

Stationary distribution of Markov chains in \mathbb{R}^d with application to global random optimization

CHANG CHUNG YU DOREA

*Departamento de Matemática, Universidade e Brasília, 70910-900 Brasília-DF, Brazil.
e-mail: cdorea@mat.unb.br*

Stationary and long-run distributions of a class of Markov chains with continuous state space $S \subset \mathbb{R}^d$ are studied. Applications to random search algorithms for global optimization are presented.

Keywords: continuous state space; global random optimization; stationary distribution

1. Introduction

We consider a Markov chain $\{X_n\}_{n \geq 0}$ that takes values on a bounded and measurable subset $S \subset \mathbb{R}^d$ with $m(S) > 0$ and whose transition probability function has the representation:

$$P(X_n \in B | X_0 = x) = P^{(n)}(B|x) = P_c^{(n)}(B|x) + P_d^{(n)}(B|x) \quad (1)$$

for $x \in S$ and $B \in \mathcal{B}^d$ (Borel subsets of \mathbb{R}^d with Lebesgue measure). The absolutely continuous part,

$$P_c^{(n)}(B|x) = \int_B p_c^{(n)}(y|x) dy \quad (2)$$

is such that $p_c^{(n)}(y|x) = 0$ if $y \notin S$; similarly, the discrete part

$$P_d^{(n)}(B|x) = \sum_{y \in B \cap S_x^{(n)}} p_d^{(n)}(y|x) \quad (3)$$

is such that $p_d^{(n)}(y|x) = 0$ for $y \notin S_x^{(n)}$ and $S_x^{(n)} \subset S$ for each $x \in S$ and $n \geq 1$.

Chains with continuous state space and the above representation arise naturally in random global optimization algorithms. A general scheme for random search algorithms can be described as follows: let $X_0 \in S$ be an initial random point and let $f: S \rightarrow \mathbb{R}$ be a function whose global optimum on S (maximum or minimum) is of interest; for each $x \in S$ let $g(y|x)$ denote a density function on \mathbb{R}^d ; if X_n is the result of the algorithm at step n then at step $n + 1$ one generates a random value Y_n according to the density g ; next X_{n+1} is taken to be Y_n with an acceptance probability $a(Y_n|X_n)$ or $X_{n+1} = X_n$ with probability $1 - a(Y_n|X_n)$. It follows that the transition function of the algorithm can be written as:

$$P(B|x) = \int_B a(y|x)g(y|x) dy + I_{(x \in B)} \left[1 - \int_S a(y|x)g(y|x) dy \right] \quad (4)$$

where I denotes the indicator function. This type of Markovian algorithm will be detailed in Section 3.

Though there is an extensive literature on general state-space Markov processes, considerable mathematical background is required to understand them. In this paper we treat Markov chains in \mathbb{R}^d by extending the notions of communicability and periodicity of states to subsets of \mathbb{R}^d . It will be shown that, as in the discrete state space, they play an important role in analysing the existence of stationary and long-run distributions,

$$\lim_{n \rightarrow \infty} P^{(n)}(B|x) = Q(B). \quad (5)$$

In Section 2 we derive conditions that will guarantee (5) and in Section 3 we apply these notions to random search algorithms studied by Dorea (1986; 1990) and to a continuous version of the ‘simulated annealing’ algorithm treated by Dekkers and Aarts (1991).

2. Continuous state space

Let $\{X_n\}_{n \geq 0}$ be a Markov process with values on S . Assuming that the n th-step transition function has the representation (1), we then have:

$$\begin{aligned} P^{(n+1)}(B|x) &= \int_S P^{(n)}(B|y)P(dy|x) \\ &= \int_S P^{(n)}(B|y)p_c(y|x) dy + \sum_{y \in S_x} P^{(n)}(B|y)p_d(y|x), \end{aligned}$$

where $S_x = \{y: y \in S, p_d(y|x) > 0\}$. Using (2) and (3), we can write

$$\begin{aligned} P^{(n+1)}(B|x) &= \int_B \left[\int_S p_c^{(n)}(z|y)p_c(y|x) dy + \sum_{y \in S_x} p_c^{(n)}(z|y)p_d(y|x) \right] dz \\ &+ \int_S \left[\sum_{z \in B \cap S_y^{(n)}} p_d^{(n)}(z|y)p_c(y|x) \right] dy + \sum_{y \in S_x} \sum_{z \in B \cap S_y^{(n)}} p_d^{(n)}(z|y)p_d(y|x). \end{aligned}$$

