

## A STUDY OF $\mathcal{Z}$ MODAL SYSTEMS

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In [10], Sobociński has shown that the addition to various S4-extensions of Zeman's formula

$$\mathbf{Z1} \quad LMp \cdot LMq \rightarrow (M(p \cdot q) \rightarrow LM(p \cdot q))$$

generates a new family that he refers to as the  $\mathcal{Z}$  modal systems. In this paper a completeness proof is given for the system  $\mathbf{Z1} = \mathbf{S4} + \mathbf{Z1}$ , and the finite model property is established. Since the system is finitely axiomatisable its decidability follows. Furthermore it is shown that each  $\mathcal{Z}$  modal system is the intersection of S5 with some system from family  $\mathcal{K}$ .

In the field of S4,  $\mathbf{Z1}$  is inferentially equivalent to

$$\mathbf{Z2} \quad L(LMp \rightarrow MLp) \vee L(Mq \rightarrow LMq)$$

the formula added to S4.4 by Schumm to obtain the system now called S4.9 (cf. Sobociński [9], p. 361)

- |               |   |  |
|---------------|---|--|
| (1)           | $LM(p \vee q) \cdot LM(\sim p \vee q) \rightarrow (M((p \vee q) \cdot (\sim p \vee q)) \rightarrow LM((p \vee q) \cdot (\sim p \vee q)))$ |  |
|               |   | $\mathbf{Z1}, p/p \vee q, q/\sim p \vee q$ |
| (2)           | $(p \vee q) \cdot (\sim p \vee q) \leftrightarrow q$  | <b>PC</b>                                  |
| (3)           | $Mq \rightarrow M((p \vee q) \cdot (\sim p \vee q))$  | (2), C2                                    |
| (4)           | $LM((p \vee q) \cdot (\sim p \vee q)) \rightarrow LMq$  | (2), C2                                    |
| (5)           | $LMp \rightarrow LM(p \vee q)$  | C2   |
| (6)           | $\sim MLp \rightarrow LM(\sim p \vee q)$  | C2   |
| (7)           | $LMp \cdot \sim MLp \rightarrow (Mq \rightarrow LMq)$   | (1), (3), (4), (5), (6), <b>PC</b>         |
| (8)           | $(LMp \rightarrow MLp) \vee (Mq \rightarrow LMq)$   | (7), <b>PC</b>                             |
| (9)           | $M(LMp \rightarrow MLp) \vee L(Mq \rightarrow LMq)$   | (8), C2                                    |
| (10)          | $M(LMp \rightarrow MLp) \rightarrow (LMp \rightarrow MLp)$  | <b>S4</b>                                  |
| (11)          | $(LMp \rightarrow MLp) \rightarrow L(LMp \rightarrow MLp)$  | $\mathbf{Z1}, \mathbf{S4}$ (cf. [1])       |
| $\mathbf{Z2}$ | $L(LMp \rightarrow MLp) \vee L(Mq \rightarrow LMq)$   | (9), (10), (11), <b>PC</b>                 |

This shows that  $\mathbf{Z1}$  contains the system  $\mathbf{S4} + \mathbf{Z2}$ . We shall subsequently establish the converse in two different ways.

Definitions and discussion of the model-theoretic concepts used below are given in Segerberg [5], [6], and [7] (cf. also the Metatheorem of [1]).

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**Proposition 1:** *A is a theorem of S4 + Z2 if and only if it is verified by every S4-frame satisfying*

$$\forall x \forall y ((xRy \rightarrow yRx) \vee \exists z (yRz \cdot \forall w (zRw \rightarrow z = w))) \quad (a)$$

*Proof: Necessity.* We leave it to the reader to check that any S4-frame satisfying (a) verifies Z2.

*Sufficiency.* Let  $A$  be any wff not derivable in S4 + Z2. Then  $A$  is false at some point  $t$  in the canonical model for S4 + Z2, and hence false at  $t$  in the submodel  $\mathbf{u}$  generated from the canonical model by  $t$ . Let  $\Psi$  be the closure under modalities of the set of all subwff of  $A$ , and  $\mathbf{u}'$  a Lemmon filtration of  $\mathbf{u}$  through  $\Psi$ . Then  $\mathbf{u}'$  is finite, reflexive, and transitive, and hence is an S4-frame (Segherberg [5], Section 3, and [6], Chapter I, Theorem 7.6). Furthermore by the Filtration Theorem ([7], p. 303)  $A$  is false in  $\mathbf{u}'$  at  $[t]$ . It remains only to show that  $\mathbf{u}'$  satisfies (a).

