# Perfect Graphs and Complex Surface Singularities with Perfect Local Fundamental Group 

Lawrence Brenton and Daniel Drucker

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#### Abstract

In this paper we introduce the term "perfect graph" to refer to those graphs which characterize resolutions of certain isolated singular points of complex surfaces. Using techniques for graphical evaluation of determinants, we reduce questions about perfect graphs to problems involving partial fraction representations of positive integers; the solutions to those Diophantine problems thus have interesting geometric interpretations.


1. Introduction and statement of results. In [5] Brieskorn gave the first examples of isolated singularities of complex $n$-varieties, $n \geq 3$, that are topologically non-singular (locally homeomorphic to the $2 n$-ball) but analytically singular. Earlier Mumford [16] had shown that this is impossible in dimension 2. In this paper we pursue the natural analogue of the Brieskorn singularities for complex surfaces, namely those singular points $x \in X$ which are homologically non-singular in the sense of being locally homeomorphic to the cone on a homology 3 -sphere. (The rational double point $E_{8}$ is the most familiar example.) This condition is equivalent to the requirement that the local fundamental group of $x$ in $X$ be a perfect group (cf., for example, [16], [17], and [19], where the topic of classifying isolated two-dimensional singularities by the group-theoretic properties of the local fundamental group is introduced and developed).

Let $x$ be an isolated singularity of a normal complex surface $X$, and let $p: \tilde{X} \rightarrow X$ be the minimal resolution of singularities. We will assume that the exceptional curve $C=p^{-1}(x)=\bigcup_{i=1}^{n} C_{i}$ is contractible, that each component $C_{i}$ is non-singular rational, and that the components meet transversally with no triple intersections. In this case the topology of the singularity is completely determined by the weighted dual intersection graph $G_{p}$ of the exceptional curve. In particular, the local fundamental group $\pi_{1}(x)$ can be computed directly from $G_{p}$ in terms of generators and relations, by the technique of Mumford [16]. Using this method it can be shown that $\pi_{1}(x)$ is perfect exactly when the intersection matrix ( $-C_{i} \cdot C_{i}$ ) has determinant 1 . Indeed, the following are necessary and sufficient conditions for a weighted graph $G$ to be the dual graph of the minimal resolution of a normal complex surface singularity whose minimal resolution is normal ("good") and whose local fundamental group is perfect:
(a) $G$ is a tree (a connected graph with no circuits).
(b) Each weight $w_{i}$ is an integer $\geq 2$.
(c) The associated intersection matrix is positive definite with determinant 1 . (Section 1 of [4] gives an elementary expository review of the geometry of complex surface
singularities as reflected by their resolutions, graphs, and local fundamental groups; [1] investigates and classifies some of the global settings in which such "perfect" singular points occur.)

These considerations motivate the following:
(1.1) Definitions. A weighted graph $G$ is perfect if it satisfies conditions (a), (b), and (c) above. A graph $G$ is perfectable if there exist integer weights $w_{i}$ for its vertices such that the resulting weighted graph $G\left(w_{1}, \cdots, w_{n}\right)$ is perfect. Such a set of weights is called a set of perfect weights. A minimal perfectable graph is a perfectable graph none of whose proper subgraphs is perfectable.

The goal of this paper is to find perfect graphs and to point out connections between perfect graphs and solutions of certain Diophantine equations of interest in number theory. Our results can be summarized as follows:
(1.2) Main Theorem. Let $G$ be any graph which is not of the form shown in Figure 1 for $n=0,1$, or 2 . Then $G$ is perfectable if and only if $G$ is a tree that contains one of the 25 minimal perfectable graphs listed in Table I at the end of this paper.

In consequence we obtain the following results for particular kinds of graphs.
(1.3) Theorem. There is no perfect weighted graph on 7 or fewer vertices. The perfect weighted graphs on 8 vertices appear in Figure 2. Of the 47 trees on 9 vertices exactly 30 are perfectable (most with several sets of perfect weights). In fact, "almost all' trees with sufficiently many vertices are perfectable; that is,

$$
\lim _{n \rightarrow \infty} \frac{\text { number of perfectable trees on } n \text { vertices }}{\text { total number of trees on } n \text { vertices }}=1
$$

The distance between two points of a graph $G$ is the number of edges in the shortest path joining them. The diameter of a graph $G$ is the maximum of the distances between


Figure 1


Figure 2
pairs of points in $G$. In the case of a tree it is the number of edges in the longest "chain" (tree without branch points) contained in $G$.
(1.4) Theorem. (a) Among ordinary and extended Dynkin diagrams of types $A, D$, and $E$, only $E_{8}$ and $\widetilde{E}_{8}$ are perfectable, and each of these graphs has only one set of perfect weights.
(b) Every tree of diameter $d \geq 29$ is perfectable except those isomorphic to $A_{n}, D_{n}$, and $\tilde{D}_{n}$.
(c) Every tree of diameter $d \geq 7$ is perfectable except $A_{n}, D_{n}, \tilde{D}_{n}$, and the graphs in Figure 3.

