Investigations in Protothetic

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In this article I present some results of five years' research into Leśniewski's protothetic.¹ I outline deductions from the axiom A_n considerably shorter than those previously known (see [9]) and I derive the laws of implication from this axiom without using the rule of extensionality.² Since this paper can best be read in the light of articles by Professor Sobociński published in this Journal (see [8], [9], and [10]), I have largely adopted his conventions of symbolism, and the following symbols in particular:

- α The rule permitting definitions of new constants (see [8], pp. 58–59).
- β The rule for distributing quantifiers (see [8], p. 59).
- 0 An informal abbreviation for (u] . u'.
- 1 An informal abbreviation for '[u]. u := . [u] . u'.

From the axiom

$$A_n[pq] :: p \equiv q := \therefore [f] \therefore f(pf(p0)) :=: [r] : f(qr) :=: q \equiv p$$

we prove the following theorems:

D1 $[p] \cdot p \equiv \operatorname{As}(p)$ $[\alpha]$ L1 $[fp] \therefore f(pf(p0)) \cdot \equiv : [r] : f(\operatorname{As}(p)r) \cdot \equiv \cdot \operatorname{As}(p) \equiv p$ $[A_n q/\operatorname{As}(p); D1]$ D2 $[pr] \therefore p \equiv r \cdot \equiv \cdot \operatorname{As}(p \equiv r) :\equiv \cdot \operatorname{Vr}(pr)$ $[\alpha]$ L2 $[pr] \cdot \operatorname{Vr}(pr)$ $[D2; D1 p/p \equiv r]$ L3 $[pr] : \operatorname{Vr}(\operatorname{As}(p)r) \cdot \equiv \cdot \operatorname{As}(p) \equiv p$ $[\operatorname{L1} f/\operatorname{Vr}; \operatorname{L2} r/\operatorname{Vr}(p0)]$

We can now establish the following four metarules of procedure:

S I (see [9], pp. 114–115) If we have in the system a thesis of the type $[x, y, \ldots] : A := B$

(with or without an initial quantifier), then we can add the corresponding thesis

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$[x, y, \ldots] : B := A.$	
a) $[x, y, \ldots] : A := B$	[The assumption]
$\mathfrak{b}) \qquad [x', y', \ldots) : A' :=.$	B' [Substitution of a to eliminate any free
	variables of shapes 'f' or 'r'.]
c) $[f, x', y',] :: f(A$	$f(A'0)) :\equiv \therefore [r] \therefore f(B'r) :\equiv B' :\equiv A'$
	$[A_n p/A', q/B'; b]$
\mathfrak{d}) [r, x', y', \ldots] \therefore Vr(B'r) = B' = A'
	[c f/Vr; L2 p/A', r/Vr(A'0)]
e) $[x', y',] : B' := .$	$A' \qquad [b; L2 p/B']$
f) $[x, y, \ldots] : B := .A$	[Substitution of e to restore the
	original variables]

S II If we have in the system two corresponding theses of types

 $[x, y, \dots] : A := B$ $[x, y, \dots] : B,$

then we can add the corresponding thesis

 $[x, y, \ldots] \cdot A.$

This follows immediately from metarule S I.

S III If we have in the system two corresponding theses of types

$$[x, y, \dots] : A := B$$
$$[x, y, \dots] \cdot F(A)$$

(where 'F(A)' represents any meaningful expression which contains at least one expression equiform with 'A' not falling within the scope of any internal quantifier), then we can add the corresponding thesis

