL WITH REAL COEFFICIENTS

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To Professor Eduard Wirsing  
with best wishes  
for his 75th birthday

**Abstract:** The length  $L(P)$  of a polynomial  $P$  is the sum of the absolute values of the coefficients. For  $P \in \mathbb{R}[x]$  the properties of  $l(P)$  are studied, where  $l(P)$  is the infimum of  $L(PG)$  for  $G$  running through monic polynomials over  $\mathbb{R}$ .

**Keywords:** length of a polynomial, unit circle.

We shall consider only polynomials with real coefficients. For such a polynomial  $P = \sum_{i=0}^d a_i x^{d-i}$  the length  $L(P)$  is defined by the formula

$$L(P) = \sum_{i=0}^d |a_i|.$$

A. Dubickas [1] has introduced the reduced length by the formula

$$\widehat{l}(P) = \inf_{G \in \widehat{\Gamma}} L(PG),$$

where

$$\widehat{\Gamma} = \left\{ \sum_{i=0}^n b_i x^{n-i} \in \mathbb{R}[x], \text{ where } b_0 = 1 \text{ or } b_n = 1 \right\}.$$

It follows, see [1], p. 3, that

$$\widehat{l}(P) = \min \{l_0(P), l_0(P^*)\},$$

where

$$l_0(P) = \inf_{G \in \Gamma_0} L(PG), \quad \Gamma_0 = \left\{ \sum_{i=0}^n b_i x^{n-i} \in \mathbb{R}[x], b_n = 1 \right\}, \quad P^* = x^{\deg P} P(x^{-1}).$$

Since polynomials with the leading coefficient 1 have a name (monic) and polynomials with the constant term 1 have no name, I prefer to work with

$$l(P) = l_0(P^*) = \inf_{G \in \Gamma} L(PG), \quad \Gamma = \left\{ \sum_{i=0}^n b_i x^{n-i} \in \mathbb{R}[x], b_0 = 1 \right\}$$

Dubickas's results about  $l_0$  translated in the language of  $l$  give the following

**Proposition A.** (Dubickas [1]) *Suppose that  $\omega, \eta, \psi \in \mathbb{R}$ ,  $\nu \in \mathbb{C}$ ,  $\bar{\nu}$  is the complex conjugate to  $\nu$ ,  $|\omega| \geq 1$ ,  $|\eta| < 1$ ,  $|\nu| < 1$ , then for every  $Q \in \mathbb{R}[x]$*

- (i)  $l(\psi Q) = |\psi|l(Q)$ ,
- (ii)  $l(x + \omega) = 1 + |\omega|$ ,
- (iii) if  $T(x) = Q(x)(x - \eta)$ , then  $l(T) = l(Q)$ ,
- (iv) if  $T(x) = Q(x)(x - \nu)(x - \bar{\nu})$ , then  $l(T) = l(Q)$ .

We shall prove the following

**Proposition.** *For all monic polynomials  $P, Q$  in  $\mathbb{R}[x]$  and all positive integers  $k$*

- (i)  $\max\{l(P), l(Q)\} \leq l(PQ) \leq l(P)l(Q)$ ,
- (ii)  $M(P) \leq l(P)$ , where  $M$  is the Mahler measure,
- (iii)  $l(P(-x)) = l(P(x))$ ,
- (iv)  $l(P(x^k)) = l(P(x))$ .

**Theorem 1.** *If  $P \in \mathbb{R}[x]$  is monic of degree  $d$  with  $P(0) \neq 0$ , then  $l(P) = \inf_{Q \in S_d(P)} L(Q)$ , where  $S_d(P)$  is the set of all monic polynomials  $Q$  over  $\mathbb{R}$  divisible by  $P$  with  $Q(0) \neq 0$  and with at most  $d + 1$  non-zero coefficients, all belonging to the field  $K(P)$ , generated by the coefficients of  $P$ .*

**Theorem 2.** *If  $P \in \mathbb{R}[x]$  has all zeros outside the unit circle, then  $l(P)$  is attained and effectively computable, moreover  $l(P) \in K(P)$  ( $l(P)$  is attained means that  $l(P) = L(Q)$ , where  $Q/P \in \Gamma$ ).*

**Corollary 1.** *If  $P \in \mathbb{R}[x]$  has no zeros on the unit circle, then  $l(P)$  is effectively computable.*

**Theorem 3.** *Let  $P, Q \in \mathbb{R}[x]$ ,  $Q$  be monic and have all zeros on the unit circle. Then for all  $m \in \mathbb{N}$*

$$l(PQ^m) = l(PQ).$$

**Theorem 4.** *If  $P \in \mathbb{R}[x]$  is monic and has all zeros on the unit circle, then  $\widehat{l}(P) = l(P) = 2$ , with  $l(P)$  attained, if and only if all zeros are roots of unity and simple.*

**Theorem 5.** *Let  $P(x) = P_0(x)(x - \varepsilon)^e$ , where  $P_0 \in \mathbb{R}[x]$ ,  $\varepsilon = \pm 1$ ,  $e \in \mathbb{N}$  and all zeros of  $P_0$  are outside the unit circle. Assume that the set  $Z$  of zeros of  $P_0$  has a subset  $Z_0$ , possibly empty, such that its elements are real of the same sign and the elements of  $Z \setminus Z_0$  are algebraically independent over  $\mathbb{Q}(Z_0)$ . Then  $l(P)$  can be effectively computed. Moreover, if  $\deg P_0 = d_0$ , then*

$$l(P) \leq \inf_{Q \in S_{d_0}(P_0)} \left\{ L(Q) + |Q(\varepsilon)| \right\}.$$

For quadratic polynomials  $P$  Theorems 2, 4 and 5 together with Proposition A (iii) and (iv) exhaust all possibilities, so that  $l(P)$  can be effectively computed. A more precise information is given by the following

**Theorem 6.** *If  $P(x) = (x - \alpha)(x - \beta)$ , where  $|\alpha| \geq |\beta| \geq 1$ , then*

$$l(P) \geq 2|\alpha|$$

with the equality attained, if and only if  $|\beta| = 1$ .

**Corollary 2.** *If  $P \in \mathbb{R}[x]$  is of degree at most two with no zeros inside the unit circle, then*

$$l(P) \in K(P).$$

**Corollary 3.** *If  $P(x) = (x - \alpha)(x - \beta)$ , where  $|\alpha| \geq |\beta| \geq 0$ , then*

$$\widehat{l}(P) = \begin{cases} |\alpha\beta|, & \text{if } |\beta| > 1, \\ 2|\alpha|, & \text{if } |\beta| = 1, \\ |\alpha| + \min\{1, |\alpha\beta|\}, & \text{if } |\alpha| > 1 > |\beta|, \\ 2, & \text{if } |\alpha| = 1, \\ 1, & \text{if } |\alpha| < 1. \end{cases}$$

**Corollary 4.** *The function  $\widehat{l}$  is not submultiplicative.*

The last corollary is of interest, because of Proposition, part (i).

The problem of computing  $l(P)$  for cubic polynomials remains open already for  $P = 2x^3 + 3x^2 + 4$ . Another open question is whether  $l(P) \in K(P)$  for all  $P \in \mathbb{R}[x]$  with no zeros inside the unit circle.

We begin with

**Proof of Proposition.** We have by definition for all monic polynomials  $R, S$  in  $\mathbb{R}[x]$

$$l(P) \leq L(PQR), \quad l(PQ) \leq L(PQRS) \leq L(PR)L(QS)$$

hence

$$l(P) \leq \inf_{R \in \Gamma} L(PQR) = l(PQ),$$

$$l(PQ) \leq \inf_{R \in \Gamma} L(PR) \inf_{S \in \Gamma} L(QS) = l(P)l(Q).$$

This proves (i). As to (ii) we have for every  $R$  in  $\mathbb{R}[x]$

$$M(R) \leq L(R)$$

(see [4]), hence

$$M(P) \leq M(PQ) \leq L(PQ),$$

thus

$$M(P) \leq \inf_{Q \in \Gamma} L(PQ) = l(P)$$

and (ii) holds. The statement (iii) follows from

$$L(P(-x)) \leq L\left(P(-x)Q(-x)(-1)^{\deg Q}\right) = L(PQ),$$

whence

$$l(P(-x)) \leq \inf_{Q \in \Gamma} L(PQ) = l(P).$$

Similarly,

$$l(P(x^k)) \leq L\left(P(x^k)Q(x^k)\right) = L(PQ),$$

whence

$$l(P(x^k)) \leq \inf_{Q \in \Gamma} L(PQ) = l(P). \tag{1}$$

Finally, if

$$P(x^k)Q(x) = \sum_{i=0}^{k-1} x^i A_i(x^k), \text{ where } A_i \in \mathbb{R}[x], \tag{2}$$

let  $A_i = Q_i P + R_i$ , where  $Q_i, R_i \in \mathbb{R}[x]$  and  $\deg R_i < \deg P$ . It follows that

$$P(x^k) \mid \sum_{i=0}^{k-1} x^i R_i(x^k)$$

and since the degree of the sum is less than that of  $P(x^k)$ ,  $R_i = 0$  ( $0 \leq i < k$ ). Let  $i$  be chosen so that  $\deg x^i A_i(x^k)$  is the greatest. It follows from (2) that  $Q_i$  is monic. Hence, by (2)

$$L\left(P(x^k)Q(x)\right) \geq L(A_i) = L(PQ_i) \geq l(P),$$

thus  $l(P(x^k)) \geq l(P)$ , which together with (1) implies (iv).

**Remark.** The above proof of (iv), simpler than author’s original proof, has been kindly suggested by A. Dubickas.

For the proof of Theorem 2 we need two lemmas

**Lemma 1.** Let  $k \geq n$ ,  $\mathbf{x} = (x_1, \dots, x_n)$ ,  $L_i(\mathbf{x})$  for  $i \leq k$  be linear forms over  $\mathbb{R}$ ;  $L_1, \dots, L_n$  linearly independent,  $a_i \in \mathbb{R}$  ( $1 \leq i \leq k$ ). Then

$$S(\mathbf{x}) = \sum_{i=1}^k |L_i(\mathbf{x}) + a_i|$$

attains its infimum.

**Proof.** Let  $L_i(\mathbf{x}) = \sum_{j=1}^n a_{ij}x_j$  ( $1 \leq i \leq k$ ),  $A = \max_{i,j \leq n} |a_{ij}|$ ,

$$D = \left| \det (a_{ij})_{i,j \leq n} \right|, \quad s = \sum_{i=1}^k |a_i|.$$

Let  $s_0$  be the infimum of  $S(\mathbf{x})$  in the hypercube (degenerated if  $s = 0$ )

$$H : \max_{1 \leq i \leq n} |x_i| \leq \frac{2n^{\frac{n-1}{2}} s A^{n-1}}{D}.$$

Since  $H$  is compact, there exists  $\mathbf{x}_0 \in H$  such  $S(\mathbf{x}_0) = s_0$ . We shall show that  $s_0 = \inf_{\mathbf{x} \in \mathbb{R}^n} S(\mathbf{x})$ . Indeed, if for some  $\mathbf{x}_1 \in \mathbb{R}^n$

$$S(\mathbf{x}_1) < s_0, \tag{3}$$

then

$$\sum_{i=1}^n |L_i(\mathbf{x}_1)| < s_0 + s \leq 2s.$$

Solving the system  $L_i(\mathbf{x}) = L_i(\mathbf{x}_1)$  ( $1 \leq i \leq n$ ) by means of Cramer's formulae and using Hadamard's inequality to estimate the relevant determinants we obtain

$$\max_{1 \leq i \leq n} |x_{1i}| < \frac{2n^{\frac{n-1}{2}} s A^{n-1}}{D},$$

hence  $\mathbf{x}_1 \in H$ , a contradiction with (3) and the definition of  $s_0$ .

**Lemma 2.** Let  $k \geq n$ ,  $\mathbf{x} \in \mathbb{R}^n$ ,  $K$  be a subfield of  $\mathbb{R}$ ,  $L_1(\mathbf{x}), \dots, L_k(\mathbf{x})$  be linear forms over  $K$ ,  $n$  of them linearly independent,  $a_i \in K$ . There exists a point  $\mathbf{x}_0 \in K^n$  in which  $S(\mathbf{x}) = \sum_{i=1}^k |L_i(\mathbf{x}) + a_i|$  attains its infimum over  $\mathbb{R}^n$  and  $L_i(\mathbf{x}_0) + a_i = 0$ , for  $n$  indices  $i = i_1, i_2, \dots, i_n$  such that  $L_{i_1}, L_{i_2}, \dots, L_{i_n}$  are linearly independent.

