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Entropy and expansion

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Abstract. Shearer's inequality bounds the sum of joint entropies of random variables in terms of the total joint entropy. We give another lower bound for the same sum in terms of the individual entropies when the variables are functions of independent random seeds. The inequality involves a constant characterizing the expansion properties of the system.

Our results generalize to entropy inequalities used in recent work in invariant settings, including the edge-vertex inequality for factor-of-IID processes, Bowen's entropy inequalities, and Bollobás's entropy bounds in random regular graphs.

The proof method yields inequalities for other measures of randomness, including covariance.

As an application, we give upper bounds for independent sets in both finite and infinite graphs.

Résumé. L'inégalité de Shearer limite la somme des entropies conjointes de variables aléatoires en termes d'entropie conjointe totale. Nous donnons une autre borne inférieure pour la même somme en termes d'entropies individuelles lorsque les variables sont des fonctions de nombres aléatoires indépendants. Le coefficient de l'inégalité caractérise les propriétés d'expansion du système.

Nos résultats se généralisent aux inégalités d'entropie utilisées dans des travaux récents dans des contextes invariants, y compris l'inégalité arête-sommet pour les processus de facteur d'IID, les inégalités d'entropie de Bowen et les limites d'entropie de Bollobás dans des graphes réguliers aléatoires.

La méthode de la preuve produit des inégalités pour d'autres mesures de l'aléatoire, y compris la covariance.

Comme application, nous donnons des bornes supérieures pour des ensembles indépendants dans des graphes finis et infinis.

1. Introduction

In recent years entropy inequalities for so-called *local algorithms* and *factor-of-IID processes* have been used to great effect in proving new results about random regular graphs and regular graphs of large girth. The method of entropy bounds actually goes back to a 1981 result of Bollobás [8] where he gave an upper bound on the independence ratio of random regular graphs. The counting argument behind that result essentially proves an *edge-vertex entropy inequality* in a special setting. Interestingly, this entropy inequality also comes up in the seminal work of L. Bowen on the f-invariant [9,10].

All previous proofs for the edge-vertex entropy inequality heavily relied on the fact that the underlying graph does not contain (short) cycles. One would think that the acyclic property of the graph is not crucial. Indeed, in this paper we present an approach that directly relates the existence of such entropy inequalities to expansion properties of the graph allowing us to vastly generalize these inequalities.

1.1. *The setup*

Throughout the paper X_v, $v \in V$, will denote a finite or infinite system of random variables obtained as follows. Suppose that we have a collection of independent random variables $Z_1, Z_2, \ldots$ that we will refer to as *random seeds*. Then for each v, X_v is obtained as a measurable function of some specified subcollection of these random seeds. The main result is a general entropy inequality involving the Shannon entropies of our random variables X_v and the joint entropies of X_v

for given subsets of V. We will see that the best coefficient for which our inequality holds can be expressed as a certain hyperedge expansion of the system.

One may think of X_v, $v \in V$, as the output of a randomized (local) algorithm on a (large) network, where V denotes the set of nodes. The algorithm is distributed among the nodes with each node having access only to a certain subcollection of the random seeds used by the algorithm, usually dictated by the restriction that each node can communicate only with close-by nodes. The entropy inequalities presented in this paper provide constraints for what can be achieved by such distributed algorithms.

As we will see, the finite version nicely complements *Shearer's inequality*, while the infinite version generalizes the afore-mentioned edge-vertex entropy inequality.

1.2. *The finite version*

For the moment let X_v, $v \in V$, denote any random variables indexed by elements of a finite set V such that each X_v takes values from a finite set. Suppose that we have a hypergraph G over the vertex set V. This simply means that we have an edge set E containing subsets e of V, that is, each hyperedge $e \in E$ is a subset $e \subseteq V$. For the sake of simplicity the reader may think of G as an ordinary graph where each edge e is a pair of vertices. One can consider the joint entropy $H(X_e)$ of the random variables X_v, $v \in e$. Given a hypergraph over V with edge set E, *Shearer's inequality* [11] says that

$$\sum_{e \in E} H(X_e) \geq k H(X_V), \tag{1}$$

where k is such that each $v \in V$ is contained by at least k hyperedges (i.e. k is the minimum degree of the hypergraph) and $H(X_V)$ denotes the joint entropy of all X_v, $v \in V$. Shearer's inequality has many elegant applications in various areas, most notably in set systems and graph homomorphisms.

We would like to complement (1) by proving another lower bound for the same sum of joint entropies but this time in terms of the individual entropies $H(X_v)$. If no additional assumption is made, then only trivial bounds can be obtained (i.e. bounds that follow from the inequalities $H(X_e) \geq H(x_v)$, $v \in e$).

To get a non-trivial inequality, we will work with random variables X_v that are measurable functions of independent random seeds with the sole restriction that each X_v "can see" only some specified subcollection of the seeds. We will prove that

$$\sum_{e \in E} H(X_e) \geq \beta \sum_{v \in V} H(X_v),$$

where β is related to the expansion properties of the hypergraph with respect to sets describing which X_v can see a given seed; see Theorem 2.4. To be more specific, but without having to go into too many details, we state our inequality here for the case when G is a regular graph.

Theorem 1 (Regular graph version). *Suppose that G is a d-regular graph with vertex set V and edge set E. Suppose that $Z_1, \ldots, Z_m$ are independent random variables and a subset $W_i \subseteq V$ is given for each $i = 1, \ldots, m$. Informally, the random value of Z_i is passed on to the vertices in W_i, and for each $v \in V$ we have a random variable X_v that is an arbitrary measurable function of the values passed on to v. More precisely, for each vertex $v \in V$ let X_v be a measurable function of those Z_i for which $v \in W_i$. Then*

$$\sum_{\{u,v\} \in E} H(X_u, X_v) \geq \beta \sum_{v \in V} H(X_v), \quad \text{where } \beta = \frac{1}{2}\left(d + \min_{1 \leq i \leq m;\, W \subseteq W_i} \frac{|\partial W|}{|W|}\right).$$

Here ∂W denotes the edge boundary *of W, that is, the set of edges with one endpoint in W and one endpoint outside W. The above minimum of $|\partial W|/|W|$ is the* edge Cheeger constant (*or* isoperimetry number) *of G with respesct to the set system $W_1, \ldots, W_m$.*

We will see that this β is the largest coefficient for which the theorem holds.

1.3. *Factor-of-IID processes*

The following question arises: can we use our method to say something if not finitely but countably many random variables X_v are given? In this case the sum of all entropies may be infinite. Therefore, to get a meaningful statement, one needs to assume some kind of invariance. A natural framework to consider is *factor-of-IID processes*.

Factors of IID are closely related to randomized local algorithms. They are also extensively studied by ergodic theory under the name of *factors of Bernoulli shifts*. In short, we start with independent and identically distributed (say $[0, 1]$ uniform) random variables Z_v, $v \in V$ (i.e. the *IID process*), then we apply a measurable function $F: [0, 1]^V \to M^V$ that is required to be Γ-equivariant for some group Γ acting on V. The collection of random variables X_v, $v \in V$, we obtain this way is called a *factor-of-IID process*. Note that the joint distribution of X_v, $v \in V$, is Γ-invariant. In this paper the state space M will usually be finite.