It follows that

$$p_c^{(n+1)}(z|x) \geq \int_S p_c^{(n)}(z|y)p_c(y|x) dy + \sum_{y \in S_x} p_c^{(n)}(z|y)p_d(y|x) \quad (6)$$

and

$$p_d^{(n+1)}(z|x) \geq \sum_{y \in S_x} p_d^{(n)}(z|y)p_d(y|x). \quad (7)$$

By expressing $P^{(n+1)}(B|x) = \int_S P(B|y)P^{(n)}(dy|x)$ we also obtain the inequalities

$$p_c^{(n+1)}(z|x) \geq \int_S p_c(z|y)p_c^{(n)}(y|x) dy + \sum_{y \in S_x^{(n)}} p_c(z|y)p_d^{(n)}(y|x)$$

and

$$p_d^{(n+1)}(z|x) \geq \sum_{y \in S_x^{(n)}} p_d(z|y)p_d^{(n)}(y|x).$$

Our main result (Theorem 1 below) shows that an irreducible and aperiodic chain possesses a long-run distribution if certain regularity conditions are met. And in this case it coincides with the unique stationary distribution. The following notation and notions will be needed:

$$p^{(n)}(y|x) = p_c^{(n)}(y|x) + p_d^{(n)}(y|x). \tag{8}$$

For $B \in \mathcal{B}^d$ let

$$B^+ = \{D: D \subset B, m(D) > 0\}, \quad \text{if } m(B) > 0 \tag{9}$$

and

$$B^+ = \{D: D \neq \emptyset, D \subset B\} \quad \text{if } B \text{ is countable.} \tag{10}$$

Let

$$S^1 = \{D: D \subset S, \text{ either } m(D) > 0 \text{ or } D \text{ is countable}\}. \tag{11}$$

Definition 1. Let $\{A, B\} \subset S^1$. We say that B is accessible from A if there exists $n_B \geq 1$ such that

$$P^{(n_B)}(B^+|x) > 0, \quad \forall x \in A, \forall B^+ \in B^+$$

(which we write as $A \xrightarrow{n_B} B$). We say that A and B are communicating subsets ($A \leftrightarrow B$) if $A \xrightarrow{n_B} B$ and $B \xrightarrow{n_A} A$.

Definition 2. We say that a chain is irreducible if there exists $\{A_1, \dots, A_k\} \subset S^1$ such that $\bigcup_{j=1}^k A_j = S$ and $A_i \leftrightarrow A_j$ for all i and j .

Definition 3. Assume $A \leftrightarrow A$. We say that d_A is the period of A if d_A is the greatest common divisor of

$$\mathcal{D}(A) = \{n: P^{(n)}(A^+|x) > 0, \forall x \in A, \forall A^+ \in A^+\}.$$

If $d_A = 1$ we say that the subset A is aperiodic.

Some immediate properties can be derived:

Proposition 1. (a) If $A \xrightarrow{n_B} B$ then, given $\epsilon_0 > 0$, there exists $A_0 \in A^+$ such that either

$$P_c^{(n_B)}(B|x) \geq \epsilon_0, \forall x \in A_0 \text{ or } P_d^{(n_B)}(B|x) \geq \epsilon_0, \forall x \in A_0, \quad (12)$$

depending on whether $m(B) > 0$ or B is countable.

(b) If $A \xrightarrow{n_B} B$ and $B \xrightarrow{n_C} C$ then $A \xrightarrow{n_B+n_C} C$.

(c) If $A \leftrightarrow B$ then they have the same period ($d_A = d_B$).

(d) If d_A is the period of A then there exists $m_A \geq 1$ such that for $m \geq m_A$ we have $md_A \in \mathcal{D}(A)$.

Proof. (a) Assume $m(B) > 0$ and let $S_x^{(n_B)} = \{y: p_d^{(n_B)}(y|x) > 0\}$. Since $B \setminus S_x^{(n_B)} \in B^+$ and $m(S_x^{(n_B)}) = 0, \forall x \in A$, we must have $P_c^{(n_B)}(B'|x) > 0, \forall x \in A$ and $\forall B' \in B^+$. Given $\epsilon_0 > 0$, define for $k \geq 1$

$$D_k = \left\{ x: x \in A, P_c^{(n_B)}(B|x) < \frac{1}{k} \right\}.$$

Since $D_{k+1} \subset D_k, \bigcap_k D_k = \emptyset$ and $m(A) < \infty$ we have $\lim_{k \rightarrow \infty} m(D_k) = 0$. Let k_0 be large enough so that $\epsilon_0 \leq 1/k_0$ and $m(A \setminus D_{k_0}) > 0$ (if $m(A) > 0$) and $A \setminus D_{k_0} \neq \emptyset$ (if A is countable). Then $A_0 = A \setminus D_{k_0} \in A^+$ and for $x \in A_0$ we have $P_c^{(n_B)}(B|x) \geq \epsilon_0$.

If B is countable clearly we have $P^{(n_B)}(B'|x) = P_d^{(n_B)}(B'|x), \forall B' \in B^+$. Now proceeding with arguments of the same type we obtain (12) by observing that $P_d^{(n_B)}(B|x) > 0, \forall x \in A$.

(b) Let $C' \in C^+$; then we have $B \xrightarrow{n_C} C'$. From (a), given $\epsilon_0 > 0$, there exists $B_0 \in B^+$ such that $P^{(n_C)}(C'|y) \geq \epsilon_0$ for $y \in B_0$. Since $B_0 \in B^+$ we also have $A \xrightarrow{n_B} B_0$. Now for $x \in A$,

$$P^{n_C+n_B}(C'|x) \geq \int_{B_0} P^{(n_C)}(C'|y)P^{(n_B)}(dy|x) \geq \epsilon_0 P^{(n_B)}(B_0|x) > 0.$$

It follows that $A \xrightarrow{n_C+n_B} C$.