**Proof by induction on  $k$ .** If  $k = 1$  we have  $n = 1$  and the assertion is trivial. Assume it is true for  $k - 1$  forms and consider the case of  $k$  forms,  $k \geq 2$ . If one of them, say  $L_k$  is identically 0, then among  $L_1, \dots, L_{k-1}$  there are  $n$  linearly independent, hence  $k - 1 \geq n$  and applying the inductive assumption to  $L_1, \dots, L_{k-1}$  we obtain the assertion. Therefore, we assume that all forms  $L_1, \dots, L_k$  are non-zero. Suppose that  $\inf S(\mathbf{x}) = S(\mathbf{x}_1)$  and  $L_i(\mathbf{x}_1) + a_i \neq 0$  for all  $i \leq k$ . Then there is an  $\varepsilon > 0$  such that  $|\mathbf{x} - \mathbf{x}_1| < \varepsilon$  implies  $\text{sgn}(L_i(\mathbf{x}) + a_i) =: \varepsilon_i$  for all  $i \leq k$ . We have

$$S(\mathbf{x}) = \sum_{i=1}^k \varepsilon_i (L_i(\mathbf{x}) + a_i) = M(\mathbf{x} - \mathbf{x}_1) + S(\mathbf{x}_1),$$

where

$$M(\mathbf{x}) = \sum_{i=1}^k \varepsilon_i L_i(\mathbf{x}).$$

If  $M \neq 0$ , then there exists a point  $\mathbf{x}_0$  with  $|\mathbf{x}_0| < \varepsilon$  and  $M(\mathbf{x}) < 0$ , hence taking  $\mathbf{x}_2 = \mathbf{x}_1 + \mathbf{x}_0$  we obtain  $S(\mathbf{x}_2) < S(\mathbf{x}_1)$ , a contradiction. Thus either  $L_{i_1}(\mathbf{x}_1) + a_{i_1} = 0$  for a certain  $i_1$ , or  $M = 0$ . In the latter case we take the point  $\mathbf{x}_2$  nearest to  $\mathbf{x}_1$  (or one of these) with  $L_{i_2}(\mathbf{x}_2) + a_{i_2} = 0$  for a certain  $i_2$ . Since the hyperplanes  $L_i(x) + a_i = 0$  either are disjoint with the ball  $|\mathbf{x} - \mathbf{x}_1| \leq |\mathbf{x}_2 - \mathbf{x}_1|$ , or are tangent to it, taking  $\langle \mathbf{x}_3, i_3 \rangle$  equal either to  $\langle \mathbf{x}_1, i_1 \rangle$  or to  $\langle \mathbf{x}_2, i_2 \rangle$  we obtain  $S(\mathbf{x}_3) = S(\mathbf{x}_1)$  and  $L_{i_3}(\mathbf{x}_3) + a_{i_3} = 0$ . Without loss of generality we may assume that  $i_3 = k$  and  $L_k$  is of positive degree in  $x_n$ . The equation  $L_k(\mathbf{x}) + a_k = 0$  is equivalent to  $x_n = C(x_1, \dots, x_{n-1}) + c$ , where  $C$  is a linear form over  $K$  and  $c \in K$ . We now apply the inductive assumption to the forms  $L'_i = L_i(x_1, \dots, x_{n-1}, C(x_1, \dots, x_{n-1}))$  and numbers  $a'_i = a_i + L_i(0, \dots, 0, c)$  ( $1 \leq i \leq k-1$ ). By virtue of the theorem about the rank of the product of matrices, the number of linearly independent among forms  $L'_i$  is  $n-1$ . By the inductive assumption there exists a point  $\mathbf{x}'_0 \in K^{n-1}$  such that  $\sum_{i=1}^{k-1} |L'_i(\mathbf{x}') + a'_i| = S'(\mathbf{x}')$  attains at  $\mathbf{x}'_0$  its infimum over  $\mathbb{R}^{n-1}$  and  $L'_i(\mathbf{x}'_0) + a'_i = 0$  for  $n-1$  indices  $i = i'_1, \dots, i'_{n-1}$  such that  $L'_{i'_1}, \dots, L'_{i'_{n-1}}$  are linearly independent. By the definition of  $L'_i$  and  $a'_i$  we have

$$S(\mathbf{x}_3) = S'(x_{3,1}, \dots, x_{3,n-1}) \geq \inf_{\mathbf{x}' \in \mathbb{R}^{n-1}} S'(\mathbf{x}') \geq \inf_{\mathbf{x} \in \mathbb{R}^n} S(\mathbf{x}) = S(\mathbf{x}_3),$$

hence

$$S'(\mathbf{x}'_0) = \inf_{\mathbf{x}' \in \mathbb{R}^{n-1}} S'(\mathbf{x}') = \inf_{\mathbf{x} \in \mathbb{R}^n} S(\mathbf{x}).$$

Moreover,  $L'_{i'_j}(\mathbf{x}'_0) + a'_{i'_j} = 0$  implies

$$L'_{i'_j}(x'_{0,1}, \dots, x'_{0,n-1}, C(\mathbf{x}'_0)) + a_{i_j} = 0$$

and the linear independence of  $L'_{i'_1}, \dots, L'_{i'_{n-1}}$  implies the linear independence of the forms  $L_{i'_1}, \dots, L_{i'_{n-1}}$ . The latter forms are also linearly independent with  $L_k$  since identity

$$L_1(x_1, \dots, x_n) = \sum_{j=1}^{n-1} c_j L_{i'_j}(x_1, \dots, x_n), \quad c_j \in \mathbb{R}$$

gives on substitution  $x_n = C(x_1, \dots, x_{n-1})$

$$0 = \sum_{j=1}^{n-1} c_j L'_{i'_j}(x_1, \dots, x_n), \quad \text{hence } c_j = 0 \quad (1 \leq j < n).$$

Taking  $\mathbf{x}_0 = (x'_{0,1}, \dots, x'_{0,n-1}, C(x'_0))$ ,  $i_j = i'_j$  ( $1 \leq j < n$ ),  $i_n = k$  we obtain the inductive assertion.

**Proof of Theorem 1.** We have by definition

$$l(P) = \inf L(PG),$$

where  $G$  runs through all monic polynomials. Let  $P = x^d + \sum_{i=1}^d a_i x^{d-i}$ ,  $G = x^n + \sum_{i=1}^n x_i x^{n-i}$ . We have

$$PG = x^{n+d} + \sum_{i=1}^{n+d} b_i x^{n+d-i},$$

where, with  $a_0 = 1$  for  $i \leq d$

$$b_i = a_i + \sum_{j=1}^{\min\{i,n\}} a_{i-j} x_j,$$

for  $i > d$

$$b_i = \sum_{j=i-d}^{\min\{i,n\}} a_{i-j} x_j.$$

Therefore,

$$l(P) = 1 + \inf_{n, \mathbf{x} \in \mathbb{R}^n} \left\{ \sum_{i=1}^d |L_i(\mathbf{x}) + a_i| + \sum_{i=d+1}^{d+n} |L_i(\mathbf{x})| \right\},$$

where

$$L_i(\mathbf{x}) = \sum_{j=\max\{1, i-d\}}^{\min\{i,n\}} a_{i-j} x_j.$$

The forms  $L_i$  satisfy the assumptions of Lemma 2. Indeed, the  $n$  forms  $L_{d+1}, \dots, L_{d+n}$  are linearly independent, since  $L_{d+1}(\mathbf{x}) = \dots = L_{d+n}(\mathbf{x}) = 0$  gives  $PG \equiv 0 \pmod{x^n}$ , hence  $G \equiv 0 \pmod{x^n}$ , i.e.  $x_1 = \dots = x_n = 0$ . Applying Lemma 2 and Proposition A (iii) with  $\eta = 0$  we obtain that for a given  $n$ ,  $PG$  with the minimal length occurs in  $S_d(P)$ .

For the proof of Theorem 2 we need

**Definition 1.** Let  $P = \prod_{s=1}^r (x - \alpha_s)^{m_s}$ , where  $\alpha_s$  are distinct and non-zero,  $m_s \in \mathbb{N}$  ( $1 \leq s \leq r$ ),  $m_1 + \dots + m_r = d$ ,  $n_0 > n_1 > \dots > n_{d-1} > n_d \geq 0$  be integers. If  $s \geq 1$ ,  $1 \leq i \leq d$ ,  $0 \leq j \leq d$ , then  $i$  can be written in the form  $i = m_1 + \dots + m_{s-1} + g$  for some  $1 \leq s \leq r$  and  $1 \leq g \leq m_s$ . We put

$$c_{ij} = \alpha_s^{n_j} \prod_{f=0}^{g-2} (n_j - n_f), \text{ where the empty product is } 1$$

and for  $\nu = 0, 1$

$$C(P; n_0, \dots, n_d) = (c_{ij})_{\substack{1 \leq i \leq d \\ 0 \leq j \leq d}}, \quad C_\nu(P; n_\nu, \dots, n_{d-1+\nu}) = (c_{ij})_{\substack{1 \leq i \leq d \\ \nu \leq j < d+\nu}}.$$

**Definition 2.**  $T_d(P) = \left\{ Q \in S_d(P) : Q = x^{n_0} + \sum_{i=1}^d b_i x^{n_i}, \text{ where } n_0 > n_1 > \dots > n_d = 0, |C_0(P; n_0, \dots, n_{d-1})| \neq 0 \neq |C_1(P; n_1, \dots, n_d)|, L(Q) \leq L(P) \right\}$ .

**Lemma 3.** We have for  $x_j \in \mathbb{C}$

$$\sum_{j=0}^d x_j x^{n_j} \equiv 0 \pmod{P} \tag{4_1}$$

if and only if

$$\sum_{j=0}^d c_{ij} x_j = 0 \quad (1 \leq i \leq d). \tag{4_2}$$

**Proof.** Clearly the condition (4<sub>1</sub>) is equivalent to

$$\sum_{j=0}^d x_j \binom{n_j}{g-1} \alpha_s^{n_j} = 0 \quad (1 \leq g \leq m_s, 1 \leq s \leq r),$$

that is to the vector equation

$$M \mathbf{x} = 0, \tag{5}$$

where  $\mathbf{x} = (x_0, x_1, \dots, x_d)^t$ ,  $M = (m_{ij})_{\substack{1 \leq i \leq d \\ 0 \leq j \leq d}}$  and if  $i = m_1 + \dots + m_{s-1} + g$ ,  $1 \leq g \leq m_s$ , then

$$m_{ij} = \binom{n_j}{g-1} \alpha_s^{n_j}. \tag{6}$$

Now define the numbers  $b_{gh}$  by the equation

$$\prod_{f=0}^{g-2} (x - n_f) = \sum_{h=1}^g b_{gh} \binom{x}{h-1} \tag{7}$$

and put for  $i = m_1 + \dots + m_{s-1} + g$ ,  $1 \leq g \leq m_s$ ,  $1 \leq j \leq d$

$$a_{ij} = \begin{cases} b_{gh} & \text{if } j = m_1 + \dots + m_{s-1} + h, 1 \leq h \leq g, \\ 0 & \text{otherwise,} \end{cases} \tag{8}$$

$$A = (a_{ij})_{1 \leq i, j \leq d}. \tag{9}$$

The matrix  $A$  is lower triangular and non-singular, since  $b_{gg} = (g-1)!$ . Hence the equation (5) is equivalent to

$$AM \mathbf{x} = 0. \tag{10}$$

However, by (6)–(9) the element in  $i$ -th row ( $1 \leq i \leq d$ ) and  $j$ -th column ( $0 \leq j \leq d$ ) of  $AM$  for  $i = m_1 + \dots + m_{s-1} + g$ ,  $1 \leq g \leq m_s$  is

$$\sum_{t=1}^d a_{it} m_{tj} = \sum_{h=1}^g b_{gh} \binom{n_j}{h-1} \alpha_s^{n_j} = \alpha_s^{n_j} \prod_{f=0}^{g-2} (n_j - n_f) = c_{ij},$$

hence (4<sub>1</sub>) is equivalent to (4<sub>2</sub>).