	$[x, y, \ldots] \cdot F(B).$	
a)	$[x, y, \ldots] : A := B$	[Assumption]
b)	$[x, y, \ldots] \cdot F(A)$	[Assumption]
c)	$[x', y', \ldots] : A' := B'$	[These substitutions of a and b
ð)	[x', y',] . F'(A')	eliminate any free variables of
		shapes 'f', 'p', 'q', or 'r'.]
e)	$[x', y', \ldots] : B' := .A'$	[S I: c]
f)	[x', y',] . As $(F'(A'))$	$[D1 \ p/F'(A'); b]$
g)	$[f, x', y', \dots] :: f(A'f(A'0)) := .$	$[r] \therefore f(B'r) :=: B' :=: A'$
		$[A_n p/A', q/B'; c]$
h)	[p, q, r, x', y',] :: Vr(As($F'(p)$	$(r)(r) :=: \operatorname{As}(F'(p)) :=: F'(q)$
	$\therefore \equiv G(px'y'\dots)(qr)$	[α]
j)	$[x', y', \ldots] \cdot G(A'x'y'\ldots)(A'G($	$A'x'y'\ldots)(A'0))$
	[ħ <i>p</i> ,	/A', q/A', r/G(A'x'y')(A'0);
	L3	p/F'(A'), r/G(A'x'y')(A'0)]
f)	$[r, x', y', \ldots] \therefore G(A'x'y'\ldots)(B$	(r) :=: B' :=: A'
		$[\mathfrak{g} f/G(A'x'y'\ldots);\mathfrak{j}]$
1)	$[r, x', y', \ldots] \cdot G(A'x'y'\ldots)(B'r$	(S II: ť; e)

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m)
$$[r, x', y', \ldots]$$
 \therefore Vr(As($F'(A')$) r) \equiv : As($F'(A')$) \equiv : $F'(B')$ n) $[x', y', \ldots]$ $F'(B')$ $[m; L2 p/As(F'(A')); f]$ o) $[x, y, \ldots]$ $F(B)$ [Substitution of n to restore the original variables]

S IV If we have in the system corresponding theses of types

 $[x, y, \dots] : A := B$ $[x, y, \dots] : F(B)$

(with 'F(B)' understood as in S III), then we can add the corresponding thesis

 $[x, y, \ldots] \cdot F(A).$

This follows immediately from S I and S III.

Aided by these rules, we prove the following theses:

L4	$[pr]: Vr(pr) := p \equiv p$	[S IV: D1; L3]
L5	$[p] \cdot p \equiv p$	[L4; L2]
L6	$[fp] \therefore f(pf(p0)) :=: [r] : f(pr) :=$	$p \equiv p$ [$A_n q/p$; L5]
D3	$[pq] \therefore q := . q \equiv p := . Fb(pq)$	[α]
D4	$[pr] : Fb(Vr(0r)p) := .Vr_1(pr)$	[α]
L7	$[r] : Fb(10) := . Vr_1(0r)$	[S III: L4 <i>p</i> /0; D4 <i>p</i> /0]
L8	$[r] \therefore Fb(10) = .1 :=: Vr_1(0r) = .$	1 [S III: L7;
		L5 $p/Fb(10)$.=. 1]
L9	$Fb(10) = .1 = :[r] : Vr_1(0r) = .1$	[L8; β]
L10	$Vr_1(0Vr_1(00)) :=: Fb(10) :=. 1$	[S IV: L9; L6 <i>f</i> /Vf ₁ , <i>p</i> /0]
L11	Fb(10) :=: Fb(10) :=. 1	[S IV: L7 <i>r</i> /Vr ₁ (00); L10]
L12	Fb(1Fb(10))	[D3 q/Fb(10), p/1; L11]
L13	[r]: Fb(1r) = .1 = 1	[L6 <i>f/Fb</i> , <i>p/</i> 1; L12]
L14	$[p] \cdot Fb(1p)$	[S II: L13 r/p; L5 p/1
L15	[p]: p := p = 1	[SII: D3 <i>p</i> /1, <i>q</i> / <i>p</i> ; L14]
L16	$[pq] \therefore q \equiv 1 = p := q \equiv p = 1$	[S III: L15 p/q ; L15 $p/q \equiv p$]

At this point we can establish the following metatheorem:

S V If we have in our system two theses of types

which are substitutions of some meaningful expression

 $[p] \cdot F(p)$

with the expressions '0' and '1' replacing the variable 'p', then we can add this latter expression to the system as a thesis in the following way:

a)	F(0)	[Assumption]
b)	F(1)	[Assumption]
c)	$[pq] \therefore F(p) = q \equiv q \equiv G(qp)$	[α]
ð)	[p]: F(p) := G(0p)	[S IV: L15 $p/F(p)$; c $q/0$]
e)	<i>G</i> (00)	[d p/0; a]

 $[\]frac{F(0)}{F(1)}$

f)	<i>G</i> (01)	[b p/1; b]
g)	G(00) = .1	$[L15 \ p/G(00); e]$
h)	G(0G(00))	[S IV: g; f]
j)	[r] : G(0r) := .1	[L6 f/G , $p/0$; \mathfrak{h}]
ť)	[p] . G(0p)	[S II: j <i>r/p</i> ; L5 <i>p/</i> 0]
ĺ)	$[p] \cdot F(p)$	[S II: b; ť].

With the help of S V, we prove the following theorems:

L17	$[pq] \therefore q \equiv p = . \equiv . 0 := : q \equiv 0 = . p$	$p = [S V: L5 p/q \equiv 0 . \equiv . 0;$
		L16 p/0]
L18	$[p] \therefore p := p \equiv 0 := 0$ [S II:	L17 q/p , $p/p \equiv 0$; L5 $p/p \equiv 0$]
L19	0 = .1 = 0 = .0 = .0 [S	III: L18 $p/1$; L17 $q/0$, $p/1 \equiv 0$]
L20	$[p] :: p := . p \equiv 0 := \therefore 0 := . 1 \equiv 0$:≡. 0 [S IV: L19; L18]
D5	$[pq] :: p := . 1 \equiv 0 := . q \equiv 0 : = .$	$q \equiv 0 ::=. Fc(pq) \qquad [\alpha]$
L21	$0 := .1 \equiv 0 := .1 := .Fc(00)$	[S IV: L15 $p/0 = .1 = 0;$
		D5 p/0, q/0]
L22	0 := .1 = 0 := .Fc(00)	[S IV: L15 $p/0$.=. 1 = 0; L21]
L23	Fc(00 = .1 = 0) [D5 p/0, q/0]	0 = .1 = 0; L20 p/0 = .1 = 0
L24	Fc(0Fc(00))	[S III: L22; L23]
L25	[r]: Fc(0r) := .1	[L6 <i>f/Fc</i> , <i>p/</i> 0; L24]
L26	Fc(00)	[S II: L25 r/0; L5 p/0]
L27	$0 = .1 \equiv 0$	[S II: L22; L26]
L28	$0 \equiv 1 = . 1 \equiv 0$	[S III: L15 p/0; L27]
L29	$1 \equiv 0 = . \equiv . 0 \equiv 1$	[S I: L28]
L30	$[p]: p \equiv 1 . \equiv . 1 \equiv p$	[S V: L28; L5 $p/1 \equiv 1$]
L31	$[p]: p \equiv 0 . \equiv . 0 \equiv p$	[S V: L5 <i>p</i> /1; L29]
L32	$[pq]: p \equiv q := . q \equiv p$	[S V: L31; L30]
L33	$[pq] \therefore q \equiv p = . \equiv . 1 : \equiv : q \equiv 1 = . p$	9 [S I: L16]
L34	$[pqr] \therefore q \equiv p \equiv . \equiv . r \equiv : q \equiv r \equiv . p$	o [S V: L17; L33]
L35	$[pqr] \therefore p \equiv q \equiv r \equiv q \equiv r \equiv q$	o [S IV: L32; L34]
L36	$[pqr] \therefore p \equiv q \equiv r \equiv p \equiv q \equiv q$	r [S IV: L32 $q/q \equiv r$; L35]
L37	$[pq] \therefore p :=: q :=. p \equiv q$	$[L36 r/p \equiv q; L5 p/p \equiv q]$
L38	$[pq] \therefore p :=: q :=. q \equiv p$	$[L36 r/q \equiv p; L32]$
L39	[pq]: p := Fb(pq)	[S IV: L38; D3]
L40	$[pqr]$: p :=: $Fb(qr)$:=. $q \equiv p$	[S III: L39 <i>p/q</i> , <i>q/r</i> ; L38]
L41	$[pqr] \therefore Fb(pFb(p0)) =: Fb(qr)$.	$\equiv . q \equiv p$
		[S III: L39 <i>q/Fb</i> (<i>p</i> 0); L40]
L42	$[pq] \therefore Fb(pFb(p0)) =: [r] : Fb(a)$	$(qr) := . q = p$ [L41; β]
L43	[fp] :: p := :: f(pf(p0)) := : [r] : f	(pr) := p = p : = p
	[L37, q/p , $p/f(pf(p0))$) $=: [r] : f(pr) = p = p; L6$
L44	[p] :: p := :: [f] :: f(pf(p0)) :=: [r]:f(pr) := p = p := p
		[L43; β]
L45	$[p] :: p := \therefore [f] \therefore f(pf(p0)) :=:$	$[r]:f(1r):=.\ 1\equiv p$
		[S IV: L15; A _n q/1]
L46	[p] :: [f] : f(pf(p0)) :=: [r] : f(pf(p(p0)) :=: [r] : f(pf(p0)) :=: [r] : f(pf(pf(p0)) :=: [r] : f(pf(pf(p0)) :=: [r] : f(pf(pf(p0)) :=: [r] : f(pf(pf(p0)) :=: [r] : f(pf(pf(pf(p0))) :=: [r] : f(pf(pf(pf(pf(pf(pf(pf(pf(pf(pf(pf(pf(pf	$1r) = .1 \equiv p : = :: [f] ::$
	f(pf(p0)) :=: [r] : f(1r) :=. 1 = p	∴≡. 1
	[L15 <i>p/f(pf(p</i> 0)) .=: $[r] : f(1r) .= .1 = p; \beta]$