An important special case, mainly because of its connection to random regular graphs and other large-girth regular graphs, is when V is the vertex set of the d-regular tree T_d equipped with the natural action of the automorpihsm group $\Gamma = \mathrm{Aut}(T_d)$. In this setting various entropy inequalities have been known for factor-of-IID processes. The simplest one is the following edge-vertex inequality:

$$H(X_u, X_v) \geq \frac{2(d - 1)}{d} H(X_v), \quad \text{where } (u, v) \text{ is an edge.} \tag{2}$$

This inequality played a central role in a couple of intriguing results recently [5,26]. All previous proofs of (2) are based on counting arguments for random regular graphs or for random permutations, and as such they heavily build on the acyclic nature of T_d, see [4,10,25]. Further inequalities and generalizations can be found in [3,16].

The new approach presented in this paper does not assume the acyclic property of the system (only certain local expansion). It provides a clean new proof for (2) and leads to numerous new versions and generalizations. Once again, we want to avoid too many technicalities at this point so we do not state our result in its most general form. Theorem 2 below is concerned with transitive graphs G satisfying some additional invariance called *unimodularity*.

There are several equivalent definitions for unimodular transitive graphs. The one most convenient for applications is the so-called *Mass-Transport Principle* saying that $\sum_{v \in V} h(o, v) = \sum_{v \in V} h(v, o)$ holds for any fixed vertex o of the vertex set V and for any function $h: V \times V \to [0, \infty)$ that is diagonally $\mathrm{Aut}(G)$-invariant (i.e. $h(u, v) = h(\varphi(u), \varphi(v))$ $\forall \varphi \in \mathrm{Aut}(G)$ $\forall u, v \in V$). For an equivalent definition that is easier to check, let $m_{u,v}$ denote the size of the orbit of v under the stabilizer of u, that is, the size of the set $\{\varphi(v) : \varphi \in \mathrm{Aut}(G); \varphi(u) = u\}$. Then a transitive G is unimodular if and only if $m_{u,v} = m_{v,u}$ for any pair of vertices.

Theorem 2. *Let G be a unimodular transitive graph with vertex set V and with degree d. Suppose that X_v, $v \in V$, is an $\mathrm{Aut}(G)$-factor-of-IID process with a finite state space. Then for an arbitrary $o \in V$ we have*

$$\frac{1}{2} \sum_{o' \sim o} H(X_o, X_{o'}) \geq \beta H(X_o), \quad \text{where } \beta = \frac{1}{2}\left(d + \inf_{W \subset V; |W| < \infty} \frac{|\partial W|}{|W|}\right).$$

Here $o' \sim o$ means that o' and o are connected by an edge in G, that is, o' runs through all neighbors of the fixed vertex o. Note that the infimum of $|\partial W|/|W|$ over finite sets W is usually called the edge Cheeger constant *of G.*

In particular, for $G = T_d$ the Cheeger constant is $d - 2$ and we get back the original edge-vertex inequality (2). See Theorem 3.3 for our general result, which is concerned with quasi-transitive group actions and joint entropies for arbitrary finite subsets.

1.4. *Proof method*

In the finite case our strategy is to reveal the independent seeds $Z_1, Z_2, \ldots, Z_m$ one by one and study how much information is gained about X_e and X_v at each step. After step i, the conditional entropy $H(X_e \mid Z_1, \ldots, Z_i)$ indicates the expected "amount of uncertainty" remaining. So the information gained at step i can be defined as the difference

$$H(X_e \mid Z_1, \ldots, Z_{i-1}) - H(X_e \mid Z_1, \ldots, Z_i).$$

It can be seen that this difference must be at least as much as the same difference for X_v provided that the vertex v is incident to e. This observation, combined with the fact that the information gain at step i is zero for vertices outside W_i, leads to Theorem 1.

As for the infinite setting, one needs to reveal the seeds in a random order and take expectation. This ensures that we have the required invariance to use the Mass-Transport Principle in the argument.

In fact, the approach described above goes beyond entropy. There are other ways of "measuring the uncertainty" in random variables. If such an uncertainty function satisfies a certain convexity criterion, then we can simply use this function – instead of entropy – in our arguments and obtain further inequalities of similar flavor. For example, we can get a bound for the correlation of neighbors in terms of the edge expansion. We will discuss this in detail in Section 4.

Telescoping sums of conditional entropies, in particular of those in random order, have been used in the literature before. The recent notion of *percolative entropy* is also based on this idea, see [2,27].

1.5. *Applications*

As we have mentioned, these inequalities can be used to get constraints for what can be achieved by factor-of-IID processes or randomized local algorithms. They can be applied to various problems; in this paper we will focus on questions regarding independent sets.

Upper bounds have been known for the density of factor-of-IID independent sets in Cayley graphs. An infinite version of the Hoffman bound provides a non-trivial bound for non-amenable graphs [22, Proposition 3.3]. The same paper explores possible improvements as well [22, Section 4]. Here, instead of using the spectral radius or other spectral parameters, we give a bound in terms of the Cheeger constant.

Theorem 3. *Suppose that G is a Cayley graph (or more generally, a unimodular transitive graph) with vertex set V and degree d. Let τ denote the normalized edge Cheeger constant*:

$$\tau = \inf_{W \subset V; |W| < \infty} \frac{|\partial W|}{d|W|}.$$

Then the density of any factor-of-IID independent set in G is at most $\varphi^{-1}(\tau)$, where φ^{-1} is the inverse of the following function:

$$\varphi(q) = \frac{-2q \log q - (1 - 2q) \log(1 - 2q)}{-q \log q - (1 - q) \log(1 - q)} - 1 \quad \text{for } q \in (0, 1/2).$$

It is easy to see that φ continuously and monotone decreasingly maps $(0, 1/2)$ onto $(0, 1)$, and hence the inverse exists.

It is, of course, very hard to assess the strength of this result without further details or examination. We will discuss the behaviour of φ^{-1} in Section 5. Note, however, that G is non-amenable if and only if $\tau > 0$, and the obtained bound $\varphi^{-1}(\tau)$ is strictly less than $1/2$ in that case. If G is the d-regular tree T_d, then we get back the original Bollobás bound: $(2 \log d)/d$ asymptotically as $d \to \infty$. The exact value of the Cheeger constant is also known for regular tessellations of the hyperbolic plane [17,19] so Theorem 3 can be applied to obtain good quantitative bounds for these lattices. Asymptotically we get the following.

Corollary 4. *Consider the regular tessellation of the hyperbolic plane by polygons with k sides and with d polygons at each vertex for some positive integers k and d satisfying $(k - 2)(d - 2) > 4$. Let $G_{d,k}$ denote the corresponding planar graph, that is, each face is a k-cycle and each degree is d. For any fixed k, the density of a factor-of-IID independent set in $G_{d,k}$ is at most*

$$\left(2 + \frac{2}{k-2} + o(1)\right) \frac{\log d}{d} \quad \text{as } d \to \infty.$$

To compare this to the spectral approach, we note that the Hoffman bound is, at best, of order $1/\sqrt{d}$ for any d-regular infinite graph, which is significantly weaker than the result above. In general, we cannot say that our bound (given by Theorem 3) is always better: the Cheeger inequalities relate the spectral radius to the Cheeger constant, and these inequalities do not exclude the possibility that the spectral approach gives slightly better results but this is probably not the typical case. Moreover, we rarely have exact values or even good estimates for the spectral radius of the adjacency operator. It tends to be a somewhat easier task to find the Cheeger constant.

