

45. Axiom Systems of Aristotle Traditional Logic. II

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In this paper, we shall give new axiom systems of Aristotle traditional logic. Some systems have been obtained by J. Lukasiewicz ([4], [5]), I. Bocheński ([1]), N. Kretzmann ([3]), and recently K. Iséki ([2]).

K. Iséki has given a method to find axiom systems. For the detail, see [2].

We use the following notations. For the categorical sentences,

- 1) Aab : Every a is b ,
- 2) Iab : At least one a is b ,
- 3) Oab : At least one a is not b ,
- 4) Eab : No a is b ,

For functors,

- 1) C : Implication, 2) N : Negation, 3) K : Conjunction.

Then we have

$$D1 \quad Eab = NIab, \quad D2 \quad Oab = NAab.$$

For moods and figures:

- 1) XY_1 : $CXabYab$,
- 2) XY_2 : $CXabYba$,
- 3) XYZ_1 : $CKXabYcaZcb$,
- 4) XYZ_2 : $CKXabYcbZca$,
- 5) XYZ_3 : $CKXabYacZcb$,
- 6) XYZ_4 : $CKXabYbcZca$.

Under these symbols, the Lukasiewicz axiom system is written in the form of

$$L1 \quad Aaa,$$

$$L2 \quad Iaa,$$

$$L3 \quad AAA_1,$$

$$L4 \quad AII_3.$$

From theses of the classical propositional calculus, we have the following deduction rules $T1$ - $T7$. We shall symbolize these rules as right sides.

$$\begin{array}{l}
 T1 \quad CK\alpha\beta\gamma \begin{array}{l} \longrightarrow CK\beta\alpha\gamma \\ \searrow CK\alpha N\gamma N\beta \\ \searrow CKN\gamma\beta N\alpha \end{array} : \alpha\beta\gamma \begin{array}{l} \longrightarrow (i) \beta\alpha\gamma, \\ \searrow (ii) \alpha N\gamma N\beta, \\ \searrow (iii) N\gamma\beta N\alpha, \end{array} \\
 T2 \quad CK\alpha\beta\gamma, C\delta\alpha \longrightarrow CK\delta\beta\gamma; \quad \alpha\beta\gamma + \delta\alpha \longrightarrow \delta\beta\gamma, \\
 T3 \quad CK\alpha\beta\gamma, C\delta\beta \longrightarrow CK\alpha\delta\gamma; \quad \alpha\beta\gamma + \delta\beta \longrightarrow \alpha\delta\gamma, \\
 T4 \quad CK\alpha\beta\gamma, C\gamma\delta \longrightarrow CK\alpha\beta\delta; \quad \alpha\beta\gamma + \gamma\delta \longrightarrow \alpha\beta\delta,
 \end{array}$$

- $T5 \quad C\alpha\beta \longrightarrow CN\beta N\alpha; \quad \alpha\beta \longrightarrow N\beta N\alpha,$
 $T6 \quad CK\alpha\beta\gamma, \quad \alpha \longrightarrow C\beta\gamma; \quad \alpha\beta\gamma + \alpha \longrightarrow \beta\gamma,$
 $T7 \quad CK\alpha\beta\gamma, \quad \beta \longrightarrow C\alpha\gamma; \quad \alpha\beta\gamma + \beta \longrightarrow \alpha\gamma.$

Theorem 1. *Under the deduction rule T1, we have the following deductively equivalent groups G1–G8:*

- $G1 \quad AAA_1 \quad AOO_2 \quad OAO_3,$
 $G2 \quad EAE_1 \quad EIO_2 \quad IAI_3,$
 $G3 \quad AAI_1 \quad AEO_2 \quad EAO_3,$
 $G4 \quad EAO_1 \quad EAO_2 \quad AAI_3,$
 $G5 \quad EIO_1 \quad EAE_2 \quad AII_3,$
 $G6 \quad AII_1 \quad AEE_2 \quad EIO_3,$
 $G7 \quad AEE_4 \quad EIO_4 \quad IAI_4,$
 $G8 \quad AAI_4 \quad AEO_4 \quad EAO_4.$

