ON MESOPRIMARY DECOMPOSITION OF MONOID CONGRUENCES

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ABSTRACT. We prove two main results concerning mesoprimary decomposition of monoid congruences, as introduced by Kahle and Miller. First, we identify which associated prime congruences appear in every mesoprimary decomposition, thereby completing the theory of mesoprimary decomposition of monoid congruences as a more faithful analogue of primary decomposition. Second, we answer a question posed by Kahle and Miller by characterizing which finite posets arise as the set of associated prime congruences of monoid congruences.

1. Introduction. A congruence is an equivalence relation on the elements of a monoid that respects the monoid operation. This paper focuses on congruences on the monoid of monomials in a polynomial ring S that arise from binomial ideals in S, that is, ideals whose generators have at most two terms. In particular, any binomial ideal $I \subset S$ identifies, up to scalar multiple, any two monomials appearing in the same binomial in I, inducing a congruence \sim_I on the monoid of monomials in S. In [3], Kahle and Miller introduce mesoprimary decompositions of binomial ideals, which are combinatorial approximations of primary decompositions constructed from the underlying congruences.

Mesoprimary decompositions are constructed in two settings: first for monoid congruences, and then for binomial ideals; both are designed to parallel standard primary decomposition in a Noetherian ring [1, Chapter 3]. At the heart of mesoprimary decomposition, for both monoid congruences and binomial ideals, lies a notion of associated objects analogous to associated prime ideals in standard primary decom-

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position. In particular, any congruence \sim has a collection of associated prime congruences, and each component in a mesoprimary decomposition for \sim has precisely one associated prime congruence. However, unlike standard primary decomposition, eliminating redundant mesoprimary components can produce decompositions in which some of the associated objects do not appear as the associated object of any component (Example 3.2).

The focus of this paper is on mesoprimary decomposition of monoid congruences, and the two main results are as follows. First, we identify the class of truly associated prime congruences (Definition 3.3), which must appear as the associated prime congruence of some component in every mesoprimary decomposition of \sim (Theorems 3.9 and 4.9), thereby completing the theory of mesoprimary decomposition of monoid congruences as a more faithful analogue of standard primary decomposition of ideals. Second, we characterize which finite posets arise as the set of associated prime congruences of a congruence, and in doing so answer [3, Problems 17.4 and 17.9].

- 2. Overview of mesoprimary decomposition of monoid congruences. In this section, we briefly review the necessary definitions and results found in [3] concerning mesoprimary decomposition of monoid congruences. The full background and theory developed in [3] is too extensive to include here; the unfamiliar reader is referred to [3] for a thorough treatment on mesoprimary decomposition, or [4] for a more algebraic overview that also includes mesoprimary decompositions of binomial ideals. Basic monoid definitions can also be found in [2].
- **2.1. Conventions.** Unless otherwise stated, Q denotes a finitely generated (equivalently, Noetherian) commutative monoid, and the nil, i.e., universally absorbing element, of Q, if it exists, is denoted $\infty \in Q$.
- **Definition 2.1.** A subset $P \subset Q$ is an *ideal* if $P + a \subset P$ for all $a \in Q$, and P is *prime* if $Q \setminus P$ is a submonoid of Q. An equivalence relation \sim on Q is a *congruence* if $a \sim b$ implies $a + c \sim b + c$ for all $a, b, c \in Q$.

Notation 2.2. Fix a congruence \sim on Q, a prime $P \subset Q$, and let $F = Q \setminus P$. We write Q_P for the localization of Q along P (that is, the

set of formal differences q-f for $q \in Q$ and $f \in F$ with the stipulation that q-f=q'-f', whenever w+q+f'=w+q'+f for some $w \in F$), and $\overline{Q}_P=Q_P/\sim$ for the quotient of Q_P modulo \sim . We denote by \overline{q} the image of $q \in Q$ in $\overline{Q}=Q/\sim$.

Definition 2.3. Fix a prime $P \subset Q$, and consider the equivalence relation \sim on Q_P relating $a \sim b$ whenever $a - b \in Q_P \setminus P_P$. Given an element $a \in Q$, the set of elements of Q lying in the same equivalence class under \sim in Q_P is called *Green's class* of a. By [3, Lemma 2.19], the set of Green's classes in Q_P is partially ordered by divisibility, and the induced preorder on elements of Q_P is called *Green's preorder*.