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L47	[p] :: p := :: [f] :: f(pf(p0)) .	$\equiv: [r] : f(1r) := .1 \equiv p := .1$
		[S IV: L45; L46]
D6 ⁴	[pq] :: p := :: [f] :: f(pf(p0))	$ =: [r] : f(qr) = q \equiv p := q $
	$\therefore \equiv . p \supset q$	[α]
L48	$[p] \cdot p \supset p$	[D6 <i>q/p</i> ; L44]
L49	$[p] \cdot p \supset 1$	[D6 q/1; L47]
L50	$[p] \ . \ 0 \supset p$	[S V: L48 p/0; L49 p/0]
L51	1⊃1.≡.1	$[L15 \ p/1 \supset 1; L48 \ p/1]$
L52	1.0.101	[S IV: L51; L48 <i>p</i> /1]
L53	0.⊃.1⊃0:≡:1.⊃.1⊃1	$[L37 \ p/0 \ . \supset . \ 1 \supset 0, \ q/1 \ . \supset . \ 1 \supset 1;$
		L50 $p/1 \supset 0$; L52]
L54	1.⊃.1⊃0:≡∴1.⊃:1.⊃.0≡	1 [S III: L15 <i>p</i> /0;
		L5 $p/1$. \supset . $1 \supset 0$]
L55	$[p] :: p : \supset . 1 \supset 0 := \therefore 1 : \supset : 1$	$\Box : 0 \equiv p$ [S V: L53; L54]
L56	1 .⊃∴ 1 .⊃: 0 .≡. 1 ⊃ 0	[L55 $p/1 \supset 0$; L48 $p/1 \supset 0$

At this point we prove the metatheorem:

S VI⁵: If we have in the system corresponding theses

$$\begin{array}{c} A \supset B \\ A \end{array}$$

then we can add to the system the thesis

B.