(c) Let $A \xrightarrow{n_B} B$ and $B \xrightarrow{n_A} A$. From (b) we have $A \xrightarrow{n_A+n_B} A$ so that $n_A + n_B \in \mathcal{D}(A)$. Now let $n_1 \in \mathcal{D}(B)$, we will show that $n_B + n_1 + n_A \in \mathcal{D}(A)$. Thus d_A divides n_1 since it divides $n_A + n_B$. It follows that $d_A \leq d_B$. Interchanging the roles of A and B , we conclude $d_A = d_B$.

Let $A' \in A^+$; since $B \xrightarrow{n_A} A'$, there exists $B_0 \in B^+$ such that $P^{(n_A)}(A'|z) \geq \epsilon_0 > 0$ on B_0 . Since $n_1 \in \mathcal{D}(B)$ we have $B \xrightarrow{n_1} B_0$, and there exists $B_1 \in B^+$ such that $P^{(n_1)}(B_0|y) \geq \epsilon_1 > 0$ on B_1 . Now for $x \in A$,

$$\begin{aligned} P^{(n_B+n_1+n_A)}(A'|x) &\geq \int_{(y \in B_1)} \int_{(z \in B_0)} P^{(n_A)}(A'|z)P^{(n_1)}(dz|y)P^{(n_B)}(dy|x) \\ &\geq \epsilon_0 \int_{(y \in B_1)} P^{(n_1)}(B_0|y)P^{(n_B)}(dy|x) \\ &\geq \epsilon_0 \epsilon_1 P^{(n_B)}(B_1|x) > 0. \end{aligned}$$

Thus $n_B + n_1 + n_A \in \mathcal{D}(A)$.

(d) Note that if $r \in \mathcal{D}(A)$ and $s \in \mathcal{D}(A)$ then, using (a), we have $r + s \in \mathcal{D}(A)$. That is, $\mathcal{D}(A)$ is closed under addition. Then there exists m_A such that $md_A \in \mathcal{D}(A)$ for $m \geq m_A$; see Doob (1953, p. 176) or Parzen (1962, p. 262). \square

Proposition 2. (a) If $A \xrightarrow{n_B} B$ then there exist $\delta_B > 0$, $A_0 \in A^+$ and $B_0 \in B^+$ such that

$$\inf \{p^{(n_B)}(y|x) : x \in A_0, y \in B_0\} \geq \delta_B. \tag{13}$$

(b) If $A \xrightarrow{n_A} A$ then there exist $\delta_A > 0$ and $A_0 \in A^+$ such that

$$\inf \{p^{(2n_A)}(y|x) : x \in A_0, y \in A_0\} \geq \delta_A. \tag{14}$$

(c) If the chain is irreducible and aperiodic, then we can decompose $S = S_c \cup S_d$ where S_d is countable and $S_c = S \setminus S_d$. Moreover, there exists $n_S \geq 1$ such that $S \xrightarrow{n_S} S_c$, and if $S_d \neq \emptyset$ we also have $S \xrightarrow{n_S} S_d$.

Proof. (a) First assume B is countable. Let $B_N = \{b_1, \dots, b_N\} \subset B$. Since $A \xrightarrow{n_B} B_N$, from (12) we can write,

$$P_d^{(n_B)}(B_N|N) \geq \delta_0 > 0 \text{ on } A_0 \in A^+. \tag{15}$$

It follows that we must have $p_d^{(n_B)}(y_j|x) \geq \delta_0/N$ for some $y_j \in B_N$. Now (13) follows using (8) and taking $B_0 = \{y_j\} \in B^+$ and $\delta_B = \delta_0/N$.

Now assume $m(B) > 0$. If A is countable let $x_0 \in A$ and $A_0 = \{x_0\} \in A^+$. Since $P_c^{(n_B)}(B|x_0) = \int_B P_c^{(n_B)}(y|x_0) dy > 0$, there exist $\delta_B > 0$ and $B_0 \in B^+$ such that on B_0 we have $p_c^{(n_B)}(y|x_0) \geq \delta_B$.

If $m(A) > 0$ then from (12),

$$P_c^{(n_B)}(B|x) \geq \delta_1 > 0 \text{ on } A_1 \in A^+. \tag{16}$$

Let $\delta_B = \delta_1/2m(B)$ and define

$$D = \{(x, y) : x \in A_1, y \in B, p_c^{(n_B)}(y|x) \geq \delta_B\}.$$

Note that for $x \in A_1$ and $D_x = \{y : (x, y) \in D\}$ we have $m(D_x) > 0$. If not, then for almost all y in D_x we have $p_c^{(n_B)}(y|x) < \delta_B$ and $P_c^{(n_B)}(B|x) \leq \delta_B m(B) \leq \delta_1/2$, which contradicts (16). Let m_2 denote the Lebesgue measure on \mathbb{R}^{2d} ; then we have $m_2(D) = \int_{A_1} m(D_x) dx > 0$. Thus there exists a rectangle $A_0 \times B_0$ with $m_2(A_0 \times B_0) > 0$, $A_0 \in A_1^+ \subset A^+$, $B_0 \in B^+$, and for $x \in A_0$ and $y \in B_0$ we have

$$p^{(n_B)}(y|x) \geq p_c^{(n_B)}(y|x) \geq \delta_B.$$

(b) From (a) there exist $A_1 \in A^+$ and $A_2 \in A^+$ such that

$$\inf \{p^{(n_A)}(y|x) : x \in A_1, y \in A_2\} \geq \delta_1 > 0.$$

By (12) there exists $A_0 \in A_2^+$ such that

$$P^{(n_A)}(A_1|y) \geq \delta_2 > 0 \text{ on } A_0.$$

Now for $x \in A_0$ and $y \in A_0$,

$$p^{(2n_A)}(y|x) \geq \int_{A_1} p^{(n_A)}(y|z)P^{(n_A)}(dz|x) \geq \delta_1 P^{(n_A)}(A_1|x) \geq \delta_1 \delta_2 > 0.$$

