

Chapter XIII

E-Recursively Enumerable Degrees

Degree theory for subsets of *E*-closed structures differs markedly from the Σ_1 admissible case. On the surface the results are similar, but the modes of argument differ considerably. Post's problem once again has a positive solution, but this time without injuries and without repeated attempts to satisfy a given requirement. The presence of Moschovakis witnesses makes all the difference. Injuries do occur in the proof of Slaman's splitting theorem.

1. Regular Sets

Let \mathcal{E} be transitive and $A, B \subseteq \mathcal{E}$. The relativization of *E*-recursiveness to *B* was introduced in Section 5.XI. In essence a new scheme,

$$\{c\}^B(x_1, \dots, x_n) = B \cap x_i \quad (c = \langle 7, n, i \rangle)$$

is added to the original six. *f* is partial *E*-recursive relative to *B* if $f \simeq \{e\}^B$ for some $e < \omega$. *D* is *E*-recursively enumerable in *p* relative to *B* if

$$D = \{x \mid \{e\}^B(x, p) \downarrow\}$$

for some e . \mathcal{E} is *E*-closed relative to *B* if

$$\{e\}^B(x) \downarrow \rightarrow \{e\}^B(x) \in \mathcal{E}$$

for all $e < \omega$ and $x \in \mathcal{E}$.

Assume \mathcal{E} is *E*-closed. A is *E*-reducible to *B* on \mathcal{E} (in symbols $A \leq_{\mathcal{E}} B$) if there exist $e < \omega$ and $p \in \mathcal{E}$ such that

- (i) $\{e\}^B(x, p) \downarrow$ for all $x \in \mathcal{E}$,
- (ii) $T_{\langle e, x, p, B \rangle} \in \mathcal{E}$ for all $x \in \mathcal{E}$, and
- (iii) $A = \hat{x}[x \in \mathcal{E} \ \& \ \{e\}^B(x, p) = 0]$.