Finally, we mention another application that shows that we can get results even in amenable graphs using local (or small set) expansion.

Theorem 5. *Let X be a factor-of-IID process with finite radius R over the integer lattice $\mathbb{Z}^n$ taking values 0 or 1. If X defines an independent set (i.e. the values of two neighboring vertices cannot be both 1), then the density of this independent set (i.e. the probability that a given vertex takes value 1) is bounded above by*

$$\frac{1}{2} - \frac{cn}{R \log(R/n)}$$

provided that $R > Cn$, where c and C are absolute constants.

Outline of the paper

In Section 2 we present and prove our entropy inequalities for a finite system of random variables, while in Section 3 we deal with the infinite case under the unimodularity condition. Section 4 explores possible generalizations, while Section 5 discusses applications. For the reader's convenience we also include an Appendix with some definitions and lemmas regarding entropy and conditional entropy.

2. Finite setting

2.1. *Expansion for hypergraphs*

Let V be a finite set and let G be a hypergraph over V, where the set of hyperedges is denoted by E. In other words, each hyperedge $e \in E$ is a subset of V. First we define a certain hyperedge expansion that will appear as a coefficient in our entropy inequality.

Definition 2.1. Given a set $W \subset V$, we define its *closure* $\overline{W}$ as the set of hyperedges incident to W:

$$\overline{W} := \{e \in E : e \cap W \neq \varnothing\}.$$

Then the *hyperedge expansion* of G with respect to the sets $W_1, \ldots, W_m \subset V$ is defined as

$$\beta := \min_{1 \leq i \leq m; W \subseteq W_i} \frac{|\overline{W}|}{|W|}.$$

More generally, we will consider hypergraphs with a nonnegative edge weight assigned to each $e \in E$, in which case we will need to replace $|\overline{W}|$ with the sum of weights in $\overline{W}$.

Remark 2.2. Note that the most standard edge expansion definition for graphs uses the *edge boundary* ∂W consisting of edges with one endpoint in W and one endpoint in $V \setminus W$. If G is a d-regular graph, then the two notions are essentially the same:

$$d|W| = 2(\#\text{edges inside } W) + |\partial W| \quad \text{and} \quad |\overline{W}| = (\#\text{edges inside } W) + |\partial W|,$$

and hence

$$\beta = \min_{1 \leq i \leq m; W \subseteq W_i} \frac{d|W| + |\partial W|}{2|W|} = \frac{1}{2}\left(d + \min_{1 \leq i \leq m; W \subseteq W_i} \frac{|\partial W|}{|W|}\right).$$

2.2. *An entropy inequality for finitely many random variables*

Let V be a finite set and X_v, $v \in V$, be random variables, each X_v taking finitely many values. For any subset $U \subseteq V$ one can consider the joint entropy of the random variables X_v, $v \in U$. This is simply defined as the Shannon entropy of the joint distribution of $X_U := (X_v)_{v \in U}$ and will be denoted by $H(X_U)$. Usually we will think of such a subset U as a hyperedge and use the notation $e \subseteq V$.

Throughout this section we will work under the following assumption.

Assumption 2.3. Let $Z_1, \ldots, Z_m$ be independent random variables. For each Z_i a subset $W_i \subseteq V$ is given, and for every $v \in V$ we assume that X_v is a measurable function of Z_i, $v \in W_i$.

An equivalent formulation would be the following. Let $\mathcal{W}$ denote the set of all subsets of all W_i. Then given countably many independent random bits, each bit is sent to X_v, $v \in W$, for some $W \in \mathcal{W}$, and every X_v is a measurable function of the bits received.

For random variables X_v obtained as above, the next theorem bounds any given sum of joint entropies in terms of the sum of the individual entropies. Note that Shearer's inequality (1) provides a lower bound for the same sum in terms of the joint entropy of all X_v.

Theorem 2.4 (Hypergraph version). *Under Assumption* 2.3 *the following inequality is true for any set* E *of hyperedges over* V:

$$\sum_{e \in E} H(X_e) \geq \beta \sum_{v \in V} H(X_v), \quad \text{where } \beta = \min_{1 \leq i \leq m; W \subseteq W_i} \frac{|\overline{W}|}{|W|}.$$

Recall that $\overline{W}$ *is defined as* $\{e \in E : e \cap W \neq \varnothing\}$.

When our hypergraph is a d-regular graph, then the coefficient β can be expressed using the standard edge expansion with respect to subsets of W_i (see Remark 2.2) so in that case we get back Theorem 1 stated in the Introduction.

Also note that in the d-regular case $\beta = d/2$ would be a trivial lower bound following from the fact that $H(X_u, X_v) \geq (H(X_u) + H(X_v))/2$. Similarly, a trivial upper bound ($\beta = d$) can be deduced using $H(X_u, X_v) \leq H(X_u) + H(X_v)$.

Finally, we state and prove the most general version of our inequality where nonnegative weights are assigned to the hyperedges. In fact, we will assign a weight to each subset $U \subseteq V$, thinking of zero weights as non-edges.

Theorem 2.5 (Weighted hypergraph version). *Suppose that to each subset* $U \subseteq V$ *a nonnegative weight* $\alpha_U \in [0, \infty)$ *is assigned. Then, under Assumption* 2.3, *the following inequality holds*:

$$\sum_{U \subseteq V} \alpha_U H(X_U) \geq \beta \sum_{v \in V} H(X_v), \quad \text{where } \beta = \min_{1 \leq i \leq m; W \subseteq W_i} \frac{\sum_{U \subseteq V; U \cap W \neq \varnothing} \alpha_U}{|W|}. \tag{3}$$

Moreover, the above inequality is sharp: β *is the largest coefficient for which* (3) *holds for any random variables satisfying Assumption* 2.3.

Proof. The idea is to reveal $Z_1, \ldots, Z_m$ one by one and look at how much information is gained on average about X_U at each step. To be more precise, for any i and U the following conditional entropy can be defined:

$$H(X_U \mid Z_1, \ldots, Z_i).$$

(See the Appendix for basic definitions regarding Shannon entropy and for how the above conditional entropy can be defined if Z_i is not discrete.) For $i = 0$ this is simply $H(X_U)$, while for $i = m$ we get 0 because $Z_1, \ldots, Z_m$ fully determine X_U. Then the average information gained at step i can be defined as:

$$h(i, U) := H(X_U \mid Z_1, \ldots, Z_{i-1}) - H(X_U \mid Z_1, \ldots, Z_i).$$

In the case of one-element sets $U = \{v\}$ we will simply write $h(i, v)$.

