

CONCERNING THE PROPER AXIOM FOR S4.04
AND SOME RELATED SYSTEMS

R. I. GOLDBLATT

This paper examines the group of modal axioms covered by the general schema

$$(X) \quad Xp \rightarrow (p \rightarrow Lp)$$

where X is an affirmative modality of S4. Familiarity is assumed with the properties of maximal-consistent sets of wff, and with the post-Henkin method of completeness proofs. Soundness proofs are left to the reader throughout.

(X) yields seven cases:

Case 1. Zeman's S4.04 axiom

$$L1 \quad LMLp \rightarrow (p \rightarrow Lp) \quad \text{cf. [5], p. 250}$$

In the field of S4, L1 is equivalent to

$$L2 \quad p \rightarrow L(MLp \rightarrow p)$$

That L1 is a consequence of L2 is easy to see. For the converse we have¹

- | | |
|---|---------------------------|
| (1) $MLp \rightarrow ML(MLp \rightarrow p)$ | C2 |
| (2) $\sim MLp \rightarrow (MLp \rightarrow p)$ | PC |
| (3) $\sim LMMLp \rightarrow ML(MLp \rightarrow p)$ | (2), C2 |
| (4) $\sim MLp \rightarrow ML(MLp \rightarrow p)$ | (3), S4, PC |
| (5) $ML(MLp \rightarrow p)$ | (1), (4), PC |
| (6) $LML(MLp \rightarrow p)$ | (5), T ⁰ |
| (7) $LML(MLp \rightarrow p) \rightarrow ((MLp \rightarrow p) \rightarrow L(MLp \rightarrow p))$ | L1, $p/MLp \rightarrow p$ |
| (8) $(MLp \rightarrow p) \rightarrow L(MLp \rightarrow p)$ | (6), (7), PC |
| (9) $p \rightarrow L(MLp \rightarrow p)$ | (8), PC |

We now present a semantic analysis that distinguishes L1 and L2 in

1. This proof is due to Professor G. E. Hughes.

systems weaker than S4. Let K be the modal logic whose rule of inference is

$$\text{Necessitation: } \frac{A}{LA}$$

and whose sole axiom is $L(p \rightarrow q) \rightarrow (Lp \rightarrow Lq)$.

For the definition of a (normal) K -model, cf. [1] pp. 56-60 (the system is called $T(C)$ in that paper). If S is any normal extension of K we define $\mathcal{P}_S = (W_S, R, V)$ to be the *canonical model* for S , where

$$\begin{aligned} W_S &= \{x \mid x \text{ is an } S\text{-maximal set of wff}\} \\ \forall x, y \in W_S, xRy &\text{ iff } \{A \mid LA \in x\} \subseteq y \\ V(p, x) &= 1 \text{ iff } p \in x, \text{ for all propositional variables } p. \end{aligned}$$

\mathcal{P}_S falsifies every non-theorem of S , so to show that S is complete with respect to a class of models satisfying a certain condition, it suffices to show that \mathcal{P}_S satisfies that condition. It is well known that S4 is characterized by the class of K -models for which R is reflexive and transitive.

Proposition 1: *If S is a normal extension of $KL2$, then \mathcal{P}_S satisfies*

$$(xRy . x \neq y) \rightarrow \forall w(yRw \rightarrow wRy) \tag{a}$$

Proof. Suppose $x, y \in W_S, xRy$ and $x \neq y$. Then there is some wff A such that $A \in x$ and $A \notin y$. Let w be such that yRw . If $LB \in w$ then $L(A \vee B) \in w$, by the K -theorem $LB \rightarrow L(A \vee B)$. Now yRw , so $ML(A \vee B) \in y$. But $A \vee B \in x$ so by **L2**, $L(ML(A \vee B) \rightarrow A \vee B) \in x$ and hence $ML(A \vee B) \rightarrow A \vee B \in y$. Thus $A \vee B \in y$ and since $A \notin y, B \in y$. We have therefore shown that $\{B \mid LB \in w\} \subseteq y$, i.e., wRy .