**Lemma 4.** We have  $\inf_{Q \in S_d(P)} L(Q) = \inf_{Q \in T_d(P)} L(Q)$ .

**Proof.** Let  $P$  be as in Definition 1. We shall prove by induction with respect to  $n$  that

$$\inf_{\substack{Q \in S_d(P) \\ \deg Q \leq n+d}} L(Q) = \inf_{\substack{Q \in T_d(P) \\ \deg Q \leq n+d}} L(Q). \tag{11}$$

If  $n = 0$  then  $Q \in S_d(P)$ ,  $\deg Q \leq n + d$  implies  $Q = P$ . We shall show that  $P \in T_d(P)$ . Otherwise  $P = x^{n_0} + \sum_{i=1}^d a_i x^{n_i}$  ( $n_i = d - i$ ) and either  $|C_0(P; n_0, \dots, n_{d-1})| = 0$ , or  $|C_1(P; n_1, \dots, n_d)| = 0$ . In the former case there exists  $[d_0, \dots, d_{d-1}] \in \mathbb{C}^d \setminus \{\mathbf{0}\}$  such that

$$\sum_{j=0}^{d-1} c_{ij} d_j = 0 \quad (1 \leq i \leq d). \tag{12}$$

By Lemma 3

$$\sum_{j=0}^{d-1} d_j x^{n_j} \equiv 0 \pmod{P} \tag{13}$$

and, since  $n_0 = d$ ,  $\sum_{j=0}^{d-1} d_j x^{n_j} = d_0 P$ ;  $d_0 \neq 0$ ,  $P(0) = 0$ , a contradiction. In the latter case, similarly, there exists  $[e_1, \dots, e_d] \in \mathbb{C}^d \setminus \{\mathbf{0}\}$  such that

$$\sum_{j=1}^d c_{ij} e_j = 0 \quad (1 \leq i \leq d). \tag{14}$$

By Lemma 3

$$\sum_{j=1}^d e_j x^{n_j} \equiv 0 \pmod{P}, \tag{15}$$

which is impossible since  $n_1 < n_0 = d = \deg P$ .

Assume now that the equality (11) holds with  $n$  replaced by  $n - 1$  and suppose that

$$\inf_{\substack{Q \in S_d(P) \\ \deg Q \leq n+d}} L(Q) = L(Q_0), \text{ where } Q_0 = x^{n_0} + \sum_{j=1}^d b_j x^{n_j} \in K(P)[x], \tag{16}$$

$$n_0 > n_1 > \dots > n_d \geq 0.$$

Clearly  $L(Q_0) \leq L(P)$ .

If  $n_0 < n + d$  the inductive assertion follows immediately from the inductive assumption. If  $n_0 = n + d$ , let  $L_i(x)$  be the linear forms defined in the proof of

Theorem 1 and  $a_i$  have the meaning of that proof, if  $i \leq d$ ,  $a_i = 0$  otherwise. We have

$$L(Q_0) = \inf_{\mathbf{x} \in \mathbb{R}^n} \sum_{i=1}^{n+d} |L_i(\mathbf{x}) + a_i|,$$

hence, by Lemma 2, the above infimum is attained in a point  $\mathbf{x}_0$ , such that for  $n$  indices  $i_1, \dots, i_n$  simultaneously  $L_{i_j}(\mathbf{x}_0) + a_{i_j} = 0$  and  $L_{i_1}, \dots, L_{i_n}$  are linearly independent. Since the system of equations  $L_{i_j}(\mathbf{x}_0) + a_{i_j} = 0$  ( $1 \leq j \leq n$ ) determines  $\mathbf{x}_0$  uniquely the coefficients of  $x^{n+d-i}$  in  $Q$ , where  $i \neq i_1, \dots, i_n$  (hence  $n+d-i = n_1, n_2, \dots, n_d$ ) are uniquely determined by the condition  $Q \equiv 0 \pmod{P}$ ,  $Q$  monic in  $\mathbb{C}[x]$ . On the other hand, if  $|C_0(P; n_0, \dots, n_{d-1})| = 0$ , then there exists  $[d_0, \dots, d_{d-1}] \in \mathbb{C}^d \setminus \{\mathbf{0}\}$  such that (12) and (13) hold again. If  $d_0 = 0$ , then

$$Q_1 := Q_0 + \sum_{i=1}^{d-1} d_j x^{n_j} \equiv 0 \pmod{P},$$

where the  $Q_1$  is again monic, contrary to the uniqueness property. If  $d_0 \neq 0$ , then by the uniqueness property

$$Q_0 = d_0^{-1} \sum_{j=1}^{d-1} d_j x^{n_j} = x^{n_{d-1}} Q_2, \quad Q_2 \in K(P)[x],$$

hence

$$L(Q_0) = L(Q_2), \quad \deg Q_2 < n + d$$

and, by the inductive assumption

$$L(Q_0) = \inf_{\substack{Q \in S_d(P) \\ \deg Q < n+d}} L(Q) = \inf_{\substack{Q \in T_d(P) \\ \deg Q < n+d}} L(Q) \geq \inf_{\substack{Q \in T_d(P) \\ \deg Q \leq n+d}} L(Q).$$

By (16) this gives (11).

If  $|C_1(P; n_1, \dots, n_d)| = 0$ , then there exists  $[e_1, \dots, e_d] \in \mathbb{C}^d \setminus \{\mathbf{0}\}$  such that (14) and (15) hold again. We have

$$Q_3 := Q_0 + \sum_{j=1}^d e_j x^{n_j} \equiv 0 \pmod{P}$$

and  $Q_3$  is again monic, contrary to the uniqueness property.

In the remaining case

$$|C_0(P; n_0, \dots, n_{d-1})| \neq 0 \neq |C_1(P; n_1, \dots, n_d)|$$

we have  $Q_0 \in T_d(P)$ , hence (11) holds again.

**Lemma 5.** *Let in the notation of Definition 1,  $i = m_1 + \dots + m_{s(i)-1} + g(i)$ ,  $1 \leq g(i) \leq m_{s(i)}$ . Then for every  $j \geq h \geq g(i) - 1$*

$$|c_{ij}| \leq |c_{ih}| \max \left\{ 1, \frac{g(i) - 1}{\log |\alpha_{s(i)}|} \right\}^{g(i)-1}. \tag{17}$$

**Proof.** For the sake of brevity, put  $s(i) = s$ ,  $g(i) = g$ . For  $g = 1$  we have  $|c_{ij}| = |\alpha_s^{n_j}| \leq |\alpha_s^{n_h}| = |c_{ih}|$ . Assume  $g > 1$ . For every  $f \leq g - 2$  the function

$$\varphi(x) = \max \left\{ 1, \frac{g - 1}{\log |\alpha_s|} \right\} |\alpha_s|^{\frac{n_h - x}{g-1}} - \frac{n_f - x}{n_f - n_h}$$

satisfies  $\varphi(n_h) \geq 0$ ,  $\varphi'(x) \leq 0$  for  $x \leq n_h$ . Hence  $\varphi(n_j) \geq 0$ ,

$$\max \left\{ 1, \frac{g - 1}{\log |\alpha_s|} \right\} |\alpha_s|^{\frac{n_h}{g-1}} (n_f - n_h) \geq |\alpha_s|^{\frac{n_j}{g-1}} (n_f - n_j)$$

and (17) follows on taking products over  $f$  from 0 to  $g - 2$ .

**Lemma 6.** *Let  $a, b, c, x \in \mathbb{R}$ ,  $a > 1$ ,  $b \geq 0$ ,  $c > 0$ ,  $x > 0$ . If*

$$a^x / x^b \leq c, \tag{18}$$

then

$$x \leq \left( \frac{2b}{e \log a} + \sqrt{\frac{b^2}{e^2 (\log a)^2} + \frac{\log c}{\log a}} \right)^2 =: \psi(a, b, c). \tag{19}$$

The function  $\psi$  is decreasing in  $a$ , increasing in  $b$  and  $c$ .

**Proof.** Put  $x = y^2$ ,  $y > 0$ . It follows from (18) that

$$y^2 \log a - 2b \log y \leq \log c$$

and, since  $\log y \leq y/e$

$$y^2 \log a - \frac{2b}{e} y \leq \log c.$$

Solving this inequality for  $y$  and squaring we obtain (19).

**Lemma 7.** *For every subset  $I$  of  $\{1, \dots, d\}$  of cardinality  $h$  we have*

$$\left| \det (c_{ij})_{\substack{i \in I \\ 1 \leq j \leq h}} \right| \leq h^{\frac{h}{2}} \prod_{i \in I} |\alpha_{s(i)}|^{n_{\max\{1, g(i)-1\}}} \prod_{f=0}^{g(i)-2} (n_f - n_h) \tag{20}$$

and

$$\left| \det (c_{ij})_{\substack{i \in I \\ 0 \leq j < h}} \right| \leq h^{\frac{h}{2}} \prod_{i \in I} |\alpha_{s(i)}|^{n_{g(i)-1}} \prod_{f=0}^{g(i)-2} (n_f - n_{h-1}). \tag{21}$$

**Proof.** For all  $i \in I$  and  $j \leq g(i) - 2$  we have  $c_{ij} = 0$ , while for  $j > g(i) - 2$

$$|c_{ij}| = |\alpha_{s(i)}|^{n_j} \prod_{f=0}^{g(i)-2} (n_f - n_h) \leq \begin{cases} |\alpha_{s(i)}|^{n_{\max\{1, g(i)-1\}}} \prod_{f=0}^{g(i)-2} (n_f - n_h) & \text{if } 1 \leq j \leq h, \\ |\alpha_{s(i)}|^{n_{g(i)-1}} \prod_{f=0}^{g(i)-2} (n_f - n_{h-1}) & \text{if } 0 \leq j < h, \end{cases}$$

hence (20) and (21) follow by Hadamard's inequality. Note that if  $g(i) > h + 1$  or  $g(i) > h$  for  $i \in I$ , then both sides of (20) or (21), respectively, are zero.

**Definition 3.** In the notation of Definition 1 and of Lemma 5 put for a positive integer  $h < d$ , positive integers  $e_1, \dots, e_h$  and a subset  $J$  of  $\{1, \dots, d\}$  of cardinality  $h + 1$  such that  $\max_{i \in J} g(i) \leq h + 1$

$$D(J; e_1, \dots, e_h) = \left| \det \left( \alpha_{s(i)}^{\sum_{\mu=j+1}^h e_\mu} \prod_{f=0}^{g(i)-2} \sum_{\nu=f+1}^j e_\nu \right) \right|_{\substack{i \in J \\ 0 \leq j \leq h}} \\ \times \prod_{i \in J} |\alpha_{s(i)}|^{-\sum_{\mu=\max\{2, g(i)\}}^h e_\mu} \prod_{i \in J} \prod_{f=0}^{g(i)-2} \left( \sum_{\nu=f+1}^h e_\nu \right)^{-1}.$$

**Definition 4.**  $D(e_1, \dots, e_h) = \max D(J; e_1, \dots, e_h)$ , where the maximum is taken over all subsets of  $\{1, \dots, d\}$  of cardinality  $h + 1$  such that  $\max_{i \in J} g(i) \leq h + 1$ .

**Remark.** The definition is meaningful, since always there exists a subset  $J$  of  $\{1, \dots, d\}$  with the required property. If for all  $i \leq d$  we have  $g(i) \leq h + 1$  this is clear and if for some  $i_0 : g(i_0) > h + 1$  we take

$$J = \{i : m_1 + \dots + m_{s(i_0)-1} < i \leq m_1 + \dots + m_{s(i_0)-1} + h + 1\}.$$

**Proof of Theorem 2.** Using the notation of Definition 1 we define the sequence  $d_1, \dots, d_d$  inductively as follows.

$$d_1 = \frac{\log(L(P) - 1)}{\log |\alpha_1|} \tag{22}$$

and, if  $d_1, \dots, d_h$  ( $d > h \geq 1$ ) are already defined, put

$$D_{h+1} = (h + 1)^{-1} h^{\frac{h}{2}} \min \{D(e_1, \dots, e_h) : 1 \leq e_i \leq d_i, D(e_1, \dots, e_h) > 0\} \tag{23}$$

(the minimum over an empty set being  $\infty$ ),  $m = \max_{1 \leq s \leq r} m_s$ ,

$$d_{h+1} = \begin{cases} \max \left\{ d_1 + \dots + d_h, \right. \\ \left. \psi \left( |\alpha_r|, m - 1, \left( \max \left\{ 2, \frac{2(m-1)}{\log |\alpha_r|} \right\} \right)^{m-1} D_{h+1}^{-1} (L(P) - 1) \right) \right\} \\ \qquad \qquad \qquad \text{if } D_{h+1} \neq \infty, \\ 0 \text{ otherwise.} \end{cases} \quad (24)$$