Defintion. Let $p_{1} \leq p_{2} \leq \cdots \leq p_{r}$ be positive integers. A graph $G$ is of type $E_{p_{1}, \cdots, p_{r}}$ if $G$ has a vertex $v_{0}$ such that $G-\left\{v_{0}\right\}$ is the disjoint union of $r$ graphs of types $A_{p_{1}, \cdots, p_{r}}$, each joined to $v_{0}$ only at a terminal vertex (cf. Figure 4).
(1.5) THEOREM. (a) There is a one-to-one correspondence between perfect weighted graphs of type $E_{p_{1}, \cdots, p_{r}}$ and solutions in reduced proper fractions $s_{i} / t_{i}$ of the equation

$$
\sum_{i=1}^{r} \frac{s_{i}}{t_{i}}+\frac{1}{\prod_{i=1}^{r} t_{i}}=n
$$

with $n$ an integer $\geq 2$ and $s_{i}, t_{i}$ positive integers for $i=1, \cdots, r$.
(b) A graph $G$ of type $E_{p_{1}, \cdots, p_{r}}$, with $r \geq 3$ and $p_{r-1} \geq 2$, is perfectable if and only if $G$ contains one of the following:

$$
E_{1,2,4} ; E_{2,2,3} ; E_{1,2,2,2} ; E_{1,1,1,2,3} ; E_{1,1,1,1,1,2,2}
$$

This last result shows one of the connections of this topic with certain problems of independent interest in number theory. For instance, putting each $s_{i}=1$ above leads to the following unresolved question in the theory of Egyptian fractions (Paul Erdös offers $\$ 100$ for a solution): Given positive integers $t_{1}, \cdots, t_{k}$, relatively prime in pairs and all $\geq 2$, do there always exist integers $n, t_{k+1}, \cdots, t_{r}$, all $\geq 2$, such that $n=\sum_{i=1}^{r}\left(1 / t_{i}\right)+1 /\left(\prod_{i=1}^{r} t_{i}\right)$ ? Connections with number theory will be discussed more


Figure 3


Figure 4
fully in Section 4.
As noted above, finding perfect weighted graphs is equivalent to finding symmetric bilinear forms $\phi: \boldsymbol{Z}^{n} \times \boldsymbol{Z}^{\boldsymbol{n}} \rightarrow \boldsymbol{Z}$ of determinant 1 , corresponding to the intersection matrices $\left(-C_{i} \cdot C_{j}\right)$. In Section 2 we present techniques for quickly calculating the determinant of the intersection matrix associated to a weighted graph directly from the graph. The results are stated in sufficient generality to apply in a wide range of settings.
2. Graphical evaluation of determinants. In this section, we describe methods of evaluating the determinant of a matrix by use of an associated graph. The initial results are useful for general sparse matrices, but do not seem to be well known. The more specialized versions are useful in our classification of perfectable graphs. (See also references [1], [4], [7], and [8].)

Let $M=\left(m_{i j}\right)$ be an $n \times n$ matrix with entries in a commutative ring $A$ with identity element 1 different from 0 . The determinant of $M$ is given by the formula

$$
\begin{equation*}
|M|=\sum_{\sigma \in S_{n}}(\operatorname{sgn} \sigma) m_{1, \sigma(1)} m_{2, \sigma(2)} \cdots m_{n, \sigma(n)} \tag{2.1}
\end{equation*}
$$

where $S_{n}$ is the set of all permutations of $\{1,2, \cdots, n\}$. Let $c=\left[i_{1}, \cdots, i_{k}\right](1 \leq k \leq n)$ denote the $k$-cycle in $S_{n}$ that cyclically permutes the distinct indices $i_{1}, \cdots, i_{k}$. (When $k=1, c$ is the identity permutation of a singleton set.) $c$ is even if $k$ is even. The weight of $c$ is the ring element $w(c)=m_{i_{1}, i_{2}} \cdots m_{i_{k-1}, i_{k}} m_{i_{k}, i_{1}} .\left(w(c)=m_{i_{1}, i_{1}}\right.$ if $k=1$.) The signed weight of $c$ is the ring element $\tilde{w}(c)=(\operatorname{sgn} c) w(c)=(-1)^{k-1} w(c)$. More generally, if $\sigma=c_{1} \cdots c_{s}$ is a product of disjoint cycles, we define $w(\sigma)=w\left(c_{1}\right) \cdots w\left(c_{s}\right)$ and $\tilde{w}(\sigma)=\tilde{w}\left(c_{1}\right) \cdots \tilde{w}\left(c_{s}\right)=(\operatorname{sgn} \sigma) w(\sigma)=(-1)^{e(\sigma)} w(\sigma)$, where $e(\sigma)$ is the number of even cycles among $c_{1}, \cdots, c_{s}$. Then (2.1) can be rewritten as

$$
\begin{equation*}
|M|=\sum_{\sigma \in S_{n}} \tilde{w}(\sigma)=\sum_{\sigma \in S_{n}}(-1)^{e(\sigma)} w(\sigma) . \tag{2.2}
\end{equation*}
$$

Define the associated graph of $M$ to be the directed graph $G=G(M)$ with $n$ vertices (labelled $1,2, \cdots, n$ ) that has a directed edge $(i, j)$ from the vertex $i$ to the vertex $j$ precisely when $m_{i j} \neq 0$. ( $i=j$ is allowed.) A circuit of length $k(1 \leq k \leq n)$ in $G$ is a $k$-cycle $c=\left[i_{1}, \cdots, i_{k}\right]$ such that $\left(i_{r}, i_{r+1}\right)$ (for $\left.1 \leq r<k\right)$ and $\left(i_{k}, i_{1}\right)$ are directed edges of $G$. A product $p=c_{1} \cdots c_{s}$ of circuits of $G$ will be called a circuit partition of $G$ if the domains of $c_{1}, \cdots, c_{s}$ form a partition of $\{1,2, \cdots, n\}$. The set of all circuit partitions of $G$ will be denoted $P$. Formula (2.2) implies that

$$
|M|=\sum_{p \in P} \tilde{w}(p)=\sum_{p \in P}(-1)^{e(p)} w(p) .
$$

This means that we can calculate $|M|$ just by looking at the graph $G$ with its directed edges labelled by the ring elements $m_{i j}$. The labelled graph $G$ uniquely determines the matrix $M$, so we can write $|G|$ for $|M|$ and obtain the graph-theoretic formula

$$
\begin{equation*}
|G|=\sum_{p \in P} \tilde{w}(p)=\sum_{p \in P}(-1)^{e(p)} w(p) . \tag{2.3}
\end{equation*}
$$

