## 156. On the Elementary Partitions of the State Set in a Multiple-Input Semiautomaton

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1. Introduction. Determination of all homomorphic images of a given semiautomaton is equivalent to constructing all admissible partitions of its state set.

For the case of a one-input semiautomaton, there exists an efficient method for the construction of all admissible partitions. This can be done easily by determining all elementary partitions [1], [2].

For the case of a multiple-input semiautomaton, it seems complicated at first sight. But, even in this case, if all elementary partitions can be constructed, we can use the same procedure as the one-input case and we can obtain all admissible partitions.

In this note, we shall give an algorithm for constructing all elementary partitions of the state set in a multiple-input semiautomaton by using known elementary partitions for the one-input case. We shall borrow many notations and terms from [1].

2. Preliminaries. Consider a semiautomaton  $A = (S, \Sigma, M)$ , where S is a set of states,  $\Sigma = \{\sigma_1, \sigma_2, \dots, \sigma_n\}$   $(n \ge 2)$  is a set of inputs, and M is a set of transition mappings.

Definition 1. Let  $\pi$  be a partition of S.  $\tilde{\pi}$  is called the admissible closure of  $\pi$  in A if and only if  $\tilde{\pi} = \prod_{i \in A} \xi_i$ , where  $\{\xi_i : i \in A\}$  is the set of all admissible partitions in A such that  $\pi \leq \xi_i (i \in A)$ .

In section 4, we shall give a method for constructing the admissible closure  $\tilde{\pi}$  of  $\pi$ .

Definition 2. An admissible partition  $\pi \neq 0$  of S in A, where 0 means the identity partition, is called elementary if and only if for every admissible partition  $\pi'$  of S in  $A, 0 \leq \pi' \leq \pi$  implies  $\pi'=0$  or  $\pi'=\pi$ .

3. Structure of elementary partitions. For the semiautomaton given in the preceding section, we shall construct following one-input semiautomata:

Put  $\Sigma_i = \{\sigma_i\}$  and  $M_i = \{\sigma_i^A\} = \{\sigma_i^{A_i}\}$  for each natural number  $i \ (1 \le i \le n)$ . Thus, we obtain the one-input semiautomata  $A_i = (S, \Sigma_i, M_i) \ (1 \le i \le n)$ .

For each semiautomaton  $A_i$   $(1 \le i \le n)$ , the set of all elementary partitions of S in  $A_i$  can be determined by the procedure introduced in [1], [2]. We denote this set by  $\mathcal{P}_i$ .

We can now prove the following theorem on the structure of an elementary partition of S in A:

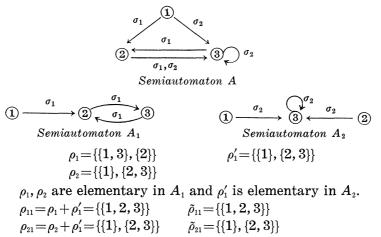
**Theorem 1.** If  $\pi$  is an elementary partition of S in A, then there exist elementary partitions  $\rho_i \in \mathcal{P}_i$   $(1 \le i \le n)$  such that  $\tilde{\rho} = \pi$ , where  $\rho = \sum_{i=1}^n \rho_i$ .

**Proof.** For each natural number i  $(1 \le i \le n)$ , we can consider  $\pi$  as an admissible partition of S in  $A_i$ . Thus, there exists an elementary partition  $\rho_i \in \mathcal{P}_i$  such that  $0 < \rho_i \le \pi$ .

Now, we can take the sum of partitions  $\rho = \sum_{i=1}^{n} \rho_i$ . Then, it is easy to see that  $0 < \rho \le \pi$ . We consider the admissible closure  $\tilde{\rho}$  of  $\rho$  in *A*. By virtue of  $0 < \rho \le \pi$  and the admissibility of  $\pi$  in *A*, we get  $0 < \tilde{\rho} \le \pi$ . Since  $\pi$  is elementary in *A*,  $\tilde{\rho}$  must be equal to  $\pi$ . Q.E.D.

**Remark.** The converse of the above theorem is not true. Indeed, there exist some elementary partitions  $\rho_i \in \mathcal{P}_i$   $(1 \le i \le n)$  such that  $\tilde{\rho}$  is not elementary in A, where  $\rho = \sum_{i=1}^{n} \rho_i$ .

