

A combinatorial proof of the Burdzy–Pitman conjecture*

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Abstract

First, we prove the following sharp upper bound for the number of high degree differences in bipartite graphs. Let (U, V, E) be a bipartite graph with $U = \{u_1, u_2, \dots, u_n\}$ and $V = \{v_1, v_2, \dots, v_n\}$. For $n \geq k > \frac{n}{2}$ we show that

$$\sum_{1 \leq i, j \leq n} \mathbb{1}\{|\deg(u_i) - \deg(v_j)| \geq k\} \leq 2k(n - k).$$

Second, as a corollary, we confirm the Burdzy–Pitman conjecture about the maximal spread of coherent and independent vectors: for $\delta \in (\frac{1}{2}, 1]$ we prove that

$$\mathbb{P}(|X - Y| \geq \delta) \leq 2\delta(1 - \delta)$$

for all random vectors (X, Y) satisfying $X = \mathbb{P}(A|\mathcal{G})$ and $Y = \mathbb{P}(A|\mathcal{H})$ for some event A and independent σ -fields \mathcal{G} and \mathcal{H} .

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1 Introduction

Let $(\Omega, \mathcal{F}, \mathbb{P})$ be a probability space. We say that a random vector (X, Y) defined on this probability space is coherent if there exist sub σ -fields $\mathcal{G}, \mathcal{H} \subset \mathcal{F}$ and an event $A \in \mathcal{F}$, such that

$$X = \mathbb{E}(\mathbb{1}_A|\mathcal{G}), \quad Y = \mathbb{E}(\mathbb{1}_A|\mathcal{H}).$$

We will also say that the joint distribution of such (X, Y) is coherent on $[0, 1]^2$. Hereinafter, we write $(X, Y) \in \mathcal{C}$ or $\mu \in \mathcal{C}$ to indicate that the vector (X, Y) or a distribution μ is coherent. By abuse of notation, \mathcal{C} will be used to denote a family of vectors and a family of distributions

As suggested in [6], a coherent vector can be interpreted as objective opinions of two autonomous experts about the odds of some random event A . In this context, we interpret \mathcal{G} and \mathcal{H} as different information sources that are available to the experts. Motivated by this application, it is natural to ask about the maximal possible spread of coherent opinions. Accordingly, Burdzy and Pal [1] proved that for any $\delta \in (\frac{1}{2}, 1]$ and

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$(X, Y) \in \mathcal{C}$ the probability $\mathbb{P}(|X - Y| \geq \delta)$ that the difference between coherent random variables exceeds a given threshold δ is bounded above by $\frac{2(1-\delta)}{2-\delta}$. They go on to show that this bound is sharp and it is attained by a random vector (X, Y) with X and Y being dependent random variables. We will write

$$\mathcal{C}_{\mathcal{I}} = \{(X, Y) \in \mathcal{C} : X \perp Y\},$$

to denote the family of those coherent vectors whose components are independent. In this paper we prove the following claim stated as a conjecture by Burdzy and Pitman in [2].

Theorem 1.1. *If $\delta \in (\frac{1}{2}, 1]$ and $(X, Y) \in \mathcal{C}_{\mathcal{I}}$ then*

$$\mathbb{P}(|X - Y| \geq \delta) \leq 2\delta(1 - \delta). \tag{1.1}$$

Moreover, the bound $2\delta(1 - \delta)$ is optimal.

In other words, Theorem 1.1 provides a sharp upper bound on the maximal spread of coherent opinions in the special case of two experts with access to independent sources of information. Let us point out that restricting δ to $(\frac{1}{2}, 1]$ does not diminish generality of the result. Consider $X' = \mathbb{1}_A$ and $Y' = \mathbb{P}(A)$ for an arbitrary event A with $\mathbb{P}(A) = \frac{1}{2}$. It is easy to see that $(X', Y') \in \mathcal{C}$. In this case, $\mathbb{P}(|X' - Y'| \geq \frac{1}{2}) = 1$. Hence, for all $\delta \in [0, \frac{1}{2}]$ the problem is trivial.