Decomposition of a determinant relative to components of the associated graph. If $G$ has connected components $G_{1}, \cdots, G_{t}$, and if $P_{1}, \cdots, P_{t}$ are the corresponding sets of circuit partitions, then $P=P_{1} \cdots P_{t}$, so

$$
\begin{equation*}
|G|=\sum \tilde{w}\left(p_{1} \cdots p_{t}\right)=\prod_{i=1}^{t}\left(\sum_{p_{i} \in P_{i}} \tilde{w}\left(p_{i}\right)\right)=\prod_{i=1}^{t}\left|G_{i}\right| \tag{2.4}
\end{equation*}
$$

where the first sum is taken over all $\left(p_{1}, \cdots, p_{t}\right) \in P_{1} \times \cdots \times P_{t}$. A directed graph is said to be strongly connected if there is a directed path from each vertex to every other vertex. Equivalently, each pair of distinct vertices can be joined by a circuit. A strongly connected component is a maximal strongly connected subgraph. Each connected component of a directed graph contains one or more strongly connected components. The strongly connected components partition the vertices of the graph, but edges that do not belong to any circuit of the graph are not in any of the strongly connected components. If $G$ has strongly connected components $G_{1}^{\prime}, \cdots, G_{s}^{\prime}$, and if $P_{1}^{\prime}, \cdots, P_{s}^{\prime}$ are the corresponding sets of circuit partitions, then $P=P_{1}^{\prime} \cdots P_{s}^{\prime}$ as before, so $|G|=\left|G_{1}^{\prime}\right| \cdots\left|G_{s}^{\prime}\right|$ by (2.3).

Expansion of a determinant relative to a vertex of the associated graph. Let $C(i)$ denote the collection of all circuits of $G$ passing through vertex $i$. Since every circuit partition of $G$ must have one factor which is a circuit passing through $i$, we have $P(G)=\{c \cdot p: c \in C(i), p \in P(G-c)\}$, where $G-c$ is obtained from $G$ by deleting all the vertices of the circuit $c$ and all the edges of $G$ incident with those vertices. By (2.3),

$$
|G|=\sum_{c \in C(i)} \sum_{p \in P(G-c)} \tilde{w}(c) \tilde{w}(p)=\sum_{c \in C(i)} \tilde{w}(c)\left[\sum_{p \in P(G-c)} \tilde{w}(p)\right] ; \text { i.e., }
$$

$$
\begin{equation*}
|G|=\sum_{c \in C(i)} \tilde{w}(c)|G-c| \tag{2.5}
\end{equation*}
$$

Let $C_{k}(i)$ denote the set of circuits of length $k$ passing through vertex $i$. Then it follows that

$$
\begin{equation*}
|G|=\sum_{1 \leq k \leq n}(-1)^{k-1} \sum_{c \in C_{k}(i)} w(c)|G-c|=m_{i i}|G-\{i\}|-\sum_{2 \leq k \leq n}(-1)^{k} \sum_{c \in C_{k}(i)} w(c)|G-c| . \tag{2.6}
\end{equation*}
$$

In the important special case where $G$ has no circuits of length $>2$,

$$
\begin{equation*}
|G|=m_{i i}|G-\{i\}|-\sum_{c \in C_{2}(i)} w(c)|G-c| . \tag{2.7}
\end{equation*}
$$

Expansion of a determinant relative to an exclusive circuit of the associated graph. A
circuit $c$ of a graph $G$ is exclusive if none of its directed edges belongs to another circuit. Let $G-E(c)$ denote the graph obtained from $G$ by deleting the directed edges (but not the vertices) of $c$. Then $P(G)=c \cdot P(G-c) \cup P(G-E(c))$, so (2.3) implies

$$
\begin{equation*}
|G|=\sum_{p \in P(G-c)} \tilde{w}(c) \tilde{w}(p)+\sum_{q \in P(G-E(c))} \tilde{w}(q)=\tilde{w}(c)|G-c|+|G-E(c)| . \tag{2.8}
\end{equation*}
$$

When $G$ has more than one exclusive circuit, successive applications of (2.8) with different exclusive circuits can be used to expand $|G|$ in terms of determinants of smaller subgraphs of $G$.

Associated graphs and determinants for a special class of symmetric matrices. We will need to compute determinants of matrices representing intersection forms, which are symmetric bilinear forms associated to resolutions of singularities of complex surfaces. In our applications, the symmetric matrices for these forms have integral entries greater than 1 along the diagonal, and nothing but 0 's and -1 's off the diagonal.

Let $M$ be an $n \times n$ symmetric matrix of the type just described. We represent $M$ by a graph on $n$ vertices, each labelled with the corresponding diagonal entry of $M$. The $i$-th such entry is the weight of the circuit of length 1 at the $i$-th vertex, so we denote it $w_{i}$ and call it the weight of the $i$-th vertex. Note: To simplify the labelling of our graphs, we omit the label $w_{i}$ when $w_{i}=2$. For distinct $i$ and $j$ in $\{1,2, \cdots, n\}$, either $m_{i j}=m_{j i}=0$ or $m_{i j}=m_{j i}=-1$. In the second case, we join the $i$-th and $j$-th vertices with a single unlabelled, undirected edge. This yields an undirected graph $G=G(M)$, some of whose vertices may be labelled with a positive integer (not 2). Conversely, any such graph $G$ together with an ordering of the vertices uniquely determines a symmetric matrix $M$ with positive integers $>1$ along its diagonal and with 0 's and -1 's elsewhere.