If we add up all $h(i, U)$ for a given U, then we get a telescoping sum and obtain that

$$H(X_U) = \sum_{i=1}^{m} h(i, U).$$

We will also need the following observation: whenever $U' \subset U$ we have $h(i, U') \leq h(i, U)$ for each i. This is clear heuristically since for larger U we should gain more information about X_U at each step. For a rigorous proof see Lemma A.1 in the Appendix. In particular, it holds for any given i and U that

$$h(i, U) \geq \max_{v \in U} h(i, v).$$

It follows that

$$\sum_{U \subseteq V} \alpha_U H(X_U) \geq \sum_{i=1}^{m} \sum_{U \subseteq V} \alpha_U \max_{v \in U} h(i, v).$$

As for the right-hand side of the inequality (3), we have

$$\beta \sum_{v \in V} H(X_v) = \beta \sum_{i=1}^{m} \sum_{v \in V} h(i, v).$$

Therefore it suffices to show that it holds for each i that

$$\sum_{U \subseteq V} \alpha_U \max_{v \in U} h(i, v) \geq \beta \sum_{v \in V} h(i, v).$$

Note that, for fixed i, $h(i, v)$ is a $V \to [0, \infty)$ function that is supported on W_i. Hence the following claim completes the proof.

Claim. *For any function $f : V \to [0, \infty)$ it holds that*

$$\frac{\sum_{U \subseteq V} \alpha_U \max_{v \in U} f(v)}{\sum_{v \in V} f(v)} \geq \min_{W \subseteq \operatorname{supp} f} \frac{\sum_{U \subseteq V; U \cap W \neq \varnothing} \alpha_U}{|W|}.$$

In other words, the fraction on the left-hand side is minimized by an indicator function $\mathbb{1}_W$.

To prove this claim we take the following unique decomposition of f into the sum of indicator functions: there exist sets $\operatorname{supp} f = A_1 \supsetneq A_2 \supsetneq \cdots \supsetneq A_k$ and positive real numbers $c_1, \ldots, c_k > 0$ such that

$$f = \sum_{j=1}^{k} c_j \mathbb{1}_{A_j}.$$

For such a decomposition one can always interchange taking maximum over a set U with taking sum:

$$\max_{v \in U} f(v) = \sum_{j=1}^{k} c_j \max_{U} \mathbb{1}_{A_j} = \sum_{j=1}^{k} c_j \mathbb{1}_{U \cap A_j \neq \varnothing}.$$

Now considering the fraction in the claim, the total contribution of the jth term to the numerator and denominator is $c_j \sum_{U \subseteq V; U \cap A_j \neq \varnothing} \alpha_U$ and $c_j |A_j|$, respectively. The claim follows by using that $\sum a_j / \sum b_j \geq \min_j a_j / b_j$ for any $a_j \geq 0, b_j > 0$.

There is a slightly different way to finish the proof after having defined the values $h(i, U)$ for a given process X. We include this ending, too, partly because it is a nice argument and partly because it easily implies that our inequality is sharp. For the sake of simplicity we first present the proof under the additional assumption that $h(i, v)$ is an integer multiple of $\log 2$ for each i and v. First we define a new set of random variables: let $\tilde{Z}_i$ consist of a finite but sufficiently large number of independent "bits" (taking values 0 and 1 with probability $1/2$ each), and let $\tilde{X}_v$ be the collection of the first $h(i, v)/\log 2$ bits of $\tilde{Z}_i$ for each $i = 1, \ldots, m$. It is clear that

$$H(\tilde{X}_v) = \sum_{i=1}^{m} h(i, v) = H(X_v)$$

and

$$H(\tilde{X}_U) = \sum_{i=1}^{m} \max_{v \in U} h(i, v) \leq \sum_{i=1}^{m} h(i, U) = H(X_U).$$

It follows that it suffices to prove (3) for this new set of random variables $\tilde{X}_v$. Now let us consider an arbitrary bit of $\tilde{Z}_i$. This bit is present in the random variables $\tilde{X}_v$, $v \in W$, for some subset $W \subseteq W_i$. Therefore its entropy ($\log 2$) will appear $\sum_{U \subseteq V; U \cap W \neq \varnothing} \alpha_U$ times on the left-hand side and $\beta |W|$ times on the right-hand side of (3). The former will be clearly larger than or equal to the latter provided that we choose β as in the theorem.

If the values $h(i, v)/\log 2$ are not necessarily integers, then we need to slightly modify the above argument. Let $\tilde{Z}_i$ be a sequence of independent discrete random variables with prescribed entropies. We can clearly choose these entropies in such a way that for any $v \in V$ the value $h(i, v)$ can be obtained as the sum of the first $k(i, v)$ of these entropies for some nonnegative integer $k(i, v)$. Then $\tilde{X}_v$ can be defined as the collection of the first $k(i, v)$ elements of $\tilde{Z}_i$, $i = 1, \ldots, m$. The rest of the proof remains essentially the same.

To see sharpness, take a subset $W \subseteq W_i$ with $(\sum_{U \subseteq V; U \cap W \neq \varnothing} \alpha_U)/|W| = \beta$, that is, W has "minimal expansion". Let X_v be the same measurable function of Z_i for all $v \in V$, and let X_v be almost surely constant for $v \notin W$. It is easy to see that (3) holds with inequality in this case. Usually there are many W with minimal expansion in which case one can consider combinations of these examples for different W. (Much as in the construction of $\tilde{X}_v$ above.) $\qquad \square$

2.3. *The general problem*

All the inequalities proved in this section, and Shearer's inequality as well, fit into the following general problem. Two weighted sums of joint entropies of finitely many random variables are given along with sets W_i describing which variables can use the same random seed. The task is to find the strongest inequality (the one with the best coefficient) between the two sums. Solving the problem in full generality is probably too much to ask for, but there might exist a common generalization of (3) and Shearer's inequality, which would already be very interesting.

In fact, it is not crucial for our approach to have individual entropies on the right-hand side of the inequality. The key observation in our proof was that the information gain is always larger for larger sets, that is, $h(i, U') \leq h(i, U)$ provided that $U' \subset U$. Note that we only used this fact in the case when U' was a one-element set $\{v\}$. If we make full use of this fact, more complicated inequalities can be obtained. For example, one could deduce an inequality between the entropies corresponding to the stars in a graph and the entropies corresponding to the edges. By a star we mean a vertex v and its neighbors $N(v)$. Given a simple graph $G = (V, E)$, a straightforward modification of the proof above yields the following star-edge entropy inequality:

$$\sum_{v \in V} H(X_{\{v\} \cup N(v)}) \geq \beta \sum_{e \in E} H(X_e), \quad \text{where } \beta = \min_{1 \leq i \leq m; F \subseteq \overline{W_i}} \frac{|\{v : \exists e \in F \text{ s.t. } v \in e\}|}{|F|}. \tag{4}$$

An infinite version of this inequality for factor-of-IID processes on T_d says that the entropy of a star is at least $d/2$ times the entropy of an edge [4,5].

3. Unimodular setting

In this section we investigate what can be said for infinitely many random variables. It turns out that our approach can be adapted provided that X_v, $v \in V$, are obtained from the seeds in a certain invariant way. We start with the simplest setup.

3.1. *Transitive graphs*

Let G be a connected infinite graph that is transitive meaning that its automorphism group $\mathrm{Aut}(G)$ acts transitively on its vertex set V. First we recall the definition of factor-of-IID processes on G. We start with an IID process over V: this is simply a collection of independent random variables Z_v, $v \in V$, each uniformly distributed on $[0, 1]$. Then we apply a measurable $\mathrm{Aut}(G)$-equivariant $[0, 1]^V \to M^V$ mapping. The result is a collection of random variables X_v, $v \in V$, each taking values in M. A process that can be obtained this way is called a factor-of-IID process on G. Note that in this paper the state space M is usually assumed to be finite.