a)	$A \supset B$	[Assumptio	n]
b)	A	[Assumptio	n]
c)	$A :\equiv :: [f] :: f(Af(A0)) :=: [r] : f(Br) :=$	$B \equiv A \therefore \equiv B$	
		[S II: D6 <i>p</i> / <i>A</i> , <i>q</i> / <i>B</i> ;	a]
ð)	[f] :: f(Af(A0)) :=: [r] : f(Br) :=. B = A	∴≡. <i>B</i> [c;	6]
e)	В	[b f/Fb; L42 p/A, q/A	B]

We then use this rule to prove

L57	$1 \supset : 0 = . 1 \supset 0$	[S VI: L56; L5 p/0]
L58	$0 = .1 \supset 0$	[S VI: L57; L5 p/0]

This completes the proof that the laws of implication can be derived from A_n without using the rule of extensionality. Similar proofs may easily be constructed for A_m , A_o , A_p , and A_q (see [8], p. 67 and [9], p. 122). On the other hand, it seems that the deductions from axiom A_l and the proof of metatheorem $A2^6$ require the use of theses of higher extensionality. Observe the crucial use of the rule for distributing quantifiers at L9, L42, L44, and L46; it is largely this rule and the unusual definition of implication at D6 which allow us to prove the laws of implication L48, L49, and L58.

It is well known that theses L32 and L36 are an adequate axiom system for the *E*-calculus⁷, and so replace the thesis A78 proved by Professor Sobociński. My research has concerned the axiom A_o , from which deductions are slightly shorter than those from A_n^8 ; even a slight simplification of the known proof of A78 required 1,205 applications of the rule of inference of \mathfrak{S}_5 , including definitions. The proof of L32 and L36 requires at most 583 applications of the rule, while the proof of the related system L34 and L59

L59 [pqr] : $r := p \equiv q :=: r := q \equiv p$

requires at most 568 applications of the rule. Hence the deductions outlined in this paper reduce to less than half its length the proof that A_0 is adequate for the *E*-calculus.

NOTES

- 1. The basic sources describing protothetic are [1], [2], and [3]. The articles [8], [9], and [10] are indispensable for serious study of the theory.
- The proof given in L37 to L58 of this paper shows that derivations from this axiom need not rely on the theoretical considerations which led to its discovery. See [7], p. 19; [8], pp. 66-67; [10], p. 146, note 83; and especially [9], p. 122: "On the other hand I am unable to prove that A_n... is a single axiom of S₅ without the application of reasonings... in which point ε plays an essential role."
- 3. See [9], pp. 115-117 and 119-120. My formulation of these rules is slightly broader, as I refer to meaningful expressions instead of functors. My deductions diverge significantly from those in [9] from definition D4 onwards.
- 4. This unusual definition differs from any of those previously known, and from any of the related definitions of conjunction. For previous definitions see [1], p. 12, and [2], D3 on p. 93 and D4 on p. 107. For definitions of conjunction see [1], p. 13; [3], pp. 138, 141, and 147; [6], pp. 202-203; and the discussion in [10], pp. 134-144. Compare the proof of L58 with the theorem A1 in [10], pp. 134-139, to which it bears a close resemblance. I do not know how to shorten the present definition of implication and still carry out the proof of L58.
- 5. This rule is quite different from the S VI of [10], p. 132.
- 6. This axiom is found in [8], p. 66, and discussed in [9], pp. 123-124. The metatheorem appears in [10], pp. 139-144.
- 7. See [4], p. 91. Initial deductions from a closely related axiom system are given in [10], p. 314.
- 8. For example, the thesis L5 can be proved from A_o in only 44 applications of the rule, but its proof from A_n appears to require at least (and certainly at most) 48 applications. Even the initial deductions proving L32 and L36 from A_o would require about 100 pages of print, so that I cannot see any possibility of publishing them.

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