We shall show that if  $Q \in T_d(P)$ ,  $Q = x^{n_0} + \sum_{j=1}^d b_j x^{n_j}$ , then

$$n_{j-1} - n_j \leq d_j \quad (1 \leq j \leq d). \quad (25)$$

We proceed by induction on  $j$ . Since  $Q \in T_d(P)$  the equation

$$\alpha_1^{n_0} + \sum_{j=1}^d b_j \alpha_1^{n_j} = 0$$

implies

$$|\alpha_1|^{n_0} \leq |\alpha_1|^{n_1} \sum_{j=1}^d |b_j| \leq |\alpha_1|^{n_1} (L(Q) - 1) \leq |\alpha_1|^{n_1} (L(P) - 1),$$

which, in view of (22) gives (25) for  $j = 1$ . Assume now that (25) holds for all  $j \leq h$  ( $h < d$ ) and consider the matrix  $(c_{ij})_{\substack{1 \leq i \leq d \\ 0 \leq j \leq h}}$  defined in Definition 1. Since  $Q \in T_d(P)$  we have

$$\text{rank } (c_{ij})_{\substack{1 \leq i \leq d \\ 0 \leq j \leq h}} = h + 1,$$

hence also

$$\text{rank } \left( c_{ij} \alpha_{s(i)}^{-n_h} \right)_{\substack{1 \leq i \leq d \\ 0 \leq j \leq h}} = h + 1.$$

Therefore, there exists a subset  $J$  of  $\{1, \dots, d\}$  of cardinality  $h + 1$  such that

$$\Delta(J) = \det \left( c_{ij} \alpha_{s(i)}^{-n_h} \right)_{\substack{i \in J \\ 0 \leq j \leq h}} \neq 0. \quad (26)$$

For every subset  $J$  with the above property consider

$$M(J) = \max_{i \in J} \left| \left( c_{i0} + \sum_{j=1}^h c_{ij} b_j \right) \alpha_{s(i)}^{-n_h} \right|.$$

Solving the system of equations

$$\left( c_{i0}x_0 + \sum_{j=1}^h c_{ij}x_j \right) \alpha_{s(i)}^{-n_h} = \left( c_{i0} + \sum_{j=1}^h c_{ij}b_j \right) \alpha_{s(i)}^{-n_h} \quad (i \in J)$$

by means of Cramer's formulae and developing the numerator according to the first column we obtain

$$1 \leq \frac{(h+1)M(J) \max \left| \det \left( c_{ij} \alpha_{s(i)}^{-n_h} \right)_{\substack{i \in I \\ 1 \leq j \leq h}} \right|}{|\Delta(J)|},$$

where the maximum is taken over all subsets  $I$  of  $J$  of cardinality  $h$ . Now, by Lemma 7, since  $|\alpha_{s(i)}| \geq 1$

$$\max \left| \det \left( c_{ij} \alpha_{s(i)}^{-n_h} \right)_{\substack{i \in I \\ 1 \leq j \leq h}} \right| \geq h^{\frac{-h}{2}} \prod_{i \in J} |\alpha_{s(i)}|^{n_{\max\{1, g(i)-1\}}} \prod_{f=0}^{g(i)-2} (n_f - n_h).$$

This gives, by Definitions 3 and 4, for every  $J$  satisfying (26)

$$M(J) \geq (h+1)^{-1} h^{-\frac{h}{2}} D(J; n_0 - n_1, \dots, n_{h-1} - n_h) > 0$$

and, since such  $J$  exist

$$\max^* M(J) \geq (h+1)^{-1} h^{-\frac{h}{2}} D(J; n_0 - n_1, \dots, n_{h-1} - n_h) < \infty,$$

where  $\max^*$  is taken over all subsets  $J$  of  $\{1, \dots, d\}$  such that  $\text{card } J = h+1$  and  $\max_{i \in J} g(i) \leq h+1$ .

By the inductive assumption and (23)

$$\max^* M(J) \geq D_{h+1} > 0,$$

thus there exists a set  $J_0 \subset \{1, \dots, d\}$  such that

$$\text{card } J_0 = h+1, \quad \max_{i \in J_0} g(i) \leq h+1 \quad \text{and}$$

$$M(J_0) \geq D_{h+1}. \tag{27}$$

On the other hand, by Lemma 3

$$c_{i0} + \sum_{j=1}^d c_{ij}b_j = 0 \quad (i \in J_0),$$

hence

$$\left| \left( c_{i_0} + \sum_{j=1}^h c_{ij} b_j \right) \alpha_{s(i)}^{-n_h} \right| \cdot |\alpha_{s(i)}|^{n_h} = \left| \sum_{j=h+1}^d c_{ij} b_j \right|. \tag{28}$$

By (27) for a certain  $i_0 \in J_0$  the left-hand side is at least  $D_{h+1} |\alpha_{s(i)}^{n_h}|$ . As to the right-hand side, replacing in Lemma 5  $h$  by  $h + 1$ , we obtain

$$\begin{aligned} \left| \sum_{j=h+1}^d c_{i_0j} b_j \right| &\leq |c_{i_0,h+1}| \left( \max \left\{ 1, \frac{g(i_0) - 1}{\log |\alpha_{s(i_0)}|} \right\} \right)^{g(i_0)-1} \sum_{j=h+1}^d |b_j| \\ &\leq |c_{i_0,h+1}| \left( \max \left\{ 1, \frac{m_{s(i_0)} - 1}{\log |\alpha_{s(i_0)}|} \right\} \right)^{m_{s(i_0)}-1} (L(P) - 1). \end{aligned} \tag{29}$$

If  $n_h - n_{h+1} \leq n_0 - n_h$ , we obtain  $n_h - n_{h+1} \leq d_1 + \dots + d_h \leq d_{h+1}$ , hence the inductive assertion holds. If  $n_h - n_{h+1} > n_0 - n_h$ , then

$$\begin{aligned} |c_{i_0,h+1}| &= |\alpha_{s(i_0)}|^{n_{h+1}} \prod_{f=0}^{g(i_0)-2} (n_f - n_{h+1}) \\ &\leq |\alpha_{s(i_0)}|^{n_{h+1}} \left( 2(n_h - n_{h+1}) \right)^{g(i_0)-1} \\ &\leq |\alpha_{s(i_0)}|^{n_{h+1}} 2^{m_{s(i_0)}-1} (n_h - n_{h+1})^{m_{s(i_0)}-1}. \end{aligned} \tag{30}$$

Combining this inequality with (28) and (29) we obtain

$$\frac{D_{h+1} |\alpha_{s(i_0)}|^{n_h - n_{h+1}}}{(n_h - n_{h+1})^{m_{s(i_0)}-1}} \leq \left( \max \left\{ 2, \frac{2(m_{s(i_0)} - 1)}{\log |\alpha_{s(i_0)}|} \right\} \right)^{m_{s(i_0)}-1} (L(P) - 1), \tag{31}$$

hence, by Lemma 6,

$$\begin{aligned} &n_h - n_{h+1} \\ &\leq \max_{1 \leq s \leq r} \psi \left( |\alpha_s|, m_s - 1, \left( \max \left\{ 2, \frac{2(m_s - 1)}{\log |\alpha_s|} \right\} \right)^{m_s-1} D_{h+1}^{-1} (L(P) - 1) \right) \\ &\leq \psi \left( |\alpha_r|, m - 1, \left( \max \left\{ 2, \frac{2(m - 1)}{\log |\alpha_r|} \right\} \right)^{m-1} D_{h+1}^{-1} (L(P) - 1) \right) \leq d_{h+1}. \end{aligned}$$

The inductive assertion being proved, it follows that

$$n_0 - n_d \leq \sum_{h=1}^d d_h.$$

However,  $n_d = 0$ , hence

$$l(P) = \inf_{Q \in U_d(P)} L(Q),$$

where  $U_d(P) = \left\{ Q \in T_d(P) : \deg Q \leq \sum_{h=1}^d d_h \right\}$ .

The set  $U_d(P)$  is finite and effectively computable, since for  $Q = x^{n_0} + \sum_{j=1}^d b_j x^{n_j} \in U_d(P)$  there are only finitely many choices for  $\langle n_0, \dots, n_d \rangle$  and for each choice the coefficients  $b_j$  are determined uniquely and are effectively computable. Moreover,  $Q \in T_d(P)$  implies  $Q \in K(P)[x]$ , hence  $L(Q) \in K(P)$ . The theorem follows.

**Proof of Corollary 1.** If  $P(x) = a_0 \prod_{i=1}^c (x - \alpha_i) \prod_{i=c+1}^d (x - \alpha_i)$ , where  $|\alpha_i| > 1$  for  $i \leq c$ ,  $|\alpha_i| < 1$  for  $i > c$ , then by Proposition A

$$l(P) = |a_0| l \left( \prod_{i=1}^c (x - \alpha_i) \right)$$

and the right hand side is effectively computable by Theorem 2.

For the proof of Theorem 3 we need two lemmas.

**Lemma 8.** If  $P_n \in \mathbb{R}[x]$ ,  $p_n, q_n \in \mathbb{N} \cup \{0\}$  ( $n = 0, 1, \dots$ ) and

$$\liminf_{n \rightarrow \infty} L(P_n(x) - P_0(x^{p_n})x^{q_n}) = 0, \tag{32}$$

then

$$\liminf_{n \rightarrow \infty} l(P_n) \leq l(P_0). \tag{33}$$

**Proof.** By definition of  $l(P_0)$  for every  $n$  there exists  $G_n$  monic such that

$$L(P_0 G_n) \leq l(P_0) + \frac{1}{n}.$$

By (32) there exists  $k_n \in \mathbb{N}$  such that  $k_n > n$  and

$$L(P_{k_n}(x) - P_0(x^{p_{k_n}})x^{q_{k_n}}) \leq \frac{1}{nL(G_n)}.$$

Hence

$$\begin{aligned} L(P_{k_n}(x)G_n(x^{p_{k_n}})) &\leq L(P_0(x^{p_{k_n}})x^{q_{k_n}}G_n(x^{p_{k_n}})) \\ &+ L((P_{k_n}(x) - P_0(x^{p_{k_n}})x^{q_{k_n}})G_n(x^{p_{k_n}})) \leq L(P_0G_n) \\ &+ L(P_{k_n}(x) - P_0(x^{p_{k_n}})x^{q_{k_n}})L(G_n) \leq l(P_0) + \frac{2}{n}, \end{aligned}$$

thus

$$l(P_{k_n}) \leq l(P_0) + \frac{2}{n}.$$

This implies (33).



**Remark.** The equality  $\lim_{n \rightarrow \infty} L(P_n - P_0) = 0$  does not imply  $\liminf_{n \rightarrow \infty} l(P_n) = l(P_0)$ , as is shown by the example  $P_n = x - \frac{n-1}{n}$ ,  $P_0 = x - 1$ , see Proposition A, (ii) and (iii).

**Lemma 9.** *Let  $Q$  be a monic polynomial, irreducible over  $\mathbb{R}$  of degree  $d \leq 2$  with the zeros on the unit circle. There exists a sequence of monic polynomials  $R_n$  such that*

$$Q^2 \mid R_n \tag{34}$$

and

$$\lim_{n \rightarrow \infty} L(R_n - x^{dn}Q) = 0. \tag{35}$$

**Proof.** It suffices to take

$$R_n = x^{x+1} - \varepsilon \left( 1 + \frac{1}{n} \right) x^n + \frac{\varepsilon^{n+1}}{n}, \quad \text{if } d = 1, Q = x - \varepsilon$$

and

$$R_n = \left( x^{n+1} - \zeta \left( 1 + \frac{1}{n} \right) x^n + \frac{\zeta^{n+1}}{n} \right) \left( x^{n+1} - \bar{\zeta} \left( 1 + \frac{1}{n} \right) x^n + \frac{\bar{\zeta}^{n+1}}{n} \right),$$

if  $d = 2$ ,  $Q = (x - \zeta)(x - \bar{\zeta})$ .

Indeed, we have for every  $\varepsilon$

$$(x - \varepsilon)^2 \mid x^{n+1} - \varepsilon \left( 1 + \frac{1}{n} \right) x^n + \frac{\varepsilon^{n+1}}{n},$$

which implies (34) and

$$L(R_n - x^{dn}Q) \leq \begin{cases} 2/n & \text{if } d = 1 \\ (8n + 4)/n^2 & \text{if } d = 2, \end{cases}$$

which implies (35).

**Proof of Theorem 3.** We proceed by induction with respect to the number  $N$  of irreducible factors of  $Q^{m-1}$  counted with multiplicities. If  $N = 1$ , then  $m = 2$ ,  $Q$  is irreducible and by Lemma 9 we have

$$PQ^2 \mid PR_n \tag{36}$$

and

$$\lim_{n \rightarrow \infty} L(PR_n - x^{dn}PQ) = 0. \tag{37}$$

By (37) and Lemma 8 we have

$$\liminf_{n \rightarrow \infty} l(PR_n) \leq l(PQ).$$

However, by Proposition (i) and (36)

$$l(PQ) \leq l(PQ^2) \leq l(PR_n),$$

hence

$$l(PQ) \leq l(PQ^2) \leq l(PQ),$$

which gives the theorem for  $N + 1$ .

