

# ADJOINT FUNCTORS AND TRIPLES<sup>1</sup>

BY

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A triple  $\mathbf{F} = (F, \eta, \mu)$  in a category  $\mathcal{A}$  consists of a functor  $F : \mathcal{A} \rightarrow \mathcal{A}$  and morphisms  $\eta : 1_{\mathcal{A}} \rightarrow F$ ,  $\mu : F^2 \rightarrow F$  satisfying some identities (see §2, (T.1)–(T.3)) analogous to those satisfied in a monoid. Cotriples are defined dually.

It has been recognized by Huber [4] that whenever one has a pair of adjoint functors  $T : \mathcal{A} \rightarrow \mathcal{B}$ ,  $S : \mathcal{B} \rightarrow \mathcal{A}$  (see §1), then the functor  $TS$  (with appropriate morphisms resulting from the adjointness relation) constitutes a triple in  $\mathcal{B}$  and similarly  $ST$  yields a cotriple in  $\mathcal{A}$ .

The main objective of this paper is to show that this relation between adjointness and triples is in some sense reversible. Given a triple  $\mathbf{F}$  in  $\mathcal{A}$  we define a new category  $\mathcal{A}^F$  and adjoint functors  $T : \mathcal{A}^F \rightarrow \mathcal{A}$ ,  $S : \mathcal{A} \rightarrow \mathcal{A}^F$  such that the triple given by  $TS$  coincides with  $\mathbf{F}$ . There may be many adjoint pairs which in this way generate the triple  $\mathbf{F}$ , but among those there is a universal one (which therefore is in a sense the “best possible one”) and for this one the functor  $T$  is faithful (Theorem 2.2). This construction can best be illustrated by an example. Let  $\mathcal{A}$  be the category of modules over a commutative ring  $K$  and let  $\Lambda$  be a  $K$ -algebra. The functor  $F = \Lambda \otimes$  together with morphisms  $\eta$  and  $\mu$  resulting from the morphisms  $K \rightarrow \Lambda$ ,  $\Lambda \otimes \Lambda \rightarrow \Lambda$  given by the  $K$ -algebra structure of  $\Lambda$ , yield then a triple  $\mathbf{F}$  in  $\mathcal{A}$ . The category  $\mathcal{A}^F$  is then precisely the category of  $\Lambda$ -modules. The general construction of  $\mathcal{A}^F$  closely resembles this example. As another example, let  $\mathcal{A}$  be the category of sets and let  $F$  be the functor which to each set  $A$  assigns the underlying set of the free group generated by  $A$ . There results a triple  $\mathbf{F}$  in  $\mathcal{A}$  and  $\mathcal{A}^F$  is the category of groups.

Let  $\mathbf{G} = (\delta, \varepsilon, G)$  be a cotriple in a category  $\mathcal{A}$ . It has been recognized by Godement [3] and Huber [4], that the iterates  $G^n$  of  $G$  together with face and degeneracy morphisms

$$G^{n+1} \rightarrow G^n, \quad G^n \rightarrow G^{n+1}$$

defined using  $\varepsilon$  and  $\delta$  yield a simplicial structure which can be used to define homology and cohomology.

Now if  $\mathbf{F}$  is a triple in  $\mathcal{A}$ , then one has an adjoint pair  $T : \mathcal{A}^F \rightarrow \mathcal{A}$ ,  $S : \mathcal{A} \rightarrow \mathcal{A}^F$  and therefore one has an associated cotriple  $\mathbf{G}$  in  $\mathcal{A}^F$ . This in turn yields a simplicial complex for every object in  $\mathcal{A}^F$ , thus paving the way for homology and cohomology in  $\mathcal{A}^F$ . In §4 we show that under suitable

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conditions this complex is a projective resolution in a suitable relative sense as developed by us in [2].

For some further developments of the ideas presented here see a forthcoming dissertation of Jon M. Beck.

### 1. Review of adjoint functors

Given a category  $\mathcal{A}$  we use the symbol  $\mathcal{A}(A, A')$  to denote the set of all morphisms  $A \rightarrow A'$  in  $\mathcal{A}$  where  $A, A'$  are objects of  $\mathcal{A}$ .

We shall use the notation

$$(1.1) \quad a : S \dashv T : (\mathcal{A}, \mathcal{B})$$

whenever  $T : \mathcal{A} \rightarrow \mathcal{B}$  and  $S : \mathcal{B} \rightarrow \mathcal{A}$  are functors and  $a$  is an isomorphism

$$a : \mathcal{A}(S, \ ) \rightarrow \mathcal{B}( \ , T)$$

of functors. Explicitly for each pair  $A \in \mathcal{A}, B \in \mathcal{B}$   $a$  yields a bijection

$$a : \mathcal{A}(S(B), A) \rightarrow \mathcal{B}(A, T(A))$$

satisfying

$$(1.2) \quad a(gfS(h)) = T(g)a(fh)$$

for

$$h : B' \rightarrow B, \quad f : S(B) \rightarrow A, \quad g : A \rightarrow A'.$$

Under the relation (1.1) the functor  $S$  is said to be the coadjoint of  $T$ , and  $T$  is said to be the adjoint of  $S$ .

Setting

$$(1.3) \quad \alpha(A) = a^{-1}(1_{T(A)}) : ST(A) \rightarrow A$$

$$(1.4) \quad \beta(B) = a(1_{T(B)}) : B \rightarrow TS(B)$$

we obtain morphisms of functors

$$\alpha : ST \rightarrow 1_{\mathcal{A}}, \quad \beta : 1_{\mathcal{B}} \rightarrow TS$$

such that the compositions

$$S \xrightarrow{S\beta} STS \xrightarrow{\alpha S} S, \quad T \xrightarrow{\beta T} TST \xrightarrow{T\alpha} T$$

are identities. Conversely we have:

$$(1.5) \quad a(f) = T(f)\beta(B) \quad \text{for } f : S(B) \rightarrow A$$

$$(1.6) \quad a^{-1}(g) = \alpha(A)S(g) \quad \text{for } g : B \rightarrow T(A).$$

We shall write  $a \sim (\alpha, \beta)$ .

Given

$$a : S \dashv T : (\mathcal{A}, \mathcal{B})$$

$$c : R \dashv Q : (\mathcal{B}, \mathcal{C})$$

we have

$$ca : SR \dashv QT : (\mathfrak{A}, \mathfrak{C})$$

where  $ca$  is the composition

$$\mathfrak{A}(SR, \quad) \xrightarrow{a} \mathfrak{B}(R, T) \xrightarrow{c} \mathfrak{C}(\quad, QT).$$

If  $a \sim (\alpha, \beta)$ ,  $c \sim (\sigma, \tau)$  then

$$ca \sim (\sigma(Q\alpha R), (S\tau T)\beta).$$

Given

$$(1.7) \quad a : S \dashv T : (\mathfrak{A}, \mathfrak{B})$$

$$(1.8) \quad a' : S' \dashv T' : (\mathfrak{A}, \mathfrak{B})$$

and given morphisms  $\varphi : S' \rightarrow S$ ,  $\psi : T \rightarrow T'$  we write

$$(1.9) \quad \varphi \dashv \psi$$

if the following diagram is commutative:

$$\begin{array}{ccc} \mathfrak{A}(S, \quad) & \xrightarrow{a} & \mathfrak{B}(\quad, T) \\ \mathfrak{A}(\varphi, \quad) \downarrow & & \downarrow B(\quad, \psi) \\ \mathfrak{A}(S', \quad) & \xrightarrow{a'} & \mathfrak{B}(\quad, T') \end{array}$$

