

## On the solution of $x^2 + dy^2 = m$

By Julius Magalona BASILLA<sup>\*)</sup>

Department of Mathematics, Sophia University

7-1, Kioicho, Chiyoda-ku, Tokyo 102-8554

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**Abstract:** A simple proof of the validity of Cornacchia's algorithm for solving the diophantine equation  $x^2 + dy^2 = m$  is presented. Furthermore, the special case  $d = 1$  is solved completely.

**Key words:** Cornacchia's algorithm; quadratic forms; diophantine equations.

In 1908, G. Cornacchia gave an algorithm for solving the diophantine equation  $x^2 + dy^2 = 4p$ , for  $p$  prime (cf. [1]). The same algorithm can be used to solve the diophantine equation

$$(1) \quad x^2 + dy^2 = m$$

where  $1 \leq d < m$ ,  $m$  may not be prime. The algorithm is briefly described as follows:

1. Put  $r_0 = m$  and  $r_1^2 \equiv -d \pmod{m}$ , where  $0 \leq r_1 \leq (m/2)$ .
2. Using Euclidean algorithm, compute  $r_{i+2} \equiv r_i \pmod{r_{i+1}}$  recursively until we arrive at  $r_k^2 < m$ .
3. If  $(m - r_k^2/d)$  is a square integer, say  $s^2$ , we get the solution  $(r_k, s)$ .

A proof of the validity of this algorithm relying on Diophantine Approximation was given by F. Morain and J.-L. Nicolas (cf. [2]). In this paper, we give a simpler proof. Moreover, we claim that if  $r_0 = m$ ,  $r_1^2 \equiv -1 \pmod{m}$ ,  $1 \leq r_1 < (m/2)$ , then  $m = r_k^2 + r_{k+1}^2$ , when  $r_{k-1}^2 > m > r_k^2$ .

**1. A simple proof.** Let  $d$  and  $m$  be integers such that  $1 \leq d < m$ .

**Lemma 1.** *If  $(x_0, y_0)$  is a primitive solution of (1), then there exists an integer  $t$ ,  $0 < t < m$ ,  $t^2 \equiv -d \pmod{m}$  such that  $(x_0, y_0) \in \langle (m, 0), (t, 1) \rangle_{\mathbf{Z}}$ .*

*Proof.* Clearly,  $\gcd(y_0, m) = 1$ . Choose  $t$ ,  $0 < t < m$  such that  $y_0 t \equiv x_0 \pmod{m}$ . We have  $0 \equiv x_0^2 + dy_0^2 \equiv y_0^2(t^2 + d) \pmod{m}$ . Thus,  $t^2 \equiv -d \pmod{m}$ . Also, for some integer  $l$ ,  $(x_0, y_0) = l(m, 0) + y_0(t, 1) \in \langle (m, 0), (t, 1) \rangle_{\mathbf{Z}}$ .  $\square$

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<sup>\*)</sup>On study leave from the Department of Mathematics, University of the Philippines, Diliman, Quezon City Philippines.

Define  $L_t := \langle (m, 0), (t, 1) \rangle_{\mathbf{Z}}$ . Clearly, if the solutions  $(x_0, y_0)$  and  $(-x_0, -y_0)$  are in  $L_t$ , the solutions  $(-x_0, y_0)$  and  $(x_0, -y_0)$  are in  $L_{m-t}$ . This suggests that, to find all the primitive solutions of (1), it is enough to consider all square roots  $t$  of  $-d$  modulo  $m$ , where  $1 \leq t \leq (m/2)$  and compute for all the vectors  $(x, y)$  in  $L_t$  with  $x^2 + dy^2 = m$ . We will discuss how to find these vectors.

**Lemma 2.** *Let  $\vec{u} = \langle u_1, u_2 \rangle$ ,  $\vec{v} = \langle v_1, v_2 \rangle$  be generators of a lattice such that  $0 \leq u_2$ ,  $0 \leq v_2$ ,  $|v_1| < |u_1|$ ,  $u_1 v_1 < 0$ . Then the vector  $\vec{w} = \langle w_1, w_2 \rangle$  with the least  $w_2$  such that  $0 < w_2$ ,  $|w_1| < |v_1|$  is given by  $\vec{w} = \vec{u} + q\vec{v}$ , where  $q = \lfloor -(u_1/v_1) \rfloor$ . Moreover,  $\langle \vec{u}, \vec{v} \rangle_{\mathbf{Z}} = \langle \vec{v}, \vec{u} + q\vec{v} \rangle_{\mathbf{Z}}$ .*

For the proof (cf. [3]).