Let us briefly describe our approach and the organization of the paper. Although there are known alternative characterizations of coherent distributions [6, 7, 9], let us quote [2]:

For reasons we do not understand well, these general characterizations seem to be of little help in establishing the evaluations of $\epsilon(\delta)$ [i.e. $\mathbb{P}(|X - Y| \geq \delta)$] discussed above, or in settling a number of related problems about coherent distributions [...].

It is our belief that this is indeed so because of the underlying combinatorial nature of these problems. Discretization and combinatorial techniques appeared already in [1, 5]. Moreover, it is a remarkable fact that the properties of two-dimensional coherent vectors are closely related to the properties of degree sequences of bipartite graphs. An intriguing example of this phenomenon can be found in [12]. Therefore, in order to take advantage of the combinatorial nature of the claim made in Theorem 1.1, we start by discussing its graph-theoretic version. More precisely, we prove the following theorem.

Theorem 1.2. *Let $G = (U, V, E)$ be a bipartite graph with an equal bipartition, i.e.*

$$U = \{u_1, u_2, \dots, u_n\}, \quad V = \{v_1, v_2, \dots, v_n\},$$

for some $n \in \mathbb{Z}_+$. For $n \geq k > \frac{n}{2}$ we have

$$\sum_{1 \leq i, j \leq n} \mathbb{1}\{|\deg(u_i) - \deg(v_j)| \geq k\} \leq 2k(n - k). \tag{1.2}$$

Note that the trivial upper bound n^2 is the best possible upper bound in the case $k \leq \frac{n}{2}$. The proof of Theorem 1.2, given in Section 2, is based on an idea similar to the spread bounding theorem of Erdős, Chen, Rousseau and Schelp – see [8, 3]. Later in the same section we provide an elementary example showing that the bound (1.2) is sharp. In Section 3 we show how to transform the Theorem 1.1 to Theorem 1.2. To this end, we make use of an appropriate sampling construction, similar in spirit to [11]. The key idea is to approximate a fixed coherent distribution with a randomly generated sequence of graphs. We then apply Theorem 1.2 to each of the graphs in the sequence and obtain (1.1) by passing to the limit.

2 Number of high degree differences in bipartite graphs

Let $G = (U, V, E)$ be a bipartite graph with an equal bipartition, that is a triplet

$$U = \{u_1, u_2, \dots, u_n\}, \quad V = \{v_1, v_2, \dots, v_n\},$$

and

$$E \subset U \times V,$$

for some fixed $n \in \mathbb{Z}_+$. Let us fix a natural number k satisfying $n \geq k > \frac{n}{2}$. Hereinafter, we denote the degree sequences of G as $(\alpha_i)_{i=1}^n$ and $(\beta_j)_{j=1}^n$, i.e., $\alpha_i = \deg(u_i)$ and $\beta_j = \deg(v_j)$ for all $1 \leq i, j \leq n$. Without loss of generality we also assume that

$$\alpha_1 \geq \alpha_2 \geq \dots \geq \alpha_n,$$

$$\beta_1 \geq \beta_2 \geq \dots \geq \beta_n.$$

We start with an observation similar to the spread bounding theorem of Erdős et al. – see [8].

Lemma 2.1. *There exist $s, t \in \{1, 2, \dots, n - k + 1\}$ such that $\alpha_s \leq \beta_{s+k-1} + k - 1$ and $\beta_t \leq \alpha_{t+k-1} + k - 1$.*

Proof. We will prove only the existence of s , as the case of t is analogous. Assume for the sake of contradiction that such a number s does not exist. Therefore, the total number of edges incident to $u_1, u_2, \dots, u_{n-k+1}$ is at least $\beta_k + \beta_{k+1} + \dots + \beta_n + k(n - k + 1)$. Observe that at least $k(n - k + 1)$ of these edges go to vertices v_1, v_2, \dots, v_{k-1} . Let us denote

$$\tilde{E} := E \cap \left(\{u_1, u_2, \dots, u_{n-k+1}\} \times \{v_1, v_2, \dots, v_{k-1}\} \right).$$