We note the following properties of adjointness of morphisms:

- (1.10) If in addition to the above we also have  $a'' : S'' \dashv T'' : (\mathfrak{A}, \mathfrak{B})$  and  $\varphi' \dashv \psi'$  for  $\varphi' : S'' \rightarrow S'$ ,  $\psi' : T' \rightarrow T''$  then  $\varphi\varphi' \dashv \psi\psi'$
- (1.11) If (1.7) and (1.8) hold and  $\varphi : S' \rightarrow S$  then there exists a unique  $\psi : T \rightarrow T'$  such that  $\varphi \dashv \psi$ . Further  $\psi$  is an isomorphism (or identity) if and only if  $\varphi$  is.
- (1.12) If (1.7), (1.8) and (1.9) hold and if  $c : R \dashv Q : (\mathfrak{B}, \mathfrak{C})$  then  $\varphi R \dashv Q\psi$  relative to the adjointness relations  $ca, ca'$ .

### 2. Triples

Let  $\mathfrak{A}$  be a category. A triple  $\mathbf{F} = (F, \eta, \mu)$  in  $\mathfrak{A}$  consists of a functor  $F : \mathfrak{A} \rightarrow \mathfrak{A}$  and of morphisms

$$\eta : 1_A \rightarrow F, \quad \mu : F^2 \rightarrow F$$

such that

$$(T.1) \quad \text{the composition } F \xrightarrow{\eta^F} F^2 \xrightarrow{\mu} F \text{ is the identity,}$$

$$(T.2) \quad \text{the composition } F \xrightarrow{F\eta} F^2 \xrightarrow{\mu} F \text{ is the identity,}$$

(T.3) the diagram

$$\begin{array}{ccc}
 F^3 & \xrightarrow{F\mu} & F^2 \\
 \mu F \downarrow & & \downarrow \mu \\
 F^2 & \xrightarrow{\mu} & F
 \end{array}$$

is commutative.

Dually a cotriple  $(\delta, \varepsilon, F)$  in  $\mathcal{A}$  is given by a triple  $(F^*, \varepsilon^*, \delta^*)$  in the dual category  $\mathcal{A}^*$ .

PROPOSITION 2.1. *Let*

$$(2.1) \quad a : S \dashv T : (\mathcal{A}, \mathcal{B})$$

with  $a \sim (\alpha, \beta)$ . Then

$$\nabla(a) = (TS, \beta, T\alpha S)$$

is a triple in  $\mathcal{B}$ . Dually

$$\Delta(a) = (S\beta T, \alpha, ST)$$

is a cotriple in  $\mathcal{A}$ .

We say that  $\nabla(a)$  is generated by (2.1) and  $\Delta(a)$  is cogenerated by (2.1).

*Proof.* Since  $(T\alpha)(\beta T) = 1_T$  we have  $(ST\alpha)(S\beta T) = 1_{ST}$  and (T.1) holds. Since  $(\alpha S)(S\beta) = 1_S$  we have  $(\alpha ST)(S\beta T) = 1_{ST}$  so that (T.2) holds. Relation (T.3) follows from the commutative diagram

$$\begin{array}{ccc}
 STST & \xrightarrow{\alpha ST} & ST \\
 ST\alpha \downarrow & & \downarrow \alpha \\
 ST & \xrightarrow{\alpha} & 1_A.
 \end{array}$$

THEOREM 2.2. *Every triple  $\mathbf{F} = (F, \eta, \mu)$  in a category  $\mathcal{A}$  admits a generator*

$$(2.2) \quad a : S \dashv T : (\mathcal{B}, \mathcal{A}).$$

Moreover, there exists a universal generator

$$(2.3) \quad \alpha^F : S^F \dashv T^F : (\mathcal{A}^F, \mathcal{A})$$

of  $\mathbf{F}$  such that for any generator (2.2) of  $\mathbf{F}$  there exists a unique functor  $L : \mathcal{B} \rightarrow \mathcal{A}^F$  such that

$$(2.4) \quad LS = S^F, \quad L\alpha = \alpha^F L.$$

These relations imply

$$(2.5) \quad T^F L = T.$$

In addition the functor  $T^F$  is faithful.

*Proof.* We define the objects of  $\mathcal{G}^F$  to be the pairs  $(A, \varphi)$  where  $A$  is an object of  $\mathcal{G}$  and  $\varphi : F(A) \rightarrow A$  is a morphism in  $A$  satisfying

$$(2.6) \quad \varphi\eta(A) = 1_A, \quad \varphi F(\varphi) = \varphi\mu(A).$$

A morphism  $[f] : (A, \varphi) \rightarrow (A', \varphi')$  in  $\mathcal{G}^F$  is given by a morphism  $f : A \rightarrow A'$  in  $\mathcal{G}$  such that

$$(2.7) \quad f\varphi = \varphi'F(f).$$

If  $[g] : (A', \varphi') \rightarrow (A'', \varphi'')$  then we define  $[g][f] = [gf]$ . This defines the category  $\mathcal{G}^F$ . The functor  $T^F : \mathcal{G}^F \rightarrow \mathcal{G}$  is given by

$$T^F(A, \varphi) = A, \quad T^F[f] = f.$$

Clearly  $T^F$  is faithful.

The functor  $S^F : \mathcal{G} \rightarrow \mathcal{G}^F$  is defined by

$$S^F(A) = (F(A), \mu(A)), \quad S^F(f) = [F(f)]$$

for  $f : A \rightarrow A'$  in  $\mathcal{G}$ .

Since  $T^F S^F = F$  we define

$$\beta^F = \eta : 1_A \rightarrow F = T^F S^F.$$

Since  $[\varphi] : (F(A), \mu(A)) \rightarrow (A, \varphi)$  is a morphism in  $\mathcal{G}^F$  we define

$$\alpha^F : S^F T^F \rightarrow 1_{\mathcal{G}^F}, \quad \alpha^F(A, \varphi) = [\varphi].$$

For each  $A$  in  $\mathcal{G}$ , the composition

$$S^F(A) \xrightarrow{S^F \beta^F} S^F T^F S^F(A) \xrightarrow{\alpha^F S^F} S^F(A)$$

becomes the composition

$$(F(A), \mu(A)) \xrightarrow{[F\eta(A)]} (F^2(A), \mu F(A)) \xrightarrow{[\mu(A)]} (F(A), \mu(A))$$

which is the identity. Similarly the composition

$$T^F(A, \varphi) \xrightarrow{\beta^F T^F} T^F S^F T^F(A, \varphi) \xrightarrow{T^F \alpha^F} T^F(A, \varphi)$$

becomes the composition

$$A \xrightarrow{\eta(A)} F(A) \xrightarrow{\varphi} A$$

which again is the identity. This yields (2.3) with  $a^F \sim (\alpha^F, \beta^F)$ . Since

$$T^F \beta^F S^F(A) = T^F \beta^F(F(A), \mu(A)) = T^F[\mu(A)] = \mu(A)$$

we have  $\nabla(a^F) = F$  so that (2.3) is a generator for  $\mathbf{F} = (F, \eta, \mu)$ .

