## ON THE LATTICE OF FACES OF THE n-CUBE<sup>1</sup>

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1. We give a simple axiom system for the lattice  $L_n$  of faces of the *n*-cube, which is independent of dimension, and we construct a partition of the lattice into a minimum number of chains, or *Dilworth partition*. This partition turns out to enjoy some notable symmetries.

We use the representation of the faces of an *n*-cube as *signed subsets* of an *n*-set, say of the set  $\{1, 2, \ldots, n\}$ . A signed subset  $A_{\sigma} = (A_1, A_2)$  is an ordered pair of disjoint subsets, where  $A_1$  is called the *positive part*, and  $A_2$  the *negative-part*. If  $B_{\sigma}$  is also a signed set, write  $A_{\sigma} \leq B_{\sigma}$  when  $A_1 \supseteq B_1$  and  $A_2 \supseteq B_2$ . Add a minimum element 0—the *improper face*—to the ordered set of signed sets, thereby making it a lattice  $L_n$ . The maximum element I of  $L_n$  is the signed set  $I = (\phi, \phi)$ . We use the terms "face" and "signed set" interchangeably.

On the lattice of signed subsets one defines diagonals  $\Delta(A_{\sigma}, \cdot)$ . For a given face  $A_{\sigma}$ , such a diagonal is a function defined on the segment  $[0, A_{\sigma}]$  of  $L_n$ , and  $\Delta(A_{\sigma}, B_{\sigma}) = C_{\sigma}$ , where  $C_1 = (A_1, B_2)$  and  $C_2 = (A_2, B_1)$ . On the improper face one sets  $\Delta(A_{\sigma}, 0) = 0$ . Geometrically, the diagonal  $\Delta(A_{\sigma}, \cdot)$  associates to each face contained in  $A_{\sigma}$  the unique opposite face inside the face  $A_{\sigma}$ . When  $A_{\sigma} = I$ , the diagonal  $\Delta(I, \cdot)$ , written  $\Delta(\cdot)$ , is a cubical analog of complementation in a Boolean algebra.

- 2. Main Theorem. Let L be a finite lattice with minimum 0 and maximum I. For every  $x \neq 0$ , let  $\Delta_x$  be a function defined on the segment [0, x] and taking values in [0, x]. Assume: (1) if  $y \leq x$ , then  $\Delta_x(\Delta_x(y)) = y$ ; (2) if  $a \leq b \leq x$ , then  $\Delta_x(a) \leq \Delta_x(b)$ ; (3) if a < x, then  $a \land \Delta_x(a) = 0$ ; (4) let a < x and b < x. Then the following two conditions are equivalent:  $\Delta_x(a) \land b < x$  and  $a \land b = 0$ . Then L is isomorphic to the lattice of faces of an n-cube for some n, and conversely.
- 3. A symmetric Dilworth partition. By Dilworth's theorem there exists a partition of  $L_n$  into  $\lfloor n/3 \rfloor$  chains. We explicitly describe one such partition, one that is invariant under the main diagonal  $\Delta$ .

Consider a signed subset as a sequence  $u_1u_2\cdots u_n$  whose digits  $u_i$  range over the alphabet  $\{x, 0, 1\}$ . Set  $u_i = x$  if the element i is unsigned; otherwise,

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 $u_i = 0$  or 1 according to whether  $u_i$  belongs to the positive or to the negative parts. We call 0 and 1 significant digits in contrast to x. A vertex is a sequence  $u_1u_2 \cdot \cdot \cdot u_n$  where no  $u_i = x$ .

The chains issuing from the vertices are constructed one step at a time, upwards from the vertices, according to the following rule:

- (1) Change the last two digits of the vertex  $u_1u_2\cdots u_n$ , according to the following table:  $11 \longrightarrow 1x$ ,  $00 \longrightarrow 0x$ ,  $10 \longrightarrow x0$ ,  $01 \longrightarrow x1$ . Denote the resulting face by  $v_1v_2\cdots v_n$ .
- (2) Change the two rightmost adjacent significant digits of  $v_1v_2\cdots v_n$  by the same rule.
- (3) Repeat until the resulting face contains no two adjacent significant digits.
- (4) Now the preceding algorithm assigns to chains issuing from vertices all faces  $u_1u_2\cdots u_n$  for  $u_1=u_2=x$ , where  $u_j$  is a significant digit for some j>2, and no two successive digits are significant.
- (5) For faces containing three successive digits x00, x01, x10 or x11 use the following bracketing algorithm: (a) Bracket x01 to (x01), etc., by inserting brackets at the appropriate places. In the resulting bracketed face, two digits are successive if they are successive in the sequence obtained after all bracketed digits have been removed. (b) Repeat the bracketing algorithm (a), for successive digits, until no unbracketed digits can be bracketed, and obtain a completely bracketed face. (c) On the set of completely bracketed faces having the same bracketed digits and the same bracket structure, remove bracketed digits and proceed as in steps (1)–(4), and then reinsert the bracketed digits in their original places. One obtains in this way a partition into chains of such a set. (d) The only remaining faces to assign to chains are those which have only x's in their unbracketed sequence, which is nonempty.

Let  $u = u_1 u_2 \cdots u_n$  be one of these remaining faces. If u = I, the maximum element, assign it arbitrarily to any chain, say, the chain containing  $0xx \cdots x$ . If  $u \neq I$ , let  $u_k$  be the first significant digit, and let  $u_m$  be the first unbracketed digit, which must be an x. Then assign u to the chain containing  $v_1 v_2 \cdots v_n$  where  $v_m = u_k$  and  $v_i = u_i$  otherwise.

The preceding technique for constructing Dilworth partitions can be extended to families of k disjoint subsets, ordered similarly to the case k=2 of the cube, and even to certain sublattices of the lattice of partitions of a set. These constructions will be presented elsewhere.

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