Proposition 2: *If S is a normal extension of $KL1$, then \mathcal{P}_S satisfies*

$$(xRy . x \neq y) \rightarrow \exists z(xRz . \forall w(zRw \rightarrow wRy)) \tag{b}$$

Proof. Suppose xRy and $x \neq y$. Let

$$z_0 = \{A \mid LA \in x\} \cup \{LMB \mid B \in y\}$$

If z_0 is not S -consistent then there are wff A_i such that $LA_i \in x$ ($1 \leq i \leq n$) and LMB_i such that $B_i \in y$ ($1 \leq i \leq m$) for which

$$\vdash_S \sim(A_1 \dots A_n . LMB_1 \dots LMB_m)$$

and hence

$$\vdash_S A \rightarrow \sim(LMB_1 \dots LMB_m) \text{ where } A = A_1 \dots A_n$$

By the K -theorem $\sim(LMp . LMq) \rightarrow \sim LM(p . q)$ we then deduce

$$\vdash_S A \rightarrow \sim LMB \text{ where } B = B_1 \dots B_m$$

Since $x \neq y$ there is some wff C such that $C \in x$ and $\sim C \in y$. Using the K -theorem $\sim(LMp \rightarrow \sim LM(p . q))$ we now deduce

$$\vdash_S A \rightarrow \sim LM(B . \sim C)$$

and so

$$\vdash_5 LA \rightarrow L \sim LM(B. \sim C)$$

i.e.,

$$\vdash_5 LA \rightarrow LML \sim (B. \sim C)$$

But

$$\vdash_K LA \leftrightarrow (LA_1 \dots LA_n)$$

so $LA \in x$, which gives $LML \sim (B. \sim C) \in x$ (1)

Now $\sim C \notin x$, so $(B. \sim C) \notin x$, hence $\sim(B. \sim C) \in y$. (2)

In the presence of L1, (1) and (2) together yield $L \sim (B. \sim C) \in x$ and so $\sim(B. \sim C) \in y$ (since xRy).

But $B \in y$ and $\sim C \in y$, so $(B. \sim C) \in y$, which contradicts the PC-consistency of y . The upshot of all this is that z_0 must be S -consistent and so, by Lindenbaum's Lemma, may be extended to an S -maximal set z . Since $\{A \mid LA \in x\} \subseteq z_0 \subseteq z$, we have xRz . Finally, suppose zRw . We want to show wRy . But if $B \notin y$ then $\sim B \in y$ and so $LM \sim B \in z$. But zRw so $M \sim B \in w$, hence $\sim M \sim B \notin w$, i.e., $LB \notin w$. This shows that wRy , and the proof is complete.

Proposition 3: *If \mathcal{P} is an S4-model, then \mathcal{P} satisfies condition (a) iff \mathcal{P} satisfies condition (b).*

Proof. If \mathcal{P} satisfies (a) then putting $y = z$ it is immediate that (b) is satisfied. Conversely, suppose xRy and $x \neq y$. We want to show that if yRt then tRy . From (b) we deduce

$$\exists z(xRz \cdot \forall w(zRw \rightarrow wRy)) \quad (1)$$

Since R is reflexive, zRz , and so by (1) zRy . Then if yRt , we have zRt by the transitivity of R . Using (1) again we deduce tRy .

From Proposition 1 it follows that S4.04 is complete with respect to the class of S4-models that satisfy condition (a). The axiom corresponding to (a) is L2. L1 corresponds (Proposition 2) to a rather more complex condition that reduces (Proposition 3) to (a) in S4-models. Our discussion would seem to indicate then that from a model-theoretic stand-point L2 is the "right" axiom for S4.04.

Case 2. Sobociński's K4 axiom

$$P1 \quad MLMp \rightarrow (p \rightarrow Lp) \quad \text{cf. [4], p. 349}$$

Proposition 4: *If S is a normal extension of K_{P1} then \mathcal{P}_S satisfies*

$$(xRy \cdot x \neq y) \rightarrow \forall z(xRz \rightarrow \exists w(zRw \cdot \forall t(wRt \rightarrow t = y))) \quad (c)$$

Proof. If xRy and $x \neq y$ then there is some wff C such that $C \in x$ and $\sim C \in y$. Let xRz and put

$$w_0 = \{A \mid LA \in z\} \cup \{LB \mid B \in y\}$$

If w_0 is not S -consistent then reasoning as in Proposition 2 we find there are wff A, B such that $LA \in z, B \in y$ and

$$\vdash_5 A \rightarrow \sim L(B \sim C).$$

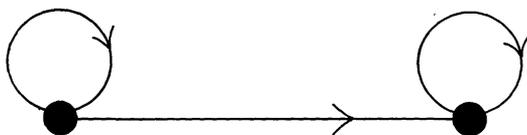
Thus

$$\vdash_5 LA \rightarrow L \sim L(B \sim C)$$

and so $LM \sim (B \sim C) \in z$. But xRz and so $MLM \sim (B \sim C) \in x$.