Proof. Under the deduction rule T1, we have

- $R1 \quad XYZ_1 \begin{cases} \longrightarrow (i) \quad XNZNY_2 \quad (T1 \text{ (ii)}), \\ \longrightarrow (ii) \quad NZYNX_3 \quad (T1 \text{ (iii)}), \end{cases}$
 $R2 \quad XYZ_2 \begin{cases} \longrightarrow (i) \quad XNZNY_1 \quad (T1 \text{ (ii)}), \\ \longrightarrow (ii) \quad YNZNX_3 \quad (T1 \text{ (iii)}, T1 \text{ (i)}), \end{cases}$
 $R3 \quad XYZ_3 \begin{cases} \longrightarrow (i) \quad NZXNY_2 \quad (T1 \text{ (ii)}, T1 \text{ (i)}), \\ \longrightarrow (ii) \quad NZYNX_1 \quad (T1 \text{ (iii)}), \end{cases}$
 $R4 \quad XYZ_4 \begin{cases} \longrightarrow (i) \quad NZXNY_4 \quad (T1 \text{ (ii)}, T1 \text{ (i)}), \\ \longrightarrow (ii) \quad YNZNX_4 \quad (T1 \text{ (iii)}, T1 \text{ (i)}). \end{cases}$

For example, we can change AAA_1 into the following thesis containing N by R1, and further can eliminate the functor N , by D1.

- $ANANA_2 \longrightarrow AOO_2 \quad (R1 \text{ (i)}, D1),$
 $NAANA_3 \longrightarrow OAO_3 \quad (R1 \text{ (ii)}, D1).$

Similarly we have

- $AOO_2 \begin{cases} \longrightarrow ANONO_1 \longrightarrow AAA_1 \quad (R2 \text{ (i)}, D1), \\ \longrightarrow NOONA_3 \longrightarrow OAO_3 \quad (R2 \text{ (ii)}, D1), \end{cases}$
 $OAO_3 \begin{cases} \longrightarrow ONONA_2 \longrightarrow AOO_2 \quad (R3 \text{ (i)}, D1), \\ \longrightarrow NOANO_1 \longrightarrow AAA_1 \quad (R3 \text{ (ii)}, D1). \end{cases}$

Therefore AAA_1 , AOO_2 , and OAO_3 are deductively equivalent each other under T1. By the same method, we have

- $EAE_1 \begin{cases} \longrightarrow ENENA_2 \longrightarrow EIO_2 \quad (R1 \text{ (i)}, D1, D2), \\ \longrightarrow NEANE_3 \longrightarrow IAI_3 \quad (R1 \text{ (ii)}, D1), \end{cases}$
 $EIO_2 \begin{cases} \longrightarrow ENONI_1 \longrightarrow EAE_1 \quad (R2 \text{ (i)}, D1, D2), \\ \longrightarrow INONE_3 \longrightarrow IAI_3 \quad (R2 \text{ (ii)}, D1, D2), \end{cases}$
 $IAI_3 \begin{cases} \longrightarrow NIINA_2 \longrightarrow EIO_2 \quad (R3 \text{ (i)}, D1, D2), \\ \longrightarrow NIANI_1 \longrightarrow EAE_1 \quad (R3 \text{ (ii)}, D1), \end{cases}$
 $AAI_1 \begin{cases} \longrightarrow ANINA_2 \longrightarrow AEO_2 \quad (R1 \text{ (i)}, D1, D2), \\ \longrightarrow NIANA_3 \longrightarrow EAO_3 \quad (R1 \text{ (ii)}, D1, D2), \end{cases}$
 $AEO_2 \begin{cases} \longrightarrow ANONE_1 \longrightarrow AAI_1 \quad (R2 \text{ (i)}, D1, D2), \\ \longrightarrow ENONA_3 \longrightarrow EAO_3 \quad (R2 \text{ (ii)}, D2), \end{cases}$
 $EAO_3 \begin{cases} \longrightarrow NOENA_2 \longrightarrow AEO_2 \quad (R3 \text{ (i)}, D2), \\ \longrightarrow NOANE_1 \longrightarrow AAI_1 \quad (R3 \text{ (ii)}, D1, D2), \end{cases}$

