

SOLUTIONS

No problem is ever permanently closed. Any comments, new solutions, or new insights on old problems are always welcomed by the problem editor.

45. [1992, 88] *Proposed by Mohammad K. Azarian, University of Evansville, Evansville, Indiana.*

Let α , β , and γ be the direction angles of a vector R . Without using Lagrange multipliers, show that

$$\cos \alpha + \cos \beta + \cos \gamma - 2 \cos \alpha \cos \beta \cos \gamma \leq \sqrt{2}.$$

Solution by the proposer.

If we let

$$\cos \alpha = \frac{a}{\sqrt{2}}, \quad \cos \beta = \frac{b}{\sqrt{2}}, \quad \cos \gamma = \frac{c}{\sqrt{2}},$$

then we need to show that

$$a + b + c - abc \leq 2.$$

Thus, it is enough to prove that

$$4(4 - (a + b + c - abc)^2) \geq 0.$$

But, using

$$a^2 + b^2 + c^2 = 2(\cos^2 \alpha + \cos^2 \beta + \cos^2 \gamma) = 2,$$

we have

$$\begin{aligned} 4(4 - (a + b + c - abc)^2) &= (2 - 2ab)(2 - 2ac)(2 - 2bc) + 4(abc)^2 \\ &= (c^2 + (a - b)^2)(b^2 + (a - c)^2)(a^2 + (b - c)^2) + 4(abc)^2. \end{aligned}$$

This completes the proof. Note that equality holds if R is perpendicular to one of the axes and is making a 45-degree angle with each of the other two axes.

46. [1992, 88] *Proposed by Robert E. Kennedy and Curtis Cooper, Central Missouri State University, Warrensburg, Missouri.*

What is the leading digit of 2^{googol} ?

Solution by Kevin Robinson, Messiah College, Grantham, Pennsylvania and the proposers.

The leading digit of 2^{googol} is 2. To see this, we observe that for any positive integer n ,

$$2^n = 10^{\lfloor n \log 2 \rfloor} \cdot 10^{\{n \log 2\}},$$

where $\lfloor \cdot \rfloor$ denotes the greatest integer function, $\{ \cdot \}$ denotes the fractional part operator, and $\log 2$ denotes the base 10 logarithm of 2. By the above identity, we have the leading digit of 2^n is the units digit of $10^{\{n \log 2\}}$. Using the computer algebra system DERIVE, the first 120 digits of $\log 2$ are

0.30102999566398119521
 37388947244930267681
 89881462108541310427
 46112710818927442450
 94869272521181861720
 40684477191430995379 .

Therefore,

$$10^{\{\text{googol} \log 2\}} = 10^{\{10^{100} \log 2\}} = 10^{.40684\dots} = 2. \dots$$

Thus, the leading digit of 2^{googol} is 2.

One incorrect solution was also received.

47. [1992, 88] *Proposed by Russell Euler, Northwest Missouri State University, Maryville, Missouri.*

Find all the solutions of

$$(x - 1)x(x + 1)(x + 2) = -1.$$

Solution I by Bob Prielipp, University of Wisconsin - Oshkosh, Oshkosh, Wisconsin; Robert L. Doucette, McNeese State University, Lake Charles, Louisiana; Kevin Robinson, Messiah College, Grantham, Pennsylvania; Joseph E. Chance, University of Texas - Pan American, Edinburg, Texas; Frank Flanigan, San Jose State University, San Jose, California; Kandasamy Muthuvel, University of Wisconsin - Oshkosh, Oshkosh, Wisconsin; Herta T. Freitag, Roanoke, Virginia; J. Sriskandarajah, University of Wisconsin Center, Richland Center, Wisconsin and the proposer.

We may construct the following collection of equivalent equalities.

$$(x - 1)x(x + 1)(x + 2) = -1$$

$$x^4 + 2x^3 - x^2 - 2x + 1 = 0$$

$$(x^2 + x - 1)^2 = 0$$

$$x = \frac{-1 + \sqrt{5}}{2} \text{ (double root) or } x = \frac{-1 - \sqrt{5}}{2} \text{ (double root).}$$

Solution II by Mohammad K. Azarian, University of Evansville, Evansville, Indiana and J. Sriskandarajah, University of Wisconsin Center, Richland Center, Wisconsin.

From the substitution

$$x = \frac{u - 1}{2},$$

we obtain

$$(x - 1)x(x + 1)(x + 2) = \left(\frac{u - 3}{2}\right)\left(\frac{u - 1}{2}\right)\left(\frac{u + 1}{2}\right)\left(\frac{u + 3}{2}\right) = -1.$$

This implies that

$$u^4 - 10u^2 + 25 = 0.$$

Thus, $u^2 = 5$ is a double root. Hence, $u = \pm\sqrt{5}$. Therefore,

$$x = \frac{\pm\sqrt{5} - 1}{2},$$

where both of these roots are of multiplicity two.

Generalized Solution I by Joseph E. Chance and Donald P. Skow (jointly), University of Texas - Pan American, Edinburg, Texas.