Definition 2.4 ([3, Definitions 2.12, 3.4, 4.7, 4.10, 7.1, 7.2, 7.7, 7.12]). Fix a congruence \sim on Q and a prime $P \subset Q$.

- (a) An element $q \in Q$ is an aide for an element $w \in Q$ and a generator $p \in P$ if
 - (i) $\overline{w} \neq \overline{q}$;
 - (ii) $\overline{w} + \overline{p} = \overline{q} + \overline{p}$; and
 - (iii) \overline{q} is maximal (under Green's preorder) in the set $\{\overline{q}, \overline{w}\}$. If q is an aide for w for each generator of P, then q is a key aide.
- (b) An element $w \in Q$ is a witness for P if it has an aide for each $p \in P$, and a key witness for P if it has a key aide. A key witness w is a cogenerator for \sim if w + p is nil modulo \sim for all $p \in P$.
- (c) The congruence \sim is P-primary if every $p \in P$ is nilpotent in \overline{Q} and every $f \in Q \setminus P$ is cancelative in \overline{Q} . A P-primary congruence \sim is mesoprimary if $\overline{a} + \overline{b} = \overline{a} + \overline{c} \neq \infty$ implies $\overline{b} = \overline{c}$ for all $\overline{a}, \overline{b}, \overline{c} \in \overline{Q}$, that is, every element of \overline{Q} is partly cancelative. The congruence \sim is coprincipal if it is mesoprimary and every cogenerator for \sim generates the same ideal in \overline{Q} .
- (d) The coprincipal component \sim_w^P of \sim cogenerated by a witness $w \in Q$ for P is the coprincipal congruence that relates $a \sim_w^P b$ if one of the following is satisfied:
 - (i) both \overline{a} and \overline{b} generate an ideal not containing \overline{w} in \overline{Q}_P ; or
 - (ii) \overline{a} and \overline{b} differ by a unit in \overline{Q}_P and $\overline{a} + \overline{c} = \overline{b} + \overline{c} = \overline{w}$ for some $\overline{c} \in \overline{Q}_P$.

A (key) witness for P may be called a (key) \sim -witness for P to specify \sim . Congruences may be called P-mesoprimary or P-coprincipal to specify P.

Theorem 2.5 ([3, Theorem 8.4]). Each congruence \sim on Q is the common refinement of the coprincipal components cogenerated by its key witnesses.

Lastly, we recall the definition of prime congruences from [3], which play the role of "associated objects" in this setting.

Definition 2.6 ([3, Definitions 5.1 and 5.2]). Fix a congruence \sim on a monoid Q, a prime ideal $P \subset Q$, and an element $q \in Q$ that is not nil modulo \sim .

- (a) Let \approx denote the congruence on Q_P that sets $a \approx b$ when
 - (i) a and b both lie in P_P , or
 - (ii) a and b both lie in $Q_P \setminus P_P$ and $\overline{a} + \overline{q} = \overline{b} + \overline{q}$. The P-prime congruence of \sim at q is $\ker(Q \to Q_P/\approx)$.
- (b) The P-prime congruence at q is associated to \sim if q is a key witness for P.

We emphasize here that, in contrast to the setting of groups, the kernel $\ker(Q \to Q')$ of a monoid homomorphism is a congruence on Q, not a submonoid of Q.

Remark 2.7. By [3, Corollary 6.7], a congruence is P-mesoprimary if and only if it is P-primary and the P-prime congruences at every non-nil element all coincide. Generally speaking, each witness w for P detects an element whose P-prime congruence differs from those in the direction(s) of P, and the coprincipal component at w distinguishes the P-prime congruence at w from those above it in the decomposition in Theorem 2.5. We direct the unfamiliar reader to [3, Example 1.3] and the accompanying graphics, which are a particularly enlightening illustration of mesoprimary decomposition at the level of congruences.

3. True witnesses of monoid congruences. Key witnesses (Definition 2.4) form a restricted class of witnesses sufficient for decomposing any monoid congruence, but the coprincipal components they cogenerate may still be redundant (Example 3.2). In this section, we further restrict to the class of true witnesses (Definition 3.3), which are still sufficient for decomposing any congruence (Theorem 3.9).

