

A recursive calculation of the Arf invariant of a link

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The Arf invariant of a knot was introduced in [11], and it can be calculated from its Alexander polynomial or its Conway polynomial [6]. The Arf invariant of a proper link (a link L is proper if $\text{lk}(K, L-K)$ is even for every component K in L , where lk means a linking number) is defined to be that of a knot which is related to it (a knot K is related to a link L if there is a smoothly and properly embedded disk with holes D in $\mathbf{R}^3 \times [0, 1]$ with $D \cap \mathbf{R}^3 \times \{0\} = K$ and $D \cap \mathbf{R}^3 \times \{1\} = -L$ [11]). K. Murasugi found a relation between the Arf invariants and the Alexander polynomials of two-component links [10]. The author showed in [9] that for some classes of proper links the Arf invariants can be expressed in terms of their Conway polynomials. See also [3].

In this paper we consider $V_L(i)$, where $V_L(t)$ is V. F. R. Jones' trace invariant [5] and $i = \sqrt{-1}$. He proposed there that one is allowed to define an Arf invariant of L as $V_L(i)$, and here we show that

$$V_L(i) = \begin{cases} (\sqrt{2})^{\#(L)-1}, & \text{if } L \text{ is proper and } \text{Arf}(L)=0, \\ -(\sqrt{2})^{\#(L)-1}, & \text{if } L \text{ is proper and } \text{Arf}(L)=1, \text{ and} \\ 0, & \text{if } L \text{ is not proper,} \end{cases}$$

where $\#(L)$ is the number of components in L , $\text{Arf}(L)$ is the Arf invariant of L , and $\sqrt{-1}$ is chosen to be $e^{(5/8) \cdot 2\pi i}$ in $V_L(i)$. This gives an answer to the Problem 12 in [2].

Using a recursive definition of $V_L(t)$ introduced by several people ([4], [8]), we can calculate the Arf invariant of any proper link recursively as follows.

DEFINITION. For any oriented link L , a numerical link type invariant $I(L)$ is defined so that it satisfies the following two axioms.

- (i) For the trivial knot O , $I(O)=1$, and
- (ii) If three links L , L' , and l are related as in Figure 1 (the other parts are identical), then

$$I(L) + I(L') = \sqrt{2} \cdot I(l).$$

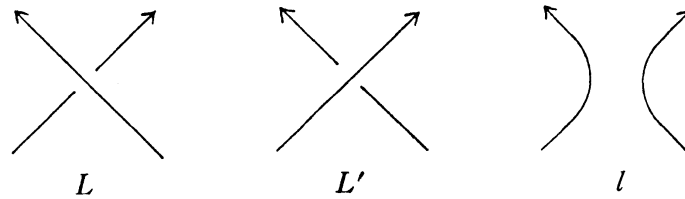


Figure 1.

REMARK. $V_L(t)$ is defined so that $1/t \cdot V_L(t) - t \cdot V_{L'}(t) = (\sqrt{t} - 1/\sqrt{t})V_l(t)$ with $V_o(t)=1$. A simple calculation shows that $I(L) = V_L(i)$, and so the above definition is well-defined. For another proof of well-definedness see [4], [8].

Then we have

THEOREM.

$$I(L) = \begin{cases} (\sqrt{2})^{\#(L)-1}, & \text{if } L \text{ is proper and } \text{Arf}(L)=0, \\ -(\sqrt{2})^{\#(L)-1}, & \text{if } L \text{ is proper and } \text{Arf}(L)=1, \text{ and} \\ 0, & \text{if } L \text{ is not proper.} \end{cases}$$

Before proving the theorem, we show the following.

LEMMA. Suppose that L, L' , and l are given as in Figure 1 and that $\#(L) = \#(L') = \#(l) - 1$. If L and L' are proper and l is not proper, then $\text{Arf}(L) \neq \text{Arf}(L')$.

PROOF. Let K, K', k_1 , and k_2 be knots in L, L' , or l as indicated in Figure 2.

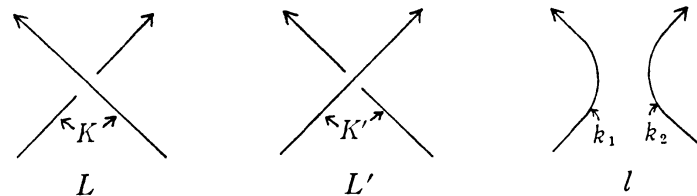


Figure 2.