To show that (2.3) has the universal property consider an arbitrary generator (2.2) of  $\mathbf{F}$ .

Given an object  $B$  in  $\mathfrak{B}$  we have  $\alpha(B) : ST(B) \rightarrow B$  and therefore

$$T\alpha(B) : FT(B) = TST(B) \rightarrow T(B).$$

We assert that  $(T(B), T\alpha(B))$  is an object in  $\mathfrak{A}^F$ . Firstly, the composition

$$T(B) \xrightarrow{\eta T} FT(B) \xrightarrow{T\alpha} T(B)$$

is the identity since  $\eta = \beta$  and  $(T\alpha)(\beta T) = 1_T$ . Secondly, from the commutative diagram

$$\begin{array}{ccc} STST & \xrightarrow{ST\alpha} & ST \\ \alpha ST \downarrow & & \downarrow \alpha \\ ST & \xrightarrow{\alpha} & 1_B \end{array}$$

we deduce

$$(T\alpha)(FT\alpha) = (T\alpha)(TST\alpha) = (T\alpha)(T\alpha ST) = (T\alpha)(\mu T).$$

Thus we may define

$$L(B) = (T(B), T\alpha(B)).$$

If  $f : B \rightarrow B'$  in  $\mathfrak{B}$  then

$$\begin{aligned} T(f)T\alpha(B) &= T(f\alpha(B)) = T(\alpha(B')ST(f)) = (T\alpha(B'))(TST(f)) \\ &= T\alpha(B')FT(f). \end{aligned}$$

Thus (2.7) holds and  $[T(f)] : L(B) \rightarrow L(B')$  is a morphism in  $\mathfrak{A}^F$ . Thus setting  $L(f) = [T(f)]$  we obtain a functor  $L : \mathfrak{B} \rightarrow \mathfrak{A}^F$ . Clearly

$$\begin{aligned} LS(A) &= (TS(A), T\alpha S(A)) = (F(A), \mu(A)) = S^F(A) \\ LS(f) &= [TS(f)] = [F(f)] = S^F(f) \end{aligned}$$

so that  $LS = S^F$ . Also

$$\begin{aligned} T^F L(B) &= T^F(T(B), T\alpha(B)) = T(B) \\ T^F L(f) &= T^F[T(f)] = T(f) \end{aligned}$$

so that  $T^F L = T$ . Further

$$\alpha^F L(B) = \alpha^F(T(B), T\alpha(B)) = [T\alpha(B)] = L\alpha(B)$$

so that  $\alpha^F L = L\alpha$ . Thus (2.4) and (2.5) hold.

To show that  $L$  is unique consider another functor  $L' : \mathfrak{B} \rightarrow \mathfrak{A}^F$  satisfying (2.4). Let  $B \in \mathfrak{B}$  and let  $L'(B) = (A, \varphi)$ . Then

$$\begin{aligned} A &= T^F(A, \varphi) = T^F L'(B) = T(B) \\ \varphi &= T^F[\varphi] = T^F \alpha^F(A, \varphi) = T^F \alpha^F L'(B) = T^F L' \alpha(B) = T\alpha(B) \end{aligned}$$

and thus  $L'(B) = L(B)$ . If  $f : B \rightarrow B'$  then both  $L(f)$  and  $L'(f)$  are morphisms  $L(B) \rightarrow L(B')$ . Since  $T^F L = T = T^F L'$  it follows that  $T^F L(f) = T^F L'(f)$ . The functor  $T^F$  being faithful it follows that  $L(f) = L'(f)$ . Thus  $L = L'$ . Since the uniqueness proof uses only (2.4) while  $L$  satisfies (2.5), it follows that (2.5) is a consequence of (2.4). This concludes the proof of 2.2.

PROPOSITION 2.3. *Let*

$$a : S \dashv T : (\mathfrak{A}, \mathfrak{B})$$

with  $a \sim (\alpha, \beta)$  and let  $\mathbf{F} = (F, \eta, \mu)$  be a triple in  $\mathfrak{A}$ . Then

$$\mathbf{F}' = (TFS, (T\eta S)\beta, (T\mu S)(TF\alpha FS))$$

is a triple in  $\mathfrak{B}$ .

*Proof.* A purely computational proof was given by Huber [4, p. 10]. The following proof is somewhat more conceptual. Let

$$c : R \dashv Q : (\mathfrak{C}, \mathfrak{A})$$

with  $c \sim (\sigma, \tau)$  be a generator of  $\mathbf{F}$ . Thus  $\mathbf{F} = (QR, \tau, Q\sigma R)$ . We then have by §1 the adjoint relationship

$$ac : RS \dashv TQ : (\mathfrak{C}, \mathfrak{B})$$

with

$$ac \sim (\sigma(R\alpha Q), (T\tau S)\beta).$$

Then by 2.1,  $ac$  generates the triple

$$\begin{aligned} \nabla(ac) &= (TQRS, (T\tau S)\beta, (TQ\sigma RS)(TQR\alpha QRS)) \\ &= (TFS, (T\eta S)\beta, (T\mu S)(TF\alpha FS)) = \mathbf{F}'. \end{aligned}$$

### 3. Adjoint triples

Let

$$(3.1) \quad a : F \dashv G : (\mathfrak{A}, \mathfrak{A}).$$

We define

$$a^n : F^n \dashv G^n : (\mathfrak{A}, \mathfrak{A})$$

for  $n = 0, 1, \dots$  inductively as follows. For  $n = 0$ ,  $F^0 = G^0 = 1$  and  $a^0$  is the identity. For  $n > 0$ ,  $a^n$  is the composition

$$\mathfrak{A}(F^n, \quad) \xrightarrow{a} \mathfrak{A}(F^{n-1}, G) \xrightarrow{a^{n-1}} \mathfrak{A}(\quad, G^n).$$

In particular  $a^1 = a$ .

If  $\mathbf{F} = (F, \eta, \mu)$  is a triple in  $\mathfrak{A}$  and  $\mathbf{G} = (\delta, \varepsilon, G)$  is a cotriple in  $\mathfrak{A}$ , then we write

$$a : \mathbf{F} \dashv \mathbf{G}$$

if (3.1) holds and if

$$\eta \dashv \varepsilon, \quad \mu \dashv \delta$$

under  $a^0, a^1$  and  $a^2$ .

**PROPOSITION 3.1.** *If  $\mathbf{F} = (F, \eta, \mu)$  is a triple in  $\mathcal{A}$  and if  $a : F \dashv G : (\mathcal{A}, \mathcal{A})$  then there exists a unique cotriple  $\mathbf{G} = (\delta, \varepsilon, G)$  such that  $a : \mathbf{F} \dashv \mathbf{G}$ .*

*Proof.* We must define  $\varepsilon$  and  $\delta$  by the conditions  $\mu \dashv \delta, \eta \dashv \varepsilon$ . Then by (1.12),  $\eta F \dashv G\varepsilon$  and by (1.10)  $\mu(\eta F) \dashv (G\varepsilon)\delta$ . Since  $\delta(\varepsilon F) = 1_F$  it follows that  $(G\varepsilon)\delta = 1_G$ . Similarly we show that  $(\varepsilon G)\delta = 1_G$  and that  $\delta(G\delta) = \delta(\delta G)$ . Thus  $G = (\delta, \varepsilon, G)$  is a cotriple as required.