Assume now that the number of irreducible factors of  $Q^{m-1}$  is  $N > 1$  and the theorem is true for the number of irreducible factors less than  $N$ . If  $m > 2$  then the number of irreducible factors of  $Q^{m-2}$  and of  $Q$  is less than  $N$ , hence applying the inductive assumption with  $P$  replaced first by  $PQ$  we obtain

$$l(PQ^m) = l(PQ^2) = l(PQ).$$

If  $m = 2$  and the number of irreducible factors of  $Q^{m-1}$  is  $N > 1$ , then  $Q$  is reducible,  $Q = Q_1Q_2$ , where  $\deg Q_i > 1$  ( $i = 1, 2$ ). The number of irreducible factors of  $Q_i$  is less than  $N$ , hence applying the inductive assumption with  $P$  replaced first by  $PQ_1^2$  and then by  $PQ_2$  we obtain

$$l(PQ^2) = l(PQ_1^2Q_2^2) = l(PQ_1^2Q_2) = l(PQ_1Q_2) = l(PQ).$$

The inductive proof is complete.

**Proof of Theorem 4.** By Theorem 3 and Proposition A (ii) we have for  $d \in \mathbb{N}$

$$l((x-1)^d) = l(x-1) = 2. \quad (38)$$

Now, let

$$P(x) = \prod_{j=1}^d (x - \exp 2\pi i r_j), \quad \text{where } r_j \in \mathbb{R}.$$

By Dirichlet's approximation theorem for every positive integer  $n$  there exists a positive integer  $p_n$  such that

$$\|p_n r_j\| \leq \frac{1}{2\pi n} \quad (1 \leq j \leq d),$$

hence

$$|\exp 2\pi i p_n r_j - 1| < \frac{1}{n}.$$

It follows that the polynomial

$$Q_n(x) = \prod_{j=1}^d (x^{p_n} - \exp 2\pi i p_n r_j)$$

satisfies

$$P \mid Q_n \quad (39)$$

and

$$L(Q_n - (x^{p^n} - 1)^d) \leq \left(2 + \frac{1}{n}\right)^d - 2^d. \tag{40}$$

Now, (39) implies by Proposition (i)

$$l(P) \leq \liminf_{n \rightarrow \infty} l(Q_n),$$

while (40), Lemma 8 and (38) imply

$$\liminf_{n \rightarrow \infty} l(Q_n) \leq l((x - 1)^d) = 2.$$

Hence  $l(P) \leq 2$ . On the other hand, if  $P \mid Q$ ,  $Q = x^n + \sum_{j=1}^n b_j x^{n-j}$ , then for a zero  $\alpha$  of  $P$  we have

$$1 = |\alpha|^n = \left| \sum_{j=1}^n b_j x^{n-j} \right| \leq \sum_{j=1}^n |b_j| = L(Q) - 1,$$

hence  $L(Q) \geq 2$ ; so

$$l(P) \geq 2,$$

which gives the first part of the theorem. In order to prove the second part assume that  $P \mid Q$ ,  $Q$  monic and  $L(Q) = 2$ . Let

$$Q = x^n + \sum_{j=1}^n b_j x^{n-j}, \quad b_n \neq 0.$$

For every zero  $\alpha$  of  $P$  we have

$$\alpha^n + \sum_{j=1}^n b_j \alpha^{n-j} = 0, \tag{41}$$

hence

$$\left| \sum_{i=1}^n b_i \alpha^{n-i} \right| = |\alpha^n| = 1 = \sum_{i=1}^n |b_i|.$$

It follows that for every  $j$  with  $b_j \neq 0$

$$\arg b_j \alpha^{n-j} = \arg b_n.$$

Since  $\arg b_i = 0$  or  $\pi$ , either  $\alpha$  is a root of unity, or  $b_j = 0$  for all  $j < n$ . However the latter case, by virtue of (41) leads to the former. Suppose now that  $\alpha$  is a multiple zero of  $P$ , hence also of  $Q$ . Then

$$n\alpha^{n-1} + \sum_{j=1}^{n-1} b_j(n-j)\alpha^{n-j-1} = 0,$$

hence

$$\left| \sum_{j=1}^{n-1} b_j(n-j)\alpha^{n-j-1} \right| = |n\alpha^n| = n > \sum_{j=1}^{n-1} |b_j|n,$$

which is impossible, since for each  $j < n$ ,  $|b_j(n-j)\alpha^{n-j-1}| \leq |b_j|n$ . Thus all zeros of  $P$  are roots of unity and simple. If this condition is satisfied, then  $P \mid x^m - 1$ , where  $m$  is the least common multiple of orders of the roots of unity in question and

$$L(x^m - 1) = 2.$$

For the proof of Theorem 5 we need seven lemmas.

**Lemma 10.** *Let  $d > 2$ ,  $I$  be a subset of  $\{1, 2, \dots, d - 1\}$  and  $J$  a subset of  $\{0, \dots, d - 2\}$  both of cardinality  $d - 2$ . Then*

$$\left| \det (c_{ij})_{\substack{i \in I \\ j \in J}} \right| \leq (d - 2)^{\frac{d-2}{2}} \prod_{i \in I} |\alpha_{s(i)}|^{n_{g(i)-1}} \prod_{f=0}^{g(i)-2} (n_f - n_{d-2}).$$

**Proof.** is similar to the proof of Lemma 7.

**Lemma 11.** *Under the assumptions of Theorem 5 we have, in the notation of Definition 1, for every  $h \leq d - e$  and  $\nu = 0, 1$*

$$D_{h\nu} = \det (c_{ij})_{\substack{1 \leq i \leq h \\ \nu \leq j < h + \nu}} \neq 0. \tag{42}$$

**Proof.** In the notation of Definition 1 we have  $|\alpha_s| > 1$  for  $s < r$ ,  $\alpha_r = \varepsilon$ ,  $m_r = e$ . Assume that  $\alpha_s \in Z_0$ , if and only if  $s \in S_0$ . In the notation of Lemma 5

$$h = m_1 + \dots + m_{s(h)-1} + g(h), \quad 1 \leq g(h) \leq m_{s(h)}.$$

If  $\nu = 0$  or  $1$ ,  $D_{h\nu} = 0$  and  $\alpha_s \in Z_0$  for all  $s \leq s(h)$ , the system of equations

$$\sum_{j=\nu}^{h-1+\nu} c_{ij}x_j = 0 \quad (1 \leq i \leq h)$$

has a solution  $\langle x_\nu, \dots, x_{h-1+\nu} \rangle \in \mathbb{R}^h \setminus \{\mathbf{0}\}$ . It follows by Lemma 3 that

$$\sum_{j=\nu}^{h-1+\nu} x_j x^{n_j} \equiv 0 \left( \text{mod } \prod_{s=1}^{s(h)-1} (x - \alpha_s)^{m_s} (x - \alpha_{s(h)})^{g(h)} \right)$$

hence the polynomial  $\sum_{j=\nu}^{h-1+\nu} x_j x^{n_j} \in \mathbb{R}[x]$  has  $h$  zeros of the same sign, counted with multiplicities. This, however, contradicts the Descartes rule of signs (see [5], Satz 12), hence, if  $\alpha_s \in Z_0$  ( $s \leq s(h)$ ) (42) holds. It also shows that  $D_{h\nu}$  as a polynomial in  $\alpha_s$  ( $s \notin S_0$ ) is not identically zero for any fixed  $\alpha_s \in Z_0$ . Since the coefficients of the polynomial in question belong to  $\mathbb{Q}(Z_0)$ , the algebraic independence of  $\alpha_s$  ( $s \notin S_0$ ) over  $\mathbb{Q}(Z_0)$  implies  $D_{h\nu} \neq 0$ .

**Lemma 12.** *Under the assumptions of Theorem 5 let  $P_0$  be of degree  $d - e$ . For all positive integers  $e_1, \dots, e_{d-e}$  there exists a unique polynomial  $Q = Q(P_0; e_1, \dots, e_{d-e})$  such that*

$$Q = x^{\sum_{\mu=1}^{d-e} e_\mu} + \sum_{j=1}^{d-e} b_j x^{\sum_{\mu=j+1}^{d-e} e_\mu}$$

and

$$Q \equiv 0 \pmod{P_0}. \tag{43}$$

Moreover,  $Q \in \mathbb{R}[x]$ .

**Proof.** For  $j = 0, \dots, d - e$  put  $n_j = \sum_{\nu=j+1}^{d-e} e_\nu$  and for  $i \leq d - e$ , let  $c_{ij}$  be defined by Definition 1 with  $P$  replaced by  $P_0$ . By Lemma 3 the congruence (43) is equivalent to

$$\sum_{j=1}^{d-e} c_{ij} b_j = -c_{i0} \quad (1 \leq i \leq d - e).$$

By Lemma 11 with  $h = d - e$  and  $\nu = 1$  the determinant of this system is non-zero, hence  $b_j$  are uniquely determined. If we replace  $c_{ij}$  by  $\bar{c}_{ij}$  we obtain the same system of equations, hence  $b_j \in \mathbb{R}$ .

**Lemma 13.** *For every positive integer  $h < d - e$  and all positive integers  $e_1, \dots, e_h$  we have in the notation of Definition 3*

$$D(\{1, \dots, h + 1\}, e_1, \dots, e_h) > 0.$$

**Proof.** We have

$$\max_{i \leq h+1} g(i) \leq \max_{i \leq h+1} i = h + 1,$$

hence  $D(\{1, \dots, h + 1\}, e_1, \dots, e_h)$  is defined. Its only factor, which could possibly vanish is

$$\det \left( \alpha_{s(i)}^{\sum_{\mu=j+1}^h e_\mu} \prod_{f=0}^{g(i)-2} \sum_{\nu=f+1}^j e_\nu \right)_{\substack{1 \leq i \leq h+1 \\ 0 \leq j \leq h}} = \det \left( c_{ij} \alpha_{s(i)}^{-n_h} \right)_{\substack{1 \leq i \leq h+1 \\ 0 \leq j \leq h}},$$

where  $n_j = \sum_{\mu=j+1}^d e_\mu$ . By (42) with  $\nu = 0$  the above determinant is non-zero.

**Definition 5.** Let the sequence  $d_i(1 \leq i \leq d-e)$  be defined inductively as follows.

$$d_1 = \frac{\log(L(P) - 1)}{\log |\alpha_1|} \tag{44}$$

and if  $d_1, \dots, d_h$  ( $h < d - e$ ) are already defined

$$D_{h+1} = (h + 1)^{-1} h^{\frac{-h}{2}} \min\{D(\{1, \dots, h + 1\}, e_1, \dots, e_h) : 1 \leq e_i \leq d_i\}, \tag{45}$$

$$d_{h+1} = \max \left\{ d_1, \dots, d_h, \right. \\ \left. \psi \left( |\alpha_{s(i)}|, m - 1, \left( \max \left\{ 2, \frac{2(m - 1)}{\log |\alpha_s(h + 1)|} \right\} \right)^{m-1} \right) D_{h+1}^{-1}(L(P) - 1) \right\}. \tag{46}$$

**Lemma 14.** For every  $Q \in T_d(P)$ ,  $Q = x^{n_0} + \sum_{j=1}^d b_j x^{n_j}$  we have

$$n_{j-1} - n_j \leq d_j \quad (1 \leq j < d). \tag{47}$$

**Proof** is by induction on  $j$ . Since  $Q \in T_d(P)$  the equation

$$\alpha_1^{n_0} + \sum_{j=1}^d b_j \alpha_1^{n_j} = 0$$

implies

$$|\alpha_1|^{n_0} \leq |\alpha_1|^{n_1} \sum_{j=1}^d |b_j| \leq |\alpha_1|^{n_1} (L(Q) - 1) \leq |\alpha_1|^{n_1} (L(P) - 1),$$

which, in view of (44) gives (47) for  $j = 1$ . Assume now that (47) holds for all  $j \leq h$  ( $h < d - 1$ ). By Lemma 11 we have

$$\Delta = \det \left( c_{ij} \alpha_{s(i)}^{-n_h} \right)_{\substack{1 \leq i \leq h+1 \\ 0 \leq j \leq h}} \neq 0.$$

Let

$$M = \max_{1 \leq i \leq h+1} \left| \left( c_{i0} + \sum_{j=1}^h c_{ij} b_j \right) \alpha_{s(i)}^{-n_h} \right|.$$