Using the axiom **P1**, a contradiction then follows exactly as in Proposition 2. Thus w_0 is S -consistent and may be extended to an S -maximal w such that zRw . Now let wRt . Then if $B \in y, LB \in w$, and since $wRt, B \in t$. Hence $y \subseteq t$ and so by the maximality of $y, y = t$.

K4 is known to be characterized by Lewis and Langford's Group II matrix (cf. [6] p. 349) and so it has two-element models that may be represented graphically as follows: (cf. [1] p. 63).



If \mathcal{P} is a reflexive model that satisfies (c) then it may be seen by elementary reasoning that if y and z are distinct from x , then $(xRz \cdot xRy) \rightarrow (x = y)$. If \mathcal{P} is also transitive and connected in the sense of [1] p. 193, then it has the form indicated above.

Case 3. When X is the improper modality, the resulting axiom is (equivalent to) $p \rightarrow Lp$ corresponding to the model condition $\forall x \forall y (xRy \rightarrow x = y)$, cf. [1] p. 214.

Case 4. $X = L$ gives a substitution-instance of a **PC**-tautology.

Case 5. $X = M$ gives

$$\mathbf{X1} \quad Mp \rightarrow (p \rightarrow Lp)$$

Proposition 5: If S is a normal extension of **KX1** then \mathcal{P}_S satisfies

$$(xRy \cdot x \neq y) \rightarrow \forall z (xRz \rightarrow z = y) \tag{d}$$

Proof. There is some wff A such that $A \in x$ and $A \notin y$. If $B \in z, A \vee B \in z$ and so $M(A \vee B) \in x$. But $A \vee B \in x$ so by **X1**, $L(A \vee B) \in x$. Thus $A \vee B \in y$, and since $A \notin y, B \in y$. This shows $z \subseteq y$, which is enough to prove $z = y$.

Given R reflexive, corresponding to $Lp \rightarrow p$, then condition (d) above reduces to

$$\forall x \forall y (xRy \rightarrow x = y)$$

which is Case 3 above. A syntactic proof is straightforward.

Case 6. $X = ML$ gives the **S4.4** axiom

$$\mathbf{R1} \quad MLp \rightarrow (p \rightarrow Lp)$$

whose corresponding model condition is

$$(xRy . x \neq y) \rightarrow \forall z(xRz \rightarrow zRy).$$

A proof is given in [3].

Case 7. $X = LM$ yields

H2 $LMp \rightarrow (p \rightarrow Lp)$

Proposition 6: *If S is a normal extension of $KH2$ then \bar{P}_S satisfies*

$$(xRy . x \neq y) \rightarrow \exists z(xRz . \forall w(zRw \rightarrow w = y)) \quad (e)$$

Proof. By a similar method to **Proposition 4** we may show

$$z_0 = \{A \mid LA \in x\} \cup \{LB \mid B \in y\}$$

is S -consistent and may be extended to an S -maximal z with the required properties. When R is reflexive, condition (e) reduces to

$$(xRy . x \neq y) \rightarrow \forall z(yRz \rightarrow y = z)$$

which was shown in [3] to be the characteristic model condition for the **K1.2** axiom

H1 $p \rightarrow L(Mp \rightarrow p)$

In [4] a proof is given that **H1** and **H2** are equivalent in the field of **S4**. The following proof shows that the equivalence holds in the field of **S2**:

- | | |
|---|-----------------------------------|
| (1) $LM(Mp \rightarrow p) \rightarrow ((Mp \rightarrow p) \rightarrow L(Mp \rightarrow p))$ | H2 $p/Mp \rightarrow p$ |
| (2) $LM(Mp \rightarrow p)$ | S2 , cf. [2], p. 140, 22.8 |
| (3) $(Mp \rightarrow p) \rightarrow L(Mp \rightarrow p)$ | (1), (2), PC |
| (4) $p \rightarrow L(Mp \rightarrow p)$ | (3), PC |

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*Victoria University of Wellington
Wellington, New Zealand*