**PROPOSITION 3.2.** *Let*

$$a : S \dashv T : (\mathcal{A}, \mathcal{B})$$

$$c : R \dashv S : (\mathcal{B}, \mathcal{A})$$

*Then in  $\mathcal{A}$  we have*

$$ca : \nabla(c) \dashv \Delta(a).$$

*Proof.* Let  $a \sim (\alpha, \beta), c \sim (\sigma, \tau)$ . Then by definition

$$\Delta(a) = (S\beta T, \alpha, ST), \quad \nabla(c) = (SR, \tau, S\sigma R)$$

By §1 we have

$$ca : SR \dashv ST.$$

Thus it suffices to show that

$$\tau \dashv \alpha, \quad S\sigma R \dashv S\beta T.$$

The relation  $\tau \dashv \alpha$  is the commutativity relation in the diagram

$$\begin{array}{ccccc} \mathcal{A}(SR, \quad) & \xrightarrow{\alpha} & \mathcal{B}(R, T) & \xrightarrow{c} & \mathcal{A}(\quad, ST) \\ & \searrow & & \swarrow & \\ \mathcal{A}(\tau, \quad) & & & & \mathcal{A}(\quad, \alpha) \\ & & \mathcal{A}(\quad, \quad) & & \end{array}$$

which follows from the commutative diagram

$$\begin{array}{ccccc} & & \mathcal{B}(R, T) & & \\ & \swarrow \alpha^{-1} & \downarrow S & \searrow c & \\ \mathcal{A}(SR, \quad) & \longleftarrow & \mathcal{A}(SR, ST) & \longrightarrow & \mathcal{A}(\quad, ST) \\ & \swarrow \mathcal{A}(SR, \alpha) & \downarrow \mathcal{A}(\tau, \alpha) & \searrow \mathcal{A}(\tau, ST) & \\ & \mathcal{A}(\tau, \quad) & \mathcal{A}(\quad, \quad) & \mathcal{A}(\quad, \alpha) & \end{array}$$

The relation  $S\sigma R \dashv S\beta T$  follows from the commutativity of the diagram

$$\begin{array}{ccccc}
 \mathfrak{A}(SR, \ ) & \xrightarrow{a} & \mathfrak{B}(R, T) & \xrightarrow{c} & \mathfrak{A}(\ , ST) \\
 \mathfrak{A}(S\sigma R, \ ) \downarrow & & \downarrow S & & \downarrow \mathfrak{A}(\ , S\beta T) \\
 \mathfrak{A}(SRSR, \ ) & \xrightarrow{ca} & \mathfrak{A}(SR, ST) & \xrightarrow{ca} & \mathfrak{A}(\ , STST).
 \end{array}$$

The commutativity in the left square follows from the commutative diagram

$$\begin{array}{ccccc}
 \mathfrak{A}(SR, \ ) & \xrightarrow{a} & \mathfrak{B}(R, T) & \xrightarrow{1} & \mathfrak{B}(R, T) \\
 \mathfrak{A}(S\sigma R, \ ) \downarrow & & \mathfrak{B}(\sigma R, T) \downarrow B(R\tau, T) & \nearrow & \downarrow S \\
 \mathfrak{A}(SRSR, \ ) & \xrightarrow{a} & \mathfrak{B}(RSR, T) & \xrightarrow{c} & \mathfrak{A}(SR, ST)
 \end{array}$$

and the commutativity of the right square is shown dually.

PROPOSITION 3.3. *Let  $\mathbf{F} = (F, \eta, \mu)$  be a triple in  $\mathfrak{A}$  with a universal generator*

$$(3.2) \quad c : R \dashv S : (\mathfrak{B}, \mathfrak{A}).$$

*Then the following properties are equivalent:*

- (i)  *$F$  has an adjoint  $F \dashv G : (\mathfrak{A}, \mathfrak{A})$ ,*
- (ii) *The triple  $\mathbf{F}$  has an adjoint cotriple  $\mathbf{F} \dashv \mathbf{G}$ ,*
- (iii) *The functor  $S$  has an adjoint*

$$(3.3) \quad a : S \dashv T : (\mathfrak{A}, \mathfrak{B}).$$

*If the above is the case, then (3.3) is a universal cogenerator for  $\mathbf{G}$ .*

*Proof.* The implication (i)  $\Rightarrow$  (ii) follows from 3.1 while the implication (iii)  $\Rightarrow$  (i) follows from 3.2. There remains to prove the implication (ii)  $\Rightarrow$  (iii) and the last statement.

Let

$$(3.4) \quad a' : S' \dashv T' : (\mathfrak{B}', \mathfrak{A})$$

be a universal cogenerator for  $\mathbf{G}$ . Assuming that  $e : \mathbf{F} \dashv \mathbf{G}$  we shall construct an isomorphism  $L : B \rightarrow B'$  of categories such that  $S'L = S$ . Once this is done we replace  $S', T', B'$  by  $S, T = T'L, B$  so that (3.4) becomes

$$(3.5) \quad a : S \dashv T : (\mathfrak{B}, \mathfrak{A})$$

which is still a universal cogenerator for  $\mathbf{G}$ .

In order to construct  $L$  we take for (3.2) and (3.4) the explicit constructions as given in §2. Thus an object of  $\mathfrak{B}$  is a pair  $(A, \varphi)$  with  $\varphi : F(A) \rightarrow A$ , satisfying

$$\varphi\eta(A) = 1_A, \quad \varphi F(\varphi) = \varphi\eta(A).$$

Dually, an object of  $\mathfrak{B}'$  is a pair  $(\psi, A)$  with  $\psi : A \rightarrow G(A)$  such that

$$\varepsilon(A)\psi = 1_A, \quad G(\psi)\psi = \delta(A)\psi.$$

Now given  $(A, \varphi) \in \mathfrak{B}$ , let  $L(A, \varphi) = (e\varphi, A)$ . Since  $\eta \dashv \varepsilon$  we have the commutative diagram

$$\begin{array}{ccc} \mathfrak{A}(F(A), A) & \xrightarrow{e} & \mathfrak{A}(A, G(A)) \\ \mathfrak{A}(\eta(A), A) & \searrow & \swarrow \mathfrak{A}(A, \varepsilon(A)) \\ & \mathfrak{A}(A, A) & \end{array}$$

so that

$$\varepsilon(A)e\varphi = \varphi\eta(A) = 1_A.$$

Since  $\mu \dashv \delta$  we have the commutative diagram

$$\begin{array}{ccc} \mathfrak{A}(F(A), A) & \xrightarrow{e} & \mathfrak{A}(A, G(A)) \\ \mathfrak{A}(\mu(A), A) & \downarrow & \downarrow \mathfrak{A}(A, \delta(A)) \\ \mathfrak{A}(F^2(A), A) & \xrightarrow{e^2} & \mathfrak{A}(A, G^2(A)) \end{array}$$

so that

$$e^2(\varphi\eta(A)) = \delta(A)e\varphi.$$

On the other hand

$$e^2(\varphi F(\varphi)) = e[e(\varphi)\varphi] = G(e\varphi)e\varphi.$$

Thus  $G(e\varphi)e\varphi = \delta(A)e\varphi$  so that  $L(A, \varphi) \in \mathfrak{B}'$ .