Solving the system of equations

$$\left( c_{i0} x_0 + \sum_{j=1}^h c_{ij} x_j \right) \alpha_{s(i)}^{-n_h} = \left( c_{i0} + \sum_{j=1}^h c_{ij} b_j \right) \alpha_{s(i)}^{-n_h} \quad (1 \leq i \leq h + 1)$$

by means of Cramer's formulae and developing the numerator according to the first column we obtain

$$1 \leq \frac{(h+1)M \max \left| \det \left( c_{ij} \alpha_{s(i)}^{-n_h} \right)_{\substack{i \in I \\ 1 \leq j \leq h}} \right|}{|\Delta|},$$

where the maximum is taken over all subsets  $I$  of  $\{1, \dots, h+1\}$  of cardinality  $h$ . Now, by Lemma 7, since  $|\alpha_{s(i)}| > 1$ , we have

$$\max \left| \det \left( c_{ij} \alpha_{s(i)}^{-n_h} \right)_{\substack{i \in I \\ 1 \leq j \leq h}} \right| \leq h^{\frac{h}{2}} \prod_{i=1}^{h+1} |\alpha_{s(i)}|^{n_{\max\{1, g(i)-1\}}} \prod_{f=0}^{g(i)-2} (n_f - n_h).$$

This gives, by Definition 3,

$$m \geq (h+1)^{-1} h^{-\frac{1}{2}} D(\{1, \dots, h+1\}, n_0 - n_1, \dots, n_{h-1} - n_h)$$

and by the inductive assumption and (45)

$$M \geq D_{h+1}. \tag{48}$$

On the other hand, by Lemma 3

$$c_{i0} + \sum_{j=1}^d c_{ij} b_j = 0 \quad (1 \leq i \leq h+1),$$

hence

$$\left| \left( c_{i0} + \sum_{j=1}^h c_{ij} b_j \right) \alpha_{s(i)}^{-n_h} \right| \cdot |\alpha_{s(i)}|^{n_h} = \left| \sum_{j=1}^h c_{ij} b_j \right|. \tag{49}$$

By (48) for a certain  $i_0 \leq h+1$  the left-hand side is at least  $D_{h+1} |\alpha_{s(i)}|^{n_h}$ . As to the right-hand side, by (29) we obtain

$$\left| \sum_{j=1}^d c_{i_0 j} b_j \right| \leq |c_{i_0, h+1}| \left( \max \left\{ 1, \frac{m_{s(i_0)}-1}{\log |\alpha_{s(i_0)}|} \right\} \right)^{m_{s(i_0)}-1} (L(P) - 1). \tag{50}$$

If  $n_h - n_{h+1} \leq n_0 - n_h$ , we obtain  $n_h - n_{h+1} \leq d_1 + \dots + d_h \leq d_{h+1}$ , hence the inductive assumption holds. If  $n_h - n_{h+1} > n_0 - n_h$ , then by (30)

$$|c_{i_0, h+1}| \leq |\alpha_{s(i_0)}|^{n_{h+1}} 2^{m_{s(i_0)}-1} (n_h - n_{h+1})^{m_{s(i_0)}-1}.$$

Combining the inequality with (49) and (50) we obtain (31), where, however,  $D_{h+1}$  has the new meaning given by (45).

It follows, by Lemma 6

$$\begin{aligned} & n_h - n_{h+1} \\ & \leq \max_{1 \leq s \leq s(h+1)} \psi \left( |\alpha_s|, m_s - 1, \left( \max \left\{ 2, \frac{2(m_s - 1)}{\log |\alpha_s|} \right\} \right)^{m_s-1} \times D_{h+1}^{-1} (L(P) - 1) \right) \\ & \leq \psi \left( \alpha_{s(h+1)}, m - 1, \left( \max \left\{ 2, \frac{2(m - 1)}{\log |\alpha_{s(h+1)}|} \right\} \right)^{m-1} D_{h+1}^{-1} (L(P) - 1) \right) \leq d_{h+1} \end{aligned}$$

and the inductive proof is complete.

**Definition 6.** Assume that, under the assumptions of Theorem 5,  $e = 1$ . Put for positive integers  $n_1 > \dots > n_{d-2} > n_{d-1}$

$$\begin{aligned}
 (c'_{ij})_{\substack{1 \leq i < d \\ 1 \leq j \leq d}} &= C(P_0; n_1, \dots, n_{d-1}, 0), \quad c'_{dj} = 1 \quad (1 \leq j \leq d), \\
 E(P_0; n_1 - n_{d-1}, \dots, n_{d-2} - n_{d-1}) \\
 &= \left| \det \left( c'_{ij} \alpha_{s(i)}^{-n_{d-1}} \right)_{1 \leq i, j < d} \right|^{-1} \prod_{i=1}^{d-1} |\alpha_{s(i)}|^{n_{g(i)} - n_{d-1}} \prod_{f=1}^{g(i)-1} (n_f - n_{d-1}). \tag{51}
 \end{aligned}$$

**Remark.**  $E(P_0; n_1 - n_{d-1}, \dots, n_{d-2} - n_{d-1})$  is well defined since  $\det(c'_{ij} \alpha_{s(i)}^{-n_{d-1}})$  is non-zero by Lemma 11 with  $h = d - 1$ ,  $\nu = 0$ . Moreover the right-hand side of (51) depends only on  $P_0$  and the differences  $n_j - n_{d-1}$  ( $1 \leq j \leq d - 2$ ).

**Lemma 15.** Assume that, under the assumptions of Theorem 5 and in the notation of Definition 1,  $e = 1$ . If for positive integers  $n_1 > \dots > n_{d-1}$  and for  $n > 1$ ,  $a \in \mathbb{R}$

$$\begin{aligned}
 n_{d-1} > \max \left\{ n_1 - n_{d-1}, \psi \left( |\alpha_{r-1}|, m - 1, d(d - 1)^{\frac{d+1}{2}} 2^{m-1} n \right. \right. \\
 \left. \left. \times E(P_0; n_1 - n_{d-1}, \dots, n_{d-2} - n_{d-1}) \max \left\{ 1, \frac{nd|a|}{nd - 1} \right\} \right) \right\}, \tag{52}
 \end{aligned}$$

then there exists a polynomial  $R \in \mathbb{R}[x]$  of degree at most  $n_1$  such that

$$P_0 \mid R(x) - a, \tag{53}$$

$$x - 1 \mid R(x), \tag{54}$$

$$L(R) < \frac{1}{n}. \tag{55}$$

**Proof.** Put

$$R(x) = \sum_{j=1}^{d-1} r_j x^{n_j} + r_d, \quad r_j \in \mathbb{C}.$$

By Lemma 3 the conditions (53) and (54) are equivalent to the following system of linear equations for  $r_j$

$$\begin{aligned}
 \sum_{j=1}^d c'_{ij} r_j &= c'_{id} a \quad (1 \leq i < d) \\
 \sum_{j=1}^d c'_{dj} r_j &= 0.
 \end{aligned} \tag{56}$$



The determinant of this system equals

$$\Delta_0 = \prod_{i=1}^d \alpha_{s(i)}^{n_{d-1}} \det \left( c'_{ij} \alpha_{s(i)}^{-n_{d-1}} \right)_{1 \leq i, j \leq d}.$$

Developing the last determinant according to the last column we obtain

$$\begin{aligned} & \det \left( c'_{ij} \alpha_{s(i)}^{-n_{d-1}} \right)_{1 \leq i, j \leq d} \\ &= \det \left( c'_{ij} \alpha_{s(i)}^{-n_{d-1}} \right)_{1 \leq i, j < d} + \sum_{k=1}^{d-1} (-1)^{k+d} c'_{kd} \alpha_{s(k)}^{-n_{d-1}} \det \left( c'_{ij} \alpha_{s(i)}^{-n_{d-1}} \right)_{\substack{i \neq k \\ j < d}}, \end{aligned}$$

hence, by (21) with  $h = d-1$ ,  $I = \{1, \dots, d\} \setminus \{k\}$  and by the condition  $\alpha_r = e = 1$

$$\begin{aligned} & \left| \Delta_0 \prod_{i=1}^d \alpha_{s(i)}^{-n_{d-1}} - \det \left( c'_{ij} \alpha_{s(i)}^{-n_{d-1}} \right)_{1 \leq i, j < d} \right| \tag{57} \\ & < (d-1)^{\frac{d+1}{2}} |\alpha_{r-1}|^{-n_{d-1}} \left( \prod_{i=1}^{d-1} |\alpha_{s(i)}|^{n_{g(i)} - n_{d-1}} \prod_{f=1}^{g(i)-1} (n_f - n_{d-1}) \right) \max_{1 \leq k < d} |c'_{kd}|. \end{aligned}$$

Since, by (52)  $n_{d-1} > n_1 - n_{d-1}$ , we have

$$\max_{1 \leq k < d} |c'_{kd}| \leq \prod_{f=1}^{m-1} n_f \leq (2n_{d-1})^{m-1}. \tag{58}$$

In view of Definition 6 the right-hand side of (57) does not exceed

$$\begin{aligned} & (d-1)^{\frac{d+1}{2}} |\alpha_{r-1}|^{-n_{d-1}} \left| \det \left( c'_{ij} \alpha_{s(i)}^{-n_{d-1}} \right)_{1 \leq i, j < d} \right| \\ & \times E(P_0; n_1 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0) 2^{m-1} n_{d-1}^{m-1}. \end{aligned}$$

Since, by (52),

$$\begin{aligned} & n_{d-1} > \psi \left( |\alpha_{r-1}|, m-1, d(d-1)^{\frac{d+1}{2}} 2^{m-1} n \right. \\ & \left. \times E(P_0; n_1 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0) \right) \end{aligned}$$

we have by Lemma 6 and (57)

$$|\Delta_0| > \left( 1 - \frac{1}{dn} \right) \left| \det (c'_{ij})_{1 \leq i, j < d} \right|, \tag{59}$$

hence by the Remark after Definition 6,  $\Delta_0 \neq 0$ . Thus the system (56) is uniquely solvable and since on replacing  $c'_{ij}$  by  $\bar{c}'_{ij}$  we obtain the same system,  $r_j$  are real ( $1 \leq j \leq d$ ).

The determinant  $\Delta_k$  obtained by substituting in  $(c'_{ij})_{1 \leq i, j \leq d}$  for the  $k$ -th column the column

$$[c'_{1d}, \dots, c'_{d-1d}, 0]^t a$$

satisfies for  $k < d$

$$\Delta_k = \pm (\det c'_{ij})_{\substack{i < d \\ j \neq k}} a,$$

hence developing the last determinant according to the last column, using Lemma 10, Definition 6 and (58) we obtain

$$\begin{aligned} |\Delta_k| &\leq |a| \sum_{l=1}^{d-1} |c'_{ld}| (d-2)^{\frac{d-2}{2}} \prod_{\substack{i=1 \\ i \neq l}}^{d-1} |\alpha_{s(i)}|^{n_{g(i)}} \prod_{f=1}^{g(i)-1} (n_f - n_{d-1}) \\ &\leq |a| (d-1)(d-2)^{\frac{d-2}{2}} (2n_{d-1})^{m-1} |\alpha_{r-1}|^{-n_{d-1}} \left| \det (c'_{ij})_{1 \leq i, j < d} \right| \\ &\quad \times E(P_0; n_1 - n_{d-1}, \dots, n_{d-2} - n_{d-1}), \end{aligned}$$

where  $(d-2)^{\frac{d-2}{2}} = 1$  for  $d = 2$ .

Since  $(d-1)(d-2)^{\frac{d-2}{2}} < (d-1)^{\frac{d+1}{2}}$  we obtain, by virtue of (52),

$$|\Delta_k| < \frac{dn-1}{d^2 n^2} \left| \det (c'_{ij})_{1 \leq i, j < d} \right|$$

hence, by (59),  $r_k = \Delta_k/\Delta_0$  satisfies

$$|r_k| < \frac{1}{dn} \quad (1 \leq k < d). \tag{60}$$

It remains to consider  $k = d$ . In this case developing  $\Delta_d$  according to the last column we obtain

$$|\Delta_d| \leq |a| \sum_{l=1}^{d-1} |c'_{ld}| \left| \det (c'_{ij})_{\substack{i \neq l \\ j < d}} \right|.$$

Using (21) with  $h = d-1$ ,  $I = \{1, \dots, d\} \setminus \{l\}$ , the condition  $\alpha_r = e = 1$  and (58) we obtain

$$\begin{aligned} |\Delta_d| &\leq |a| (d-1)^{\frac{d+1}{2}} (2n_{d-1})^{m-1} |\alpha_{r-1}|^{-n_{d-1}} \\ &\quad \times \prod_{i=1}^{d-1} |\alpha_{s(i)}|^{n_{g(i)}} \prod_{f=1}^{g(i)-1} (n_f - n_{d-1}) \leq (d-1)^{\frac{d+1}{2}} 2^{m-1} n_{d-1}^{m-1} |\alpha_{r-1}|^{-n_{d-1}} \\ &\quad \times \left| \det (c'_{ij})_{1 \leq i, j < d} \right| E(P_0; n_1 - n_{d-1}, \dots, n_{d-2} - n_{d-1}). \end{aligned}$$

Again, by virtue of (52) and of Lemma 6,

$$|\Delta_d| < \frac{dn - 1}{d^2 n^2} \left| \det (c'_{ij})_{1 \leq i, j < d} \right|,$$

hence  $r_d = \Delta_d/\Delta_0$  satisfies

$$|r_d| < \frac{1}{dn}.$$

It follows now from (60) that

$$L(R) = \sum_{k=1}^d |r_k| < \frac{1}{n},$$

which proves (55).

