

ON A RECENT ALLOTMENT OF PROBABILITIES TO
OPEN AND CLOSED SENTENCES

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Probabilities, though frequently allotted to closed sentences, have rarely been allotted to open ones. The recent scheme by Kemeny, Mirkil, Snell, and Thompson in *Finite Mathematical Structures* for allotting probabilities to sentences of the form ' $f(x) = a$ ' is therefore of considerable interest.¹ It has, however, a shortcoming which I should like to discuss here and, possibly, remedy.

Let ' f ' be a functional constant, ' x ' an individual variable, and ' a ' an individual constant; let U be the (finite) set of values of ' x '; and let A be the subset of U whose members satisfy ' $f(x) = a$ '. Kemeny *et al.* then take the probability of ' $f(x) = a$ ' to be $m(A)$, where $m(A)$ is the measure (in some appropriate sense of the word 'measure') of A .² Their scheme is attractive enough and mirrors to some extent what mathematicians understand by the probability of a set.³ Kemeny *et al.* are careful, of course, to restrict it to open sentences of the form ' $f(x) = a$ '. Consider, however, a closed sentence of the kindred form ' $f(b) = a$ ', where ' b ' is an individual constant. Since ' $f(b) = a$ ' does not contain any occurrence of ' x ', it would normally be held to be satisfied by every member of U when true, by none when false. One would accordingly expect Kemeny *et al.* to take the probability of ' $f(b) = a$ ' to be 1 when ' $f(b) = a$ ' is true, 0 when ' $f(b) = a$ ' is false. Yet in their scheme for allotting probabilities to closed sentences, a scheme I shall go into below, they let the probability of a closed sentence equal 1 only when the sentence is logically true, 0 only when it is logically false.⁴ ' $f(b) = a$ ' being neither logically true nor logically false, its probability must therefore differ by that scheme from either one of 1 and 0, a disturbing enough result.

The difficulty becomes even more acute when the calculus, call it C , to whose sentences probabilities are allotted is a simple applied predicate calculus of the first order with identity.

Assume indeed that a set D of individuals has been singled out as the domain of C , a member of D paired with each individual constant W of C as the individual designated by W , and a class of ordered n -tuples of members of D paired with each n -adic predicate constant F of C as the extension of

F. Assume next that an assignment *Asst* of members of *D* to the individual variables of *C* is said to satisfy a sentence *S* of *C* under the following circumstances:

D1. (a) Let *S* be of the form $F(W_1, W_2, \dots, W_n)$, where *F* is an *n*-adic predicate constant and W_1, W_2, \dots, W_n are *n* individual constants or variables. If the ordered *n*-tuple made up of the members of *D* respectively designated by or assigned by *Asst* to W_1, W_2, \dots, W_n belongs to the extension of *F*, then *Asst* satisfies *S*;

(b) Let *S* be of the form $W_1 = W_2$, where W_1 and W_2 are two individual constants or variables. If the members of *D* respectively designated by or assigned by *Asst* to W_1 and W_2 are the same, then *Asst* satisfies *S*.

(c) Let *S* be of the form $\sim (S')$. If *Asst* does not satisfy *S'*, then *Asst* satisfies *S*.

(d) Let *S* be of the form $(S') \supset (S'')$. If *Asst* does not satisfy *S'* or *Asst* satisfies *S''*, then *Asst* satisfies *S*.

(e1) Let *S* be of the form $(\forall W) (S')$, where *W* is an individual variable and *W* is not free in *S'*. If *Asst* satisfies *S'*, then *Asst* satisfies *S*.

(e2) Let *S* be of the form $(\forall W) (S')$, where *W* is an individual variable and *W* is free in *S'*. If *Asst* satisfies *S'* and every assignment of members of *D* to the individual variables of *L* which is like *Asst* except for the member of *D* it assigns to *W* also satisfies *S'*, then *Asst* satisfies *S*.⁵ Assume finally that probabilities are allotted to the sentences of *C* as follows:

D2. Let *S* be a sentence of *C*. Then the probability of *S* equals $m(Asst_S)$, where $Asst_S$ is the set of assignments of members of *D* to the individual variables of *C* which satisfy *S* and $m(Asst_S)$ is the measure (in some appropriate sense of the word 'measure') of $Asst_S$.

It is clear that if the sentence *S* in D1-D2 is allowed to be closed as well as open, if a closed sentence *S* of *C* is taken, as usual, to be true when *S* is satisfied by all assignments of members of *D* to the individual variables of *C*, and if a closed sentence *S* of *C* is taken, as usual again, to be false when *S* is not true, then the probability of a closed sentence *S* of *C* will automatically equal 1 or 0, when *S* is true, 0 when *S* is false. If, on the other hand, the sentence *S* in D1-D2 is presumed to be open, then other probabilities besides 1 and 0 may consistently be allotted to the closed sentences of *C*. Note, however, that with the sentence *S* in D1-D2 thus presumed to be open, the probability of $(W = W) \supset (S)$, where *W* is an individual variable of *C* and *S* is a closed sentence of *C*, will nonetheless equal 1 or 0, 1 when *S* is true, 0 when *S* is false.⁶ One could therefore not allot probabilities other than 1 and 0 to *S* without thereby allotting different probabilities to *S* and $(W = W) \supset (S)$. But *S* and $(W = W) \supset (S)$ are logically equivalent. One could therefore not allot probabilities other than 1 and 0 to *S* without thereby allotting different probabilities to logically equivalent sentences, a distressing result.⁷

I would accordingly recommend that (1) the Kemeny *et al.* scheme for allotting probabilities to open sentences of the form ' $f(x) = a$ ' be made to cover as well closed sentences of the form ' $f(b) = a$ '. I would also

recommend, when it comes to allotting probabilities to the sentences of C , that (2) the sentence S in $D1$ - $D2$ be suffered to be open as well as closed, or that (3) the probability of a closed sentence S of C be taken to be that of the open sentence $(W = W) \supset (S)$ we just considered.⁸ In all three cases the probability of a closed sentence, be it of the form ' $f(b) = a$ ' or any other, would equal 1 or 0, and in the last two cases the requirement that equivalent sentences be allotted equal probabilities would be met.