Let  $[f] : (A, \varphi) \rightarrow (A', \varphi')$  be a morphism in  $\mathfrak{B}$ . Then by definition  $f : A \rightarrow A'$  and  $\varphi'F(f) = f\varphi$ . Then

$$G(f)e\varphi = e(f\varphi) = e(\varphi'F(f)) = (e\varphi')f$$

so that  $[f] : (e\varphi, A) \rightarrow (e\varphi', A')$  is a morphism in  $\mathfrak{B}'$ . This yields the required functor  $L$ .

#### 4. Relations with projective classes

We rapidly review some of the notions discussed in [2] and needed here.

Let  $\mathfrak{A}$  be a pointed category. A *sequence* in  $\mathfrak{A}$  is a diagram

$$(*) \quad A' \xrightarrow{f} A \xrightarrow{g} A''$$

such that  $gf = 0$ . The sequence  $(*)$  is *exact* if  $f$  admits a factorization  $f = kl$  where  $k$  is a kernel of  $g$  and  $l$  is an epimorphism.

Let  $\mathfrak{E}$  be a class of sequences  $(*)$  in  $\mathfrak{A}$ . An object  $P$  of  $\mathfrak{A}$  is  $\mathfrak{E}$ -projective if for every sequence  $(*)$  in  $\mathfrak{E}$  the sequence

$$(**) \quad A(P, A') \rightarrow A(P, A) \rightarrow A(P, A'')$$

is exact in the category of pointed sets. The class  $\mathfrak{E}$  is called *projective* if the following two conditions hold: (1) If  $(*)$  is sequence in  $\mathfrak{A}$  such that  $(**)$  is exact for every  $\mathfrak{E}$ -projective object  $P$ , then  $(*)$  is in  $\mathfrak{E}$ ; (2) For every  $g : A \rightarrow A''$  in  $\mathfrak{A}$  there exists a sequence  $(*)$  in  $\mathfrak{E}$  in which  $A'$  is  $\mathfrak{E}$ -projective.

If the class of all exact sequences in  $\mathcal{A}$  is projective then  $\mathcal{A}$  is called *projectively perfect*.

Let  $\mathcal{A}$  be a pointed category with kernels and  $\mathbf{G} = (\delta, \varepsilon, G)$  a cotriple in  $\mathcal{A}$ , where the functor  $G : A \rightarrow A$  is pointed.

We verify that for any  $A, B \in \mathcal{A}$  the morphism

$$(4.1) \quad \alpha(G(B), \varepsilon(A)) : \alpha(G(B), G(A)) \rightarrow \alpha(G(B), A)$$

is surjective. Indeed, let  $\varphi : G(B) \rightarrow A$ . Then

$$\varepsilon(A)G(\varphi)\delta(B) = \varphi\varepsilon(G(B))\delta(B) = \varphi.$$

By [2, Ch. I, §6] this implies that the cotriple  $\mathbf{G}$  determines a projective class  $\mathfrak{G}$  in  $\mathcal{A}$  as follows. A sequence

$$A' \xrightarrow{f} A \xrightarrow{g} A''$$

(with  $gf = 0$ ) is in  $\mathfrak{G}$  provided the sequence

$$\alpha(G(B), A') \rightarrow \alpha(G(B), A) \rightarrow \alpha(G(B), A'')$$

is exact for every  $B \in \mathcal{A}$ . The  $\mathfrak{G}$ -projective objects are the objects  $G(B)$ ,  $B \in \mathcal{A}$  and their retracts. The class  $\mathfrak{G}$  is exact if and only if  $\varepsilon(A)$  is an epimorphism for every  $A \in \mathcal{A}$ .

The kernel functor of  $\varepsilon$

$$0 \rightarrow K \xrightarrow{\kappa} G \xrightarrow{\varepsilon} 1_{\mathcal{A}}$$

leads to a sequence

$$(4.2) \quad \dots \rightarrow GK^3 \xrightarrow{d_3} GK^2 \xrightarrow{d_2} GK \xrightarrow{d_1} G \xrightarrow{\varepsilon} 1_{\mathcal{A}} \rightarrow 0$$

where  $K^0 = 1_{\mathcal{A}}$ ,  $K^n = KK^{n-1}$  and  $d_n : GK^n \rightarrow GK^{n-1}$  is the composition

$$GK^n \xrightarrow{\varepsilon K^n} K^n \xrightarrow{\kappa K^{n-1}} GK^{n-1}.$$

Applied to any object  $A$  in  $\mathcal{A}$  the sequence (4.2) yields a sequence in  $\mathfrak{G}$  and since  $GK^n(A)$  is  $\mathfrak{G}$ -projective, there results a  $\mathfrak{G}$ -projective resolution of  $A$ , called the *canonical resolution* (relative to the cotriple  $\mathbf{G}$ ).

Now consider the morphisms

$$\varepsilon^i : G^{n+1} \rightarrow G^n, \quad \delta^i : G^{n+1} \rightarrow G^{n+2} \quad i = 0, 1, \dots, n$$

defined as follows

$$\varepsilon^i = G^i \varepsilon G^{n-i}, \quad \delta^i = G^i \delta G^{n-i}.$$

The “face” operators  $\varepsilon^i$  and the “degeneracy” operators  $\delta^i$  satisfy the usual simplicial identities so that there results a simplicial functor  $\tilde{G}$  with  $G^{n+1}$  in degree  $n$ . If the category  $\mathcal{A}$  is preadditive then we may construct the boundary operator

$$\partial_n : G^{n+1} \rightarrow G^n$$

by setting

$$\partial_n = \sum_{i=0}^n (-1)^i \varepsilon^i.$$

There results a complex

$$(4.3) \quad \dots \rightarrow G^4 \xrightarrow{\partial_3} G^3 \xrightarrow{\partial_2} G^2 \xrightarrow{\partial_1} G^1 \xrightarrow{\varepsilon} 1_A \rightarrow 0$$

with  $\varepsilon$  as augmentation. This is the *standard complex* of the cotriple  $\mathbf{G}$ .

**PROPOSITION 4.1.** *Let  $\mathcal{A}$  be a preadditive category with kernels, and let  $\mathbf{G} = (\delta, \varepsilon, G)$  be a cotriple in  $\mathcal{A}$  with  $G$  an additive functor. If for each  $A \in \mathcal{A}$  the morphism  $G\kappa(A)$  is a kernel of  $G\varepsilon(A)$ , then for each  $A \in \mathcal{A}$ , the sequence (4.3) applied to  $A$  is a  $\mathcal{G}$ -projective resolution of  $A$ .*

*Proof.* Since  $(G\varepsilon)\delta = 1_G$  it follows from the hypothesis that

$$0 \rightarrow GK \xrightarrow{G\kappa} G^2 \xrightarrow{G\varepsilon} G \rightarrow 0$$

is a split exact sequence. Since the functor  $G$  is additive, it follows that

$$0 \rightarrow G^n K \xrightarrow{G^n \kappa} G^{n+1} \xrightarrow{G^n \varepsilon} G^n \rightarrow 0$$

is split exact.

