

5. PARTIAL CONSERVATIVITY

A sentence φ is Γ -conservative over T if for every Γ sentence θ , if $T + \varphi \vdash \theta$, then $T \vdash \theta$. In this chapter we study this phenomenon for its own sake. Results on Γ -conservativity are, however, also very useful in many contexts, in particular in connection with interpretability (see Chapters 6 and 7).

Our task in this chapter is to develop general methods for constructing partially conservative sentences satisfying additional conditions such as being nonprovable in a given theory.

We assume throughout that $PA \dashv T$. The results of this chapter do not depend on the assumption that T is reflexive.

A first example of a Π_1 -conservative sentence is given in the following:

Theorem 1. $\neg Con_T$ is Π_1 -conservative over T .

Proof. Suppose θ is Π_1 and

$$(1) \quad T + \neg Con_T \vdash \theta.$$

From (1) we get $PA \vdash Pr_T(\neg\theta) \rightarrow Pr_T(Con_T)$, whence

$$(2) \quad PA \vdash Pr_T(\neg\theta) \rightarrow \neg Con_{T+\neg Con_T}.$$

By provable Σ_1 -completeness,

$$(3) \quad PA \vdash \neg\theta \rightarrow Pr_T(\neg\theta).$$

By Corollary 2.2,

$$(4) \quad PA + Con_T \vdash Con_{T+\neg Con_T}.$$

Combining (2), (3), (4) we get $PA \vdash \neg\theta \rightarrow \neg Con_T$ and so by (1), $T \vdash \theta$. ■

By Corollary 2.4, Theorem 1 provides us with an example of a (Σ_1) sentence φ which is Π_1 -conservative over T and nontrivially so, i.e. such that $T \not\vdash \varphi$, even if T is not Σ_1 -sound.

If φ is Γ -conservative over T and ψ is Γ^d , then clearly φ is Γ -conservative over $T + \psi$. Also note that if T is Σ_1 -sound and π is Π_1 , then π is Σ_1 -conservative over T iff π is true iff $T + \pi$ is consistent.

Let us now try to construct a sentence φ which is nontrivially Γ -conservative over T . Thus, given that

$$(1) \quad T + \varphi \vdash \theta,$$

where θ is Γ , we want to be able to conclude that $T \vdash \theta$. This follows if (1) implies that

$$(2) \quad T + \neg\theta \vdash \varphi.$$

The natural way to ensure that (1) implies (2) is to let φ be a sentence saying of itself that there is a false Γ sentence (namely θ) which φ implies in T . Thus, let φ be such that

$$(3) \quad PA \vdash \varphi \leftrightarrow \exists u(\Gamma(u) \wedge Pr_{T+\varphi}(u) \wedge \neg Tr_\Gamma(u)),$$

where $\Gamma(x)$ is a PR binumeration of the set of Γ sentences. Then (1) implies (2).