**Lemma 16.** Assume, under the assumptions of Theorem 5, that  $\varepsilon = e = 1$ . Then

$$l(P) \leq \inf_{Q \in S_{d-1}(P_0)} \left\{ L(Q) + |Q(1)| \right\}.$$

**Proof.** Let

$$Q = x^{q_0} + \sum_{j=1}^{d-1} b_j x^{q_j},$$

where  $q_0 > q_1 > \dots > q_{d-1}$  and we may assume  $q_{d-1} = 0$ . By Lemma 15 with  $a = Q(1)$ ,  $n_j = n_{d-1} + q_j$  ( $1 \leq j < d$ ), if

$$n_{d-1} > \max \left\{ q_1, \psi \left( |\alpha_{r-1}|, m-1, d(d-1)^{\frac{d+1}{2}} 2^{m-1} n \right) \right. \\ \left. \times E(P_0; q_1, \dots, q_{d-2}) \max \left\{ 1, \frac{nd}{nd-1} |Q(1)| \right\} \right\}$$

there exists a polynomial  $R \in \mathbb{R}[x]$  of degree at most  $n_1$  satisfying (53)–(55). We consider the polynomial

$$S(x) = Q(x)x^{n_{d-1}} + R(x) - Q(1).$$

It follows from (53)–(54) that

$$P_0 \mid S, \quad x - 1 \mid S, \quad \text{thus } P \mid S$$

and since  $S$  is monic

$$l(P) \leq L(S).$$

On the other hand, by (55)

$$L(S) \leq L(Q) + |Q(1)| + \frac{1}{n}.$$

Since  $n$  is arbitrary, the lemma follows.

**Proof of Theorem 5.** Since, by Proposition (iii),  $l(P(-x)) = l(P(x))$ , we may assume that  $\varepsilon = 1$  and, by virtue of Theorem 3, that  $e = 1$ . Thus Lemmas 15 and 16 are applicable. The second part of the theorem follows from Lemma 16. In order to prove the first part we shall show that for every  $n > 1$

$$0 \geq l(P) - \min^* \min \{L(Q(P; n_0, \dots, n_{d-1}, 0)), L(Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)) + |Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)(1)|\} > -\frac{1}{n}, \tag{61}$$

where the  $\min^*$  is taken over all integers  $n_0 > \dots > n_{d-1} > 0$  such that

$$n_{j-1} - n_j \leq d_j \quad (1 \leq j < d) \tag{62}$$

$$\begin{aligned} n_{d-1} \leq & \psi \left( |\alpha_{r-1}|, m-1, d(d-1)^{\frac{d+1}{2}} 2^{m-1} n \right. \\ & \times E(P_0; n_1 - n_{d-1}, \dots, n_{d-2} - n_d) \\ & \left. \times \max \left\{ 1, \frac{nd}{nd-1} |Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)(1)| \right\} \right) \end{aligned} \tag{63}$$

and, in the notation of Definition 1

$$|C_1(P; n_1, \dots, n_{d-1}, 0)| \neq 0. \tag{64}$$

The condition (64) implies that there is a unique polynomial

$$Q = x^{n_0} + \sum_{j=1}^{d-1} b_j x^{n_j} + b_d$$

divisible by  $P$ , denoted in (61) by  $Q(P; n_0, \dots, n_{d-1}, 0)$ . Similarly  $Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)$  is the unique polynomial

$$Q = x^{n_0 - n_{d-1}} + \sum_{j=1}^{d-1} b_j x^{n_j - n_{d-1}}$$

divisible by  $P_0$ . The inequality

$$l(P) \leq \min^* L(Q(P; n_0, \dots, n_{d-1}, 0))$$

is clear and the inequality

$$\begin{aligned} l(P) \leq & \min^* \{L(Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)) \\ & + |Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)(1)|\} \end{aligned}$$

follows from Lemma 16. This shows the first of inequalities (61). In order to prove the second one we notice that by Lemmas 4 and 14

$$l(P) = \inf L(Q(P; n_0, \dots, n_{d-1}, 0)), \tag{65}$$

where  $\langle n_0, \dots, n_{d-1} \rangle$  runs through all strictly decreasing sequences of  $d$  positive integers satisfying (62) and (64). If (63) is satisfied then, clearly

$$L(Q(P; n_0, \dots, n_{d-1}, 0)) \geq \min^* L(Q(P; n_0, \dots, n_{d-1}, 0)) \tag{66}$$

and, if not, then by Lemma 15 there exists a polynomial  $R \in \mathbb{R}[x]$  of degree at most  $n_1$  such that (53)–(55) hold with

$$a = Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)(1).$$

Then the polynomial

$$\begin{aligned} S(x) &= Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0) x^{n_{d-1}} + R(x) \\ &\quad - Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)(1) \end{aligned}$$

is monic, satisfies

$$P \mid S(x)$$

and, by (64),

$$S(x) = Q(P; n_0, \dots, n_{d-1}, 0). \tag{67}$$

By (55)

$$\begin{aligned} L(S) &> L(Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)) \\ &\quad + |Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)(1)| - \frac{1}{n}. \end{aligned} \tag{68}$$

The formulae (66)–(68) imply

$$\begin{aligned} &L(Q(P; n_0, \dots, n_{d-1}, 0)) \\ &\geq \min^* \min \left\{ L(Q(P; n_0, \dots, n_{d-1}, 0)), L(Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)) \right. \\ &\quad \left. + |Q(P_0; n_0 - n_{d-1}, \dots, n_{d-2} - n_{d-1}, 0)(1)| \right\} - \frac{1}{n} \end{aligned}$$

for all sequences  $\langle n_0, \dots, n_{d-1} \rangle$  satisfying (64), hence by (65) the second of the inequalities (61) follows. The conditions (62) and (63) are for a given  $n$  satisfied by only finitely many sequences  $\langle n_0, \dots, n_{d-1} \rangle$ , since

$$n_j - n_{d-1} \leq \sum_{\mu=j+1}^{d-1} d_\mu$$

and all such sequences can be effectively determined, hence  $l(P)$  can be effectively computed.

For the proof of Theorem 6 we need

**Definition 7.** For  $\alpha, \beta$  in  $\mathbb{C}$  and  $n > m > 0$

$$Q_n(\alpha, \beta) = \begin{cases} \frac{\alpha^n - \beta^n}{\alpha - \beta} & \text{if } \alpha \neq \beta, \\ n\alpha^{n-1} & \text{if } \alpha = \beta, \end{cases}$$

$$E_{n,m}(\alpha, \beta) = \left| \frac{Q_n(\alpha, \beta)}{Q_m(\alpha, \beta)} \right| + |\alpha\beta|^m \left| \frac{Q_{n-m}(\alpha, \beta)}{Q_m(\alpha, \beta)} \right|,$$

$$F_{n,m}(x, \beta) = x^n - \beta^n + |\beta|^m x^m (x^{n-m} - \beta^{n-m}) - (2x - 1)(x^m - \beta^m).$$

**Lemma 17.** In the notation of Definition 2, if  $P(x) = (x - \alpha)(x - \beta)$ ,  $\alpha\beta \neq 0$ , then all elements of  $T_2(P)$  are of the form

$$x^n - \frac{Q_n(\alpha, \beta)}{Q_m(\alpha, \beta)} x^m + (\alpha\beta)^m \frac{Q_{n-m}(\alpha, \beta)}{Q_m(\alpha, \beta)} = F_{n,m}(x; \alpha, \beta), \tag{69}$$

where  $n > m > 0$ ,  $Q_m(\alpha, \beta)Q_{n-m}(\alpha, \beta) \neq 0$ .

**Proof.** Let an element  $Q$  of  $T_2(P)$  be  $x^n + Ax^m + B$ , where  $n > m > 0$ . By Lemma 3 the condition  $Q \equiv 0 \pmod{P}$  is equivalent to

$$c_{i0} + c_{i1}A + c_{i2}B = 0 \quad (i = 1, 2), \tag{70}$$

where  $c_{ij}$  are given in Definition 1 for  $\alpha_1 = \alpha$ ,  $\alpha_2 = \beta$ , hence

$$\begin{aligned} c_{10} &= \alpha^n, \quad c_{11} = \alpha^m, \quad c_{12} = 1; \\ c_{20} &= \beta^n, \quad c_{21} = \beta^m, \quad c_{22} = 1, \quad \text{if } \beta \neq \alpha; \\ c_{20} &= 0, \quad c_{21} = (m - n)\beta^m, \quad c_{22} = -n, \quad \text{if } \beta = \alpha. \end{aligned}$$

Since  $Q \in T_2(P)$  we have

$$|C_0(P; n, m)| \neq 0 \neq |C_1(P; n, m)|,$$

hence  $Q_m(\alpha, \beta)Q_{n-m}(\alpha, \beta) \neq 0$ . Solving the system (70) we obtain for  $Q$  the form (69).

**Lemma 18.** If  $\beta \in \mathbb{R}$ ,  $|\beta| \geq 1$ , then for all positive integers  $n > m$  and all integers  $k \geq 0$  we have

$$G_{n,m,k}(\beta) = \frac{1}{k!} \frac{d^k}{dx^k} F_{n,m}(x, \beta)|_{x=|\beta|} \geq 0$$

and if  $|\beta| > 1$  for  $k = 0$  or 1

$$\inf_{n > m} G_{n,m,k}(\beta) > 0.$$

**Proof.** Consider first the case  $\beta > 0$ . For  $k = 0$  we have  $G_{n,m,k}(\beta) = 0$ . For  $k \geq 1$  we have

$$\begin{aligned} G_{n,m,k}(\beta) &= \binom{n}{k} \beta^{n-k} + \binom{n}{k} \beta^{n+m-k} - \binom{m}{k} \beta^{n+m-k} \\ &\quad - 2 \binom{m+1}{k} \beta^{m-k+1} + \binom{m}{k} \beta^{m-k} + 2 \binom{1}{k} \beta^{m-k+1} \\ &= \beta^{m-k} \left( \binom{n}{k} \beta^{n-m} + \binom{n}{k} \beta^n - \binom{m}{k} \beta^n \right. \\ &\quad \left. - 2 \binom{m+1}{k} \beta + \binom{m}{k} + 2 \binom{1}{k} \beta \right). \end{aligned}$$

The expression in the parenthesis is non-negative, since for  $\beta = 1$  it is equal to

$$2 \binom{n}{k} - 2 \binom{m+1}{k} + 2 \binom{1}{k} \geq 0$$

and its derivative with respect to  $\beta$  is

$$\begin{aligned} &\binom{n}{k} (n-m) \beta^{n-m-1} + \left( \binom{n}{k} - \binom{m}{k} \right) n \beta^{n-1} - 2 \binom{m+1}{k} + 2 \binom{1}{k} \\ &\geq \binom{m+1}{k} + \left( \binom{m+1}{k} - \binom{m}{k} \right) (m+1) - 2 \binom{m+1}{k} + 2 \binom{1}{k} \\ &= (k-1) \binom{m+1}{k} + 2 \binom{1}{k} \geq 2 \binom{1}{k}. \end{aligned}$$

It follows that

$$G_{n,m,k}(\beta) \geq 2(\beta - 1)$$

and the obtained lower bound, independent of  $n, m$  is positive for  $\beta > 1$ . Consider now the case  $\beta < 0$ . We distinguish four cases according to the parity of  $n, m$ .

If  $n \equiv m \equiv 0 \pmod{2}$ , then

$$G_{n,m,k}(\beta) = G_{n,m,k}(|\beta|)$$

and the case reduces to the former.