We define morphisms

$$\tau_n : GK^n \rightarrow G^{n+1}$$

by induction as follows:  $\tau_0 = 1_G : G \rightarrow G$ ; for  $n > 0$ ,  $\tau_n$  is the composition

$$GK^n \xrightarrow{\tau_{n-1}K} G^n K \xrightarrow{G^n \kappa} G^{n+1}.$$

We verify that for  $n > 0$

$$\tau_{n-1} d_n = \varepsilon^0 \tau_n, \quad \varepsilon^i \tau_n = 0 \quad \text{if } i > 0$$

There results a commutative diagram

$$(4.4) \quad \begin{array}{ccccccc} \dots & \rightarrow & GK^3 & \xrightarrow{d_3} & GK^2 & \xrightarrow{d^2} & GK & \xrightarrow{d_1} & G \\ & & \downarrow \tau_3 & & \downarrow \tau_2 & & \downarrow \tau_1 & & \downarrow \tau_0 \\ \dots & \rightarrow & G^4 & \xrightarrow{\partial_3} & G^3 & \xrightarrow{\partial_2} & G^2 & \xrightarrow{\partial_1} & G \end{array}$$

The upper row is the canonical resolution. To show that the lower row also is a resolution (with augmentation  $\varepsilon : G \rightarrow 1_{\mathcal{A}}$ ) it suffices to show that the morphism (4.4) is a homotopy equivalence of complexes. To verify this it suffices to show that the upper row of (4.4) is the normalized subcomplex of the simplicial complex  $\tilde{G}$  with  $\tau$  as inclusion (for a neat exposition see [1, §3]). To do this it suffices to show that  $\tau_n(A) : GK^n(A) \rightarrow G^{n+1}(A)$  is the simultaneous kernel of  $\varepsilon^i(A) : G^{n+1}(A) \rightarrow G^n(A)$  for  $i = 1, 2, \dots, n$ . This means that if  $f : C \rightarrow G^{n+1}(A)$  is such that  $\varepsilon^i(A)f = 0$  for  $i = 1, 2, \dots, n$ , then

there exists a unique  $g : C \rightarrow GK^n(A)$  such that  $\tau_n g = f$ . For  $n = 0$  this is clear. We now assume that  $n > 0$  and proceed by induction.

Since  $G^n \kappa : G^n K \rightarrow G^{n+1}$  is the kernel of  $\varepsilon^n = G^n \varepsilon : G^{n+1} \rightarrow G^n$ , therefore,  $f$  admits a unique factorization

$$(4.5) \quad C \xrightarrow{f'} G^n K(A) \xrightarrow{G^n \kappa} G^{n+1}(A).$$

Let  $B = K(A)$  and consider the commutative diagrams for  $i = 1, 2, \dots, n - 1$ :

$$\begin{array}{ccc} C \xrightarrow{f'} G^i(B) & \xrightarrow{\varepsilon^i(B)} & G^{i-1}(B) \\ & \searrow f & \downarrow G^i \kappa \\ & & G^{i+1}(A) \xrightarrow{\varepsilon^i(A)} G^i(A). \end{array}$$

We have  $(G^{n-1} \kappa)(\varepsilon^i(B))f' = 0$  and since  $G^{n-1} \kappa$  is a kernel it follows that  $\varepsilon^i(B)f' = 0$  for  $i = 1, 2, \dots, n - 1$ . Thus, by the inductive hypothesis,  $f'$  admits a unique factorization

$$C \xrightarrow{g} GK^{n-1}(B) \xrightarrow{\tau_{n-1}(B)} G^n(B).$$

Combining this with the factorization (4.5) of  $f$  we obtain a unique factorization

$$C \xrightarrow{g} GK^n(A) \xrightarrow{\tau_n(A)} G^{n+1}(A)$$

of  $f$ , as required.

An alternative proof of 4.1 may be given as follows. Denote by  $\hat{G}$  the complex (4.3) with the augmentation included (i.e. with  $1_A$  in degree  $-1$  and with  $\partial_0 = \varepsilon$ ). Next show that the complex  $\hat{G}(G(A))$  is contractible (i.e. is split exact). This is done by defining  $s_n : G^{n+1} \rightarrow G^{n+2}$  by  $s_{-1} = 0, s_n = (-1)^{n-1} G^n \delta$  for  $n > 0$ . Then a calculation shows that

$$\partial_{n+1} s_n + s_{n-1} \partial_n = 1_{G^{n+1}}.$$

It follows that for every  $B \in \mathcal{A}$  we have

$$(4.6) \quad H\mathcal{A}(G(B), \hat{G}(G(A))) = 0.$$

Next observe that the sequence

$$0 \rightarrow K(A) \rightarrow G(A) \rightarrow A \rightarrow 0$$

is in  $\mathcal{G}$  while the sequences

$$0 \rightarrow G^n K(A) \rightarrow G^{n+1}(A) \rightarrow G^n(A) \rightarrow 0$$

are split exact for  $n > 0$ . There results an exact sequence of complexes

$$0 \rightarrow \mathcal{A}(G(B), \hat{G}K(A)) \rightarrow \mathcal{A}(G(B), \hat{G}G(A)) \rightarrow \mathcal{A}(G(B), \hat{G}(A)) \rightarrow 0.$$

Thus (4.6) implies an isomorphism

$$H_n \mathfrak{A}(G(B), \hat{G}(A)) \approx H_{n-1} \mathfrak{A}(G(B), \hat{G}K(A)).$$

Since  $H_{-2} \mathfrak{A}(G(B), \hat{G}(A)) = 0$  for all  $A, B \in \mathfrak{A}$ , it follows inductively that  $H_n \mathfrak{A}(G(B), \hat{G}(A)) = 0$  for all  $n$ . Thus  $\hat{G}(A)$  is in  $\mathfrak{G}$ , as required.

This proof has the disadvantage of not exhibiting the canonical complex as the normalized standard complex.

### 5. Properties of universal generators

Let  $\mathbf{F} = (F, \eta, \mu)$  be a triple in a category  $\mathfrak{A}$  and let

$$(5.1) \quad \alpha^F : S^F \dashv T^F : (\mathfrak{A}^F, \mathfrak{A}), \quad \alpha^F \sim (\alpha^F, \beta^F),$$

be the universal generator of  $\mathbf{F}$ . From the explicit construction given in §2 it is clear that if  $\mathfrak{A}$  is a pointed (or additive) category and if  $F$  is a pointed (or additive) functor, then  $\mathfrak{A}^F$  is a pointed (or additive) category and the functors  $S^F$  and  $T^F$  are pointed (or additive).

**PROPOSITION 5.1.** *If  $\mathfrak{A}$  is a pointed category with kernels and if the functor  $F$  is pointed, then  $\mathfrak{A}^F$  is a pointed category with kernels and the functor  $T^F$  preserves and reflects kernels (i.e.  $g$  is a kernel of  $f$  in  $\mathfrak{A}^F$  if and only if  $T^F g$  is a kernel of  $T^F f$  in  $\mathfrak{A}$ ).*