Although this paper is focused on monoid congruences, examples are easiest to state using the language of binomial ideals; thus, first we give a few basic definitions. In what follows, k denotes an arbitrary field.

Definition 3.1. A binomial in $\mathbb{k}[Q]$ is an element of the form $\mathbf{t}^a - \lambda \mathbf{t}^b$ where $a, b \in Q$ and $\lambda \in \mathbb{k}$. An ideal $I \subset \mathbb{k}[Q]$ is binomial (respectively, monomial) if it can be generated by binomials (respectively, monomials). The congruence \sim_I on Q induced by a binomial ideal $I \subset \mathbb{k}[Q]$ sets $a \sim_I b$ whenever $\mathbf{t}^a - \lambda \mathbf{t}^b \in I$ for some nonzero $\lambda \in \mathbb{k}$.

Example 3.2. Let

$$I = \langle x^3 - xy^2, x^3(z-1), x^2y - y^3, y^3(w-1), x^4, y^4 \rangle \subset \mathbb{k}[x, y, z, w].$$

Its congruence \sim_I on $Q=\mathbb{N}^4$ is depicted in Figure 1, projected onto the xy-plane. The congruence \sim_I is P-primary for $\mathfrak{m}_P=\langle x,y\rangle$ and has five Green's classes of key witnesses, namely, those containing the monomials x^2, y^2, x^3, y^3 and x^3y , respectively. Indeed, x^2 and y^2 are each key aides for the other, wx^3 is a key aide for x^3, zy^3 is a key aide for y^3 and x^3y has nil as a key aide. Of these, x^2 and y^2 yield redundant components in the coprincipal decomposition for \sim_I in Theorem 2.5, and the remaining three comprise a mesoprimary decomposition for \sim_I with no redundant components.

Definition 3.3. Fix a congruence \sim on Q, a prime $P \subset Q$, and an element $w \in Q$.

- (a) A P-cover congruence of w is the P-prime congruence at a non-nil element w + p for some generator p of P.
- (b) The discrete testimony of w at P is the set $T_P(w)$ of P-cover congruences of w. The discrete testimony of w is suspicious if the common refinement of the P-cover congruences in the testimony coincides with the P-prime congruence at w.

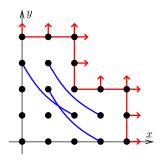


FIGURE 1. A congruence \sim on \mathbb{N}^4 with key witnesses whose coprincipal components are redundant, projected onto the xy-plane.

- (c) We say w is a true witness if
 - (i) w is maximal among \sim -witnesses for P, or
 - (ii) the discrete testimony of w is not suspicious.
- (d) A *P*-prime congruence \approx is *truly associated* to \sim if it is the *P*-prime congruence at a true \sim -witness for *P*.

Example 3.4. Conditions (i) and (ii) in Definition 3.3 (c) are both necessary. Indeed,

$$I_1 = \langle x^2 - xy, xy - y^2 \rangle \subset \mathbb{k}[x, y]$$

induces a congruence on \mathbb{N}^2 with two witnesses for the maximal prime P, both of which are maximal among witnesses for P, but neither of which has suspicious testimony since \overline{Q} has no nil element. Additionally, the congruence induced by

$$I_2 = \langle z^4 - 1, x(z-1), y(z^2 - 1), x^2, xy, y^2 \rangle \subset \mathbb{k}[x, y, z]$$

has three witnesses for the maximal prime P, one of which (the origin) has suspicious testimony but is not maximal among witnesses for P. The congruences induced by I_1 and $I_2 \cap \mathbb{k}[x,y]$ are depicted in Figure 2.

Remark 3.5. Character witnesses [3, Definition 16.3] are the binomial ideal analogues of true witnesses, except that their testimony is computed by intersecting ideals instead of refining congruences. In general, however, character witnesses need not be true, and true witnesses

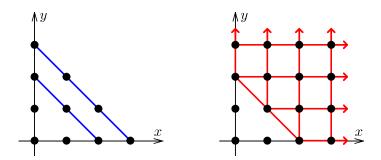


FIGURE 2. Congruences for the ideals I_1 (left) and I_2 (right) in Example 3.4.

nesses need not be character. Additionally, Corollary 3.8 states that true witnesses are key, a fact that fails for character witnesses; see [3, Examples 16.5–16.7] for a demonstration of this behavior.