(c) Let $\{A_i, \dots, A_k\}$ satisfy Definition 2. Since the chain is aperiodic, from Proposition 1 there exist m_1, \dots, m_k such that for $r \geq \max\{m_1, \dots, m_k\}$ we have $r \in \mathcal{L}(A_j)$ for $j = 1, \dots, k$.

Since $m(S) > 0$ not all A_j 's can be countable. Without loss of generality, assume that A_1, \dots, A_ℓ have positive measure ($\ell \leq k$). Let $S_c = \bigcup_{r=1}^{\ell} A_r$ and $B \in S_c^+$. Then for some j we have $B_j = B \cup A_j \in A_j^+$. For $i = 1, \dots, k$ let n_i be such that $A_i \xrightarrow{n_i} B_j$. Since $m(B_j) > 0$, from the proof of (a) we have, for $i = 1, \dots, k$,

$$\inf \{p_c^{(n_i)}(y|x): x \in A'_i, y \in B'_j(i)\} \geq \delta_i > 0, \tag{17}$$

where $A'_i \in A_i^+$ and $B'_j(i) \in B_j^+$.

Now take n_S large enough so that $n_S - n_i \in \mathcal{L}(A_i)$ for $i = 1, \dots, k$. From (2) and (17) we have, for $z \in A'_i$,

$$P_c^{(n_i)}(B'_j(i)|z) \geq \delta_i m(B'_j(i)) > 0. \tag{18}$$

Let $\delta_B = \min_{1 \leq i \leq k} \{\delta_i m(B'_j(i))\}$ and $x \in S$. Then $x \in A_i$ for some i , and we have

$$P^{(n_S)}(B|x) \geq P^{(n_S)}(B'_j(i)|x) \geq \int_{A'_i} P_c^{(n_i)}(B'_j(i)|z)P^{(n_S-n_i)}(dz|x).$$

From (18) and the fact that $n_S - n_i \in \mathcal{L}(A_i)$, we have

$$P^{(n_S)}(B|x) \geq \delta_B P^{(n_S-n_i)}(A'_i|x) > 0. \tag{19}$$

Since (19) holds for all $B \in S_c^+$, we have $S \xrightarrow{n_S} S_c$.

If $\ell = k$ then the proof is completed by taking $S_d = \emptyset$. If not, let $S_d = \bigcup_{r=\ell+1}^k A_r$. Let $B \in S_d^+$ and $\ell + 1 \leq j \leq k$ such that $B_j = B \cap A_j \in A_j^+$. The proof is exactly the same, except that p_c is replaced by p_d in (17), and P_c and $m(B'_j(i))$ by P_d and $\|B'_j(i)\|$ in (18) (where $\|\cdot\|$ denotes cardinality of the set). And we have $S \xrightarrow{n_S} S_d$. □

Remark 1. Let $A \xrightarrow{n_B} B$. Then from the proofs of Propositions 1 and 2 we also have the following:

(a) If $m(B) > 0$ then

$$P_c^{(n_B)}(B'|x) > 0, \forall x \in A, \forall B' \in B^+, \tag{20}$$

$$\inf \{p_c^{(n_B)}(y|x): x \in A_0, y \in B_0\} \geq \delta_B > 0, \tag{21}$$

with $A_0 \in A^+$ and $B_0 \in B^+$.

(b) If B is countable then (20) and (21) hold with P_d and p_d in place of P_c and p_c , respectively.

(c) If $A \xleftrightarrow{n_A} A$ and $m(A) > 0$, then

$$\inf \{p_c^{(2n_A)}(y|x): x \in A_0, y \in A_0\} \geq \delta_A > 0, \tag{22}$$

with $A_0 \in A^+$. If A is countable we have (22) with p_d in place of p_c .

Our next result requires the following condition:

Condition 1. If $m(E_k) \rightarrow 0$ as $k \rightarrow \infty$ then:

$$\lim_{k \rightarrow \infty} P_c(E_k|x) < 1, \quad \text{uniformly on } x. \tag{23}$$

Note that for each $x \in S$ we always have $P_c(E_k|x) \xrightarrow{k} 0$. Condition 1 requires that the convergence be uniform on S . Also if $p_c(y|x) \leq K < \infty$ is bounded then (23) holds trivially, since $P_c(E_k|x) \leq Km(E_k)$.

Theorem 1. If a chain is irreducible and aperiodic, and if Condition 1 is satisfied, then it possesses a long-run distribution

$$\lim_{n \rightarrow \infty} P^{(n)}(B|x) = Q(B), \quad \forall B \in \mathcal{B}^d, \tag{24}$$

where Q is a probability on $(\mathbb{R}^d, \mathcal{B}^d)$.