We have just shown that $|\tilde{E}| \geq k(n - k + 1)$. On the other hand, we clearly have

$$|\tilde{E}| \leq (k - 1)(n - k + 1),$$

which is a contradiction. □

Proof of Theorem 1.2. For $1 \leq i, j \leq n$, let us call (i, j) an \mathcal{A} -pair if $\alpha_i \geq \beta_j + k$. Analogously, let us call (i, j) a \mathcal{B} -pair if $\beta_j \geq \alpha_i + k$. Since $k > \frac{n}{2}$, we have $\alpha_i > \frac{n}{2}$ for all \mathcal{A} -pairs (i, j) and $\alpha_i < \frac{n}{2}$ for all \mathcal{B} -pairs (i, j) . As a consequence, there exists an $i_0 \in \{1, 2, \dots, n + 1\}$ such that:

1. $i \leq i_0 - 1$ for any \mathcal{A} -pair (i, j) ,
2. $i \geq i_0$ for any \mathcal{B} -pair (i, j) .

Analogously, there exists $j_0 \in \{1, 2, \dots, n + 1\}$ such that:

3. $j \leq j_0 - 1$ for any \mathcal{B} -pair (i, j) ,
4. $j \geq j_0$ for any \mathcal{A} -pair (i, j) .

Observe that by Lemma 2.1,

5. for any \mathcal{A} -pair (i, j) either $i < s$ or $j > s + k - 1$,
6. for any \mathcal{B} -pair pair (i, j) either $j < t$ or $i > t + k - 1$.

We will now show that conditions 1–6 imply that the total number of \mathcal{A} -pairs and \mathcal{B} -pairs is at most $2k(n - k)$. Let us fix $i_0, j_0 \in \{1, 2, \dots, n + 1\}$. First, we will show that it is sufficient to consider only s and t such that $s, t \in \{1, n - k + 1\}$ because these values of s and t are optimal in the sense that they maximize the total number of pairs (i, j) fulfilling all conditions 1–6.

Note that the variable s appears only in the 5-th condition and thus the value of s is not relevant for bounding the number of \mathcal{B} -pairs. Moreover, observe that if $i_0 \leq n - k + 1$, then for $s = n - k + 1$ condition 5 is automatically fulfilled and thus $s = n - k + 1$ is an optimal value. Similarly, if $j_0 \geq k + 1$, then for $s = 1$ condition 5 is also automatically fulfilled and $s = 1$ is an optimal value. Finally, let us assume that $i_0 \geq n - k + 2$ and $j_0 \leq k$. In this case, the restrictions imposed by condition 5 remove exactly $(i_0 - s)(s + k - j_0)$ additional pairs. Therefore, as the last expression is a concave function of $s \in [1, n - k + 1]$, it is minimized in one of the endpoints. Hence we may assume that $s = 1$ or $s = n - k + 1$, as desired. Analogously, we show that $t = 1$ or $t = n - k + 1$ is optimal. There are four possible cases now:

- a. $s = 1, t = n - k + 1$. We have $j \geq k + 1$ for all \mathcal{A} -pairs and $j \leq n - k$ for all \mathcal{B} -pairs (i, j) . Thus any i participates in at most $n - k$ of \mathcal{A} -pairs and in at most $n - k$ of \mathcal{B} -pairs. Therefore, since a fixed vertex can not participate in both types of pairs, every i participates overall in at most $n - k$ pairs. As a consequence, the total number of pairs does not exceed $n(n - k) < 2k(n - k)$.
- b. $s = n - k + 1, t = 1$. This case is symmetric to the previous one.
- c. $s = 1, t = 1$. We have $j \geq k + 1$ for all \mathcal{A} -pairs and $i \geq k + 1$ for all \mathcal{B} -pairs (i, j) . Let us denote $a := \max(k + 1, j_0)$ and $b := \max(k + 1, i_0)$. Then the total number of \mathcal{A} -pairs is bounded by $(n - a + 1)(b - 1)$, while the total number of \mathcal{B} -pairs is at most $(n - b + 1)(a - 1)$. Notice, that for $a, b \in [k + 1, n + 1]$ the sum

$$S := (n - a + 1)(b - 1) + (n - b + 1)(a - 1),$$

is bilinear and it is maximized at one of four corners. For $a = b = k + 1$, we get $S = 2k(n - k)$. For, say $a = n + 1$, we get $S = n(n - b + 1) \leq n(n - k) < 2k(n - k)$.