Proposition 3.6 and Corollary 3.7 each give an equivalent condition for identifying true witnesses that will be useful in proving Theorem 4.9.

Proposition 3.6. Fix a congruence \sim on Q and a witness w for P. The discrete testimony of w is not suspicious if and only if w has a key aide that is nil or generates the same ideal as w in Q_P .

Proof. If w has ∞ as a key aide, then its discrete testimony is empty. If w has a key aide w' in its Green's class in Q_P , then each prime congruence in its discrete testimony identifies w and w', and thus, so does their common refinement. Either way, the discrete testimony of w is not suspicious.

Now, suppose that the discrete testimony of w is not suspicious and that $\infty \in Q$ is not a key aide. The set $T_P(w)$ is thus nonempty, and the common refinement of the prime congruences in $T_P(w)$ relates some distinct u and v outside of P that are not related under the prime congruence \approx at w. Since u and v are distinct elements outside of P, the Green's class of w in Q_P must contain some other element w'. Choosing w' so that w + u = w' + v implies w + p = w' + p for each $p \in P$, making w' a key aide for w.

Corollary 3.7. Fix a congruence \sim on Q and a witness w for P. The element w is a true witness if and only if

- (i) w is maximal among \sim -witnesses for P, or
- (ii) w has a key aide that generates the same ideal as w in Q_P .

Corollary 3.8. Every true \sim -witness is a key \sim -witness, and every truly associated prime congruence of \sim is associated to \sim .

We are now ready for the main result of this section. Theorem 3.9 shows that, when constructing an induced coprincipal decomposition for a given congruence, it suffices to consider true witnesses. In particular, any component in the decomposition given in Theorem 2.5, cogenerated by a non-true witness is redundant, and can be omitted.

Theorem 3.9. Every congruence on Q is the common refinement of the coprincipal components cogenerated by its true witnesses.

Proof. Since primary decomposition of monoid congruences commutes with localization [3, Theorem 3.12], it suffices to assume that $Q = Q_P$. Fix a congruence \sim on Q and a key \sim witness w for P that is not true. By Theorem 2.5, we can write $\sim = \bigcap_i \sim_i$ as the common refinement of the congruences cogenerated by the key \sim -witnesses. In order to prove that the congruence \sim_w^P is redundant in this decomposition, it suffices to produce, for $q, q' \in Q$ not identified under \sim_w^P , a component $\sim_j \neq \sim_w^P$ not identifying q and q'.

First, suppose that q and q' lie in distinct Green's classes in Q. Since w is not true, it is not maximal; thus, some maximal witness v for P lies above w. The nil class of \sim_v^P is properly contained in the nil class of \sim_w^P ; hence, q and q' are both not nil under \sim_v^P . Furthermore, outside of its nil class, \sim_v^P does not relate any elements that lie in separate Green's classes. In particular, \sim_v^P does not relate q and q'.

Next, suppose that q and q' lie in the same Green's class in Q. Since q and q' are not both nil modulo \sim_w^P , there exists a $u \in Q$ such that q + u and q' + u are in the same Green's class as w. Furthermore, any component that does not relate q + u and q' + u will not relate q and q'; thus, upon replacing q with q + u and q' with q' + u, it suffices to assume u = 0 and q' = w. As w is not a true witness, q is not a key

aide for w. Therefore, $w+p \approx q+p$ for some generator $p \in P$. This means that some component \sim_j does not relate w+p and q+p, and thus does not relate w and q, as desired.

4. Irredundant mesoprimary decompositions of congruences. In this section, we prove that each truly associated prime congruence of a given congruence \sim appears as the associated prime congruence of some mesoprimary component in every mesoprimary decomposition for \sim (Theorem 4.9). As a consequence, we prove that any congruence with no embedded associated monoid primes possesses both a unique minimal mesoprimary decomposition and a unique irredundant mesoprimary decomposition (Corollary 4.12). Making statements about "all" mesoprimary decompositions necessitates some mild restrictions, see Remark 4.4 and [3, Example 8.2].