*Proof.* Let  $[f] : (A, \varphi) \rightarrow (A'', \varphi'')$  be a morphism in  $\mathfrak{A}^F$  and let  $g : A' \rightarrow A$  be a kernel of  $f$  in  $\mathfrak{A}$ . Since

$$f\varphi^F(g) = \varphi''F(f)F(g) = \varphi''F(fg) = 0,$$

it follows that there exists a unique  $\varphi' : F(A') \rightarrow A'$  such that  $g\varphi' = \varphi^F(g)$ . Since

$$g\varphi'\eta(A') = \varphi^F(g)\eta(A') = \varphi\eta(A)g = g$$

and since  $g$  is a monomorphism, it follows that  $\varphi'\eta(A') = 1_{A'}$ . Similarly since

$$\begin{aligned} g\varphi'F(\varphi') &= \varphi^F(g)F(\varphi') = \varphi^F(\varphi)F^2(g) = \varphi\mu(A)F^2(g) \\ &= \varphi^F(g)\mu(A') = g\varphi'\mu(A') \end{aligned}$$

we have  $\varphi'F(\varphi') = \varphi'\mu(A')$ . Thus  $(A', \varphi')$  is in  $\mathfrak{A}^F$  and

$$[g] : (A', \varphi') \rightarrow (A, \varphi).$$

Now let  $[h] : (A_1, \varphi_1) \rightarrow (A, \varphi)$  be a morphism in  $B$  such that  $[f][h] = 0$ . Then  $fh = 0$  and there is a unique morphism  $k : A_1 \rightarrow A'$  such that  $h = gk$ . Then

$$gk\varphi_1 = h\varphi_1 = \varphi^F(h) = \varphi^F(g)F(k) = g\varphi'F(k)$$

and therefore  $k\varphi_1 = \varphi'F(k)$ . Thus  $[k] : (A_1, \varphi_1) \rightarrow (A', \varphi')$ , and  $[h] = [g][k]$ . Thus  $[g]$  is a kernel of  $[f]$  and the proposition is established.

**PROPOSITION 5.2.** *If  $\mathfrak{A}$  is a pointed category with cokernels and the functor  $F$*

is pointed and preserves cokernels, then the category  $\mathcal{A}^F$  has cokernels and the functor  $T^F$  preserves and reflects cokernels.

*Proof.* Let  $[g] : (A', \varphi') \rightarrow (A, \varphi)$  and let  $f : A \rightarrow A''$  be a cokernel of  $g$  in  $\mathcal{A}$ . Then  $F(f)$  is a cokernel of  $F(g)$ . Since  $f\varphi F(g) = fg\varphi' = 0$  it follows that there exists a unique  $\varphi'' : F(A'') \rightarrow A''$  such that  $f\varphi = \varphi''F(f)$ . Since

$$\varphi''\eta(A'')f = \varphi''F(f)\eta(A) = f\varphi\eta(A) = f$$

it follows that  $\varphi''\eta(A'') = 1_{A''}$ . Similarly since

$$\begin{aligned} \varphi''F(\varphi'')F^2(f) &= \varphi''F(f)F(\varphi) = f\varphi F(\varphi) = f\varphi\mu(A) \\ &= \varphi''F(f)\mu(A) = \varphi''\mu(A'')F^2(f) \end{aligned}$$

and since  $F^2(f)$  is an epimorphism, it follows that  $\varphi''F(\varphi'') = \varphi''\delta(A'')$ . Thus  $(A'', \varphi'')$  is in  $\mathcal{A}^F$  and  $[f] : (A, \varphi) \rightarrow (A'', \varphi'')$ . Now let

$$[h] : (A, \varphi) \rightarrow (A_1, \varphi_1)$$

be such that  $[h][g] = 0$ . Then  $hg = 0$  and there is a unique morphism  $k : A'' \rightarrow A_1$  such that  $h = kf$ . Then

$$k\varphi''F(f) = kf\varphi = h\varphi = \varphi_1F(h) = \varphi_1F(k)F(f)$$

implies that  $k\varphi'' = \varphi_1F(k)$ . Thus  $[k] : (A'', \varphi'') \rightarrow (A_1, \varphi_1)$  and  $[h] = [k][f]$ . Thus  $[f]$  is a cokernel of  $[g]$  and the proposition is established.

The above two propositions and known facts about abelian categories imply

**PROPOSITION 5.3.** *If the category  $\mathcal{A}$  is abelian and the functor  $F$  is additive and preserves cokernels, then the category  $\mathcal{A}^F$  is abelian and the functor  $T^F$  preserves and reflects exact sequences.*

Now assume that the category  $\mathcal{A}$  and the functor  $F$  are pointed, and let  $\mathcal{E}$  be a projective class in the category  $\mathcal{A}$ . By the adjoint theorem for projective classes [2, Ch. II, §2], there results in  $\mathcal{A}^F$  a projective class  $\mathcal{E}^F = (T^F)^{-1}\mathcal{E}$ . Explicitly a sequence

$$(A', \varphi') \xrightarrow{[f]} (A, \varphi) \xrightarrow{[g]} (A'', \varphi'')$$

is in  $\mathcal{E}^F$  if and only if

$$A' \xrightarrow{f} A \xrightarrow{g} A''$$

is in  $\mathcal{E}$ . The  $\mathcal{E}^F$  projective objects are the retracts of objects  $S^F(A) = (F(A), \mu(A))$  where  $A$  is  $\mathcal{E}$ -projective. Since the functor  $T^F$  is faithful, it follows [2, Ch. II, §2] that if the class  $\mathcal{E}$  is exact then the class  $\mathcal{E}^F$  also is exact.

In particular, if the category  $\mathcal{A}$  is projectively perfect, and  $\mathcal{E}_1$  is the class of all exact sequences in  $\mathcal{A}$ , then  $\mathcal{E}_1^F$  is the class of all exact sequences in the category  $\mathcal{A}^F$ , which is therefore projectively perfect.

If in  $\mathcal{A}$  we take  $\mathcal{E}_0$  to be the class of all split exact sequences, there results a projective class  $\mathcal{E}_0^F$  in  $\mathcal{A}^F$ . The  $\mathcal{E}_0^F$ -projective objects are the retracts of ob-

jects  $S^F(A) = (F(A), \mu(A))$  where  $A$  is any object of  $\mathfrak{A}$ . This class  $\mathfrak{E}_0^F$  may also be arrived at in a different way as follows. The relation (5.1) induces in  $A^F$  a cotriple

$$\mathbf{G} = \Delta(a) = (S\beta T, \alpha, ST)$$

where the superscript  $F$  has been omitted. This cotriple defines a projective class  $\mathfrak{G}$  in  $A^F$ . The  $\mathfrak{G}$ -projective objects of  $\mathfrak{A}^F$  are the retracts of objects  $ST(A, \varphi) = S(A) = (F(A), \mu(A))$  where  $(A, \varphi) \in \mathfrak{A}^F$ . Since the composition

$$S \xrightarrow{S\beta} STS \xrightarrow{\alpha S} S$$

is the identity, it follows that for any  $A \in \mathfrak{A}$ ,  $S(A)$  is a retract of  $STS(A)$ . Thus  $S(A)$  is  $\mathfrak{G}$ -projective. It follows that the  $\mathfrak{E}_0^F$ -projective and the  $\mathfrak{G}$ -projective objects coincide and thus  $\mathfrak{E}_0^F = \mathfrak{G}$ . In particular, the canonical resolution yields an  $\mathfrak{E}_0^F$ -projective resolution for every object of  $\mathfrak{A}^F$ .

If the category  $\mathfrak{A}$  is preadditive and the functor  $F$  is additive then the category  $\mathfrak{A}^F$  also is preadditive and the functors  $S, T$  and  $G = ST$  are additive. If further  $\mathfrak{A}$  has kernels, then  $\mathfrak{A}^F$  has kernels. We shall show that the conditions of 4.1 are satisfied and therefore the standard complex for the cotriple  $\mathbf{G}$  yields  $\mathfrak{E}_0^F$ -projective resolutions.