If  $n \equiv 0, m \equiv 1 \pmod{2}$ , then

$$G_{n,m,0}(\beta) = 2(|\beta|)^m (|\beta|^n - 2|\beta| + 1) \geq 2(|\beta| - 1)^2$$

and the obtained lower bound, independent of  $n, m$  is positive for  $|\beta| > 1$ . Further, for  $k \geq 1$

$$\begin{aligned} G_{n,m,k}(\beta) &= \binom{n}{k} |\beta|^{n-k} + \binom{n}{k} |\beta|^{n+m-k} + \binom{m}{k} |\beta|^{n+m-k} \\ &\quad - 2 \binom{m+1}{k} |\beta|^{m-k+1} + \binom{m}{k} |\beta|^{m-k} - 2 \binom{1}{k} |\beta|^{m-k+1} \\ &= |\beta|^{m-k} \left( \binom{n}{k} |\beta|^{n-m} + \binom{n}{k} |\beta|^n + \binom{m}{k} |\beta|^n - 2 \binom{m+1}{k} |\beta| \right. \\ &\quad \left. + \binom{m}{k} - 2 \binom{1}{k} |\beta| \right). \end{aligned}$$

The expression in the parenthesis is non-negative, since

$$\binom{n}{k} |\beta|^{n-m} + \binom{n}{k} |\beta|^n \geq 2 \binom{m+1}{k} |\beta|$$

and

$$\binom{m}{k} |\beta|^n + \binom{m}{k} \geq \binom{m}{k} (|\beta|^2 + 1) \geq 2 \binom{1}{k} |\beta|.$$

If  $n \equiv 1$ ,  $m \equiv 0 \pmod{2}$ , then

$$G_{n,m,k}(\beta) \geq G_{n,m,k}(|\beta|)$$

and the case reduces to the already considered one.

Finally, if  $n \equiv m \equiv 1 \pmod{2}$ , then

$$G_{n,m,0}(\beta) = 2|\beta|^m (|\beta|^{n-m} - 2|\beta| + 1) \geq 2(|\beta| - 1)^2$$

and the obtained lower bound, independent of  $n, m$  is positive for  $|\beta| > 1$ .

Further, for  $k \geq 1$

$$\begin{aligned} G_{n,m,k}(\beta) &= \binom{n}{k} |\beta|^{n-k} + \binom{n}{k} |\beta|^{n+m-k} - \binom{m}{k} |\beta|^{n+m-k} \\ &\quad - 2 \binom{m+1}{k} |\beta|^{m-k+1} + \binom{m}{k} |\beta|^{m-k} - 2 \binom{1}{k} |\beta|^{m-k+1} \\ &= |\beta|^{m-k} \left( \binom{n}{k} |\beta|^{n-m} + \binom{n}{k} |\beta|^n - \binom{m}{k} |\beta|^n - 2 \binom{m+1}{k} |\beta| \right. \\ &\quad \left. + \binom{m}{k} - 2 \binom{1}{k} |\beta| \right). \end{aligned}$$

The expression in the parenthesis is non-negative, since for  $|\beta| = 1$  it is equal to

$$2 \binom{n}{k} - 2 \binom{m+1}{k} - 2 \binom{1}{k} \geq 2 \binom{m+2}{k} - 2 \binom{m+1}{k} - 2 \binom{1}{k} \geq 0$$

and its derivative with respect to  $|\beta|$  is

$$\begin{aligned} &\binom{n}{k} (n-m) |\beta|^{n-m-1} + \binom{n}{k} n - |\beta|^{n-1} - \binom{m}{k} n |\beta|^{n-1} - 2 \binom{m+1}{k} - 2 \binom{1}{k} \\ &\geq \binom{n}{k} (n-m) + \binom{n}{k} n - \binom{m}{k} n - 2 \binom{m+1}{k} - 2 \binom{1}{k} \\ &\geq 2 \binom{m+2}{k} + \left( \binom{m+2}{k} - \binom{m}{k} \right) (m+2) - 2 \binom{m+1}{k} - 2 \binom{1}{k} \geq 0. \end{aligned}$$

**Proof of Theorem 6.** Consider first the case of  $\alpha, \beta$  real. Since, by Proposition (iii),  $l(P(-x)) = l(P(x))$ , we may assume that  $\alpha > 0$ , hence  $\alpha \geq |\beta|$ .



By the Taylor formula we have in the notation of Lemma 18

$$(\alpha^m - \beta^m)(E_{n,m}(\alpha, \beta) - 2\alpha + 1) = \sum_{k=0}^n G_{n,m,k}(\beta)(\alpha - |\beta|)^k,$$

hence, by the said lemma,

$$(\alpha^m - \beta^m)(E_{n,m}(\alpha, \beta) - 2\alpha + 1) \geq 0 \tag{71}$$

and, if  $\alpha > |\beta| > 1$

$$\inf_{n>m} (\alpha^m - \beta^m)(E_{n,m}(\alpha, \beta) - 2\alpha + 1) > 0. \tag{72}$$

If  $\alpha \neq \pm\beta$  then (71) gives

$$E_{n,m}(\alpha, \beta) \geq 2\alpha - 1,$$

hence by Lemma 17

$$\inf_{Q \in T_2(P)} L(Q) \geq 2\alpha$$

and by Lemma 4,

$$l(P) \geq 2\alpha. \tag{73}$$

Now, if  $\beta = -1$ , then  $L(P) = 2\alpha$ , hence  $l(P) \leq 2\alpha$  and, by (73),  $l(P) = 2\alpha$ . If  $\beta = 1$ , then by Theorem 5 with  $P_0 = x - \alpha$

$$l(P) \leq L(P_0) + |P_0(1)| = 1 + \alpha + \alpha - 1 = 2\alpha$$

and, by (73),  $l(P) = 2\alpha$  again.

If  $\alpha > |\beta| > 1$ , then by (72)

$$\inf_{\substack{n>m \\ m < m_0}} E_{n,m}(\alpha, \beta) > 2\alpha - 1 \tag{74}$$

for every  $m_0$ . Choose now

$$m_0 = \frac{\log 4\alpha - \log(\alpha - |\beta|)}{\log |\beta|}.$$

Then for  $m \geq m_0 : E_{n,m}(\alpha, \beta) \geq |\alpha\beta|^m \frac{\alpha - |\beta|}{2\alpha^m} \geq 2\alpha$  and, by (74)

$$\inf_{n>m} E_{n,m}(\alpha, \beta) > 2\alpha - 1.$$

Using, as above, Lemmas 17 and 4 we obtain

$$l(P) > 2\alpha.$$

If  $\alpha = -\beta$ , then  $P(x) = x^2 - \alpha^2$  and by Proposition (iv) and Proposition A (ii)

$$l(P) = l(x - \alpha^2) = 1 + \alpha^2 \begin{cases} = 2\alpha & \text{if } \alpha = 1, \\ > 2\alpha, & \text{otherwise.} \end{cases}$$

If  $\alpha = \beta$ , then

$$E_{n,m}(\alpha, \beta) - 2\alpha + 1 = \frac{n\alpha^{n-m} + (n-m)\alpha^n}{m} - 2\alpha + 1.$$

The right-hand side is equal to  $2(n-m)/m > 0$  for  $\alpha = 1$  and its derivative with respect to  $\alpha$  is

$$\frac{n(n-m)}{m} (\alpha^{n-m-1} + \alpha^{n-m}) - 2 > \alpha - 1.$$

For  $\alpha = \beta = 1$ ,  $l(P) = 2 = 2\alpha$ , by Theorem 4; otherwise

$$\inf_{n>m} E_{n,m}(\alpha, \alpha) > 2\alpha - 1$$

and, by Lemmas 17 and 4,  $l(P) > 2\alpha$ .

Consider now the case, where  $\alpha, \beta$  are complex conjugate:

$$\alpha = |\alpha|e^{2i\varphi}, \quad \beta = |\alpha|e^{-2i\varphi}, \quad \varphi \in \left(0, \frac{\pi}{2}\right), \quad |\alpha| > 1$$

(the case  $|\alpha| = 1$  is settled by Theorem 4). Then

$$E_{n,m}(\alpha, \beta) = |\alpha|^{n-m} \left| \frac{\sin n\varphi}{\sin m\varphi} \right| + |\alpha|^n \left| \frac{\sin(n-m)\varphi}{\sin m\varphi} \right|,$$

where, by virtue of the condition  $Q_m(\alpha, \beta) \neq 0$  we have  $\sin m\varphi \neq 0$ . Since

$$|\sin m\varphi| \leq |\sin n\varphi| + |\sin(n-m)\varphi| \tag{75}$$

we have

$$E_{n,m}(\alpha, \beta) \geq |\alpha|^{n-m} \geq |\alpha|^2$$

unless  $n-m=1$ . In this final case we have, by (75)

$$\left| \frac{\sin n\varphi}{\sin m\varphi} \right| \geq 1 - \left| \frac{\sin \varphi}{\sin m\varphi} \right|$$

and by the well known inequality

$$\left| \frac{\sin \varphi}{\sin m\varphi} \right| \geq \frac{1}{m}.$$

Hence

$$\begin{aligned} E_{n,m}(\alpha, \beta) &\geq |\alpha| \left( 1 - \left| \frac{\sin \varphi}{\sin m\varphi} \right| \right) + |\alpha|^{m+1} \left| \frac{\sin \varphi}{\sin m\varphi} \right| \\ &\geq |\alpha| + \frac{|\alpha|^{m+1} - |\alpha|}{m} \geq |\alpha| + |\alpha|(|\alpha| - 1) = |\alpha|^2, \end{aligned}$$

where in the middle we have used Bernoulli's inequality. It follows, by Lemma 17, that  $L(Q) \geq 1 + |\alpha|^2$  for every  $Q \in T_2(P)$ , hence, by Lemma 4,

$$l(P) \geq 1 + |\alpha|^2 > 2|\alpha|.$$

**Proof of Corollary 2.** If  $\deg P = 1$ , then  $l(P) \in K(P)$  follows from Proposition A. If  $P = a(x - \alpha)(x - \beta)$ , where  $|\alpha| \geq |\beta| > 1$ , then, by Theorem 2,  $l(P)$  is attained and by Theorem 1,  $l(P) \in K(P)$ . If  $P = a(x - \alpha)(x - \beta)$ , where  $|\beta| = 1$ , then, by Theorem 6,  $l(P) = 2|a\alpha|$ . Since either  $|\alpha| = 1$  or  $\alpha \in \mathbb{R}$ ,  $l(P) \in K(P)$  follows.

**Proof of Corollary 3.** If, in the notation of the Corollary,  $|\beta| > 1$ , then, by Proposition A,  $l(P^*) = |\alpha\beta|$  and, by Proposition (ii)  $l(P) \geq |\alpha\beta|$ , thus  $\widehat{l}(P) = |\alpha\beta|$ . If  $|\alpha| > 1 = |\beta|$ , then, by Proposition (iii) and Theorem 6,  $l(P^*) = 2|\alpha| = l(P)$ , thus  $\widehat{l}(P) = 2|\alpha|$ . If  $|\alpha| > 1 > |\beta|$ , then, by Proposition A,  $l(P^*) = 1 + |\alpha|$ ,  $l(P) = |\alpha\beta|(1 + |\beta|^{-1})$ , hence  $\widehat{l}(P) = |\alpha| + \min\{1, |\alpha\beta|\}$ . If  $|\alpha| = 1 = |\beta|$ , then by Theorem 6,  $l(P) = l(P^*) = 2$ . If  $|\alpha| = 1 > |\beta|$ , then, by Proposition A,  $l(P) = 2$ , by Theorem 6,  $l(P^*) = |\alpha\beta|2|\beta|^{-1} = 2$ , thus  $\widehat{l}(P) = 2$ . Finally, if  $|\alpha| < 1$ , then by Proposition A,  $l(P) = 1$ , by Proposition (ii)  $l(P^*) \geq 1$ , thus  $\widehat{l}(P) = 1$ .

**Proof of Corollary 4.** If  $|\alpha| > 1 > |\beta| > 0$  we have  $\widehat{l}(x - \alpha) = |\alpha|$ ,  $\widehat{l}(x - \beta) = 1$ ,

$$\widehat{l}((x - \alpha)(x - \beta)) = |\alpha| + \min\{1, |\alpha\beta|\} > |\alpha|.$$

**Note added in proof.** An apparently similar problem has been considered in [2] and [3]. However, the restriction of  $G$  in the definition of  $l(P)$  to polynomials with integer coefficients makes a great difference, shown by the fact, clear from Lemma 17 above, that no analogue of Lemma 2 of [2] or Lemma 3 of [3] holds in our case.

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