Definition 4.1 ([3, Definition 8.1]). An expression $\sim = \bigcap_i \approx_i$ of a congruence \sim as a common refinement of mesoprimary congruences is a *mesoprimary decomposition* if, for each \approx_i with associated prime P_i , the P_i -prime congruences of \sim and \approx_i at each cogenerator for \approx_i coincide. This decomposition is key if every cogenerator for every \approx_i is a key witness for \sim .

Remark 4.2. Theorems 2.5 and 3.9 both yield key mesoprimary decompositions.

Definition 4.3. A mesoprimary decomposition $\sim = \bigcap_i \sim_i$ is:

- (a) induced if each \sim_i is a common refinement of coprincipal components;
- (b) minimal if \sim_i and \sim_j have distinct associated prime congruences for $i \neq j$; or
- (c) irredundant if no \sim_i can be omitted.

Remark 4.4. The coprincipal component \sim_w^P of a congruence \sim at a witness w for P is determined by the congruence \sim . More precisely, it is the finest coprincipal congruence with cogenerator w that can appear in a mesoprimary decomposition for \sim . As such, for the purpose of minimality, we restrict our attention to induced mesoprimary decompositions. Indeed, if the induced condition is

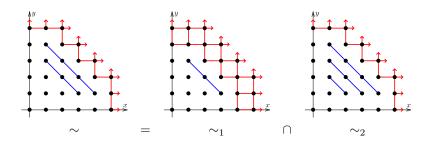


FIGURE 3. The non-induced mesoprimary decomposition from Example 4.5.

relaxed, coprincipal components whose cogenerator is a non-key \sim -witness need not be redundant, see Example 4.5.

Example 4.5. The ideal $I = \langle x^3y - x^2y^2, x^2y^2 - xy^3, x^5, y^5 \rangle$ is the intersection of $I_1 = \langle x^3y - x^2y^2, x^2y^2 - xy^3, x^4, y^4 \rangle$ and $I_2 = \langle x^2y - xy^2, x^5, y^5 \rangle$. Their congruences \sim , \sim_1 and \sim_2 , respectively, are depicted in Figure 3. Both \sim and \sim_2 are coprincipal with cogenerator (4,1), but \sim_2 is not the coprincipal component cogenerated by (4,1) since it also identifies (2,1) and (1,2). As such, this mesoprimary decomposition is not induced. Additionally, \sim_1 is cogenerated by a non-key non-character witness for \sim , but neither component of this mesoprimary decomposition can be omitted.

An important observation is that any witness whose discrete testimony is not suspicious must appear as a cogenerator in every mesoprimary decomposition. Note the absence of "induced" here; we do indeed mean *every* mesoprimary decomposition. This fact is recorded in Lemma 4.6, which serves as the foundation for Theorem 4.9.

Lemma 4.6. Fix a mesoprimary decomposition $\sim = \bigcap_i \sim_i$ and a \sim -witness w for P. If the discrete testimony of w is not suspicious, then w is a cogenerator for some \sim_i .

Proof. Let \approx denote the *P*-prime congruence at w, and let \approx_i denote the prime congruence associated to \sim_i for each i. From Proposition 3.6, either w has ∞ as a key aide, or w has a key aide w' that is Green's equivalent to w in the localization Q_P . If w has ∞ as a key aide, then

it is a cogenerator for \sim ; thus, any mesoprimary component \sim_i under which w is not nil also has w as a cogenerator.

Alternatively, suppose that w has a key aide w' in the same Green's class as w in Q_P . Since $w \nsim w'$, some mesoprimary component \sim_i does not relate w and w'. Neither w nor w' is nil under \sim_i , but for each generator p of P, the prime congruence at w+p relates w and w'. This means that each w+p must be nil under \sim_i since \sim_i is mesoprimary; hence, w is a cogenerator for \sim_i .

The symmetry in Example 4.7, which also appeared as [3, Example 2.19], demonstrates that Lemma 4.6 cannot be generalized to arbitrary true witnesses since eliminating all redundancy sometimes requires making arbitrary choices. That said, Lemma 4.8 demonstrates that the phenomenon in Example 4.7 is the only possible obstruction.