For such a graph $G$, the terms "connected" and "strongly connected" are synonymous. Every circuit of length $>1$ has signed weight -1 , so (2.6) and (2.8) become

$$
\begin{equation*}
|G|=w_{i}|G-\{i\}|-\sum_{k>1} \sum_{c \in C_{k}(i)}|G-c| \tag{2.9}
\end{equation*}
$$

and

$$
\begin{equation*}
|G|=|G-E(c)|-|G-c| . \tag{2.10}
\end{equation*}
$$

Example. We compute the determinant of the graph $G$ in Figure 5.
Using (2.3) is impractical because the number of circuit partitions is too large. It is better to break up the calculation into several easier calculations by use of (2.9)


Figure 5


Figure 6
and/or (2.10) in combination with (2.4).
First note that recursive use of (2.9) shows that a chain of $k$ vertices has determinant $k+1$. Now it follows from (2.9) that the cyclic graph or "necklace" formed by joining a vertex to the two ends of a chain of $k$ vertices (with $k \geq 2$ ) has determinant 0 .

If we apply (2.9) and (2.4) to the original graph $G$ at $v$, we get $|G|=$ $2 \cdot 8 \cdot 4-4 \cdot 4 \cdot 4-8 \cdot 3-8 \cdot 3-8-8=-64$. The last two terms arise from deletion of the two circuits of length 4 through $v$.

A better method is to apply (2.10) to the exclusive circuit of length 2 represented by the undirected edge $e$ of $G$. There are only two terms, as pictured in Figure 6. Thus $|G|=8 \cdot 0-4 \cdot 4 \cdot 4=-64$.

When $G$ is a tree, (2.9) and (2.10) become

$$
\begin{equation*}
|G|=w_{i}|G-\{i\}|-\sum_{c \in C_{2}(i)}|G-c| \tag{2.11}
\end{equation*}
$$

and

$$
\begin{equation*}
|G|=\left|G_{1}\right|\left|G_{2}\right|-\left|G_{1}^{\prime} \| G_{2}^{\prime}\right| \tag{2.12}
\end{equation*}
$$

where $G-E(c)=G_{1} \cup G_{2}$ and $G-c=G_{1}^{\prime} \cup G_{2}^{\prime}$ (canonical disjoint unions).
It is useful to recast (2.11) in notation that emphasizes what remains of the graph rather than what was deleted. For clarity, we now use subscripted $v$ 's rather than integers to name vertices.

Proposition. Let $v_{0}$ be a vertex with weight $w_{0}$ in a weighted tree $\tilde{G}$, and set $G=\tilde{G}-\left\{v_{0}\right\}$. Let $v_{1}, \cdots, v_{r}$ be the vertices joined to $v_{0}$ in $\tilde{G}$. For $i=1, \cdots, r$, let $G_{i}$ denote the component of $G$ that contains $v_{i}$, and put $G_{i}^{\prime}=G_{i}-\left\{v_{i}\right\}$. Then

$$
\begin{equation*}
|\tilde{G}|=w_{0} \prod_{i=1}^{r}\left|G_{i}\right|-\sum_{i=1}^{r}\left|G_{i}^{\prime}\right| \prod_{j \neq i}\left|G_{j}\right| \tag{2.13}
\end{equation*}
$$

so

$$
\begin{equation*}
w_{0}=\sum_{i=1}^{r} \frac{\left|G_{1}^{\prime}\right|}{\left|G_{i}\right|}+\frac{|\tilde{G}|}{\prod_{i=1}^{r}\left|G_{i}\right|} . \tag{2.14}
\end{equation*}
$$

(2.15) Application. Say that $v_{0}$ is a terminal vertex if exactly one other vertex of $\tilde{G}$ is joined to $v_{0}$. In that case $G_{1}=G$, so we write $G^{\prime}$ for $G_{1}^{\prime}$ and (2.13) becomes

$$
\begin{equation*}
|\tilde{G}|=w_{0}|G|-\left|G^{\prime}\right| \tag{2.16}
\end{equation*}
$$

In practice, we start with a tree $G$ and join a new vertex $v_{0}$ to a vertex $v_{i}$ of $G$ to form a larger tree $\tilde{\boldsymbol{G}}$. We try to choose $w_{0}$ so that $1 \leq|\widetilde{G}| \leq|G|$. Repetition of the process often produces a graph with determinant 1.

We also note for later use that if $G$ is positive definite (that is, if the associated matrix of $G$ is positive definite) and $|\tilde{G}|>0$, then $\tilde{G}$ is positive definite. This follows from the fact [18, p. 250] that a symmetric $n \times n$ matrix $M$ is positive definite if and only if the upper left $k \times k$ submatrix of $M$ has positive determinant for $k=1,2, \cdots, n$.
3. Perfect graphs. We will now apply these ideas to the special graphs under consideration in this paper.
(3.1) Lemma. If $G$ is a perfectable graph, then any tree containing $G$ is also perfectable.

Proof. Every tree $\tilde{G}$ containing $G$ can be constructed from $G$ by successively adjoining vertices $v_{n+1}, \cdots, v_{m}$ to $G$, each by means of a single edge. In this way we obtain a chain of trees $G=\widetilde{G}_{n} \subset \widetilde{G}_{n+1} \subset \cdots \subset \widetilde{G}_{m}=\widetilde{G}$. For $k=n+1, \cdots, m$, let $u_{k}$ be the vertex of $\widetilde{G}_{k-1}$, to which $v_{k}$ was joined in forming $\tilde{G}_{k}$. Set $\widetilde{G}_{k}^{\prime}=\tilde{G}_{k-1}-\left\{u_{k}\right\}$.

