

# 1. The Basics

## 1.1 Preliminaries and Notation

We assume that the reader is familiar with the basic definitions and results normally found in a first course in mathematical logic. Specifically, we will freely use the concepts of a *first-order language*, a *structure* or *model* in that language, and the *satisfaction relation* between models and formulas. We also assume that the reader knows the Compactness and Omitting Types Theorems, and can carry out an elimination of quantifiers argument for a specific theory such as dense linear orders without endpoints or divisible abelian groups. In this first section we will review some of these results as a way of setting our notation and viewpoint and jogging the student's memory.

### Notation. (Model Theory)

- A first-order language is denoted by  $L$ ,  $L'$ ,  $L_0$ , etc. The cardinality of a language  $L$ ,  $|L|$ , is simply the cardinality of the set of nonlogical symbols of  $L$ .
- Formulas are denoted by lower case Greek letters. Writing  $\varphi(v_0, \dots, v_n)$  indicates that the free variables in  $\varphi$  are in  $\{v_0, \dots, v_n\}$ . If  $t_0, \dots, t_n$  are terms in the language,  $\varphi(t_0, \dots, t_n)$  is the formula obtained by substituting  $t_i$  for  $v_i$ . A sentence is a formula with no free variables.
- We use  $\mathcal{M}$  or  $\mathcal{N}$ , decorated with various subscripts and superscripts, to denote a model or structure in a first-order language. The universe of, e.g.,  $\mathcal{M}_0$ , is  $M_0$ . Elements of the universe are denoted by lower case letters such as  $a$ ,  $b$ ,  $c$ , etc. If  $X$  is an element of the language in which  $\mathcal{M}$  is a model  $X^{\mathcal{M}}$  denotes the interpretation of  $X$  in  $\mathcal{M}$ .
- Given models  $\mathcal{M}$  and  $\mathcal{N}$  in a language  $L$ , a function  $f : M \rightarrow N$  is an *isomorphism* if  $f$  is a bijection and for all symbols  $X \in L$ ,  $f(X^{\mathcal{M}}) = X^{\mathcal{N}}$ . When there is an isomorphism of  $\mathcal{M}$  onto  $\mathcal{N}$  we write  $\mathcal{M} \cong \mathcal{N}$  and say  $\mathcal{M}$  and  $\mathcal{N}$  are isomorphic. An *automorphism* of  $\mathcal{M}$  is an isomorphism of  $\mathcal{M}$  onto itself. For  $\mathcal{M}$  a model,  $\text{Aut}(\mathcal{M})$  denotes the automorphism group of  $\mathcal{M}$ .
- A *theory* in the language  $L$  is a consistent set of sentences of  $L$ . A set of sentences need not be complete in order to be called a theory. (A set of sentences is *consistent* if it has a model. A theory is *complete* if all