The scheme Kemeny *et al.* used to allot probabilities to closed sentences is roughly as follows. A (finite) set U of so-called logical possibilities is assumed to be given; a subset A of U , consisting of all the members of U which are not precluded (so to speak) by a closed sentence S , is then paired with S as the so-called truth-set of S ; the probability of S is finally taken to be $m(A)$, where $m(A)$ is the measure (in some appropriate sense of the word 'measure') of the truth-set A of S . It is clear from my brief description of the scheme that the probability of a closed sentence S , where S is neither logically true nor logically false, must be other than 1 or 0.⁹

This allotment of probabilities, reminiscent of Carnap's *Logical Foundations of Probabilities*, is at variance with the one I just recommended.¹⁰ It might nonetheless be retained alongside mine if a distinction which Kemeny *et al.* ignore were drawn, that between the truth-value of a sentence and an estimate of the truth-value of a sentence. The probability I have just allotted to a closed sentence S coincides with the truth-value of S . The one Kemeny *et al.* allot, on the other hand, to such a sentence has all the earmarks of an estimate of the truth-value of S . But there is room in inductive logic for estimates of truth-values as well as for truth-values. The Kemeny *et al.* allotment of probabilities to closed sentences might thus be retained—*subject to the above reinterpretation*—alongside mine. It might even be extended in one interesting direction.

Whereas, indeed, the probability Kemeny *et al.* allot to a closed sentence has all the earmarks of an estimate of a truth-value, the one they allot to an open sentence has all these of a truth-value in a generalized sense of the word 'truth-value'. As a matter of fact, when the members of the domain D of such a calculus as C are paired in a one-to-one fashion with the individual constants of C , the probability allotted by $D2$ to an open sentence S of C proves to be a weighted average of the truth-values (in the usual sense of the word) of the so-called instances of S .¹¹ There being room in inductive logic for estimates of the truth-values of open sentences as well as of closed ones, the Kemeny *et al.* allotment of probabilities to closed sentences might therefore be extended to cover open sentences as well.

I cannot recount here the various steps to be taken in carrying out that extension.¹² Once they are taken, however, a sentence S (be it closed or open) comes to be allotted two probabilities: one, the truth-value of S either in the traditional sense or in a generalized sense of the word, has a statistical flavor of its own and, for that reason, might be dubbed *the statistical probability of S*; the other, an estimate of the truth-value of S , has

an inductive flavor of its own and, for that reason, might be dubbed *the inductive probability of S*.

Kemeny *et al.* ignore, I said, the distinction I urged above between a truth-value and an estimate of a truth-value. They also ignore as a result the distinction I urge here between a statistical probability and an inductive one.¹³ They do so, however, at the price, as I hope to have shown, of allotting different probabilities to the two equivalent sentences S and $(W = W) \supset (S)$.¹⁴

NOTES

[1] See John G. Kemeny, Hazelton Mirkil, J. Laurie Snell, and Gerald L. Thompson, *Finite Mathematical Structures*, Englewood Cliffs, N. J.: Prentice-Hall, Inc., 1959, chapter 2, section 4, and chapter 3, section 1.

[2] For a definition of the word 'measure', see *Finite Mathematical Structures*, p. 113.

[3] Kemeny *et al.* transfer indeed to the defining condition of a set, understood here as an open sentence, the probability which mathematicians normally allot to the set of elements satisfying that condition.

[4] See *Finite Mathematical Structures*, chapter 2, section 3, and chapter 3, sections 1 and 2.

[5] Some additions, which the reader can easily supply, should be made to the text if C contained functional constants as well as individual and predicate constants.

[6] A new definition of the phrase ' S is true', which the reader can easily supply, would of course be required here.

[7] The same result would hold if C were a simple applied predicate calculus of the first order without identity. The open sentence $((F(W)) \supset (F(W))) \supset (S)$, where F is, say, the alphabetically first predicate constant of C , could then serve in lieu of $(W = W) \supset (S)$.

[8] A different scheme for allotting what I shall call below statistical probabilities will be found in the author's "On chances and estimated chances of being true," *Revue Philosophique de Louvain*, vol. 57 (Mai 1959), pp. 225-239.

[9] See references in footnote 4.

[10] See Rudolph Carnap, *Logical Foundations of Probability*, Chicago: The University of Chicago Press, 1950, chapter V.

[11] This point is made in the author's "On chances and estimated chances of being true," theorem T3.12, for a special weighting of the truth-values of the instances of S . It can be made for all weighting of the truth-values in question once D3.1(b), (c1), and (c2) are suitably generalized.

[12] See the author's "On chances and estimated chances of being true," where the extension in question is carried out for a family of calculi C .

[13] In *A Philosopher Looks at Science*, Princeton, N. J.: D. Van Nostrand Company, Inc., 1959, chapter 4, Kemeny distinguishes between two kinds of probabilities. The distinction he draws there, however, is not reproduced in *Finite Mathematical Structures*.

[14] This paper was given by title at the 1960 International Congress for Logic, Methodology and Philosophy of Science, Stanford University, Stanford, California.

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