Let  $[x]$  be any point in  $\mathbf{u}'$  and suppose there is some point  $[y]$  in  $\mathbf{u}'$  such that

$$[x] R' [y] \text{ and not } [y] R' [x] \text{ in } \mathbf{u}' \quad (b)$$

We have to prove the second disjunct of (a).

Since  $\mathbf{u}'$  is finite, reflexive, and transitive there is some final cluster  $Y$  in  $\mathbf{u}'$  that either contains  $[y]$  or succeeds  $[y]$  ([5], p. 19). From (b) it follows that  $[x]$  precedes  $Y$  and so  $Y$  is a non-initial final cluster. Now if we can show that  $Y$  is a simple final cluster, i.e., it consists of a single element with no alternative except itself, then the proof will be complete.

To get a contradiction we suppose  $Y$  is a proper cluster and therefore contains at least two distinct elements, say  $[z]$  and  $[w]$ . Since  $[z] \neq [w]$  there is some wff  $B$  in  $\Psi$  such that

$$B \text{ is true in } \mathbf{u}' \text{ at } [z], \text{ but false at } [w] \quad (c)$$

Since the relation  $R'$  is universal over  $Y$ , it follows from (c) that

$$MB \text{ is true, and } LB \text{ is false, in } \mathbf{u}' \text{ at every point in } Y \quad (d)$$

Now if  $zRu$  in  $\mathbf{u}$ ,  $[z] R' [u]$  in  $\mathbf{u}'$ . But  $Y$  is final, so  $[u] \in Y$ , whence by (d),  $MB$  is true but  $LB$  is false in  $\mathbf{u}'$  at  $[u]$ . But  $MB$  and  $LB$  are in  $\Psi$ , so by the Filtration Theorem  $MB$  is true and  $LB$  is false in  $\mathbf{u}$  at  $u$ . This shows that  $LMB$  is true and  $MLB$  is false in  $\mathbf{u}$  at  $z$ , hence

$$(LMB \rightarrow MLB) \text{ is false at } z \quad (e)$$

Now the model  $\mathbf{u}$  is transitive and generated by  $t$ , and so

$$tRu, \text{ for all } u \text{ in } \mathbf{u} \quad (f)$$

Then (e) and (f) together yield

$$L(LMB \rightarrow MLB) \text{ is false in } \mathbf{u} \text{ at } t \quad (g)$$

As in [6], Chapter II, Lemma 2.1 we can construct a Boolean combination  $C$  of members of  $\Psi$  such that

$$C \text{ is true in } \mathbf{u} \text{ at } u \text{ iff } [u] \notin Y \quad (h)$$

Now from (b),  $[x] \notin Y$ , so by (h),  $C$  is true in  $\mathbf{u}$  at  $x$ . Thus by (f)

$$MC \text{ is true at } t \tag{i}$$

Now if  $zRu$ ,  $[u] \in Y$  (see above) so by (h),  $C$  is false at  $u$ . Hence  $MC$  is false at  $z$ , so by (f)

$$LMC \text{ is false at } t \tag{j}$$

The reflexivity of  $\mathbf{u}$ , together with (i) and (j) show that

$$L(MC \rightarrow LMC) \text{ is false at } t \tag{k}$$

But (g) and (k) contradict the fact that every substitution-instance of  $Z2$  is true at every point in  $\mathbf{u}$ , and in particular at  $t$ . This ends our proof.

It is an easy matter to check that  $Z1$  is verified by every  $S4$ -frame satisfying condition (a) above, and so by Proposition 1 is derivable in  $S4 + Z2$ . This, with our earlier result establishes that  $Z1 = S4 + Z2$ .

Corollary ([5], section 3)  $Z1$  has the finite model property.

*Proof.* The model  $\mathbf{u}'$  in Proposition 1 has at most  $2^{4n}$  elements, where  $n$  is the number of subwff of the non-theorem  $A$  that it falsifies ( $Z1$  has the same fourteen distinct modalities as  $S4$ ).