If $G$ is perfectable, fix a set of perfect weights $w_{1}, \cdots, w_{n}$. Assign weights $w_{k}$ to the vertices $v_{k}(k=n+1, \cdots, m)$ by defining $w_{k}=\left|\widetilde{G}_{k}^{\prime}\right|+1$. Then repeated use of Application (2.15) above shows that each weighted tree $\tilde{G}_{k}$ is positive definite with determinant 1. Hence each $\tilde{G}_{\boldsymbol{k}}$ is perfect. In particular, $\widetilde{G}=\widetilde{G}_{m}$ is perfect.

In view of Lemma (3.1), finding all perfectable graphs is equivalent to finding all minimal perfectable graphs. Our main result is that the graphs of Table I are minimal perfectable. To show this we must first check that each of these trees $G$ is positive definite and has determinant 1 . These properties can easily be verified recursively by use of (2.16). For example, consider entry (13) of Table I. We construct this weighted graph one vertex at a time, using the formula $\left|\widetilde{G}_{k+1}\right|=w_{k+1}\left|\widetilde{G}_{k}\right|-\left|\widetilde{G}_{k-1}\right|$ to compute the determinants. (See Figure 7.) $\tilde{G}_{k}$ is the subgraph spanned by $v_{1}, \cdots, v_{k}$. The determinants $\left|\widetilde{G}_{1}\right|, \cdots,\left|\widetilde{G}_{9}\right|$ are $12,23,34,45,179,1381,692,3$, and 1 . Since each determinant is positive and the last determinant is 1 , the graph is perfect. The other 24 examples are checked similarly.


Figure 7

To show that each of these perfectable graphs $G$ is minimally perfectable we must show that for each proper subgraph $G^{\prime}$ of $G$ there do not exist weights $w_{i} \geq 2$ for which $G^{\prime}$ is perfect. On the integer lattice $Z^{n}$, write $\left(x_{1}, \cdots, x_{n}\right)<\left(y_{1}, \cdots, y_{n}\right)$ if $x_{i} \leq y_{i}$ for all $i$ and $x_{i_{0}}<y_{i_{0}}$ for at least one index $i_{0}$. If $G$ is any graph on vertices $v_{1}, \cdots, v_{n}$, we have the mapping $|G|: \boldsymbol{Z}^{n} \rightarrow \boldsymbol{Z}$ whose value at $\left(w_{1}, \cdots, w_{n}\right)$ is the determinant of the weighted graph $G\left(w_{1}, \cdots, w_{n}\right)$.
(3.2) Lemma. Let $G$ be a tree on vertices $v_{1}, \cdots, v_{n}$ and let $\bar{w}=\left(w_{1}, \cdots, w_{n}\right)$ be a point of $\boldsymbol{Z}^{n}$ at which $G\left(w_{1}, \cdots, w_{n}\right)$ is positive definite. Then if $\bar{y}=\left(y_{1}, \cdots, y_{n}\right)>\bar{w}$, $G\left(y_{1}, \cdots, y_{n}\right)$ is also positive definite and $\left|G\left(y_{1}, \cdots, y_{n}\right)\right|>\left|G\left(w_{1}, \cdots, w_{n}\right)\right|$.

Proof. The first assertion is obvious, since the intersection matrix of $G(\bar{y})$ is the sum of the intersection matrix of $G(\bar{w})$ and the diagonal matrix $D(\bar{y}-\bar{w})$ whose diagonal entries are the non-negative integers $y_{i}-w_{i}$ and whose off-diagonal entries are zero. Since $G(\bar{w})$ is positive definite and $D(\bar{y}-\bar{w})$ is positive semi-definite, $G(\bar{y})$ is positive definite.

As for the determinant, (2.13) implies that, for each $i$,

$$
|G|\left(w_{1}, \cdots, w_{n}\right)=w_{i} \cdot\left|G-\left\{v_{i}\right\}\right|\left(w_{1}, \cdots, \hat{w}_{i}, \cdots, w_{n}\right)-\left(\text { terms that do not involve } w_{i}\right),
$$

where ^ means "omit this entry". Since $G-\left\{v_{i}\right\}$ is positive definite at $\left(w_{1}, \cdots, \hat{w}_{i}, \cdots, w_{n}\right),|G|$ is a strictly increasing linear function of $w_{i}$. Hence $|G|(\bar{y})>|G|(\bar{w})$ as claimed.

In fact, more is true.
(3.3) Lemma. Let $v_{0}$ be a vertex in a weighted tree $\tilde{G}$ and set $G=\tilde{G}-\left\{v_{0}\right\}$. Then the ratio $|\boldsymbol{G}| /|\widetilde{G}|$ strictly decreases as $\left(w_{0}, w_{1}, \cdots, w_{n}\right)$ increases with respect to $<$ in $\boldsymbol{Z}^{n+1}$.

Proof (induction on $n$ ). If $n=0$ the claim is just that $1 / w_{0}>1 / y_{0}$ whenever $w_{0}<y_{0}$. Now let $n>0$ and suppose the assertion to be true for all smaller trees. Let $v_{1}, \cdots, v_{r}$ be the vertices joined to $v_{0}$ in $\widetilde{G}$, and for $i=1, \cdots, r$, let $G_{i}$ be the component of $G$ that contains $v_{i}$. By (2.14),

$$
\frac{|\tilde{G}|}{|G|}=w_{0}-\sum_{i=1}^{r} \frac{\left|G_{i}^{\prime}\right|}{\left|G_{i}\right|},
$$

where $G_{i}^{\prime}=G_{i}-\left\{v_{i}\right\} .|\tilde{G}| /|G|$ clearly increases as a function of $w_{0}$, and by the induction hypothesis, each term $\left|G_{i}^{\prime}\right| /\left|G_{i}\right|$ strictly decreases as a function of ( $w_{1}, \cdots, w_{n}$ ). Thus $|\tilde{G}| /|\mathrm{G}|$ strictly increases as a function of $\left(w_{0}, w_{1}, \cdots, w_{n}\right)$.