Such a process is said to be a *block factor* if it can be obtained in a way that each (or, equivalently, any given) X_v is influenced only by finitely many Z_u's. In other words, there exists a finite radius R such that X_v is a function of Z_u, $u \in B_R(v)$, where $B_R(v)$ denotes the ball of radius R around v w.r.t. the graph distance in G. Block factors are also called *finite-radius factors*. For a more detailed introduction to factors of IID, see [3, Section 2].

We will also need to assume *unimodularity* for G, which is equivalent to the *Mass-Transport Principle*, see Section 1.3 for definitions and [1] for details.

Theorem 3.1. *Let G be a unimodular transitive graph with vertex set V and with degree d. Suppose that X_v, $v \in V$, is a block factor-of-IID process with radius R and with a finite state space. For an arbitrary fixed vertex $o \in V$, let $B_R = B_R(o)$ denote the R-ball around o. Then we have*

$$\frac{1}{2} \sum_{o' \sim o} H(X_o, X_{o'}) \geq \beta H(X_o), \quad \text{where } \beta = \frac{1}{2} \left(d + \min_{W \subseteq B_R} \frac{|\partial W|}{|W|} \right).$$

Theorem 2 stated in the Introduction follows immediately considering the fact that any factor-of-IID process is the weak limit of block factors. A short argument proving this fact for transitive graphs can be found in [18, Proof of Proposition 4.4].

Proof of Theorem 3.1. Let Z_u, $u \in V$, denote an IID process on the vertex set V of G. Suppose that the process X_v, $v \in V$, is obtained as an $\mathrm{Aut}(G)$-equivariant function of the IID process in such a way that each X_v is the function of Z_u, $u \in B_R(v)$.

Definition 3.2. A function $f \colon V \times V \to [0, \infty]$ is said to be diagonally Γ-invariant if

$$f(\gamma u, \gamma v) = f(u, v) \quad \forall u, v \in V \ \forall \gamma \in \Gamma.$$

We say that the *Mass-Transport Principle* holds for a Γ-action with respect to a measure μ on the orbit classes if for all diagonally Γ-invariant function f it holds that

$$\mathbb{E}_{\mu(o)} \sum_{v \in V} f(o, v) = \mathbb{E}_{\mu(o)} \sum_{v \in V} f(v, o).$$

This is actually equivalent to the following statement on the orbit sizes of the stabilizers:

$$\frac{|\operatorname{Stab}(o)o'|}{|\operatorname{Stab}(o')o|} = \frac{\mu(o')}{\mu(o)} \quad \text{for any pair } o, o'.$$

Next we introduce the weight function $\alpha \colon \mathcal{P}(V) \to [0, \infty)$ that assigns nonnegative real numbers to subsets of V. We require α to satisfy the following properties:

- α is Γ-invariant: $\alpha_U = \alpha_{\gamma U}$ for any $U \subset V$ and $\gamma \in \Gamma$;
- α is supported on finite sets: $\alpha_U = 0$ if U is not finite;
- α is "locally finitely supported": for any $v \in V$ there are finitely many U such that $v \in U$ and $\alpha_U = 0$.

Now we are in a position to state the most general form of our theorem in the infinite setting. We started with the transitive graph version (Theorem 3.1) because both the statement and the proof are much easier to digest in that special case. The proof itself is a straightforward modification of that of Theorem 3.1, and we leave the details to the reader.

Theorem 3.3. *Let V, Γ, μ, α be as above, and let Z_u, $u \in V$, be an IID process on G. Suppose that X_v, $v \in V$, is a Γ-factor of Z with finite state space such that for any vertex u the value of Z_u influences only the labels X_v, $v \in W_u$, for some finite set $W_u \subset V$. Then*

$$\mathbb{E}_{\mu(o)} \sum_{U : o \in U} \frac{\alpha_U}{|U|} H(X_U) \geq \beta \mathbb{E}_{\mu(o)} H(X_o), \quad \text{where } \beta = \min_{u \in V; W \subseteq W_u} \frac{\sum_{U \subseteq V; U \cap W \neq \varnothing} \alpha_U}{|W|}.$$

3.3. *Corollaries*

We mention a couple of consequences of Theorem 3.3 in this section.

The first special case concerns factors of IID over Cayley graphs of finitely-generated groups.

Theorem 3.4. *Let Γ be an infinite group with a finite symmetric generating set S. We denote the unit element of Γ by e. Suppose that X_γ, $\gamma \in \Gamma$, is a Γ-factor-of-IID process over (the Cayley diagram of) Γ, where X_γ takes values from a finite set. Given arbitrary nonnegative real numbers $\alpha_\sigma \geq 0$, $\sigma \in S$, with $\alpha_\sigma = \alpha_{\sigma^{-1}}$ we have the following inequality:*

$$\sum_{\sigma \in S} \alpha_\sigma H(X_e, X_\sigma) \geq \beta H(X_e), \quad \text{where } \beta = \inf_{W \subset \Gamma, |W| < \infty} \frac{\sum_{\sigma \in S} \alpha_\sigma |W \cup \sigma W|}{|W|}.$$

Remark 3.5. When Γ is a free group and the weights are all equal, we get back an inequality that was known by Bowen [10] as the fact the so-called *f-invariant* is nonnegative for factors of the Bernoulli shift. See also [16, Theorem 2.3], where this inequality is explicitly stated.

To see that Theorem 3.4 is indeed a special case of Theorem 3.3, we set $V = \Gamma$ and consider the natural left action of Γ on itself. It is a well-known fact that this is a unimodular transitive action. Using α_σ we define our weight function $\alpha \colon \mathcal{P}(V) \to [0, \infty)$ as follows: for two-element sets of the form $U = \{\gamma, \gamma\sigma\}$, $\gamma \in \Gamma$, $\sigma \in S$, we assign $\alpha_U := \alpha_\sigma$, and we set $\alpha_U = 0$ for any other set U. This yields Theorem 3.4 for block factors. By taking weak limit, we get the inequality for any factor-of-IID process.

The other special case we briefly mention concerns factor-of-IID processes on T_d. The edge-vertex inequality (2) compares the entropy of a vertex to the joint entropy of two neighboring vertices. One could compare the vertex entropy to the entropy $H(X_{U_0})$ for any finite set $U_0 \subset V(T_d)$, and ask for the largest coefficient $c = c(U_0)$ such that

$$H(X_{U_0}) \geq c H(X_v)$$

holds for any factor-of-IID process X. Setting $\Gamma = \mathrm{Aut}(T_d)$, $V = V(T_d)$, and $\alpha(U) = 1$ whenever U can be mapped to U_0 via an automorphism of T_d, we can apply Theorem 3.3 to get the above inequality with some coefficient c. It is actually not straightforward to see that the obtained coefficient is optimal. We omit the details here but it can be shown that the coefficient Theorem 3.3 provides is actually equal to

$$c = c(U_0) = \inf_R \frac{|B_R(U_0)|}{|B_R(v)|}, \quad \text{where } B_R(U_0) \text{ is the } R\text{-neighborhood of } U_0.$$

It is easy to construct factor-of-IID processes for which the ratio $H(X_{U_0})/H(X_v)$ tends to the above coefficient, see the example described in [16, Section 5.1], showing that this coefficient is indeed the best possible.