Let \bar{k} be a knot obtained from $l - k_1$ after a fusion, \bar{l} be the resulting two-component link obtained from l , and \bar{L} and \bar{L}' be the corresponding knots obtained from L and L' respectively.

We will show that $\text{lk}(k_1, \bar{k})$ is odd. Since L is proper, $0 \equiv \text{lk}(K, L - K) \equiv \text{lk}(k_1, l - k_1) + \text{lk}(k_2, l - k_2) \pmod{2}$. Thus we have $\text{lk}(k_1, \bar{k}) \equiv 1 \pmod{2}$ since otherwise $\text{lk}(k_1, l - k_1) \equiv \text{lk}(k_2, l - k_2) \equiv 0 \pmod{2}$ and l cannot be non-proper.

Now it follows from Theorem 10.7 in [7] (see also Lemma 3.1 in [12]) that $\text{Arf}(L) + \text{Arf}(L') \equiv \text{Arf}(\bar{L}) + \text{Arf}(\bar{L}') \equiv \text{lk}(k_1, \bar{k}) \equiv 1 \pmod{2}$. Thus $\text{Arf}(L) \neq \text{Arf}(L')$, completing the proof.

PROOF OF THE THEOREM. If L, L' , and l are given as in Figure 1, then we write $L=L' \oplus l$ and also $L'=L \oplus l$. Continuing this, we can write $L=L_1 \oplus L_2 \oplus \dots \oplus L_m$ (here we omit parentheses), where L_j is a trivial link ($j=1, 2, \dots, m$) [7]. We define $d(L)$ to be the minimum number of such m 's ($d(L) \geq 1$).

We will induct on $d(L)$. If $d(L)=1$, L is a trivial link.

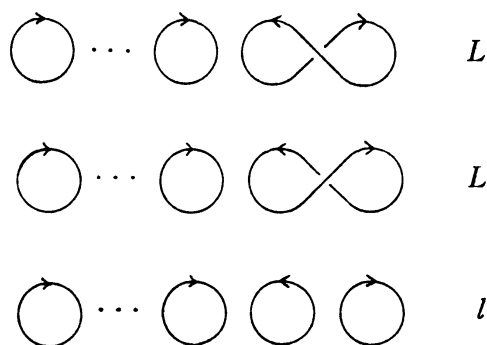


Figure 3.

Figure 3 and a simple induction will show that $I(L)=(\sqrt{2})^{\#(L)-1}$ in this case, while $\text{Arf}(L)=0$.

Now suppose that the theorem is proved for every link L' with $d(L') < m$ and consider a link L with $d(L)=m$. We may assume that L, L' , and l are as in Figure 1 and that $d(L') < m$ and $d(l) < m$. There are two cases.

Case I. Suppose that $\#(L)=\#(L')=\#(l)-1$.

(A) First assume that L is proper. Then L' is also proper. If l is proper, then $\text{Arf}(L)=\text{Arf}(L')=\text{Arf}(l)$. So from the inductive hypothesis $I(L)=\sqrt{2} \cdot \pm(\sqrt{2})^{\#(l)-1} - (\pm(\sqrt{2})^{\#(L')-1}) = \pm(\sqrt{2})^{\#(L)-1}$ according to whether $\text{Arf}(L)$ is 0 or 1. If l is not proper, then from the above lemma $\text{Arf}(L) \neq \text{Arf}(L')$. So $I(L) = \mp(\sqrt{2})^{\#(L)-1}$ according to whether $\text{Arf}(L)$ is 1 or 0.

(B) Next assume that L is non-proper. Then L' is also non-proper. It is easily shown that l is non-proper and so $I(L)=0$.

Case II. Suppose that $\#(L)=\#(L')=\#(l)+1$.

(A) Assume that L is proper. Then L' is non-proper and l is proper. Since $\text{Arf}(L)=\text{Arf}(l)$, $I(L)=\sqrt{2} \cdot \pm(\sqrt{2})^{\#(l)-1} = \pm(\sqrt{2})^{\#(L)-1}$ according to whether $\text{Arf}(L)$ is 0 or 1.

(B) Assume that L is non-proper. If L' is proper, then l is proper and $\text{Arf}(L')=\text{Arf}(l)$. So $I(L)=\sqrt{2} \cdot I(l) - I(L')=0$. If L' is non-proper, then it is easily proved that l is also non-proper and so $I(L)=0$.

Now the proof is complete.

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