To determine whether any particular graph is perfectable or not is now a finite calculation (perhaps a lengthy one if $G$ is complicated). If $G, v_{0}, G_{i}$, and $G_{i}^{\prime}$ are defined as above, then it is clear from Lemma (3.3) that for each $i$ the function $\left|G_{i}^{\prime}\right| /\left|G_{i}\right|$ achieves a maximum value $M_{i}$ on the part of $Z^{n}$ where $G_{i}$ is positive definite. Hence we have the bound

$$
w_{0}=\sum_{i=1}^{r} \frac{\left|G_{i}^{\prime}\right|}{\left|G_{i}\right|}+\frac{1}{\prod_{i=1}^{r}\left|G_{i}\right|} \leq \sum_{i=1}^{r} M_{i}+1
$$

for the weight $w_{0}$ on $v_{0}$. Since there is a similar bound for each weight, only finitely many choices of weights $w_{0}, \cdots, w_{n}$ need be checked. In practice it may not be easy to determine the maxima $M_{i}$ if $G_{i}$ is a large and complicated graph, but shortcuts of a number-theoretic nature are often available. For example, for most of the graphs in question, judicious use of (2.14) shows that many of the weights must be quite small (often 2 is the only possibility), and must be chosen such that, for all choices of $v_{0}$ and of components $G_{i}, G_{j}$ of $G-\left\{v_{0}\right\}$, the determinants $\left|G_{i}\right|$ and $\left|G_{j}\right|$ are coprime. If all but two weights have been determined, the last two weights must satisfy a quadratic equation whose coefficients are the determinants of various subgraphs and which has at most finitely many solutions (often none) in integers.

We have carried out these calculations for each of the graphs in Table II at the end of this paper, with this result:

## (3.4) Proposition. None of the graphs in Table II is perfectable.

It is now easy to complete the proof of the minimal perfectability of the graphs in Table I: each proper subgraph of a graph in Table I is contained in one of the non-perfectable graphs in Table II, so by Lemma (3.1) it must also be non-perfectable.

Likewise the proofs of Theorems (1.2), (1.3), (1.4), and part (b) of (1.5) are completed by verifying that every graph described in these theorems either contains a graph from Table I, and so is perfectable, or else is contained in a graph from Table II, and so is not perfectable. In particular, every tree not of the type shown in Figure 1 with $n=0,1$, or 2 , is accounted for. Also, for any perfectable graph $G$, a finite computational search suffices to find all sets of perfect weights. This was done, for instance, for the 8 -vertex graphs listed in Theorem (1.3). (The first part of Theorem (1.5) will be proved in the next section.)

As for the last assertion of Theorem (1.3), the existence of a single perfectable graph is sufficient to prove that

$$
\lim _{n \rightarrow \infty} \frac{\text { number of perfectable graphs on } n \text { vertices }}{\text { total number of trees on } n \text { vertices }}=1
$$

since for any given tree $G_{0}$, almost all trees with sufficiently many vertices contain $G_{0}$. Indeed, our results show that for each $n>31$, all but at most 5 trees on $n$ vertices are


Figure 8
perfectable, the 5 exceptions being the Dynkin diagrams $A_{n}$ and $D_{n}$, the extended Dynkin diagram $\tilde{D}_{n-1}$, and possibly one or both of the graphs in Figure 8.
4. Open questions and connections with number theory. The equation $w_{0}=$ $\sum\left(\left|G_{i}^{\prime}\right| /\left|G_{i}\right|\right)+1 /\left(\prod\left|G_{i}\right|\right)$ makes it clear that finding perfect graphs is equivalent to finding solutions of certain Diophantine equations. As a particularly interesting example, we now give a proof of part (a) of Theorem (1.5).
(4.1) Proposition. There is a one-to-one correspondence between perfect weighted graphs of the form $E_{p_{1}, \cdots, p_{r}}$ and solutions in integers of the equation

$$
w_{0}=\sum_{i=1}^{r} \frac{A_{i}}{B_{i}}+\frac{1}{\prod_{i=1}^{r} B_{i}}
$$

with $w_{0} \geq 2, B_{i} \geq 2,0<A_{i}<B_{i}$, and $\left(A_{i}, B_{i}\right)=1$.
First we need a lemma.
(4.2) Lemma. Given relatively prime positive integers $A<B$, there is a unique weighted chain $G$ (cf. Figure 9) such that $|G|=B,\left|G^{\prime}\right|=A$, and $w_{i} \geq 2$ for all $i$.

Proof (induction on $B$ ). If $B=2$, then $A=1$, and the unique solution is that $G$ consists of one vertex of weight 2 and $G^{\prime}$ is empty. Now let $B>2$ and assume that the result is true for all smaller numbers. Given $0<A<B$ with $(A, B)=1$, if $A=1$, then again one solution is for $G$ to be a vertex of weight $B$ and $G^{\prime}$ an empty graph. This solution is unique since no non-empty chain $G^{\prime}$ with all weights $\geq 2$ can have determinant 1.