Example 4.7. Let $I = \langle x^2 - xy, xy - y^2 \rangle \subset \mathbb{k}[x, y]$. The congruence \sim_I has two associated primes, namely, \emptyset and the maximal ideal P. Theorem 2.5 produces the coprincipal decomposition

$$I = \langle x^2 - xy, xy - y^2 \rangle = \langle x^2, y \rangle \cap \langle x, y^2 \rangle \cap \langle x - y \rangle.$$

The first two components are P-primary, and the third is \emptyset -primary. Either, but not both, of the first two components can be omitted without affecting the intersection, although each is cogenerated by a true witness for \sim_I .

Lemma 4.8. Fix a congruence \sim , a key \sim -witness w for P, and a key aide w' for w. If w is a maximal witness for P, then every mesoprimary decomposition $\sim = \bigcap_i \sim_i has$ a component with w or w' as a cogenerator.

Proof. Suppose that w is maximal among \sim -witnesses for P. Since primary decomposition of congruences commutes with localization by [3, Theorem 3.12], it suffices to replace Q with Q_P so that P is maximal. If w' is nil, then w is a cogenerator for \sim ; thus, it is a cogenerator for any P-primary component \sim_i under which it is not nil. If, instead, w' lies in the same Green's class as w in Q_P , then we are finished by Lemma 4.6. Lastly, assume that w' is not nil and lies in a different Green's class in Q_P . Since $w \sim w'$, some component \sim_i separates w and w'. Localization Q at any prime P' properly contained

in P identifies w and w' since w + p = w' + p for any $p \in P \setminus P'$. This means that any P'-primary component also identifies w and w'. Hence, \sim_i must be P-primary. Since w is maximal among witnesses for P, it is either a cogenerator for \sim_i or nil modulo \sim_i ; the latter implies that w' is a cogenerator for \sim_i . In either case, the proof is complete. \square

Theorem 4.9. Fix a congruence \sim , a true \sim -witness w for a prime P, and let \approx denote the P-prime congruence at w.

- (a) *If*
 - (i) the discrete testimony of w is not suspicious, or
 - (ii) the P-prime congruence at some non-nil key aide w' for w equals ≈,

then \approx appears as the associated prime congruence of some component in every mesoprimary decomposition $\bigcap_i \sim_i$ of \sim .

(b) If w satisfies neither (i) nor (ii), then the component in the coprincipal decomposition in Theorem 3.9 with cogenerator w is redundant.

Proof. If the discrete testimony of w is not suspicious, then apply Lemma 4.6. On the other hand, if w has a key aide w', whose prime congruence is also \approx , then by Lemma 4.8, one of w and w' must appear as a cogenerator of some component \sim_i . This proves part (a).

Next, fix $a, b \in Q$ with $a \nsim b$. The proof of Theorem 2.5 at the source [3, Theorem 8.4] implies that there is a prime $P \subset Q$ and $u \in Q$ such that (after possibly swapping a and b) a + u is a key witness with key aide b + u. If a + u has suspicious discrete testimony, then by Proposition 3.6, it does not have nil as a key aide. Thus, b + u is also a key witness for P. If, additionally, a + u and b + u have distinct P-prime congruences, then, since a + u and b + u have identical discrete testimony, the discrete testimony of b + u is not suspicious. Since $a \nsim_{b+u}^P b$, this completes the proof.

Corollary 4.10. Fix a mesoprimary decomposition $\sim = \bigcap_i \sim_i$. Each truly associated prime congruence in of \sim is associated to some component \sim_i , and any component whose associated prime congruence is not truly associated to \sim is redundant.

We conclude this section by characterizing the minimal and irredundant mesoprimary decompositions of congruences with no embedded associated monoid primes.

Theorem 4.11. Fix a mesoprimary decomposition $\sim = \bigcap_i \sim_i$. If P is a minimal associated prime of \sim , then every true witness w of P is a cogenerator of some component.

Proof. Let ≈ denote the *P*-prime congruence at w, and let ≈ $_i$ denote the prime congruence associated to \sim_i for each i. If $P = \emptyset$, then, since P is associated to \sim , some component \sim_i is P-primary, and in fact, $\sim_i = \approx$. Now, assume that P is nonempty. Once again, after localizing at P, assume that P is maximal. Since P is a minimal associated prime, \sim is P-primary by [3, Corollary 4.21]. Since w is true, either it is a maximal witness for P, in which case it has ∞ as a key aide, or its testimony is not suspicious. In either case, we are done by Lemma 4.6.