$$\begin{array}{l}
EAO_1 \longrightarrow ENONA_2 \longrightarrow EAO_2 \quad (R1 \text{ (i)}, D2), \\
 \searrow \longrightarrow NOANE_3 \longrightarrow AAI_3 \quad (R1 \text{ (ii)}, D1, D2), \\
EAO_2 \longrightarrow ENONA_1 \longrightarrow EAO_1 \quad (R2 \text{ (i)}, D2), \\
 \searrow \longrightarrow ANONE_3 \longrightarrow AAI_3 \quad (R2 \text{ (ii)}, D1, D2), \\
AAI_3 \longrightarrow NIANA_2 \longrightarrow EAO_2 \quad (R3 \text{ (i)}, D1, D2), \\
 \searrow \longrightarrow NIANA_1 \longrightarrow EAO_1 \quad (R3 \text{ (ii)}, D1, D2), \\
EIO_1 \longrightarrow ENONI_2 \longrightarrow EAE_2 \quad (R1 \text{ (i)}, D1, D2), \\
 \searrow \longrightarrow NOINE_3 \longrightarrow AII_3 \quad (R1 \text{ (ii)}, D1, D2), \\
EAE_2 \longrightarrow ENENA_1 \longrightarrow EIO_1 \quad (R2 \text{ (i)}, D1, D2), \\
 \searrow \longrightarrow ANENE_3 \longrightarrow AII_3 \quad (R2 \text{ (ii)}, D1), \\
AII_3 \longrightarrow NIANI_2 \longrightarrow EAE_2 \quad (R3 \text{ (i)}, D1), \\
 \searrow \longrightarrow NIINA_1 \longrightarrow EIO_1 \quad (R3 \text{ (ii)}, D1, D2), \\
AII_1 \longrightarrow ANINI_2 \longrightarrow AEE_2 \quad (R1 \text{ (i)}, D1), \\
 \searrow \longrightarrow NIINA_3 \longrightarrow EIO_3 \quad (R1 \text{ (ii)}, D1, D2), \\
AEE_2 \longrightarrow ANENE_1 \longrightarrow AII_1 \quad (R2 \text{ (i)}, D1), \\
 \searrow \longrightarrow ENENA_3 \longrightarrow EIO_3 \quad (R2 \text{ (ii)}, D1, D2), \\
EIO_3 \longrightarrow NOENI_2 \longrightarrow AEE_2 \quad (R3 \text{ (i)}, D1, D2), \\
 \searrow \longrightarrow NOINE_1 \longrightarrow AII_1 \quad (R3 \text{ (ii)}, D1, D2), \\
AEE_4 \longrightarrow NEANE_4 \longrightarrow IAI_4 \quad (R4 \text{ (i)}, D1), \\
 \searrow \longrightarrow ENENA_4 \longrightarrow EIO_4 \quad (R4 \text{ (ii)}, D1, D2), \\
EIO_4 \longrightarrow NOENI_4 \longrightarrow AEE_4 \quad (R4 \text{ (i)}, D1, D2), \\
 \searrow \longrightarrow INONE_4 \longrightarrow IAI_4 \quad (R4 \text{ (ii)}, D1, D2), \\
IAI_4 \longrightarrow NIINA_4 \longrightarrow EIO_4 \quad (R4 \text{ (i)}, D1, D2), \\
 \searrow \longrightarrow ANINI_4 \longrightarrow AEE_4 \quad (R4 \text{ (ii)}, D1), \\
AAI_4 \longrightarrow NIANA_4 \longrightarrow EAO_4 \quad (R4 \text{ (i)}, D1, D2), \\
 \searrow \longrightarrow ANINA_4 \longrightarrow AEO_4 \quad (R4 \text{ (ii)}, D1, D2), \\
AEO_4 \longrightarrow NOANE_4 \longrightarrow AAI_4 \quad (R4 \text{ (i)}, D1, D2), \\
 \searrow \longrightarrow ENONA_4 \longrightarrow EAO_4 \quad (R4 \text{ (ii)}, D2), \\
EAO_4 \longrightarrow NOENA_4 \longrightarrow AEO_4 \quad (R4 \text{ (i)}, D2), \\
 \searrow \longrightarrow ANONE_4 \longrightarrow AAI_4 \quad (R4 \text{ (ii)}, D1, D2).
\end{array}$$

Therefore the proof is complete.