If $A>1$, let $w_{0}$ be the unique integer $\geq 2$ for which $B<w_{0} A<A+B$. Put $A^{\prime}=w_{0} A-B$. Then $0<A^{\prime}<A<B$, and $\left(A^{\prime}, A\right)=1$, so by the induction hypothesis there is a unique weighted chain $G^{\prime}$ (cf. Figure 10) with $\left|G^{\prime}\right|=A$ and $\left|G^{\prime \prime}\right|=A^{\prime}$. But then for this choice of weights the graph $G$ in Figure 9 has $\left|G^{\prime}\right|=A$ and $|G|=w_{0}\left|G^{\prime}\right|-\left|G^{\prime \prime}\right|=B$ as required. $G$ is unique since if $\bar{G}, \bar{G}^{\prime}$ is another solution, with weights $\bar{w}_{0}, \cdots, \bar{w}_{\bar{n}}$, then $B=|\bar{G}|=\bar{w}_{0}\left|\bar{G}^{\prime}\right|-\left|\bar{G}^{\prime \prime}\right|=\bar{w}_{0} A-\left|\bar{G}^{\prime \prime}\right|$. But $\bar{w}_{2}, \cdots, \bar{w}_{\bar{n}} \geq 2$ implies that $\left|\bar{G}^{\prime \prime}\right|<A$. (To see this, note that the chain $A_{\bar{n}}\left(\bar{w}_{1}, \cdots, \bar{w}_{\bar{n}}\right)$ is positive definite if $\bar{w}_{i} \geq 2$ for all $i$, so by Lemma (3.3), $\left|\bar{G}^{\prime \prime}\right| / A=\left|\bar{G}^{\prime \prime}\right| /\left|\bar{G}^{\prime}\right|<\left|A_{\bar{n}-1}(2, \cdots, 2)\right| /\left|A_{\bar{n}}(2, \cdots, 2)\right|=\bar{n} /(\bar{n}+1)<1$.) Thus $\left|\bar{G}^{\prime \prime}\right|=A^{\prime}$ and $\bar{w}_{0}=w_{0}$, so uniqueness of the solution for $A, A^{\prime}$ implies that $G=\bar{G}$.

Remark. The proof shows that the unique solution $\left\{w_{1}, \cdots, w_{n}\right\}$ is just the set of integers that appear in the continued fraction expansion


Figure 9


Figure 10

$$
\left[\left[w_{0}, w_{1}, \cdots, w_{n}\right]\right]=w_{0}-\frac{1}{w_{1}-\frac{1}{w_{2}-\frac{1}{w_{n-1}-\frac{1}{w_{n}}}}}
$$

for $B / A$ (cf. Brieskorn [6, Section 2.4] for example).
We now prove Proposition (4.1). Let $G$ be a perfect weighted graph of type $E_{p_{1}, \cdots, p_{r}}$ whose central vertex $v_{0}$ is joined to terminal vertices $v_{1}, \cdots, v_{r}$ of graphs $G_{1}, \cdots, G_{r}$ of types $A_{p_{1}}, \cdots, A_{p_{r}}$ respectively. Then by (2.14) the weight of $v_{0}$ satisfies

$$
w_{0}=\sum_{i=1}^{r} \frac{A_{i}}{B_{i}}+\frac{1}{\prod_{r=1}^{r=1} B_{i}},
$$

where $B_{i}=\left|G_{i}\right|$ and $A_{i}=\left|G_{i}-\left\{v_{i}\right\}\right|$. Clearly $\left(A_{i}, B_{i}\right)=1$, as is seen by clearing denominators. Also $A_{i} / B_{i}<1$ since if all weights on $G_{i}$ are 2 then $A_{i} / B_{i}=p_{i} /\left(p_{i}+1\right)<1$, and $A_{i} / B_{i}$ is a decreasing function of its weights by Lemma (3.3).

Conversely, let $w_{0}=\sum_{i=1}^{r}\left(A_{i} / B_{i}\right)+1 /\left(\prod_{i=1}^{r} B_{i}\right)$ be a solution to this Diophantine equation, with $w_{0} \geq 2, B_{i} \geq 2$, and $A_{i}<B_{i}$ for all $i$. By Lemma (4.2), for each $i$ there exists a unique weighted chain $G_{i}$ as in Figure 9 with $\left|G_{i}\right|=B_{i}$ and $\left|G_{i}^{\prime}\right|=A_{i}$. Then the graph $G$ of type $E_{p_{1}, \cdots, p_{r}}$ whose central vertex $v_{0}$ has weight $w_{0}$ and whose arms are $G_{1}, \cdots, G_{r}$, with $v_{0}$ joined to $G_{i}$ at the vertex of $G_{i}-G_{i}^{\prime}$, is the required perfect graph. This completes the proof.

In a similar fashion, given any special type of graph we can identify the Diophantine equation that must be solved to produce the perfect weights. In particular, we will do this for the graphs of the type shown in Figure 1 (with no restrictions on $n$ ) for which we do not know all minimal solutions.

Given a rational number $A / B$, and Egyptian fraction expansion for $A / B$ is a decomposition of the form

$$
\frac{A}{B}=\sum_{i=1}^{N} \frac{1}{w_{i}}
$$

with the $w_{i}$ 's distinct positive integers. It is well known that every positive rational number can be so expressed, and in many different ways. Indeed, papers such as [10], [11], [2], and [3] either prove this fact in an especially nice way or give algorithms for producing such expansions with particular features, such as a minimal number of summands.
(4.3) Lemma. Let $G$ be the weighted graph in Figure 11 with all weights $\geq 2$. Then $|G|=1$ if and only if the following equation is satisfied:

$$
\begin{equation*}
\frac{A}{B}=\sum_{i=1}^{k} \frac{1}{x_{i}}+\frac{1}{B \prod_{i=1}^{k} x_{i}} \tag{4.4}
\end{equation*}
$$

where $A / B$ has the "continued fraction" expansion

$$
\left[\left[w_{0}, w_{1}, \cdots, w_{n}, w_{n+1}-\sum_{j=1}^{l} \frac{1}{y_{j}}\right]\right] .
$$