Corollary 4.12. Any congruence \sim on Q with no embedded associated monoid primes has a unique irredundant induced coprincipal decomposition and a unique induced mesoprimary decomposition. In particular, this holds when \sim is primary.

Proof. Theorem 3.9 produces the unique induced coprincipal decomposition since omitting any component yields an expression that cannot decompose \sim by Theorem 4.11. Furthermore, replacing any set of components with their common refinement whenever they share an associated prime congruence results in a minimal mesoprimary decomposition by [3, Proposition 6.9].

5. Posets of associated mesoprimes. Here, we answer a question posed by Kahle and Miller. It is known that any poset occurs as the set of associated primes of a monomial ideal; as such, the question is posed only for primary congruences so that the nilpotent directions of the associated prime congruence, i.e., the "monomial part" of an ideal inducing the congruence, all coincide.

Problem 5.1 ([3, Problem 17.4]). Characterize the posets of associated prime congruences of primary congruences.

Theorem 5.5 provides a full, albeit unsatisfying, answer to Problem 5.1, as stated. The issue is that, in the constructed congruence, most of the witnesses are incomparable under the divisibility poset of Q. In view of this, we introduce the prime congruence poset (Definition 5.6), which only renders associated prime congruences comparable if they occur at comparable elements under divisibility in \overline{Q}_P . Surprisingly, the prime congruence poset has no further restrictions than the poset of truly associated prime congruences (Theorem 5.8).

Note that the content of this section also answers [3, Problem 17.9], see Remark 5.9.

Definition 5.2. The poset of truly associated prime congruences of a congruence \sim is

$$MesoAss(\sim) = \{ \approx \text{ truly associated to } \sim \},$$

partially ordered by refinement.

Lemma 5.3. For any primary congruence \sim , MesoAss(\sim) has a unique minimum.

Proof. The prime congruence \approx at the origin refines the prime congruence at every non-nil element, and any nilpotent element that is maximal among those with prime congruence \approx is a true witness by Theorem 3.6. As such, $\approx \in \text{MesoAss}(\sim)$.

Proposition 5.4. Fix a finite subposet $\Omega = \{p_0, \ldots, p_d\}$ of the power set of $\{1, \ldots, n\}$ and distinct primes $a_0, \ldots, a_n \in \mathbb{Z}$. For $0 \le i \le d$, let

$$b_i = \prod_{j \in p_i} a_j$$
 and $I_i = \langle y^{b_i} - 1 \rangle \subset \Bbbk[y].$

- (a) The posets
 - (i) $\{b_0, \ldots, b_d\}$, ordered by divisibility, and
 - (ii) $\{I_0, \ldots, I_d\}$, ordered by reverse containment, each coincide with Ω .
- (b) No ideal I_i equals the intersection of a collection of ideals in $\{I_0, \ldots, I_d\} \setminus \{I_i\}$.

Proof. This follows from the fact that $(y^c-1) \mid (y^{c'}-1)$ if and only if $c \mid c'$.

Theorem 5.5. Suppose that $\Omega = \{p_0, p_1, \dots, p_d\}$ is a finite subposet of the power set of $\{1, \dots, n\}$ with $p_0 = \emptyset$. Let $I_0, \dots, I_d \subset \mathbb{k}[z]$ denote the ideals from Proposition 5.4. If

$$I = I_0 + x_1 I_1 + \dots + x_d I_d + \langle x_1, \dots, x_d \rangle^2 \subset \mathbb{k}[x_1, \dots, x_d, y],$$

then the poset MesoAss(\sim_I) is isomorphic to Ω .

Proof. For each $i \in \{1, \ldots, d\}$, the monoid element corresponding to x_i is a key witness for \sim_I with associated prime congruence induced by I_i , and the prime congruence at the origin is the congruence included by I_0 . As such, $\operatorname{MesoAss}(\sim_I) = \{\sim_0, \sim_1, \ldots, \sim_d\}$ is isomorphic to the poset Ω .