If the element  $[t]$  of Proposition 1 is contained in a final cluster then, since  $[t]$  generates  $\mathbf{u}'$ , the underlying frame of  $\mathbf{u}'$  consists of a single cluster and therefore verifies  $S5$ . If, on the other hand,  $[t]$  is not contained in a final cluster then every final cluster in  $\mathbf{u}'$  is non-initial and therefore by the proof of Proposition 1 is simple. But a finite  $S4$ -frame in which every final cluster is simple is a frame for  $K1$  (Seegerberg [5]). Thus  $\mathbf{u}'$  is either an  $S5$ -model or a  $K1$ -model and since it falsifies the non- $Z1$ -theorem  $A$ ,  $A$  must be a non-theorem of either  $S5$  or  $K1$ , and hence of  $S5 \cap K1$ . But  $Z1$  is contained in both  $S5$  and  $K1$ , and so we conclude that  $Z1 = S5 \cap K1$ .

This intersection result, and the fact that  $Z1$  is a consequence in  $S4$  of  $Z2$ , may alternatively be deduced from Theorem 2 of Halldén [2] (cf. also Theorem 2 of Kripke [3]). This theorem shows that if  $S$  is a system containing  $PC$  and having modus ponens as its only primitive rule of inference, and if  $A$  and  $B$  are two formulae with no propositional variable in common, then

$$(S + (A \vee B)) = (S + A) \cap (S + B).$$

Now Simons [8] has given an axiomatisation of  $S4$  that has modus ponens as its only primitive rule. Furthermore, the same system results when the same extra axioms are added to this basis and to Lewis' basis for  $S4$ . This observation, and Halldén's theorem show immediately that  $S4 + Z2$  is the intersection of  $S4 + L(Mq \rightarrow LMq)$  and  $S4 + L(LMp \rightarrow MLP)$ , i.e., of  $S5$  and  $K1$ . But  $Z1$  is easily derivable in both  $S5$  and  $K1$ , whence it is derivable in  $S4 + Z2$ .

**Proposition 2.** *Let  $S$  be a system contained in  $S5$  and obtained by adding some axioms  $S$  (which may be a finite conjunction of wff) to  $S4$ . Then  $(S + Z1) = S5 \cap (S + K1)$ .*

*Proof:* Since  $Z1$  is a theorem of  $S5$ , and deducible from  $K1$ , it is immediate that  $(S + Z1) \subseteq S5 \cap (S + K1)$ . For the converse we use again an axiomatisation of  $S4$  with modus ponens the only primitive rule. Letting  $B = L(Mq \rightarrow LMq)$  and  $K = L(LMp \rightarrow MLP)$  we see by Halldén's theorem that the system  $S5 \cap (S + K1)$  can be axiomatised as  $S4 + (B \vee (S.K))$ , where  $B$  and  $(S.K)$  have no variable in common. Now if  $A$  is a theorem of this system then by the Deduction Theorem there are finitely many instances  $(B_i \vee (S_i . K_i))$ ,  $(1 \leq i \leq n)$ , such that  $(B_1 \vee (S_1 . K_1)) \dots (B_n \vee (S_n . K_n)) \rightarrow A$  is derivable in  $S4$ . Hence  $(B_1 \vee S_1) . (B_1 \vee K_1) \dots (B_n \vee S_n) . (B_n \vee K_n) \rightarrow A$  and so  $S_1 . (B_1 \vee K_1) \dots S_n . (B_n \vee K_n) \rightarrow A$  is derivable in  $S4$ . But each conjunct in the antecedent of this last formula is derivable in  $S + Z1$ , and therefore so is  $A$ . QED

From Proposition 2 we read off the following connections between  $Z$  and  $K$  modal systems.

$$\begin{aligned} Z1 &= S5 \cap K1 \\ Z2 &= S5 \cap K1.2 \\ Z3 &= S5 \cap K1.1 \\ Z4 &= S5 \cap K2 \\ Z5 &= S5 \cap K2.1 \\ Z6 &= S5 \cap K3 \\ Z7 &= S5 \cap K3.1 \\ Z8 &= S5 \cap K3.2 \\ Z9 &= S5 \cap K4 = S4.9^1 \end{aligned}$$