- d. $s = n - k + 1, t = n - k + 1$. This case is analogous to c.

Hence we have shown that Theorem 1.2 holds in all cases. This ends the proof. □

We end this section with an example showing that the upper bound $2k(n - k)$ in (1.2) cannot be improved. Note that a straightforward modification of this example shows that $2\delta(1 - \delta)$ in (1.1) is also sharp.

Example 2.2. Consider $n, k \in \mathbb{Z}_+$, with $n \geq k > \frac{n}{2}$. Let $G_{n,k} = (U, V, E)$, where $U = V = \{1, 2, \dots, n\}$ and

$$E = \{(u, v) \in U \times V : \max(u, v) \leq k\}.$$

We clearly have

$$\sum_{1 \leq i, j \leq n} \mathbb{1}\{|\deg(u_i) - \deg(v_j)| \geq k\} = 2k(n - k).$$

Moreover, one can check that inequality (1.2) becomes an equality exactly for those graphs G that are isomorphic to $G_{n,k}$ or to its complement $\overline{G}_{n,k}$. This follows easily from the proof of Theorem 1.2 and we leave the details to interested reader.

3 Proof of the Burdzy–Pitman conjecture

By $\mathcal{C}_{\mathcal{I}}(m)$ we denote the set of $(X, Y) \in \mathcal{C}_{\mathcal{I}}$ such that both X and Y take at most m different values.

Proposition 3.1. *Let (X, Y) be coherent and independent, and let m be a positive integer. Then there exists $(X_m, Y_m) \in \mathcal{C}_{\mathcal{I}}(m)$, such that $|X - X_m| \leq \frac{1}{m}$ and $|Y - Y_m| \leq \frac{1}{m}$, almost surely.*

The proof of the above Proposition can be found in [4, 1]. In what follows, fix any $\delta \in (\frac{1}{2}, 1]$.

Proposition 3.2. *To prove Theorem 1.1 it is enough to verify it for all $(X, Y) \in \mathcal{C}_{\mathcal{I}}(m)$, $m \geq 1$.*

Proof. Fix $(X, Y) \in \mathcal{C}_{\mathcal{I}}$ and choose (X_m, Y_m) as in Proposition 3.1. By the triangle inequality we get

$$\mathbb{P}(|X - Y| \geq \delta) \leq \mathbb{P}(|X_m - Y_m| \geq \delta - 2/m).$$

Thus, assuming that Theorem 1.1 is true for all $(X, Y) \in \cup_{m=1}^{\infty} \mathcal{C}_{\mathcal{I}}(m)$, for m large enough so that $\delta - 2/m > 1/2$, we obtain

$$\mathbb{P}(|X - Y| \geq \delta) \leq 2(\delta - 2/m)(1 - \delta + 2/m).$$

Letting $m \rightarrow \infty$ completes the proof. □

We are now able to prove our main result.

Proof of Theorem 1.1. Fix $(X, Y) \in \cup_{m=1}^{\infty} \mathcal{C}_{\mathcal{I}}(m)$. There exists a probability space $(\Omega, \mathcal{F}, \mathbb{P})$, independent sub σ -fields $\mathcal{G}, \mathcal{H} \subset \mathcal{F}$ and an event $A \in \mathcal{F}$, such that $X = \mathbb{E}(\mathbb{1}_A | \mathcal{G})$ and $Y = \mathbb{E}(\mathbb{1}_A | \mathcal{H})$. Furthermore, for some $N, M \in \mathbb{Z}_+$, we may suppose that X takes values x_1, x_2, \dots, x_N on sets G_1, G_2, \dots, G_N and Y takes values y_1, y_2, \dots, y_M on sets H_1, H_2, \dots, H_M . We can also assume without loss of generality that

$$\mathcal{G} = \sigma(G_1, G_2, \dots, G_N),$$

$$\mathcal{H} = \sigma(H_1, H_2, \dots, H_M).$$

For $1 \leq i \leq N$ and $1 \leq j \leq M$, let $p_i = \mathbb{P}(G_i)$, $q_j = \mathbb{P}(H_j)$ and

$$\rho_{i,j} = \frac{\mathbb{P}(G_i \cap H_j \cap A)}{\mathbb{P}(G_i \cap H_j)}.$$