Proof. The proof requires several steps and uses some of the techniques found in Doob (1953).

(a) Since the chain is irreducible and aperiodic, by Proposition 2 there exists $S_c \in S^+$ such that $S \xrightarrow{n_S} S_c$ (also $S_c \xrightarrow{n_S} S_c$). From (22) there exist $\delta_1 > 0$ and $S'_c \in S_c^+$ such that

$$\inf \{p_c^{(2n_S)}(y|x) : x \in S'_c, y \in S'_c\} \geq \delta_1. \tag{25}$$

From (20) we have

$$P_c^{(n_S)}(S'_c|x) > 0, \quad \forall x \in S. \tag{26}$$

Let $E_k = \{x : P_c^{(n_S)}(S'_c|x) < 1/k\}$; then, by (26) and the fact that $m(S) < \infty$, we have $m(E_k) \rightarrow 0$. From Condition 1, there exist $\epsilon_0 > 0$ and k_0 such that

$$P_c(E_{k_0}|x) \leq 1 - \epsilon_0, \quad \forall x \in S. \tag{27}$$

Since $P_c^{(n_S)}(S'_c|z) \geq 1/k_0$ for $z \in S \setminus E_{k_0}$, using (6) and (27) we can write, for $x \in S$,

$$\begin{aligned} P_c^{(n_S+1)}(S'_c|x) &\geq \int_{S \setminus E_{k_0}} P_c^{(n_S)}(S'_c|z)P(dz|x) \\ &\geq \frac{1}{k_0} P(S \setminus E_{k_0}|x) \geq \frac{\epsilon_0}{k_0}. \end{aligned} \tag{28}$$

Now take $D = S_c^+$ (thus $m(D) > 0$), $n_D = 3n_S + 1$ and $\delta_D = \delta_1 \epsilon_0 / k_0$. Then, using (6), (25) and (28), we have for $y \in D$ and $x \in S$,

$$\begin{aligned} p_c^{(3n_S+1)}(y|x) &\geq \int_D p_c^{(2n_S)}(y|z)p_c^{(n_S+1)}(z|x) dz \\ &\geq \delta_1 P_c^{(n_S+1)}(S'_c|x) \geq \delta_D > 0. \end{aligned}$$

Thus there exist $\delta_D > 0$, $n_D \geq 1$ and $D \in S^+$ such that

$$\inf \{P_c^{(n_D)}(y|x) : x \in S, y \in D\} \geq \delta_D. \tag{29}$$

(b) Let D and δ_D satisfy (29) and $\epsilon_D = \delta_D m(D)$, then

$$|P^{(kn_D)}(B|x) - P^{(kn_D)}(B|y)| \leq (1 - \epsilon_D)^k \tag{30}$$

$\forall B \in \mathcal{B}^d, \forall x \in S, \forall y \in S$ and $k \geq 1$.

From (1) and (29) we have

$$P^{(n_D)}(B|x) \geq \int_{B \cap D} P_c^{(n_D)}(y|x) dy \geq \delta_D m(B \cap D)$$

and

$$P^{(n_D)}(B^c|x) \geq \delta_D m(B^c \cap D) = \epsilon_D - \delta_D m(B \cap D).$$

It follows that for $x \in S$,

$$\delta_D m(B \cap D) \leq P^{(n_D)}(B|x) \leq 1 - \epsilon_D + \delta_D m(B \cap D). \tag{31}$$

Using inequality (31) with y in place of x , we can write

$$P^{(n_D)}(B|x) - P^{(n_D)}(B|y) \leq 1 - \epsilon_D.$$

Interchanging the roles of x and y , we obtain

$$|P^{(n_D)}(B|x) - P^{(n_D)}(B|y)| \leq 1 - \epsilon_D. \tag{32}$$

For $k \geq 2$, let

$$L(dz; x, y, k) = P^{((k-1)n_D)}(dz|x) - P^{((k-1)n_D)}(dz|y) \tag{33}$$

$$U = (L(dz; x, y, k) \geq 0) \text{ and } V = (L(dz; x, y, k) < 0).$$

And we can write

$$P^{(kn_D)}(B|x) - P^{(kn_D)}(B|y) = \int_U P^{(n_D)}(B|z)L(dz; x, y, k) + \int_V P^{(n_D)}(B|z)L(dz; x, y, k).$$

From (31) we have

$$\int_U P^{(n_D)}(B|z)L(\cdot) \leq (1 - \epsilon_D + \delta_D m(B \cap D)) \int_U L(\cdot)$$

and

$$\int_V P^{(n_D)}(B|z)L(\cdot) \leq \delta_D m(B \cap D) \int_V L(\cdot).$$

Since $\int_U L(\cdot) + \int_V L(\cdot) = 0$, we have

$$P^{(kn_D)}(B|x) - P^{(kn_D)}(B|y) \leq (1 - \epsilon_D) \int_U L(\cdot). \tag{34}$$

If $k = 2$ we have, from (32),

$$\int_U L(\cdot) = P^{(n_D)}(U|x) - P^{(n_D)}(U|y) \leq 1 - \epsilon_D.$$

Thus

$$P^{(2n_D)}(B|x) - P^{(2n_D)}(B|y) \leq (1 - \epsilon_D)^2.$$

Induction arguments and (34) give us (30).