**Lemma 3.** *Let  $r_0, r_1$  be positive integers. Construct the finite sequences  $\{r_i\}$ ,  $\{q_i\}$ ,  $\{P_i\}$ , and  $\{Q_i\}$  as follows:*

$$\begin{aligned} r_i &= q_i r_{i+1} + r_{i+2}, & q_i &= \left\lfloor \frac{r_i}{r_{i+1}} \right\rfloor \\ P_{-1} &= 0; P_0 = 1; & P_{i+1} &= q_i P_i + P_{i-1} \\ Q_{-1} &= 1; Q_0 = 0; & Q_{i+1} &= q_i Q_i + Q_{i-1} \end{aligned}$$

for  $0 \leq i \leq n-1$ , where  $r_n = \gcd(r_0, r_1)$  and  $r_{n+1} = 0$ . Then, for  $0 \leq i \leq n$ ,

$$(2) \quad r_0 = P_i r_i + P_{i-1} r_{i+1},$$

$$(3) \quad r_{i+1} = (-1)^i (P_i r_1 - Q_i r_0).$$

The proof is by induction on  $i$ .

Putting  $r_0 = m$  and  $r_1^2 \equiv -d \pmod{m}$ , from (3) we get,

$$(4) \quad r_i^2 + dP_{i-1}^2 \equiv 0 \pmod{m}$$

for  $0 \leq i \leq n+1$ .

**Proposition 4.** *Let  $r_0 = m$  and  $r_1 = t$  where  $t^2 \equiv -d \pmod{m}$ . Construct the sequence  $\{r_i\}$ ,*

$\{P_i\}$  as in Lemma 3. The diophantine equation (1) has a solution in  $L_t$  if and only if  $dP_{k-1}^2 < m$ , when  $r_{k-1}^2 > m > r_k^2$ .

*Proof.* ( $\Rightarrow$ ) Let  $(x_0, y_0)$  be a solution of (1) in  $L_t$ . Without loss of generality, we can change the sign of  $x_0$  and assume  $y_0 > 0$ . If  $r_{k-1}^2 > m > r_k^2$ , we have  $|x_0|^2 < x_0^2 + dy_0^2 = m < r_{k-1}^2$ . That is  $|x_0| < |r_{k-1}|$ .

Put  $\vec{u}_0 = (-m, 0)$  and  $\vec{u}_1 = (t, 1)$ . By Lemma 2, the vector  $\vec{u}_2 = (-r_2, q)$  is the vector  $\vec{w} = (w_1, w_2)$  with the least  $w_2 > 0$  such that  $|w_1| < |t|$ . Note that the pair  $\vec{u}_1, \vec{u}_2$  satisfies the premises of Lemma 2. Thus we can apply the Lemma repeatedly. Inductively, we can show that the vector  $\vec{w} = (w_1, w_2)$  with the least  $w_2 > 0$  such that  $|w_1| < r_{i-1}$  is  $\vec{u}_i = ((-1)^{i-1}r_i, P_{i-1})$ .

In particular, the vector  $\vec{w} = (w_1, w_2)$  with the least  $w_2$  such that  $|w_1| < |(-1)^{k-2}r_{k-1}|$  is  $((-1)^{k-1}r_k, P_{k-1})$ . Since  $|x_0| < r_{k-1}$ , it follows that  $P_{k-1} \leq y_0$  and hence,  $dP_{k-1}^2 \leq dy_0^2 < m$ .

( $\Leftarrow$ ) From (4), we have  $dP_{k-1}^2 \equiv -r_k^2 \equiv m - r_k^2 \pmod{m}$ . We get the solution  $(r_k, P_{k-1})$ .  $\square$

We now consider the special case  $d = 1$ .

**Proposition 5.** Let  $t^2 \equiv -1 \pmod{m}$ ,  $0 < t < (m/2)$ . Set  $r_0 = m$  and  $r_1 = t$  and construct the finite sequence  $\{r_i\}$ ,  $r_i = q_i r_{i+1} + r_{i+2}$ , for  $0 \leq i \leq n-1$ , where  $r_0 > r_1 > \dots > r_n = 1 > r_{n+1} = 0$ . If  $r_{k-1}^2 > m > r_k^2$  then  $m = r_k^2 + r_{k+1}^2$ .