4. Beyond entropy: Possible generalizations

The goal of this section is to explore what properties of entropy were crucial in our arguments to see if there are any other quantities (assigned to random variables) that our proofs would work for.

4.1. *Uncertainty functions*

Let us consider the space $\mathcal{M}$ of probability measures on some measurable space $(A, \mathcal{A})$. This space is clearly closed under convex combination. Therefore we can talk about the convexity and concavity of $\mathcal{M} \to \mathbb{R}$ functions.

Definition 4.1. We say that a function $\delta \colon \mathcal{M} \to \mathbb{R}$ is an *uncertainty function* if δ is concave and $\delta(\nu) = 0$ holds for any measure ν concentrated on one point (i.e. unit mass).

The concavity of δ essentially means that the δ-value of a mixture distribution must be greater than or equal to the expectation of the δ-values.

To keep things simple, here we restrict our attention to the finite simple graph case: let G be a finite graph with vertex set V and edge set E. Suppose that X_v, $v \in V$, are $(A, \mathcal{A})$-valued random variables satisfying Assumption 2.3. This time, however, X_v does not need to be discrete. Let δ be an uncertainty function for probability measures over A as described above. Now we can talk about the uncertainty corresponding to a vertex: the δ-value of (the distribution of) the random variable X_v. To be able to talk about the uncertainty corresponding to an edge $e = (u, v)$ as well, we need to choose an uncertainty function Δ for probability measures over $A \times A$. Of course, δ and Δ should be compatible in some sense. It turns out that the only assumption we need is the following.

Assumption 4.2. Let δ and Δ be uncertainty functions over A and $A \times A$, respectively. Furthermore, for a measure μ on $A \times A$, let $\pi_1(\mu)$ and $\pi_2(\mu)$ denote the marginals of μ. We assume that

$$\Delta - \delta \circ \pi_1 \quad \text{and} \quad \Delta - \delta \circ \pi_2 \quad \text{are both concave.} \tag{8}$$

Under this assumption, essentially the same proof works as in Theorem 2.5. For the sake of simplicity, let us assume that the random seeds $Z_1, \ldots, Z_m$ are discrete. Then for any event $Z_1 = z_1, \ldots, Z_i = z_i$ with positive probability, we can consider the distribution of X_v conditioned on this event and take the δ-value of this conditional distribution. By taking expectation in $z_1, \ldots, z_i$ we get the expected uncertainty of the vertex v after revealing the first i seeds. We denote this by $\delta(i, v)$. Since δ is assumed to be concave, it is easy to see that $\delta(i, v)$ is monotone decreasing in i, and hence for any fixed i, $h(i, v) := \delta(i - 1, v) - \delta(i, v)$ is a nonnegative $V \to [0, \infty)$ function (supported on W_i).

Similarly, for an edge $e = (u, v)$, we define $\Delta(i, e)$ by the expectation of the Δ-value of the joint distribution of (X_u, X_v) conditioned on the first i seeds. Then we define $h(i, e)$ as $\Delta(i - 1, e) - \Delta(i, e)$. The crucial fact our proof relied on was that

$$h(i, e) \geq \max\big(h(i, u); h(i, v)\big),$$

which easily follows from (8). The rest of the proof remains the same yielding the following result.

Theorem 4.3. *Suppose that $G = (V, E)$ is a finite graph, and X_v, $v \in V$, are $(A, \mathcal{A})$-valued random variables satisfying Assumption 2.3. Furthermore, let δ and Δ be uncertainty functions as in Assumption 4.2. Then*

$$\sum_{e \in E} \Delta(X_e) \geq \beta \sum_{v \in V} \delta(X_v), \quad \text{where } \beta = \min_{1 \leq i \leq m; \, W \subseteq W_i} \frac{|\overline{W}|}{|W|}.$$

Recall that $\overline{W}$ is defined as $\{e \in E : e \cap W \neq \varnothing\}$.

As for the infinite setting, Theorem 2 can be generalized similarly.

When A is finite, we can use entropy for both δ and Δ, and we get back our original inequality. Note that conditional entropy can be obtained as the expectation of entropies of conditional distributions; see the Appendix. Therefore $\delta(i, v)$ and $\Delta(i, e)$ in this case are the same conditional entropies that came up in the original proof.

4.2. *Correlation bound*

Another natural way to measure uncertainty of a distribution is variance. Next we study what Theorem 4.3 gives if we set $\delta(v)$ to be the variance of v.

Suppose that $A = \mathbb{R}$ equipped with the Borel σ-algebra. We will work under the assumption that the random variables X_v, $v \in V$, have finite variance. For a real-valued random variable v of finite variance, $\delta(v)$ simply denotes the variance of v:

$$\delta(v) = \int_{\mathbb{R}} x^2 \, dv(x) - \left(\int_{\mathbb{R}} x \, dv(x) \right)^2.$$

It is easy to see that this δ is indeed concave. (Note that the first term is linear in v and hence does not influence convexity/concavity.)

A natural candidate for Δ, the uncertainty function over $\mathbb{R} \times \mathbb{R}$, would be the covariance of the two coordinates. Some constant multiples of the variances of the marginals should be added to the covariance so that (8) holds true. More precisely, let μ be a probability measure on $\mathbb{R} \times \mathbb{R}$ with marginals $\mu_1 = \pi_1(\mu)$ and $\mu_2 = \pi_2(\mu)$. We want to find an uncertainty function of the following form:

$$\Delta(\mu) = a\big(\delta(\mu_1) + \delta(\mu_2)\big) + b\left(\int_{\mathbb{R} \times \mathbb{R}} xy \, d\mu(x, y) - \left(\int_{\mathbb{R}} x \, d\mu_1(x) \right)\left(\int_{\mathbb{R}} y \, d\mu_2(y) \right) \right).$$

It is easy to see that (8) is satisfied if and only if

$$a > 1 \quad \text{and} \quad |b| \leq 2\sqrt{a(a-1)}.$$

Setting $b = \pm 2\sqrt{a(a-1)}$, we get the following from Theorem 4.3 for d-regular graphs:

$$\pm 2\sqrt{a(a-1)} \sum_{(u,v) \in E} \mathrm{cov}(X_u, X_v) \leq d(a - \beta/d) \sum_{v \in V} \mathrm{var}(X_v).$$

Easy calculations show that the best inequality is obtained by setting $a = \frac{\beta/d}{2\beta/d - 1}$, and we get the following result.

Theorem 4.4. *Let G be a d-regular graph with vertex set V and edge set E. The random variables X_v, $v \in V$, and the coefficient β are as in Theorem 1. If $\mathrm{var}(X_v) < \infty$ for each $v \in V$, then*

$$\left| \sum_{(u,v) \in E} \mathrm{cov}(X_u, X_v) \right| \leq \sqrt{\beta(d - \beta)} \sum_{v \in V} \mathrm{var}(X_v).$$

This means that if each X_v has the same variance, then the average correlation is at most $2\sqrt{\beta(d - \beta)}/d$. Note that the same bound could be deduced for factor-of-IID processes over unimodular transitive graphs with $\beta = (d + \psi)/2$ where ψ is the edge Cheeger constant. In this infinite transitive graph setting, however, there is another way to prove this bound by combining the following known results.

