Global Notational Conventions

Objects of type 0: $a, b, \ldots, e, i, j, \ldots, w$ Objects of type 1: α, β, γ, δ, ε f, g, h, E, F, \ldots, K A, B, C, D, M, N, O, W P, Q, \ldots, V Objects of type 2: E, F, ..., K A, B, C, D, M, N, W P,Q,...,V Objects of type 3: **E**, **F**, . . . , **K** A, B, C, D, M, N, W $\mathbb{P}, \mathbb{Q}, \ldots, \mathbb{V}$ Objects of type 4: $\mathcal{E}, \mathcal{F}, \ldots, \mathcal{K}$ 9.2....V Other: κ, λ, μ, ν, π, ρ, σ, τ, υ Γ,Λ Ľ T A. B. C M. N. U **x**, y, z

X, Y, Z

 $\varphi, \psi, \chi, \theta$

natural numbers = elements of ω

total unary functions $\omega \to \omega$ = elements of " ω partial functions " $\omega \to \omega$ sets of natural numbers = subsets of ω k-ary relations on ω = subsets of " ω

partial functionals ${}^{k,l}\omega \to \omega$ sets of (total unary) functions = subsets of ${}^{\omega}\omega$ (k, l)-ary relations = subsets of ${}^{k,l}\omega$

partial functionals ${}^{k,l,l'}\omega \to \omega$ sets of total unary functionals = subsets of ${}^{(m_{\omega})}\omega$ (k, l, l')-ary relations = subsets of ${}^{k,l,l'}\omega$

partial functionals ${}^{k,l,l',l''}\omega \rightarrow \omega$ (k, l, l', l'')-ary relations = subsets of ${}^{k,l,l',l''}\omega$

ordinal numbers inductive operators formal language formal theory formulas of a formal language structures, models arbitrary objects arbitrary sets arbitrary functions

The last three categories are often subject to local conventions

In most instances a bold-face letter denotes a finite sequence of objects of the type denoted by the light-face letter, Exceptions: Σ , Π , Δ , ∇ , κ , ω_1 [-].