

INDEPENDENT NECESSARY CONDITIONS FOR FUNCTIONAL
COMPLETENESS IN m -VALUED LOGIC

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A function f is *functionally complete* in m -valued logic if the set of functions which can be defined explicitly from f is exactly the set of all functions of m -valued logic. A *Sheffer function* is a two-place functionally complete function. Post [1] and Webb [2], among others, have identified some Sheffer functions in m -valued logic. Martin [3] identified four properties (i.e., proper substitution, co-substitution, proper closing, and \vdash -closing), the absence of which are necessary conditions for functional completeness. In this paper, we will prove that co-substitution implies proper substitution; or with respect to our necessary conditions, if f does not have the proper substitution property, then f does not have the co-substitution property. Consequently, the co-substitution property can be discarded from our set of necessary conditions for functional completeness. Finally, we show that the remaining three necessary conditions are independent.

Theorem If $f(p, q)$ is a two-place function satisfying the co-substitution property, then $f(p, q)$ satisfies the proper substitution property as well.

Proof: Let $K = \{1, 2, 3, \dots, m\}$ be the set of m truth values, and D be a decomposition of K into \hat{m} disjoint non-empty classes, $2 \leq \hat{m} < m$. We will say $i \sim j (D)$ if i and j are elements of the same class, $i, j \in K$. Further, let \hat{D} be the decomposition of the two-dimensional space K^2 such that $(p, q) \sim (r, s) (\hat{D})$ if and only if $p \sim r (D)$ and $q \sim s (D)$. Let f satisfy the co-substitution law of Martin; that is, for any $h, i, j, k \in K$, whenever $f(h, i) \sim f(j, k) (D)$, then $h \sim j (D)$ or $i \sim k (D)$.

Assume there exist $(a, b) \sim (c, d) (\hat{D})$ such that $f(a, b) \not\sim f(c, d) (D)$. There are $\hat{m} - 1$ classes of \hat{D} , we will call them $C_1, C_2, \dots, C_{\hat{m}-1}$, such that if $(w_i, x_i) \in C_i$ then $a \not\sim w_i (D)$, $b \not\sim x_i (D)$, $c \not\sim w_i (D)$, and $d \not\sim x_i (D)$. Further, if $(w_i, x_i) \in C_i$ and $(w_j, x_j) \in C_j$, $i \neq j$, then $w_i \not\sim w_j (D)$ and $x_i \not\sim x_j (D)$. Since f satisfies co-substitution, $f(a, b) \not\sim f(w_1, x_1) (D)$ and $f(c, d) \not\sim f(w_1, x_1) (D)$. Further, $f(a, b) \not\sim f(w_2, x_2) (D)$, $f(c, d) \not\sim f(w_2, x_2) (D)$ and $f(w_1, x_1) \sim f(w_2, x_2) (D)$. Continuing, we reach the case that $f(w_{\hat{m}-1}, x_{\hat{m}-1})$

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cannot be specified without violating the co-substitution property. Therefore, our original assumption is contradicted; that is, that $(a, b) \sim (c, d) (\hat{D})$ and $f(a, b) \not\sim f(c, d) (D)$. But $(a, b) \sim (c, d) (\hat{D})$ implies $f(a, b) \sim f(c, d) (D)$ is exactly the proper substitution property. Q.E.D.

Remark: Obviously, if D is the trivial decomposition consisting of m classes, then every one-place function trivially satisfies co-substitution. However, as is well known, one-place functions cannot be functionally complete, so this case is not considered.

It remains to show that the three remaining properties, i.e., proper substitution, proper closing, and \dagger -closing are indeed independent. This, we easily do by means of examples. First, we review the definitions of \dagger -closing and proper closing from [3]. Let $\dagger(p)$ be a one-place function which satisfies the following:

$$\begin{aligned} \dagger^m(j) &= j, 1 \leq j \leq m \\ \dagger^i(j) &\neq j, 1 \leq i \leq m - 1, 1 \leq j \leq m. \end{aligned}$$

Then $f(p, q)$ is \dagger -closing if there exists a $\dagger(p)$ such that, for all i, j , there exists a k such that $f(\dagger^i(p), \dagger^j(p)) = \dagger^k(p)$. A function $f(p, q)$ is proper closing if some non-empty proper subset of the m truth values is closed under $f(p, q)$.

Example 1: $f_1(p, q)$ has the proper substitution property, but is not proper closing and is not \dagger -closing.

			q			
			1	2	3	4
$f_1:$	p	1	3	4	4	4
		2	3	4	3	3
		3	4	3	2	1
		4	4	3	2	1

Example 2: $f_2(p, q)$ is proper closing, but not \dagger -closing and does not have the proper substitution property.

			q			
			1	2	3	4
$f_2:$	p	1	3	2	1	1
		2	4	3	2	4
		3	1	2	4	3
		4	2	4	4	3

Example 3: $f_3(p, q)$ is \dagger -closing, but not proper closing and does not have the proper substitution property.

		q				
		1	2	3	4	
$f_3:$	p	1	2	4	3	1
	2	2	2	3	1	4
	3	1	1	3	4	2
	4	3	3	2	4	1

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