(c) For $k \geq 1, m \geq 1$ and $x \in S$, we have

$$|P^{(kn_D+m)}(B|x) - P^{(kn_D)}(B|x)| \leq (1 - \epsilon_D)^k. \tag{35}$$

Since $\int_S P^{(m)}(dy|x) = 1$ and $P^{(kn_D+m)}(B|x) = \int_S P^{(kn_D)}(B|y)P^{(m)}(dy|x)$, we can write

$$P^{(kn_D+m)}(B|x) - P^{(kn_D)}(B|x) = \int_S [P^{(kn_D)}(B|y) - P^{(kn_D)}(B|x)]P^{(m)}(dy|x).$$

and (35) follows from (30).

(d) $P^{(n)}(B|x)$ is a Cauchy sequence by (35). For $B \in \mathcal{B}^d$ let $Q(B) = \lim_{n \rightarrow \infty} P^{(n)}(B|x)$, which is independent of x by (30). It is easy to verify that Q is σ -additive on \mathcal{B}^d and since $Q(S) = 1$ it is a probability on $(\mathbb{R}^d, \mathcal{B}^d)$. \square

Remark 2. (a) Under the hypothesis of Theorem 1 the long-run distribution Q necessarily has an absolutely continuous part. Note that from (29) we have $p_c^{(n_D)}(y|x) \geq \delta_D > 0, \forall y \in D$ and $\forall x \in S$ with $m(D) > 0$. And from (6) for $y \in D, x \in S$,

$$p_c^{(n_D+1)}(y|x) \geq \int_S p_c^{(n_D)}(y|z)P(dz|x) \geq \delta_D.$$

Thus for $D' \in D^+$ we have

$$\lim_{n \rightarrow \infty} P_c^{(n)}(D'|x) \geq \delta_D m(D').$$

(b) Our next theorem shows that the results of Theorem 1 hold if we assume the following condition:

Condition 1'. if $m(E_k) \rightarrow 0$ then $\lim_{n \rightarrow \infty} P_d(E_k|x) = 0$ uniformly on S .

Theorem 1'. Assume that the chain is irreducible and aperiodic with $S_d \neq \emptyset$. Then (24) holds if Condition 1' is satisfied.

Proof. From Proposition 2(c), if the chain is irreducible and aperiodic then $S = S_c \cap S_d$ with S_d countable and $S_c = S \setminus S_d$. Since $S_d \neq \emptyset$ there exists $n_S \geq 1$ with $S \xrightarrow{n_S} S_d$, and by (22) there exist $S'_d \in S_d^+$ and $\delta_1 > 0$ such that

$$\inf \{p_d^{(2n_S)}(y|x): x \in S'_d, y \in S'_d\} \geq \delta_1.$$

And by (20) we have $P_d^{(n_S)}(S'_d|x) > 0, \forall x \in S$.

Let $E_k = \{x: P_d^{(n_S)}(S'_d|x) < 1/k\}$; then $m(E_k) \rightarrow 0$. From Condition 1', given $\epsilon_0 > 0$, there exists k_0 such that $P_d(E_{k_0}^c|x) \geq \epsilon_0$ for $x \in S$. From (7) we have

$$\begin{aligned} P_d^{(n_S+1)}(S'_d|x) &\geq \sum_{z \in (E_{k_0}^c \cap S_x)} P_d^{(n_S)}(S'_d|z) p_d(z|x) \\ &\geq \frac{1}{k_0} P_d(E_{k_0}^c|x) \geq \frac{\epsilon_0}{k_0}. \end{aligned}$$

Now let $D = S'_d$, $n_D = 3n_S + 1$ and $\delta_D = \delta_1 \epsilon_0 / k_0$ and we have for, $y \in D$ and $x \in S$,

$$\begin{aligned} P_d^{(3n_S+1)}(y|x) &\geq \sum_{z \in D} P_d^{(2n_S)}(y|z) P_d^{(n_S+1)}(z|x) \\ &\geq \delta_1 P_d^{(n_S+1)}(S'_d|x) \geq \frac{\delta_1 \epsilon_0}{k_0} = \delta_D > 0. \end{aligned}$$

Thus there exist $\delta_D > 0$, $n_D \geq 1$ and $D \in S_d^+$ such that

$$\inf \{P_d^{(n_D)}(y|x): x \in S, y \in D\} \geq \delta_D. \tag{36}$$

It follows that for $B \in B^d$

$$P^{(n_D)}(B|x) \geq \sum_{(y \in B \cap D \cap S_x^{(n_D)})} p_d^{(n_D)}(y|x) \geq \delta_D \|D \cap B\|$$

and

$$\delta_D \|D \cap B\| \leq P^{(n_D)}(B|x) \leq 1 - \delta_D \|D\| + \delta_D \|D \cap B\|.$$

Since $0 < \delta_D \|D\| < 1$, using the same arguments as in Theorem 1 we obtain (24). □

Theorem 2. Let $\{E_1, \dots, E_k\} \subset S^1$ be mutually communicating and aperiodic subsets of S . For $E = \bigcup_{i=1}^k E_i$, let $F = S \setminus E$. Assume that $F \neq \emptyset$, $m(E) > 0$. Condition 1 holds and that for some r and n_F we have $F \xrightarrow{n_F} E_r$. Then the chain has a long-run distribution.