Definition 5.6. Fix a primary congruence \sim on Q. Given $q \in Q$, let \approx_q denote the P-prime congruence of \sim at q. The prime congruence poset $(\Omega(\sim), \preceq)$ consists of

- the set $\Omega(\sim)$ of pairs (q, \approx_q) for non-nil $q \in \overline{Q}_P$ modulo the equivalence relation generated by relating (a, \approx_a) and $([a + b], \approx_{a+b})$ whenever $\approx_a = \approx_{a+b}$, and
- the partial ordering \leq under which $(a, \approx_a) \leq (b, \approx_b)$ whenever $\langle a \rangle \supset \langle b \rangle$.

Lemma 5.7. Fix a primary congruence \sim . The poset relation of \sim is an equivalence relation which coarsens \sim , and the order \leq on $\Omega(\sim)$ is a partial order.

Proof. The important observation is that, when $\langle a \rangle \supset \langle b \rangle$ for non-nil $a, b \in \overline{Q}_P$, the prime congruence \approx_a at a coarsens the prime congruence \approx_b at b. This implies

- (i) (a, \approx_a) and (b, \approx_b) are identified in $\Omega(\sim)$ whenever a and b lie in the same Green's class in \overline{Q}_P , and
- (ii) if $\langle a \rangle \supset \langle b \rangle$ and $\approx_a = \approx_b$, then the prime congruence \approx_c at any c satisfying $\langle a \rangle \supset \langle c \rangle \supset \langle b \rangle$ agrees with \approx_a and \approx_b .

As such, any pairs (a, \approx_a) and (b, \approx_b) identified in $\Omega(\sim)$ do indeed satisfy $\approx_a = \approx_b$.

At this point, checking that \preceq is a partial order is straightforward. Clearly, \preceq is reflexive, and transitivity of \preceq follows from transitivity of Green's preorder on Q and the transitivity of the equivalence relation defining $\Omega(\sim)$. Lastly, if $(a, \approx_a) \preceq (b, \approx_b)$ and $(b, \approx_b) \preceq (a, \approx_a)$, then the observations in the above paragraph imply that $\approx_a = \approx_b$, meaning that (a, \approx_a) and (b, \approx_b) are identified in $\Omega(\sim)$. This completes the proof.

Theorem 5.8. Fix a finite subposet $\Omega = \{p_0, p_1, \dots, p_d\}$ of the power set of $\{1, \dots, n\}$ with $p_0 = \emptyset$. Let I_0, \dots, I_d denote the ideals from Proposition 5.4, and define

$$M = \langle x_1^2, \dots, x_d^2 \rangle + \langle x_i x_j : p_i, p_j \ incomparable \rangle$$
 and $B = \langle x_i x_j - x_i x_k : p_i \supset p_j \ and \ p_i \supset p_k \rangle$. The ideal
$$I = B + x_1 I_1 + \dots + x_d I_d + M \subset \mathbb{k}[x_1, \dots, x_d, y]$$
 has $\Omega(\sim_I)$ isomorphic to Ω .

Proof. The only monomials in the variables x_1, \ldots, x_d that lie outside of I are either of degree 1 or have the form $x_i x_j$ for $p_i \supset p_j$ (in particular, I contains every monomial of total degree 3). The only prime congruences that occur are induced by I_0, \ldots, I_d ; I_0 induces the prime congruence at the origin, and I_i for $i \geq 1$ induces the prime congruence at the elements corresponding to the monomials $\{x_i, x_i x_j : p_i \supseteq p_j\}$. The binomials generating B ensure that this set has exactly two distinct elements modulo I, the larger of which corresponds to the unique true witness, whose associated prime congruence is induced by I_i . Divisibility among the nonzero monomials modulo I ensures that $\Omega(\sim)$ is isomorphic to Ω .

Remark 5.9. In general, the set of associated prime congruences, as well as the prime congruence poset, can differ if different classes of witnesses are used in place of true witnesses, e.g., they may have different cardinalities. However, every witness for every congruence constructed in Theorems 5.5 and 5.8 is true. This means that, if one relaxes the problem to allow prime congruences at any more general class of wit-

nesses, then the resulting poset would be the same. Consequently, the content of this section also answers [3, Problem 17.9], the analogue of Problem 5.1 for binomial ideals. Indeed, upon referencing [3, Definitions 10.4 and 12.1], it can easily be verified that each ideal I defined in Theorems 5.5 or 5.8 decomposes as an intersection of mesoprimary ideals, whose poset of associated mesoprimes is also isomorphic to the given poset Ω .

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