- The average correlation is bounded by the spectral radius of the adjacency operator divided by d. This follows from the results of [6] on the covariance structure of factor-of-IID processes.
- By a Cheeger-tpye inequality [23, Theorem 6.7], the spectral radius is bounded by

$$\sqrt{d^2 - \psi^2} = \sqrt{(d + \psi)(d - \psi)} = 2\sqrt{\beta(d - \beta)}.$$

4.3. *Specific problems*

It is an intriguing question whether there are other general uncertainty functions that satisfy Assumption 4.2 and lead to useful inequalities. Instead of looking for further general inequalities, another approach is that, given a specific application, we try to find the best functions δ and Δ to use. In fact, there is hope that this approach might provide the optimal bound for factor-of-IID independent sets.

In the setting of independent sets, each vertex can have two states $(0, 1)$ and vertices of state 1 cannot be neighbors. So the state space is $A = \{0, 1\}$, and δ can be described by a concave function $f : [0, 1] \to \mathbb{R}$ with $f(0) = f(1) = 0$. Furthermore, the measures μ in question are concentrated on $A \times A \setminus \{(1, 1)\} = \{(0, 0); (0, 1); (1, 0)\}$. Such a measure μ can be represented by a point (x, y) of the triangle $T = \{(x, y) : x, y \geq 0; x + y \leq 1\} \subset \mathbb{R}^2$, where x, y, and $1 - x - y$ are the probabilities of $(1, 0)$, $(0, 1)$, and $(0, 0)$, respectively. Then Δ corresponds to a function $F : T \to \mathbb{R}$ with $F(0, 0) = F(0, 1) = F(1, 0) = 0$ such that the functions $F(x, y) - f(x)$ and $F(x, y) - f(y)$ are both concave on T. For d-regular graphs, if each vertex is included in the independent set with probability q, then we get that $F(q, q) \geq (2\beta/d) f(q)$. In particular, we get the following for factor-of-IID independent sets in T_d, in which setting we have $\beta = d - 1$.

Lemma 4.5. *Let q be the (maximum) density of a factor-of-IID independent set over the d-regular tree T_d. By T we denote the triangle $\{(x, y) : x, y \geq 0; x + y \leq 1\} \subset \mathbb{R}^2$. Suppose that $f : [0, 1] \to \mathbb{R}$ is a concave function with $f(0) = f(1) = 0$ and $F : T \to \mathbb{R}$ is such that $F(0, 0) = F(0, 1) = F(1, 0) = 0$ and $F(x, y) - f(x)$ and $F(x, y) - f(y)$ are both concave on T. Then the following holds true:*

$$F(q, q) \geq \frac{2(d - 1)}{d} f(q).$$

This gives the following upper bound for the maximum density q:

$$\inf\{q : F(q, q) < (2(d - 1)/d) f(q)\}.$$

So we need to find functions f, F such that the above infimum is as small as possible. Solving this convex optimization problem, even numerically for a given d, would be very interesting. The question to determine the independence ratio of a random d-regular graph and the closely related problem of finding the maximum density of a factor-of-IID independent set on T_d have been thoroughly studied; see Section 5 for details. Note that if we use entropy for δ and Δ, then we get the well-known Bollobás bound. In recent years there has been a lot of activity, especially regarding the case $d = 3$: see [12,13,20] for lower bounds and [7,21,24] for upper bounds. Finding the optimal f and F may lead to a breakthrough in this problem. The above optimization problem can also be written as an LP problem in terms of the second derivatives of f and F. Its dual solution provides necessary and sufficient conditions for a factor-of-IID independent set to reach the obtained bound. Therefore the bound must be optimal provided that these conditions can be satisfied with an arbitrarily small error.

5. Applications

Entropy inequalities can be used to acquire bounds for the size of different combinatorial structures (such as minimum or maximum cuts) on graphs. Bounds for independent sets have the most extensive literature. In this section we demonstrate how the new inequalities presented in this paper lead to results concerning random independent sets.

One of the first results in this direction is due to Bollobás who gave an upper bound for the independence ratio of random d-regular graphs [8]. The counting argument behind this result essentially proves the edge-vertex entropy inequality in this special setting. The upper bound obtained is asymptotically sharp: $2 \log(d)/d$ as $d \to \infty$. The lower bound is due to Frieze and Łuczak [14]. Later Gamarnik and Sudan showed [15] that so-called local algorithms cannot produce independent sets of this optimal size on large-girth graphs. Then in [26] the edge-vertex entropy inequality was used in a non-trivial way to prove that independent sets provided by local algorithms can have density at most $\log(d)/d$ (asymptotically as $d \to \infty$). This is equivalent to saying that factor-of-IID independent sets on T_d have density at most $(1 + o(1)) \log(d)/d$.

5.1. *Random independent sets and fractional colorings*

The graph parameter *fractional chromatic number* can be defined as the inverse of the maximum "homogeneous density" of a random independent set. More precisely, let G be a finite graph and let q be the largest real number with the property

that there exists a random independent set I on G such that any given vertex lies in I with probability at least q. Then q coincides with the inverse $1/\chi_f(G)$ of the fractional chromatic number.

It is natural to consider random independent sets that are obtained by local algorithms; one could define the corresponding *local fractional chromatic number*. We will treat the finite and infinite settings simultaneously.

Infinite setting: given a unimodular transitive graph and a finite radius R, we consider radius-R factor-of-IID processes X such that X_v takes values 0, 1 and $\{v : X_v = 1\}$ is an independent set, that is, $X_u = X_v = 1$ cannot happen for neighboring vertices u, v. The probability q that $X_v = 1$ is the same for each v due to invariance.

Finite setting: given a d-regular finite graph G and subsets $W_1, \ldots, W_m$ of the vertex set V, we consider random variables X_v, $v \in V$, as in Assumption 2.3 such that $\{v : X_v = 1\}$ is an independent set. Let q be such that $X_v = 1$ with probability at least q for each vertex $v \in V$. One can easily modify such an algorithm in a way that each of these probabilities is actually equal to q.

5.2. *Bounds via entropy inequalities*

Our goal is to find an upper bound for q in terms of the edge expansion of G. For each vertex v we clearly have

$$H(X_v) = -q \log q - (1-q) \log(1-q).$$

Moreover, for any pair of neighboring vertices u, v:

$$H(X_u, X_v) = -2q \log q - (1-2q) \log(1-2q).$$

Let

$$\varphi(q) = \frac{H(X_u, X_v)}{H(X_v)} - 1 = \frac{-2q \log q - (1-2q) \log(1-2q)}{-q \log q - (1-q) \log(1-q)} - 1.$$

It is easy to see that this function is continuous and monotone decreasing on $(0, 1/2)$ with $\lim_{q \to 0+} \varphi(q) = 1$ and $\lim_{q \to 1/2-} \varphi(q) = 0$. Straightforward calculations show that the behavior of φ around 0 and $1/2$ is as follows:

$$\varphi(q) = 1 + \frac{q}{\log q} + O\left(\frac{q}{(\log q)^2}\right);$$

$$\varphi(1/2 - t) = \frac{-2t \log t}{\log 2} + O(t).$$

Therefore one can define a monotone decreasing inverse function $\varphi^{-1} : [0, 1] \to [0, 1/2]$ such that asymptotically at 0 and 1 we have

$$\varphi^{-1}(x) - \frac{1}{2} \sim \frac{\log 2}{2} \frac{x}{\log x} \quad \text{as } x \to 0;$$

$$\varphi^{-1}(1-x) \sim -x \log x \quad \text{as } x \to 0.$$

For an explicit upper bound near 0 and 1, one can use the following estimates for φ^{-1}:

$$\varphi^{-1}(x) \leq \begin{cases} \frac{1}{2} - \frac{\log 2}{2} \frac{x}{-\log x + 2 \log(-\log x)} & \text{if } 0 < x < 1/30, \\ -(1-x) \log(1-x) & \text{if } 2/3 < x < 1. \end{cases}$$