*Proof.* Construct the sequence  $\{P_i\}$  as in Lemma 3. We get the following relations.

$$\begin{aligned} m = r_0 &= P_n r_n + P_{n-1} r_{n+1} = P_n \\ r_{n-1} &= q_{n-1} r_n + r_{n+1} = q_{n-1} \geq 2 \\ m = P_n &= q_{n-1} P_{n-1} + P_{n-2} > 2P_{n-1}, \end{aligned}$$

since  $n \geq 2$  and  $P_{n-2} \neq 0$ . Also

$$\begin{aligned} 1 = r_n &= (-1)^{n-1} (P_{n-1} r_1 - Q_{n-1} r_0) \\ &\equiv (-1)^{n-1} P_{n-1} t \pmod{m} \\ t &\equiv (-1)^n P_{n-1} \pmod{m}. \end{aligned}$$

It follows that  $n$  must be even say  $n = 2k$ , and  $t = P_{n-1}$ . Inductively, we can show  $P_{n-i} = r_i$  for  $0 \leq i \leq n+1$ . From Lemma 3, we have  $m = r_0 = P_k r_k + P_{k-1} r_{k+1} = r_k^2 + r_{k+1}^2$ . Observe that  $r_{k-1}^2 = (q_{k-1} r_k + r_{k+1})^2 > r_k^2 + r_{k+1}^2 = m > r_k^2$ . This proves the proposition.  $\square$

**2. A proof of uniqueness.** We now show that for a particular square root  $t$  modulo  $m$  of  $-d$ ,

the only primitive solution of (1) belonging to the lattice  $\langle (m, 0), (t, 1) \rangle_{\mathbf{Z}}$  is  $(r_k, P_{k-1})$ . In the case  $d = 1$ , because of symmetry we have two:  $(r_k, r_{k+1})$  and  $(r_{k+1}, r_k)$ .

Assume that  $(r_k, P_{k-1})$  is the solution of (1) obtained by applying Cornacchia's algorithm. When  $d = 1$ , it is clear that  $dP_{k+1}^2 \geq m$ .

**Lemma 6.** If  $d > 1$ , then  $dP_k^2 \geq m$ .

*Proof.* Suppose that  $dP_k^2 < m$ . As in the proof of Proposition 4, we get  $r_{k+1}^2 + dP_k^2 = m$ . From (2), we have

$$\begin{aligned} m = r_0 &= P_k r_k + P_{k-1} r_{k+1} \\ &\leq \frac{r_k^2 + P_k^2}{2} + \frac{r_{k+1}^2 + P_{k-1}^2}{2} \\ &< \frac{r_k^2 + dP_{k-1}^2 + r_{k+1}^2 + dP_k^2}{2} = m. \end{aligned}$$

We have a contradiction.  $\square$

**Proposition 7.** Let  $d > 1$  and  $t^2 \equiv -d \pmod{m}$ . If  $(x, y)$ ,  $y > 0$  is a solution of (1) such that  $(x, y) \in L_t$ , then  $(x, y) = ((-1)^{k-1}r_k, P_{k-1})$ .

*Proof.* Since  $x^2 + dy^2 = m$  it follows that  $|x|^2 < m$  and  $dy^2 < m$ . Thus,  $|x| < r_{k-1}$ . If  $|x| < r_k < r_{k-1}$ , then  $y \geq P_k$ , by minimality of  $P_k$ . Thus  $dy^2 \geq dP_k^2 \geq m$ .

If  $|x| > r_k$ , since  $|x| < r_{k-1}$ , then  $y \geq P_{k-1}$  by minimality of  $P_{k-1}$ . Then  $x^2 + dy^2 > r_k^2 + dP_{k-1}^2 = m$ .  $\square$

**Proposition 8.** Let  $t^2 \equiv -1 \pmod{m}$ . If  $(x, y)$ ,  $y > 0$  is a solution of  $x^2 + y^2 = m$  such that  $(x, y) \in L_t$ , then  $(x, y) = ((-1)^{k-1}r_k, r_{k+1})$  or  $(x, y) = ((-1)^k r_{k+1}, r_k)$ .

The proof is similar to that of Proposition 7.

This proves that using Cornacchia's algorithm on all the square-root  $t$  modulo  $m$  of  $-d$ , we can find all the primitive solutions of (1).

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