

# ON THE DIAMETER OF A GRAPH

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A graph  $G_n$  consists of a set of  $n$  nodes, some pairs of which are joined by a single edge. The degree of a node  $x$  is the number  $d(x)$  of edges joining it to other nodes. A graph is connected if it cannot be represented as the union of disjoint smaller graphs. The diameter of a connected graph  $G_n$  is the least integer  $k$  such that any pair of nodes of  $G_n$  can be joined by a sequence of at most  $k$  edges, consecutive ones of which have a node in common.

In what follows,  $n$  and  $k$  denote integers satisfying the inequality  $n - 1 > k > 2$ . Let  $g(n, k)$  be the least integer  $r$  such that if  $d(x) \geq r$  for every node  $x$  of a connected graph  $G_n$ , then the diameter of  $G_n$  is at most  $k$ . The object of this note is to prove the following result (here  $[x]$  denotes the largest integer not exceeding  $x$ ).

**THEOREM.**

$$g(n, k) = \begin{cases} \left[ \frac{n}{t} \right] & \text{if } k = 3t - 4, \\ \left[ \frac{n-1}{t} \right] & \text{if } k = 3t - 3, \\ \left[ \frac{n-2}{t} \right] & \text{if } k = 3t - 2. \end{cases}$$

*Proof.* Let us suppose that  $k = 3t - 3$ , where  $t$  is an integer greater than 1. We first show that

$$g(n, k) \leq \left[ \frac{n-1}{t} \right].$$

To accomplish this we assume the contrary, namely, that there exists a connected graph  $G_n$ , the degree of each of whose nodes is at least  $\left[ \frac{n-1}{t} \right]$  and whose diameter exceeds  $k$ . From this we shall deduce a contradiction.

It is easy to see that by introducing additional edges, we can transform  $G_n$  into a graph  $G'_n$  of the form

$$\langle 1 \rangle \text{ --- } \langle a_1 \rangle \text{ --- } \langle a_2 \rangle \text{ --- } \cdots \text{ --- } \langle a_{3t-3} \rangle \text{ --- } \langle 1 \rangle$$

(here  $\langle j \rangle$  denotes a (nonempty) complete subgraph with  $j$  nodes and  $\binom{j}{2}$  edges, and two nodes in different indicated subgraphs are joined by an edge if and only if the two subgraphs are adjacent in our representation), in such a way that the diameter of  $G'_n$  is  $k + 1$  and the degree of each node of  $G'_n$  is at least  $\left[ \frac{n-1}{t} \right]$ .

Consider such a graph. Clearly, the degree of the first node on the left is  $a_1$ . The degree of any node in  $\langle a_3 \rangle$  is  $a_2 + (a_3 - 1) + a_4$ . Continuing in this fashion,

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we obtain the following set of inequalities, since by hypothesis the degree of every node is at least  $\left\lceil \frac{n-1}{t} \right\rceil$ .

$$\begin{aligned} a_1 &\geq \left\lceil \frac{n-1}{t} \right\rceil, \\ a_2 + (a_3 - 1) + a_4 &\geq \left\lceil \frac{n-1}{t} \right\rceil, \\ a_5 + (a_6 - 1) + a_7 &\geq \left\lceil \frac{n-1}{t} \right\rceil, \\ &\dots \\ a_{3t-7} + (a_{3t-6} - 1) + a_{3t-5} &\geq \left\lceil \frac{n-1}{t} \right\rceil, \\ a_{3t-3} &\geq \left\lceil \frac{n-1}{t} \right\rceil. \end{aligned}$$

Upon adding and simplifying, we conclude that

$$n - t - a_{3t-4} \geq t \left\lceil \frac{n-1}{t} \right\rceil > n - 1 - t.$$

But this is false, since  $a_{3t-4} \geq 1$ ; hence, the above upper bound for  $g(n, k)$  is established.

If we can exhibit a connected graph  $G_n$  the degree of each of whose nodes is at least  $\left\lceil \frac{n-1}{t} \right\rceil - 1$  and whose diameter is  $k+1$ , then it will follow that

$$g(n, k) \geq \left\lceil \frac{n-t}{t} \right\rceil.$$

To do this, consider a graph of diameter  $k+1$  whose structure is as illustrated above. Let

$$a_2 = a_4 = a_5 = a_7 = a_8 = \dots = a_{3t-5} = a_{3t-4} = 1.$$

If the remaining nodes can be distributed in such a way that  $a_1$  and  $a_{3t-3}$  are at least as large as  $\left\lceil \frac{n-1}{t} \right\rceil - 1$ , and each of the remaining  $t-2$  integers  $a_i$  is positive and at least as large as  $\left\lceil \frac{n-1}{t} \right\rceil - 2$ , then it is clear that the degree of each node in the resulting connected graph will be at least  $\left\lceil \frac{n-1}{t} \right\rceil - 1$ . Upon counting the number of nodes needed in the  $3t-1$  classes, we see that these requirements can be satisfied if

$$(2t-1) + 2 \left( \left\lceil \frac{n-1}{t} \right\rceil - 1 \right) + (t-2) \left( \left\lceil \frac{n-1}{t} \right\rceil - 2 \right) = t \left\lceil \frac{n-1}{t} \right\rceil + 1 \leq n$$

when  $k \leq n - 4$ , or if  $3t - 1 \leq n$  when  $k = n - 3$  or  $n - 2$ . But these inequalities are certainly true, and hence the lower bound for  $g(n, k)$  holds.

The required result for the case  $k = 3t - 3$  now follows if we combine these bounds for  $g(n, k)$ . The proofs for the remaining two cases involve only minor modifications in the above argument, and therefore we omit them.

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