Indeed, the exact sequence

$$0 \rightarrow K \xrightarrow{\kappa} G \xrightarrow{\epsilon} 1_{\mathfrak{A}^F} \rightarrow 0$$

is the exact sequence

$$0 \rightarrow K \xrightarrow{\kappa} ST \xrightarrow{\alpha} 1_{\mathfrak{A}^F} \rightarrow 0.$$

Since  $T$  preserves exact sequences, it follows that

$$(5.2) \quad 0 \rightarrow TK \xrightarrow{T\kappa} TST \xrightarrow{T\alpha} T \rightarrow 0$$

is exact. Since  $(T\alpha)(\beta T) = 1_T$  it follows that the sequence (5.2) is split exact and therefore since  $S$  is additive

$$0 \rightarrow STK \xrightarrow{ST\kappa} STST \xrightarrow{ST\alpha} ST \rightarrow 0$$

is exact. Thus the sequence

$$0 \rightarrow GK \xrightarrow{G\kappa} G^2 \xrightarrow{G\epsilon} G \rightarrow 0$$

is exact, as required.

### 6. Examples

Let  $K$  be a commutative ring and  $\mathfrak{A}$  the category of  $K$ -modules. Then  $\mathfrak{A}$  is an abelian category.  $\mathfrak{A}$  is also a  $K$ -category, so that  $\mathfrak{A}(A, A')$  is again an object of  $\mathfrak{A}$ . The tensor product  $A \otimes B$  over  $K$  yields a functor  $\mathfrak{A} \times \mathfrak{A} \rightarrow \mathfrak{A}$

and we have the natural isomorphism

$$(6.1) \quad a : \mathcal{G}(\Lambda \otimes B, A) \rightarrow \mathcal{G}(B, \mathcal{G}(\Lambda, A)).$$

We shall also employ the standard identifications

$$(6.2) \quad K \otimes A = A, \quad \mathcal{G}(K, A) = A$$

Let  $\Lambda$  be a  $K$ -algebra. Then we have morphisms

$$(6.3) \quad \bar{\eta} : K \rightarrow \Lambda, \quad \bar{\mu} : \Lambda \otimes \Lambda \rightarrow \Lambda$$

satisfying the usual identities. There results a triple  $\mathbf{F} = (F, \eta, \mu)$  where  $F = \Lambda \otimes$

$$\eta(A) = \bar{\eta} \otimes A, \quad \mu(A) = \bar{\mu} \otimes A.$$

We also have a cotriple  $\mathbf{G} = (\delta, \varepsilon, G)$  where  $G = \mathcal{G}(\Lambda, \quad)$

$$\varepsilon(A) = \mathcal{G}(\bar{\eta}, A) : \mathcal{G}(\Lambda, A) \rightarrow \mathcal{G}(K, A) = A$$

and  $\delta(A)$  is the composition

$$\mathcal{G}(\Lambda, A) \xrightarrow{\mathcal{G}(\bar{\mu}, A)} \mathcal{G}(\Lambda \otimes \Lambda, A) \xrightarrow{a} \mathcal{G}(\Lambda, \mathcal{G}(\Lambda, A)).$$

The relation (6.1) yields  $a : F \dashv G$ . Further it is easy to verify that  $\eta \dashv \varepsilon$  and  $\mu \dashv \delta$ . Thus we have  $a : \mathbf{F} \dashv \mathbf{G}$ .

Let  ${}_{\Lambda}M$  be the category of left  $\Lambda$ -modules,  $T : {}_{\Lambda}M \rightarrow \mathcal{G}$  the usual "forgetful" functor and let  $S, R : A \rightarrow {}_{\Lambda}M$  be defined by  $S = \Lambda \otimes, S' = \mathcal{G}(\Lambda, \quad)$ . Then the relation (6.1) induces adjointness relations

$$\begin{aligned} a_1 : S \dashv T : ({}_{\Lambda}M, \mathcal{G}) \\ a_2 : T \dashv R : (\mathcal{G}, {}_{\Lambda}M) \end{aligned}$$

which are respectively the universal generator for  $\mathbf{F}$  and the universal co-generator for  $\mathbf{G}$ , in agreement with 3.3.

Using theorems of Watts [5] it is easy to show that (up to isomorphisms) the  $K$ -algebras yield all the triples  $\mathbf{F}$  and all the cotriples  $\mathbf{G}$  in  $\mathcal{G}$  such that  $F$  preserves cokernels and (arbitrary) coproducts (i.e. direct sums) while  $G$  preserves kernels and (arbitrary) products.

A  $K$ -coalgebra  $\Lambda$  is given by morphisms

$$(6.4) \quad \varepsilon : \Lambda \rightarrow K, \quad \bar{\delta} : \Lambda \rightarrow \Lambda \otimes \Lambda$$

satisfying the usual identities. There results a cotriple  $\mathbf{G} = (\delta, \varepsilon, G)$  in  $\mathcal{G}$  where  $G = \Lambda \otimes$

$$\varepsilon(A) = \varepsilon \otimes A, \quad \delta(A) = \bar{\delta} \otimes A.$$

We also have a triple  $\mathbf{F} = (F, \eta, \mu)$  where  $F = \mathcal{G}(\Lambda, \quad)$

$$\eta(A) = \mathcal{G}(\varepsilon, A) : A = \mathcal{G}(K, A) \rightarrow \mathcal{G}(\Lambda, A)$$

and  $\mu(A)$  is the composition

$$\mathcal{G}(\Lambda, \mathcal{G}(\Lambda, A)) \xrightarrow{a^{-1}} \mathcal{G}(\Lambda \otimes \Lambda, A) \xrightarrow{\mathcal{G}(\bar{\delta}, A)} A(\Lambda, A).$$

We still have the relations  $a : G \dashv F$ ,  $\varepsilon \dashv \eta$  and  $\delta \dashv \mu$ , so that in a sense we have  $\mathbf{G} \dashv \mathbf{F}$ . However  $\mathbf{G}$  being a cotriple and  $\mathbf{F}$  being a triple, 3.3 no longer applies. Indeed, the construction of the universal generator for  $\mathbf{F}$  yields the category  ${}^{\Lambda}M$  of left comodules (i.e.  $K$ -modules  $A$  with a morphism  $A \rightarrow \Lambda \otimes A$  satisfying suitable identities) while the construction of the universal cogenerator for  $\mathbf{G}$  yields the category  ${}^{\Lambda}M^*$  of  $\Lambda$ -contramodules (i.e.  $K$ -modules  $A$  with a morphism  $\mathcal{G}(\Lambda, A) \rightarrow A$  satisfying suitable identities) [see 2, Ch. III, §5]. These categories are in general distinct except when  $\Lambda$  is  $K$ -projective and finitely generated over  $K$ , in which case  ${}^{\Lambda}M$  and  ${}^{\Lambda}M^*$  are both isomorphic with the category  $\hat{\Lambda}M$  where  $\hat{\Lambda} = \mathcal{G}(\Lambda, K)$  is the  $K$ -algebra dual to the coalgebra  $\Lambda$ .

Again it can be shown that (up to isomorphisms) the  $K$ -coalgebras yield all the triples  $\mathbf{F}$  and all the cotriples  $\mathbf{G}$  in which  $F$  preserves kernels and (arbitrary) products while  $G$  preserves cokernels and (arbitrary) coproducts.

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