Now let τ denote the edge expansion (w.r.t. the R-ball or $W_1, \ldots, W_m$) normalized by d:

$$\tau = \min_{W \subseteq B_R} \frac{|\partial W|}{d|W|} \quad \text{and} \quad \tau = \min_{1 \leq i \leq m; W \subseteq W_i} \frac{|\partial W|}{d|W|}$$

in the infinite and finite setting, respectively. Then our entropy inequalities (Theorem 1 and Theorem 3.1) say that

$$\frac{H(X_u, X_v)}{H(X_v)} - 1 \geq \tau.$$

It follows that $\varphi(q) \geq \tau$ must hold. Equivalently, we get the upper bound

$$q \leq \varphi^{-1}(\tau).$$

In the infinite setting we get Theorem 3 stated in the Introduction. Corollary 4 is concerned with the special case when our graph corresponds to a regular tessellation of the hyperbolic plane. The Cheeger constant of the graph $G_{d,k}$ is known [17,19] to be

$$(d-2)\sqrt{1 - \frac{4}{(k-2)(d-2)}}.$$

If k is fixed and $d \to \infty$, then we have $\tau = 1 - (2 + 2/(k-2))/d + o(1)$. Using the asymptotics of φ^{-1} at $\tau = 1$, Corollary 4 follows.

Furthermore, we can get something non-trivial even over amenable graphs. Let G be the standard Cayley graph of $\mathbb{Z}^n$ which is a d-regular transitive graph for $d = 2n$. Although G is amenable, the normalized edge expansion τ with respect to a finite-radius ball B_R is at least cnR^{-1} for some positive constant c, and hence Theorem 5 follows. (Note that one can easily prove – simply by using that labels of vertices at distance $2R + 1$ are independent – the upper bound $1/2 - c/R$ for any n. Therefore the bound given in Theorem 5 is interesting when R is sub-exponential in n. Also note that it is easy to construct a factor-of-IID independent set on $\mathbb{Z}^n$ with radius R and density $1/2 - cn/R$.)

Lastly, we mention one immediate corollary of the finite version.

Theorem 5.1. *Let G be a random d-regular bipartite graph on $2N$ vertices. Then for any $\delta > 0$ there exists a positive ε such that the following holds true for G with high probability as N goes to infinity: a randomized algorithm on G for which each seed influences at most $1 - \delta$ fraction of the vertices cannot produce an independent set of homogeneous density at least $1/2 - \varepsilon$ (that is, at least one vertex will be included in the independent set with probability less than $1/2 - \varepsilon$).*

Appendix

Entropy and conditional entropy

Let X be a discrete random variable taking m distinct values with probabilities $p_1, \ldots, p_m$. Then the *Shannon entropy* of X is defined as

$$H(X) := \sum_{i=1}^{m} -p_i \log(p_i).$$

The *joint entropy* of finitely many random variables is the entropy of their joint distribution.

One can define the *conditional entropy* of X conditioned on Y by $H(X|Y) = H(X, Y) - H(Y)$, where $H(X, Y)$ is the joint entropy of X and Y. This conditional entropy can be expressed as as the expectation (in Y) of the entropy of the (conditional) distribution of X conditioned on Y, that is,

$$H(X \mid Y) = \sum_{j=1}^{n} \mathbb{P}(Y = y_j) \sum_{i=1}^{m} -\mathbb{P}(X = x_i \mid Y = y_j) \log \mathbb{P}(X = x_i \mid Y = y_j),$$

where $x_1, \ldots, x_m$ and $y_1, \ldots, y_n$ denote the values taken by X and Y, respectively. In other words, if f_i denotes the mapping $y \mapsto \mathbb{P}(X = x_i \mid Y = y)$, then

$$H(X \mid Y) = \mathbb{E} \sum_{i=1}^{m} -f_i(Y) \log f_i(Y).$$

This second definition of conditional entropy can be generalized even when the random variable is not discrete: for an event A the mapping $y \mapsto \mathbb{P}(A \mid Y = y)$ needs to be replaced by the conditional expectation $\mathbb{E}(\mathbb{1}_A \mid Y)$, which is a measurable function of Y.

Lemma A.1. *Let X_1, X_2, Y_1, Y_2 be random variables such that X_1 and X_2 take finitely many values. Then*

$$H(X_1, X_2 \mid Y_1) - H(X_1, X_2 \mid Y_1, Y_2) \geq H(X_1 \mid Y_1) - H(X_1 \mid Y_1, Y_2).$$

Proof. We will refer to the left-hand side and the right-hand side of the above inequality as LHS and RHS, respectively. Let us notice that

$$H(X_1, X_2 \mid Y_1) = H(X_1 \mid Y_1) + H(X_2 \mid X_1, Y_1);$$

$$H(X_1, X_2 \mid Y_1, Y_2) = H(X_1 \mid Y_1, Y_2) + H(X_2 \mid X_1, Y_1, Y_2).$$

By taking the difference of the above equations we get

$$\text{LHS} = \text{RHS} + \underbrace{H(X_2 \mid X_1, Y_1) - H(X_2 \mid X_1, Y_1, Y_2)}_{\geq 0} \geq \text{RHS}$$

Note that the above argument works even if Y_1 and Y_2 are not necessarily discrete. $\square$

Next we state a weighted version of Shearer's inequality and we include a short proof.

Theorem A.2 (Shearer's inequality, weighted version). *Let X_v, $v \in V$, be finitely many random variables. Suppose that for each $U \subseteq V$ a nonnegative weight α_U is given. Then*

$$\sum_{U \subseteq V} \alpha_U H(X_U) \geq \lambda H(X_V), \quad \text{where } \lambda = \min_{v \in V} \sum_{v \in U \subseteq V} \alpha_U.$$

Proof. We may assume that $V = \{1, \ldots, n\}$ for some positive integer n. For a sequence $1 \leq k_1 < \cdots < k_r \leq n$ we decompose the joint entropy $H(X_U)$ corresponding to the subset $U = \{k_1, \ldots, k_r\}$ in the following way:

$$H(X_U) = \sum_{i=1}^{r} H(X_{k_i} \mid X_{k_1}, \ldots, X_{k_{i-1}}) \geq \sum_{i=1}^{r} H(X_{k_i} \mid X_1, \ldots, X_{k_{i-1}}).$$

It follows that

$$\sum_{U \subseteq V} \alpha_U H(X_U) \geq \sum_{k=1}^{n} \underbrace{\sum_{k \in U \subseteq V} \alpha_U}_{\geq \lambda} H(X_k \mid X_1, \ldots, X_{k-1}) \geq \lambda H(X_V).$$

$\square$

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