Then by independence we have $\mathbb{P}(G_i \cap H_j) = p_i q_j$ and

$$x_i = \sum_{j=1}^M q_j \rho_{i,j}, \quad 1 \leq i \leq N, \tag{3.1}$$

$$y_j = \sum_{i=1}^N p_i \rho_{i,j}, \quad 1 \leq j \leq M. \tag{3.2}$$

First, we show how to construct a sequence of bipartite graphs $G_n = (U_n, V_n, E_n)$ with $|U_n| = |V_n| = n$, such that:

(C1) there are $p_i n + O(n^{3/4})$ vertices in U_n of degree $x_i n + O(n^{3/4})$, $i = 1, 2, \dots, N$,

(C2) there are $q_j n + O(n^{3/4})$ vertices in V_n of degree $y_j n + O(n^{3/4})$, $j = 1, 2, \dots, M$,

where by $O(n^{3/4})$ we denote any quantity bounded in magnitude by $Cn^{3/4}$ for some constant $C < \infty$ independent of n, N, M, i and j .

Fix $n \geq 1$ and choose n independent points u_1, u_2, \dots, u_n in the initial space Ω (distributed according to \mathbb{P}) and for $1 \leq i \leq n$ denote $\alpha_i = s$ if $u_i \in G_s$. In other words, $(\alpha_1, \alpha_2, \dots, \alpha_n)$ is an i.i.d. sample from the set $\{1, 2, \dots, N\}$ with weights p_1, p_2, \dots, p_N , respectively. We can think about this sample as a randomly generated sequence of labels. Let $A_s = \sum_{i=1}^n \mathbb{1}_{\{\alpha_i=s\}}$ be the number of labels equal to s , $1 \leq s \leq N$. Observe that A_s is the sum of n independent Bernoulli random variables. Hence, by Hoeffding’s inequality [10], we have

$$\mathbb{P}(|A_s - np_s| \geq nr) \leq 2 \cdot e^{-2nr^2},$$

for all positive r . Consequently, setting $r = n^{-1/4}$ we get

$$\mathbb{P}(|A_s - np_s| \geq n^{3/4}) \leq 2 \cdot e^{-2\sqrt{n}}.$$

Thus, for large n , with high probability we have $|A_s - np_s| < n^{3/4}$ simultaneously for all $1 \leq s \leq N$.

Analogously, we choose points v_1, v_2, \dots, v_n and generate an i.i.d. sample $(\beta_1, \beta_2, \dots, \beta_n)$ from the set $\{1, 2, \dots, M\}$ with weights q_1, q_2, \dots, q_M . If $B_t = \sum_{j=1}^n \mathbb{1}_{\{\beta_j=t\}}$ for $1 \leq t \leq M$ then, for large n , with high probability $|B_t - nq_t| < n^{3/4}$ for all t simultaneously.

Given the points $(u_i)_{i=1}^n$ and $(v_j)_{j=1}^n$ and the corresponding labels $(\alpha_i)_{i=1}^n$ and $(\beta_j)_{j=1}^n$, we will generate independently edges of a random bipartite graph (U_n, V_n, E_n) , where $U_n = \{u_1, u_2, \dots, u_n\}$ and $V_n = \{v_1, v_2, \dots, v_n\}$. The subscripts on $\mathbb{P}_{\alpha, \beta}$ and $\mathbb{E}_{\alpha, \beta}$ will denote conditioning on $(\alpha_i)_{i=1}^n$ and $(\beta_j)_{j=1}^n$.

1. Generate independent indicator random variables $Z_{i,j}$ for $1 \leq i, j \leq n$ satisfying

$$\mathbb{P}_{\alpha, \beta}(Z_{i,j} = 1) = 1 - \mathbb{P}_{\alpha, \beta}(Z_{i,j} = 0) = \rho_{\alpha_i, \beta_j},$$

2. for $1 \leq i, j \leq n$, set $(u_i, v_j) \in E_n$ iff $Z_{i,j} = 1$.