Let k be an integer. Find all the solutions of

$$(x - k)(x - k + 1)(x - k + 2)(x - k + 3) = -1.$$

Let

$$r = (x - k + 1)(x - k + 2).$$

Then

$$r = (x - k)^2 + 3(x - k) + 2.$$

Hence,

$$r - 2 = (x - k)^2 + 3(x - k) = (x - k)(x - k + 3).$$

Thus, the original equation becomes

$$r(r - 2) = -1.$$

This implies $r = 1$. Hence,

$$(x - k)^2 + 3(x - k) + 1 = 0$$

and so

$$x - k = \frac{-3 \pm \sqrt{5}}{2}.$$

Therefore,

$$x = \frac{-3 \pm \sqrt{5} + 2k}{2}.$$

If $k = 1$, we have the solution to problem 47.

Generalized Solution II by Frank Flanigan, San Jose State University, San Jose, California.

Let

$$P(x) = (x + a_1)(x + a_2)(x + b_1)(x + b_2) + k.$$

Suppose $a_1, a_2, b_1, b_2,$ and k are elements of a field F . Assume

$$a_1 + a_2 = b_1 + b_2 = s.$$

Now, if $a = a_1a_2, b = b_1b_2$ and $y = x^2 + sx$, then

$$\begin{aligned} P(x) &= (x^2 + sx + a)(x^2 + sx + b) + k \\ &= (y + a)(y + b) + k \\ &= y^2 + (a + b)y + (ab + k) \\ &= (y - \sigma)(y - \tau) \\ &= (x^2 + sx - \sigma)(x^2 + sx - \tau) \end{aligned}$$

It is important that the roots σ, τ of the quadratic in y will lie in the original scalar field F of characteristic $\neq 2$ if and only if the discriminant

$$\Delta = \Delta_y = (a - b)^2 - 4k$$

is a square in F . This is decisive for the irreducibility of $P(x)$ in the ring $F[x]$ of polynomials.

Problem 47 can now be solved by taking $F = \mathbb{R}$ and $a_1 = 0, a_2 = 1, b_1 = -1, b_2 = 2,$ and $k = 1$.

Also solved by N. J. Kuenzi, University of Wisconsin - Oshkosh, Oshkosh, Wisconsin; J. Sriskandarajah, University of Wisconsin Center, Richland Center, Wisconsin (second solution) and Donald P. Skow, University of Texas - Pan American, Edinburg, Texas. Frank Flanigan commented that his general solution to $P(x) = 0$ can be extended to the case when $P(x)$ involves the product of six, rather than four, suitable linear factors.

48. [1992, 89] *Proposed by Alvin Beltramo (student), Central Missouri State University, Warrensburg, Missouri.*

A standard deck of 52 cards is shuffled and then the cards are displayed, one at a time. Before each card is displayed, a person with a perfect memory guesses what each card is. How many cards can this person expect to guess correctly?

Solution I by Jenhua Tao, Central Missouri State University, Warrensburg, Missouri; Robert L. Doucette, McNeese State University, Lake Charles, Louisiana; K. L. D. Gunawardena, University of Wisconsin - Oshkosh, Oshkosh, Wisconsin; N. J. Kuenzi, University of Wisconsin - Oshkosh, Oshkosh, Wisconsin and Kevin Robinson, Messiah College, Grantham, Pennsylvania.

Define X_i to be the Bernoulli random variable for the i th card, i.e., $X_i = 0$ if the person guesses the i th card incorrectly and $X_i = 1$ if the person guesses the i th card correctly. Then

$$X = X_1 + X_2 + \cdots + X_{52}$$

is the number of cards this person guesses correctly. Thus,

$$E(X) = E(X_1) + E(X_2) + \cdots + E(X_{52}).$$

Since this person has perfect memory,

$$P(X_i = 1) = \frac{1}{52 - i + 1}$$

and

$$E(X_i) = \frac{1}{52 - i + 1}.$$

So,

$$E(X) = \frac{1}{52} + \frac{1}{51} + \frac{1}{50} + \cdots + 1 \doteq 4.538.$$

Generalized Solution I by K. L. D. Gunawardena, University of Wisconsin - Oshkosh, Oshkosh, Wisconsin.

Let $X_i = 1$ if the i th card is guessed correctly and $X_i = 0$ otherwise. Then

$$S = \sum_{i=1}^{52} X_i$$

denotes the total number of cards guessed correctly. Now suppose the person does not have a perfect memory, but he can remember the last k cards displayed prior to the displaying of the i th card. Thus,

$$P(X_i = 1) = \begin{cases} \frac{1}{53-i}, & 1 \leq i \leq k; \\ \frac{1}{52-k}, & k+1 \leq i \leq 52. \end{cases}$$

Therefore,

$$E(S) = \sum_{i=1}^k \frac{1}{53-i} + \sum_{i=k+1}^{52} \frac{1}{52-k} = 1 + \sum_{i=1}^k \frac{1}{53-i}.$$

Note that $k = 51$ corresponds to a person with a perfect memory.

Generalized Solution II by the proposer. Suppose there are n cards, numbered 1 to n in the deck. Then the expected number of correct guesses by a person with a perfect memory is

$$\sum_{k=1}^n \frac{1}{k}.$$

But,

$$\sum_{k=1}^n \frac{1}{k} = \log n + \gamma + O\left(\frac{1}{n}\right),$$

where \log is the natural logarithm and γ is Euler's constant. Thus, the expected number of correct guesses for a person with perfect memory, in a deck of n cards numbered 1 to n is asymptotic to $\log n$.

Also solved by the proposer. One incorrect solution was also received. Kevin Robinson and the proposer noted that the exact value is

$$\frac{14063600165435720745359}{3099044504245996706400}.$$