Proof. Since $m(E) > 0$ we may assume $m(E_i) > 0$ for $i = 1, \dots, \ell$ and E_i countable for $i = \ell + 1, \dots, k$. Let $E_c = \bigcup_{i=1}^{\ell} E_i$ and $E_d = \bigcup_{i=\ell+1}^k E_i$.

First, we will show that there exists $n'_F \geq 1$ such that

$$F \xrightarrow{n'_F} E_c \text{ and } F \xrightarrow{n'_F} E_d \quad (\text{if } E_d \neq \emptyset). \tag{37}$$

Since the E_i are communicating and aperiodic subsets we can take m large enough so that $E_r \xrightarrow{m} E_i$ for $i = 1, \dots, k$. Since $F \xrightarrow{n_F} E_r$ we have (37) by setting $n'_F = n_F + m$ and using Proposition 1.

Using aperiodicity again, there exists $n_E \geq 1$ such that

$$S \xrightarrow{n_E} E_c \text{ and } S \xrightarrow{n_E} E_d \quad (\text{if } E_d \neq \emptyset). \tag{38}$$

Now $m(E_c) > 0$ and $E_c \xrightarrow{n_E} E_c$. Using exactly the same type of argument as in the proof of Theorem 1, we show that there exist $D \in E^+$, $\delta_D > 0$ and $n_D \geq 1$ such that

$$\inf \{p_c^{(n_D)}(y|x) : x \in S, y \in D\} \geq \delta_D.$$

Following the proof of Theorem 1, we have (24). □

3. Applications

Consider the problem of estimating the global minimum of $f: S \rightarrow \mathbb{R}$, that is,

$$y_{\min} = \min_{x \in S} \{f(x)\} \text{ or } S_{\min} = \{x : x \in S, f(x) = y_{\min}\}. \tag{39}$$

Assume that S is bounded with $m(S) > 0$, the global minimum y_{\min} is finite, f is continuous in a neighbourhood of each minimum point $x_{\min} \in S_{\min}$ and the minimum points are interior points of S .

The following random search algorithm will be used: let $X_0 \in S$ be an initial random point; for each $x \in S$ let $g(\cdot|x)$ be a density function on \mathbb{R}^d ; for $k \geq 0$ let X_k denote the value of the algorithm at step k ; at step $k + 1$ a random value Y_k is generated according to the density $g(\cdot|X_k)$ and we define

$$X_{k+1} = \begin{cases} Y_k & \text{with probability } a(Y_k|X_k) \\ X_k & \text{with probability } 1 - a(Y_k|X_k). \end{cases}$$

It follows that the Markov chain $\{X_n\}_{n \geq 0}$ has the transition probability function given by $P(B|x) = P_c(B|x) + P_d(B|x)$, with

$$P_c(B|x) = \int_B p_c(y|x) dy, \quad p_c(y|x) = a(y|x)g(y|x) \tag{40}$$

and

$$P_d(B|x) = \sum_{y \in B \cap \{x\}} p_d(y|x), \quad p_d(x|x) = 1 - \int_S p_c(y|x) dy, \tag{41}$$

and $p_d(y|x) = 0$ if $y \neq x$.

Note that the second step transition is given by

$$P^{(2)}(B|x) = \int_S P(B|y)p_c(y|x) dy + P(B|x)p_d(x|x),$$

and writing $P(B|x) = \int_B p_c(y|x) dy + I_{(x \in B)} p_d(x|x)$ we have

$$P_c^{(2)}(B|x) = \int_B \left[\int_S p_c(z|y)p_c(y|x) dy + p_d(z|x)p_c(z|x) + p_c(z|x)p_d(x|x) \right] dz$$

and

$$P_d^{(2)}(B|x) = I_{(x \in B)} p_d^2(x|x).$$

In general we have

$$p_d^{(n)}(y|x) = p_d^n(y|x)$$

and

$$p_c^{(n)}(y|x) = \int_S p_c^{(n-1)}(y|z)p_c(z|x) dz + p_d^{n-1}(y|y)p_c(y|x) + p_d(y|y)p_c^{(n-1)}(y|x).$$

Note that, in this case, inequality (6) is strict and we have equality in (7). Two types of algorithms will be analysed.

Algorithm 1. Take $g(y) = g(y|x)$ independent of x and the acceptance probability to be $a(y|x) = I_{(f(y) \leq f(x))} I_{(y \in S)}$.

Algorithm 2. Take $a(y|x) = \min \{1, \exp \{-c(f(y) - f(x))\}\}$, where $c > 0$ is a constant.

For Algorithm 2 we assume the same type of hypothesis as in Dekkers and Aarts (1991) (but weaker relative to the objective function f and the set of minimum points S_{\min}): (i) if $m(A) > 0$ then $\int_A g(y|x) dy > 0, \forall x \in S$; (ii) if $m(E_k) \xrightarrow{k} 0$ then $\int_{E_k} g(y|x) dy \xrightarrow{k} 0$ uniformly on x ; and (iii) $\int_S g(y|x) dy = 1$ for all $x \in S$ and $g(y|x) = g(x|y)$.

