

ERRATA FOR RAYMOND D. GUMB AND KAREL LAMBERT'S *DEFINITIONS IN
NONSTRICT POSITIVE FREE LOGIC*,
Modern Logic, vol. 7 (1997), pp. 25–55.

p. 25, line 13: delete this line
and, after . . . gov

p. 25, line 20: $\text{haskell} \Leftarrow \text{Haskell}$

p. 27, line 18: $\text{bound}'' \Leftarrow \text{bound variable}''$

p. 27, line $\omega - 5$: $\text{Farmer} \Leftarrow , \text{Farmer}$

p. 28, line $\omega - 9$: $\text{Bencxivenga} \Leftarrow \text{Bencivenga}$

p. 28, line $\omega - 2$: $\text{haskell} \Leftarrow \text{Haskell}$

p. 28, line $\omega - 1$: $\text{Thomason} \Leftarrow \text{Thompson}$

p. 29, line 9:
regardless of whether there is a j ($1 \leq j \leq n$) such that $x_j = \text{err}$. Only non-

p. 29, line 12: $\text{haskell} \Leftarrow \text{Haskell}$

p. 30, line 8: $\text{nonlogical} \Leftarrow \text{logical}$

p. 30, line 16: $, \text{and are also} \Leftarrow \text{are also}$

p. 31, lines 5–7:

$$[t = t' \rightarrow (A \rightarrow A[t' / t])$$

$$[t = t' \rightarrow (A[t' / t] \rightarrow A)$$

$$\text{err} = t \leftrightarrow \neg \exists X (X = t)$$

p. 32, line 7–8:

nonempty set, $D_M = \{\mathbf{err}\}$, $\mathbf{err} \notin \mathbf{D}_M$, and I_M is an interpretation function mapping symbols into appropriate objects. \mathbf{D}_M is the *inner domain* of M ,

p. 32, line 15:

$$4. I(f^n(t_1, \dots, t_n)) = I(f^n)(I(t_1), \dots, I(t_n)) \in \mathbf{D} \cup D,$$

p. 34, line 2: definitions \Leftarrow definitions

p. 35, line 11:

$$\forall X_1 \dots \forall X_n \forall Y (f^n(X_1, \dots, X_n) = Y \leftrightarrow A(X_1, \dots, X_n, Y))$$

p. 35, lines 15–16:

$$A(X_1, \dots, X_n, Y),$$

3. the only nonlogical symbols occurring in $A(X_1, \dots, X_n, Y)$ are

p. 35, lines 19–20:

$$\forall X_1 \dots \forall X_n \exists Y A(X_1, \dots, X_n, Y)$$

$$\forall X_1 \dots \forall X_n \forall Y' (A(X_1, \dots, X_n, Y) \wedge (A(X_1, \dots, X_n, Y') \rightarrow Y = Y'))$$

p. 36, line 7:

$$\forall X_1 \dots \forall X_n \forall Y (C(X_1, \dots, X_n, Y) \rightarrow$$

$$(f^n(X_1, \dots, X_n) = Y \leftrightarrow (A(X_1, \dots, X_n, Y))))$$

p. 36, lines 15–17:

$$\forall X_1 \dots \forall X_n (C(X_1, \dots, X_n, Y) \rightarrow \exists Y A(X_1, \dots, X_n, Y))$$

$$\forall X_1 \dots \forall X_n \forall Y \forall Y' C(X_1, \dots, X_n, Y) \rightarrow$$

$$A(X_1, \dots, X_n, Y) \wedge (A(X_1, \dots, X_n, Y') \rightarrow Y = Y')$$

p. 36, lines 22–23:

$$\forall X_1 \dots \forall X_n \forall Y (f^n (X_1, \dots, X_n) = Y \leftrightarrow$$

$$(C(X_1, \dots, X_n) \rightarrow A(X_1, \dots, X_n, Y)) \wedge$$

$$(\neg C(X_1, \dots, X_n) \rightarrow Y = 0))$$

p. 37, $\omega - 11$:

$$\exists Y A(a_1, \dots, a_n, Y) \vee A(a_1, \dots, a_n, \text{err})$$

p. 38, $\omega - 15$ to $\omega - 14$: inner and outer domains coincide \Leftrightarrow outer but not the inner domain is empty.

p. 41, line 12:

then, for some term t , $\neg A(t)$, $\exists X (X = t) \in S$.

p. 42, lines 7-8:

cation for Padoa's method. Let p^n be an n -ary predicate that does not occur in the sentence $A(a_1, \dots, a_n)$ and

p. 42, line 10:

The predicate p is called *fully and explicitly definable* in

p. 42, line 15:

T . The predicate q^n is *implicitly definable* in T if

p. 42, line 18: and $\& \Leftrightarrow \&$

p. 42, line ω :

$$(ID) (T \cup T_q \vdash \forall^* a (p(a) \leftrightarrow q(a)))$$

p. 43, line 8:

$$1. (\{\exists X (X = a)\} \cup T \cup \{\neg p(a)\}) \cup (\{\exists X (X = a)\} \cup T_q \cup \{q(a)\})$$

p. 43, line 10:

$$2. (T \cup \{\neg p(\text{err})\}) \cup T_q \cup \{q(\text{err})\}$$

p. 44, line 17: $((\Leftarrow ($

p. 44, line 18: $((\Leftarrow ($

p. 44, line 20:

$$\forall^* a \{ \text{zero}(a) = 0 \} \equiv \forall X \{ \text{zero}(X) = 0 \} \wedge \text{zero}(\text{err}) = 0 \rightarrow$$

p. 44, line 22:

$$\forall X \forall Y \{ \text{zero}(X) = Y \leftrightarrow Y = 0 \} \wedge \forall X \{ \text{zero}(X) = \text{err} \leftrightarrow \text{err} = 0 \} \wedge$$

p. 45, line 9:

$$\forall^* a_1 \dots \forall^* a_n \forall^* b \{ f^n(a_1, \dots, a_n) = b \leftrightarrow A(a_1, \dots, a_n, b) \} \Rightarrow$$

p. 45, line 11:

$$\exists Y A(a_1, \dots, a_n, Y) \rightarrow b = Y \} \wedge \{ \neg \exists Y A(a_1, \dots, a_n, Y) \rightarrow b = \text{err} \}$$

p. 45, line $\omega - 4$:

$$\forall^* a_1 \dots \forall^* a_n \{ \neg (E!(a_1) \wedge \dots \wedge E!(a_n)) \rightarrow f^n(a_1, \dots, a_n) = \text{err} \}$$

p. 45, line $\omega - 2$:

$$\forall^* a_1 \dots \forall^* a_n \{ \neg (E!(a_1) \wedge \dots \wedge E!(a_n)) \rightarrow \neg p^n(a_1, \dots, a_n) \}$$

p. 46, lines 4–6:

$$\forall X_1 \dots \forall X_n \{ p^n(X_1, \dots, X_n) \leftrightarrow A(X_1, \dots, X_n) \} \Rightarrow$$

$$\forall^* a_1 \dots \forall^* a_n \{ p^n(a_1, \dots, a_n) \leftrightarrow E!(a_1) \wedge \dots \wedge E!(a_n) \wedge A(a_1, \dots, a_n) \}$$

p. 46, lines 10–13:

$$\begin{aligned} & \forall X_1 \dots \forall X_n \forall Y (f^n (X_1, \dots, X_n) = Y \leftrightarrow A(X_1, \dots, X_n, Y)) \Rightarrow \\ & \forall^* a_1 \dots \forall^* a_n \forall^* b (f^n (a_1 \dots a_n) = b \leftrightarrow \\ & \quad (E! (a_1) \wedge \dots \wedge E! (a_n) \rightarrow A(a_1, \dots, a_n, b)) \wedge \\ & \quad (\neg (E! (a_1) \wedge \dots \wedge E! (a_n)) \rightarrow b = \text{err})) \end{aligned}$$

p. 47, line 3:

$$(E! (a) \wedge E! (b) \wedge E! (c) \wedge b = 0 \rightarrow a = b \times c) \wedge$$

p. 47, line 12:

$$\forall^* a_1 \dots \forall^* a_n (p^n (a_1, \dots, a_n) \leftrightarrow A(a_1, \dots, a_n))$$

p. 47, line $\omega - 5$: (\Leftarrow [

p. 48, line 12: , and \Leftarrow ,

p. 48, line $\omega - 16$:

$$\forall X_1 \dots \forall X_n \forall Y (f^n (X_1, \dots, X_n) = Y \leftrightarrow A(a_1, \dots, a_n, Y)) \Rightarrow$$

p. 48, line $\omega - 2$:
 $\vdash_T \forall X \forall Y \forall Y' (A(X, Y) \wedge A(X, Y') \rightarrow Y = Y')$. However, the second, third,

p. 49, line 3:

$$\vdash_T \forall X (A(X, \text{err}) \rightarrow (\exists Y A(X, Y) \rightarrow \forall Y A(X, Y))) \wedge$$

p. 49, line $\omega - 11$:

$$\forall X \forall Y (\zeta (X) = Y \leftrightarrow Y = 0)$$

p. 51, line 18: 1979 \Leftarrow 1997

p. 51, line $\omega - 12$: I am \Leftarrow We are

p. 51, line $\omega - 5$:

this, consider in standard first-order logic that $\{\forall X A(X)\} \cup \{\forall Y \neg A(Y)\}$ is

p. 53, $\omega - 4$ to $\omega - 3$: *Festschrift . . . forthcoming* \Leftarrow *Das weite Spektrum der analytischen Philosophie*, (Berlin, de Gruyter), 135–149.