Hence, $Z_{i,j} = \mathbb{1}_{\{(u_i, v_j) \in E_n\}}$. For $1 \leq i \leq n$,

$$\mathbb{E}_{\alpha, \beta} \deg(u_i) = \mathbb{E}_{\alpha, \beta} \left(\sum_{j=1}^n Z_{i,j} \right) = \sum_{t=1}^M B_t \rho_{\alpha_i, t} = \sum_{t=1}^M (nq_t + O(n^{3/4})) \rho_{\alpha_i, t},$$

and hence, by (3.1),

$$\mathbb{E}_{\alpha, \beta} \deg(u_i) = nx_{\alpha_i} + O(n^{3/4}). \tag{3.3}$$

Similarly, for $1 \leq j \leq n$, by (3.2) we get

$$\mathbb{E}_{\alpha, \beta} \deg(v_j) = ny_{\beta_j} + O(n^{3/4}). \tag{3.4}$$

We apply Hoeffdings’s inequality again to obtain

$$\mathbb{P}_{\alpha, \beta} \left(|\deg(u_i) - \mathbb{E}_{\alpha, \beta} \deg(u_i)| \geq n^{3/4} \right) \leq 2 \cdot e^{-2\sqrt{n}}, \tag{3.5}$$

and

$$\mathbb{P}_{\alpha, \beta} \left(|\deg(v_j) - \mathbb{E}_{\alpha, \beta} \deg(v_j)| \geq n^{3/4} \right) \leq 2 \cdot e^{-2\sqrt{n}}, \tag{3.6}$$

for all $i, j \in \{1, 2, \dots, n\}$. Note that the concentration rates (3.5) and (3.6) are exponential in \sqrt{n} . Thus, since n is large, with high probability all these inequalities hold simultaneously. Then, by (3.3) and (3.4), we have $\deg(u_i) = nx_{\alpha_i} + O(n^{3/4})$ and $\deg(v_j) = ny_{\beta_j} + O(n^{3/4})$ for all $i, j \in \{1, 2, \dots, n\}$ with high probability. This, together

with bounds on $(A_s)_{s=1}^N$ and $(B_t)_{t=1}^M$, proves that (deterministic) G_n 's satisfying conditions (C1)-(C2) exist for large n .

In what follows, we add additional subscripts and write $u_i^{(n)}$ and $v_j^{(n)}$ for generic elements of U_n and V_n , respectively. We can now write

$$\begin{aligned} \mathbb{P}(|X - Y| \geq \delta) &= \sum_{\substack{1 \leq i \leq N \\ 1 \leq j \leq M}} \mathbb{1}_{\{|x_i - y_j| \geq \delta\}} \cdot p_i q_j \\ &= \lim_{n \rightarrow \infty} \frac{1}{n^2} \sum_{\substack{1 \leq i \leq N \\ 1 \leq j \leq M}} \mathbb{1}_{\{|nx_i - ny_j| \geq n\delta\}} \cdot \left(p_i n + O(n^{3/4})\right) \left(q_j n + O(n^{3/4})\right). \end{aligned} \quad (3.7)$$

By the triangle inequality

$$|nx_{\alpha_i} - ny_{\beta_j}| \leq |\deg(u_i^{(n)}) - \deg(v_j^{(n)})| + 2 \cdot O(n^{3/4}),$$

for all $i, j \in \{1, 2, \dots, n\}$. This and (C1)-(C2) imply that we can bound the right hand side of (3.7) by

$$\leq \limsup_{n \rightarrow \infty} \frac{1}{n^2} \sum_{1 \leq i, j \leq n} \mathbb{1}_{\{|\deg(u_i^{(n)}) - \deg(v_j^{(n)})| \geq n\delta - 2O(n^{3/4})\}}.$$

Finally, applying Theorem 1.2 to bipartite graphs G_n , we obtain

$$\leq \limsup_{n \rightarrow \infty} \frac{1}{n^2} \cdot 2 \left(n\delta - 2O(n^{3/4})\right) \left(n - n\delta + 2O(n^{3/4})\right) = 2\delta(1 - \delta),$$

which ends the proof. □

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