There is some interest in the relationship between  $Z2$  and the formula

$$\Delta \quad (LMp \rightarrow MLP) \vee (Mq \rightarrow LMq)$$

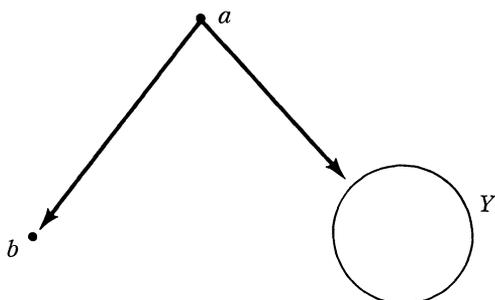
which is obviously derivable from  $Z2$ . In showing that  $Z2$  was a consequence of  $Z1$  we first derived  $\Delta$  from  $Z1$  and then used the  $S4.01$  axiom

$$\Gamma4 \quad (LMp \rightarrow MLP) \rightarrow L(LMp \rightarrow MLP)$$

which is also derivable from  $Z1$  in  $S4$ , to obtain  $Z2$ . Thus  $\Delta$  and  $Z2$  are equivalent in the field of  $S4.01$  and therefore in the field of every proper  $S4$ -extension except  $S4.02$  and  $S4.04$  (cf. [1]). However in  $S4$  itself  $\Delta$  is weaker than  $Z2$ . This follows from consideration of the matrix  $\mathfrak{M}11$  of Sobociński [9], p. 350. By the methods of Lemmon [4], section 5,  $\mathfrak{M}11$  is seen to be the matrix representation of the four-element reflexive frame that can be displayed graphically as

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1. A similar, and independent, proof that the systems  $Z9$  and  $S4.9$  are identical has been obtained by K. Fine (cf. *Notre Dame Journal of Formal Logic*, vol. XIII (1972), p. 118).



with the circle  $Y$  denoting a two-element cluster. Now  $(LMp \rightarrow MLp)$  cannot be falsified at  $a$  or  $b$ , since each of these points has an alternative (*viz*  $b$ ) that has no alternative except itself (Segerberg [5], p. 18). On the other hand  $(Mq \rightarrow LMq)$  cannot be falsified on the cluster  $Y$ . Thus the frame, and hence the matrix must verify  $\Delta$ . But this frame does not satisfy condition (a) of Proposition 1 and so there will be some assignment on it that falsifies  $Z2$ . In fact for  $p = 6, q = 6, L(M6 \rightarrow LM6) \vee L(LM6 \rightarrow ML6) = L(5 \rightarrow 13) \vee L(13 \rightarrow 16) = L9 \vee L4 = 9 \vee 12 = 9$ . Thus the addition to  $S4$  of  $\Delta$  yields a new system properly contained in  $Z1$ .

Concerning other matrices of Sobociński [9] we make the following comments:

- 1)  $\mathfrak{M}11$  also rejects  $S4.01$  (*cf.* [1]). But  $\mathfrak{M}8$  verifies  $S4.01$  while rejecting  $\Delta$ . For  $p = 6, q = 4, (M4 \rightarrow LM4) \vee (LM6 \rightarrow ML6) = (4 \rightarrow 8) \vee (1 \rightarrow 8) = 5 \vee 8 = 5$ . Thus  $S4 + \Delta$  is independent of  $S4.01$ .
- 2)  $\mathfrak{M}5$  rejects  $S4.02$  and  $S4.04$  but verifies  $K1$  and hence  $\Delta$ . But  $\mathfrak{M}8$  verifies  $S4.02$  and  $S4.04$ , so these two systems are independent of  $S4 + \Delta$ .
- 3) Since  $\mathfrak{M}11$  verifies  $S4.04$  and  $\Delta$  but rejects  $Z2$  we obtain two new systems,  $S4.02 + \Delta$  and  $S4.04 + \Delta$ , properly contained in  $Z3$  and  $Z2$  respectively. The matrix  $\mathfrak{M}4$  verifies  $S4.02$  and  $\Delta$  but rejects  $S4.04$ , showing that  $S4.02 + \Delta$  is a proper subsystem of  $S4.04 + \Delta$ .

I conjecture that  $S4 + \Delta$  is characterised by the class of  $S4$ -frames satisfying

$$\forall x \forall y ((xRy \rightarrow yRx) \vee \exists z (xRz \cdot \forall w (zRw \rightarrow z = w))),$$

or alternatively by the class of finite  $S4$ -frames in which every non-final cluster is succeeded by at least one simple final cluster.

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