We will show that the hypothesis of Theorem 2 is satisfied and the long-run distribution is given by

$$Q(B) = \int_B \alpha e^{-c(f(y) - y_{\min})} dy \quad \text{with } \alpha^{-1} = \int_S e^{-c(f(y) - y_{\min})} dy. \tag{42}$$

For $\epsilon > 0$ define

$$\eta(\epsilon) = \{x: x \in S, |x - x_0| \leq \epsilon \text{ for some } x_0 \in S_{\min}\}. \tag{43}$$

Let $y_{\min}(\epsilon) = \inf \{f(x): x \in S \setminus \eta(\epsilon)\}$ and

$$B(\epsilon) = \eta(\epsilon) \cap \{x: x \in S, f(x) \leq y_{\min}(\epsilon)\}. \tag{44}$$

Since f is continuous in a neighbourhood of each minimum point we have $m(B(\epsilon)) > 0$. We will show that $B(\epsilon) \xleftrightarrow{1} B(\epsilon)$ and $S \setminus B(\epsilon) \xrightarrow{1} B(\epsilon)$. This, together with (ii), verifies the conditions of Theorem 2. Thus the long-run distribution exists and coincides with the unique stationary distribution. To prove (42) it is enough to show that the stationary density is given by $q(y) = \alpha \exp \{-c(f(y) - y_{\min})\}$. And this can be done by verifying that q satisfies

$$q(y) = \int_S p_c(y|x)q(x) dx + q(y)p_d(y|y).$$

To prove $B(\epsilon) \xleftrightarrow{1} B(\epsilon)$, first note that $f(y) - f(x) \leq y_{\min}(\epsilon) - y_{\min}$ for $x \in B(\epsilon)$ and $y \in B(\epsilon)$. It follows that $a(y|x) \geq \delta_\epsilon = \exp \{-c(y_{\min}(\epsilon) - y_{\min})\}$. Now let $B' \in B^+(\epsilon)$ and $x \in B(\epsilon)$; then by (40) and (i) we have

$$P_c(B'|x) = \int_{B'} a(y|x)g(y|x) dy \geq \delta_\epsilon \int_{B'} g(y|x) dy > 0.$$

To prove that $S \setminus B(\epsilon) \xrightarrow{1} B(\epsilon)$, note that for $z \in S \setminus B(\epsilon)$ and $y \in B(\epsilon)$ we have $f(y) \leq f(z)$ so that $a(y|z) = 1$. And by (i),

$$P(B'|z) \leq \int_{B'} g(y|z) dy > 0, \quad \forall B' \in B^+(\epsilon).$$

As for Algorithm 1, we assume that x_0 is the unique minimum point and that $g(y) > 0$ in a neighbourhood of x_0 . An atypical situation arises: $S_{\min} \xleftrightarrow{1} S_{\min}$ but S_{\min} is not accessible from any other subset of S (for all $n \geq 1$ we have $P^{(n)}(\{x_0\}|x)$ equal to 0 if $x \neq x_0$ and equal to 1 if $x = x_0$). Now let $B(\epsilon)$ be defined by (44) and $\epsilon > 0$ small enough so that $g(y) > 0$ on $B(\epsilon)$. Then we can show that $S \setminus B(\epsilon) \xrightarrow{1} B(\epsilon)$. In this case one can prove directly that the long-run distribution Q is the probability mass at x_0 . Note that for all $n \geq 1$ and $\epsilon > 0$ we have $P^{(n)}(B(\epsilon)|x_0) = 1$. And for $x \neq x_0$ and $q_\epsilon = \int_{B(\epsilon)} g(y) dy$ we have $P(B^c(\epsilon)|x) = 1 - q_\epsilon$. Using induction arguments it is easy to show that, for $x \neq x_0$,

$$P^{(n)}(B^c(\epsilon)|x) = \int_{B^c(\epsilon)} P^{(n-1)}(B^c(\epsilon)|y)P(dy|x) = (1 - q_\epsilon)^n. \quad (45)$$

From (45), if $\eta(\epsilon)$ is an ϵ -neighbourhood of x_0 , we have

$$\lim_{n \rightarrow \infty} P^{(n)}(\eta(\epsilon)|x) = 1, \quad \forall x \in S.$$

It follows that $X_n \rightarrow x_0$ in probability and $Q(\{x_0\}) = 1$.

Acknowledgement

The research for this paper was partially supported by CNPq-Brasil.

References

- Dekkers, A. and Aarts, E. (1991). Global optimization and simulated annealing. *Math. Programming*, **50**, 367–393.
- Doob, J.L. (1953). *Stochastic Processes*. New York: Wiley.
- Dorea, C.C.Y. (1986). Limiting distribution for random optimization methods. *SIAM J. Control Optim.*, **24**, 76–82.
- Dorea, C.C.Y. (1990). Stopping rules for a random optimization method. *SIAM J. Control Optim.*, **28**, 841–850.
- Parzen, E. (1962). *Stochastic Processes*. San Francisco: Holden-Day.

Received October 1995 and revised July 1996.