Proof. Expand $|G|$ about the vertex $w_{0}$ and apply (2.14) to obtain

$$
w_{0}=\sum_{i=1}^{k} \frac{1}{x_{i}}+\frac{\left|G_{2}\right|}{\left|G_{1}\right|}+\frac{|G|}{\left|G_{1}\right| \prod_{i=1}^{k} x_{i}},
$$

where $G_{1}$ and $G_{2}$ are the graphs in Figure 12. As in the remark following Lemma (4.2), it is easy to check by induction that

$$
\frac{\left|G_{2}\right|}{\left|G_{1}\right|}=\left[\left[w_{1}, \cdots, w_{n}, w_{n+1}-\sum_{j=1}^{l} \frac{1}{y_{j}}\right]\right] .
$$

The assertion now follows by putting $B=\left|G_{1}\right|$ and $A=w_{0} B-\left|G_{2}\right|$.
For $n \geq 3$ a complete set of minimal solutions to this equation is represented by graphs (20), (22), (23), (24), and (25) of Table I. For instance, the example (22) corresponds to the solution

$$
[[2,2,2,2,2,3]]=\frac{1}{2}+\frac{1}{3}+\frac{1}{5}+\frac{1}{7}+\frac{1}{179}+\frac{1}{24323}+\frac{1}{11 \cdot 2 \cdot 3 \cdot 5 \cdot 7 \cdot 179 \cdot 24323}
$$

To complete our analysis of perfectable graphs, then, we must find all minimal solutions for $n=0,1$, and 2 . Some solutions that may be minimal are represented by the seven


Figure 11


Figure 12

Table I. Minimal Perfectable Graphs.
Each graph is pictured with a set of perfect weights. (Unlabelled vertices have weight 2.) The choice of perfect weights is not unique in general.
(1)

(3)

(4) $\longrightarrow$
(6)

(7)

(8)

(9)

(10)

(11)

(12)


Table I. (Continued).
(13)


(21)

(22)

(24)

(25)


Table II. Examples of Non-perfectable Graphs.



weighted graphs of Table III.
To illustrate the role of Egyptian fractions in problems of this kind we will show how we determined the weights for the fourth example of Table III. (Examples 1, 2, 3, and 7 are similar.) Suppose that we wish to find perfect weights for the graph in Figure 13. By (4.4) we have

$$
\frac{A}{B}=\sum_{i=1}^{N} \frac{1}{x_{i}}+\frac{1}{B \prod_{i=1}^{N} x_{i}}
$$

with


Figure 13

Table III. Some Special Perfect Graphs.
(Unlabelled vertices have weight 2).
(1)

(2)

(3)

(4)

(5)


Table III. (Continued).
(6)

(7)


$$
\frac{A}{B}=[[2,2,2, y]]=\frac{4 y-3}{3 y-2}=1+\frac{1}{3}-\frac{1}{3(3 y-2)} .
$$

Thus if one of the $x_{i}$, say $x_{N}$, is equal to 3 , we have

$$
1=\sum_{i=1}^{N-1} \frac{1}{x_{i}}+\frac{1}{3(3 y-2)}+\frac{1}{3(3 y-2) \prod_{i=1}^{N-1} x_{i}} .
$$

That is, we seek a solution to the unit fraction equation

$$
\begin{equation*}
1=\sum_{i=1}^{N+1} \frac{1}{n_{i}} \tag{4.5}
\end{equation*}
$$

with the additional constraints that $n_{N}=3(3 y-2) \equiv 3 \bmod 9$ and that $n_{N+1}=\prod_{i=1}^{N} n_{i}$. For then the perfect weights are $x_{i}=n_{i}$ for $i=1, \cdots, N-1, x_{N}=3$, and $y=\left(n_{N}+6\right) / 9$.

Now the problem (4.5) of expressing 1 as the sum of unit fractions has been much studied and has a substantial literature. (See for example the bibliographies in [9] and [12].) In particular, in [4] we considered the condition $n_{N+1}=\prod_{i=1}^{N} n_{i}$ in some detail, and, by computer search techniques, obtained a list of solutions for small $N . N=9$ is the smallest value of $N$ for which we have a solution $\left(n_{1}, \cdots, n_{N+1}\right)$ that satisfies the extra condition $n_{N} \equiv 3 \bmod 9$ (but $n_{N}>3$ so that $y=\left(n_{N}+6\right) / 9 \geq 2$ ). The solution is

$$
\begin{gather*}
2,5,7,11,17,157,961,4398619,8687184244716671,  \tag{4.6}\\
75467170101653548887992820605571,
\end{gather*}
$$

5695293763151911320400374304363730155668749225304912374335630470.

These numbers give the perfect weights for the graph of Example 4.


Figure 14
It is worth noting that from a purely number-theoretic viewpoint the most interesting of our unsolved cases is the star-shaped graph in Figure 14. For this graph, Equation (4.4) is simply

$$
\begin{equation*}
w_{0}=\sum_{i=1}^{N} \frac{1}{x_{i}}+\frac{1}{\prod_{i=1}^{N} x_{i}} \tag{4.7}
\end{equation*}
$$

(cf. (4.5)), which we wish to solve in integers $w_{0}, x_{i}$, all $\geq 2$. In [4] we show that there is no solution for $N \leq 58$. No solution to (4.7) is known for any $N$.

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Mathematics Department<br>Wayne State University<br>Detroir, MI 48202<br>U.S.A.