Theorem 2. *The Aristotle traditional logic is given by the following set:*

L1, L2, a thesis of G1, and a thesis of G5.

Proof. An axiom system of J. Lukasiewicz [5] is given by *L1, L2, L3, and L4*. *L3* is a thesis of *G1*, and *L4* is a thesis of *G5*. Therefore, by Theorem 1, we have Theorem 2.

Theorem 3. *The Aristotle traditional logic is given by the following set:*

L1, L2, a thesis of G1, and a thesis of G7.

Proof.

- L1* *Aaa,*
- L2* *Iaa,*
- G1* *AAA₁ ~ AOO₂ ~ OAO₃,*
- G7* *IAI₄ ~ EIO₄ ~ AEE₄.*

From the rule of term order changes (see, [2]), we have

$$\begin{array}{l}
 RT1 \quad XYZ_1 \begin{array}{l} \longrightarrow (i) \tilde{X}YZ_2, \\ \longrightarrow (ii) X\tilde{Y}Z_3, \end{array} \\
 RT2 \quad XYZ_2 \begin{array}{l} \longrightarrow (i) \tilde{X}YZ_1, \\ \longrightarrow (ii) X\tilde{Y}Z_4, \end{array} \\
 RT3 \quad XYZ_3 \begin{array}{l} \longrightarrow (i) \tilde{X}YZ_4, \\ \longrightarrow (ii) X\tilde{Y}Z_1, \end{array} \\
 RT4 \quad XYZ_4 \begin{array}{l} \longrightarrow (i) \tilde{X}YZ_1, \\ \longrightarrow (ii) X\tilde{Y}Z_3, \end{array}
 \end{array}$$

where \tilde{X} means the order change of terms contained in X .

In $T6$, put $\alpha = Iaa$, $\beta = Aab$, $\gamma = Iba$, then by $G7$ and $L2$, we have $AabIba$. That is:

$$\begin{array}{l}
 IaaAI_4 + Iaa \longrightarrow 1, (T6, G7, L2) \\
 1 \quad AI_2. \\
 IAaaI_4 + Aaa \longrightarrow 2, (T7, G7, L1) \\
 2 \quad II_2. \\
 AI_2 \longrightarrow 3, (T5, 1, D1, D2) \\
 3 \quad EO_2. \\
 II_2 \longrightarrow 4, (T5, 2, D1) \\
 4 \quad EE_2. \\
 EIO_4 + AI_2 \longrightarrow 5, (T3, G7, 1, RT4 (ii)) \\
 5 \quad EAO_2 \sim AAI_3. \\
 AAaaI_3 + Aaa \longrightarrow 6, (T7, 5, L1) \\
 6 \quad AI_1. \\
 AI_1 \longrightarrow 7, (T5, 6, D1, D2) \\
 7 \quad EO_1. \\
 IAI_4 + II_2 \longrightarrow 8, (T2, G7, 2, RT4 (i)) \\
 8 \quad IAI_3. \\
 IAbbI_3 + Abb \longrightarrow 9, (T7, 8, L1) \\
 9 \quad II_1. \\
 II_1 \longrightarrow 10, (T5, 9, D1) \\
 10 \quad EE_1. \\
 EIO_4 + EE_2 \longrightarrow 11, (T2, G7, 4, RT4 (i)) \\
 11 \quad EIO_3. \\
 EIO_3 + II_2 \longrightarrow 12, (T3, 11, 2, RT3 (ii)) \\
 12 \quad EIO_1 \sim AII_3 \sim EAE_2.
 \end{array}$$

$L1$, $L2$, $L3$ (one of $G1$), and 12 are axioms by J. Lukasiewicz, I. Bocheński and K. Iséki. Therefore the proof is complete.