Example. Let  $A = (\{1, 2, 3\}, \{\sigma_1, \sigma_2\}, M)$  be a semiautomaton whose transition graph is the following:



 $\tilde{\rho}_{_{21}}$  is elementary in *A*, but  $\tilde{\rho}_{_{11}}$  is not so.

**Theorem 2.** Let  $\rho_i$   $(1 \le i \le n)$  be partitions such that  $\rho_i \in \mathcal{P}_i$  and put  $\rho = \sum_{i=1}^{n} \rho_i$ . If there exist no partitions  $\rho'_i \in \mathcal{P}_i$   $(1 \le i \le n)$  such that  $\tilde{\rho}' < \tilde{\rho}$   $(\rho' = \sum_{i=1}^{n} \rho'_i)$ , then  $\tilde{\rho}$  is elementary in A.

**Proof.** Suppose  $\tilde{\rho}$  not to be elementary in A under the above assumption. Then, there exists an elementary partition  $\pi$  in A such that  $0 < \pi < \tilde{\rho}$ . In this case, from Theorem 1, there exist elementary partitions  $\rho'_i \in \mathcal{P}_i$   $(1 \le i \le n)$  such that  $\tilde{\rho}' = \pi$ , where  $\rho' = \sum_{i=1}^n \rho'_i$ . Consequently, we get  $0 < \tilde{\rho}' = \pi < \tilde{\rho}$ . But, this is a contradiction. Q.E.D.

4. Computation of  $\tilde{\pi}$ . Let  $\pi$  be a partition of S. For each natural number p, we construct inductively the partition  $\pi^{(p)}$  of S,

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starting with  $\pi^{(0)} = \pi$ . The construction method of  $\pi^{(p)}$   $(p \ge 1)$  from  $\pi^{(p-1)}$  is as follows:

Let  $\pi^{(p-1)} = \{B_1, B_2, \dots, B_m\}$  be a partition of S, where each of  $B_i$ 's is a block of  $\pi^{(p-1)}$ .

- (i) For each pair of numbers s, t  $(1 \le s \le n, 1 \le t \le m)$ , compute the set  $B_{st} = B_t \sigma_s^A$ .
- (ii) For each pair of numbers  $i, j (1 \le i, j \le m)$ , check whether  $B_i \sim B_j$ , according to the following definition:

 $B_i \sim B_j$  if and only if  $B_i = B_j$ , or there exist some numbers s, t ( $1 \leq s \leq n, 1 \leq t \leq m$ ) such that  $B_i \cap B_{st} \neq \emptyset, B_j \cap B_{st} \neq \emptyset$ .

(iii) For each pair of numbers  $i, j \ (1 \le i, j \le m)$ , check whether  $B_i \approx B_j$ , according to the following definition:

 $B_i \approx B_j$  if and only if there exists some sequence of numbers  $i=i_0, i_1, i_2, \dots, i_u=j$  such that  $B_{i_w} \sim B_{i_{w+1}}$  (w=0, 1, 2, ..., u-1).

- (iv) For each natural number i  $(1 \le i \le m)$ , compute the set  $\overline{B}_i = \bigcup_{B_i \approx B_i} B_j$ .
- (v) Let  $\pi^{(p)}$  be the partition of S whose set of all blocks is  $\{\overline{B}_i; 1 \le i \le m\}$ .

From the following procedure, the admissible closure  $\tilde{\pi}$  of  $\pi$  can be determined:

- (vi) Find a number q such that  $\pi^{(q)} = \pi^{(q-1)}$ .
- (vii) Put  $\tilde{\pi} = \pi^{(q)}$ .

5. Algorithm. We can now give the following algorithm for constructing all elementary partitions of the state set S in a semi-automaton  $A = (S, \Sigma, M)$   $(\Sigma = \{\sigma_1, \sigma_2, \dots, \sigma_n\})$ :

- (i) For each natural number i  $(1 \le i \le n)$ , construct the one-input semiautomaton  $A_i = (S, \Sigma_i, M_i)$ .
- (ii) For each natural number  $i \ (1 \le i \le n)$ , construct the set of all elementary partitions of S in  $A_i$ , i.e.,  $\mathcal{P}_i$ .
- (iii) Construct the following set:

$$\mathcal{P} = \{ \rho; \rho = \sum_{i=1}^{n} \rho_i, \rho_i \in \mathcal{P}_i \}.$$

(iv) Construct the following set:

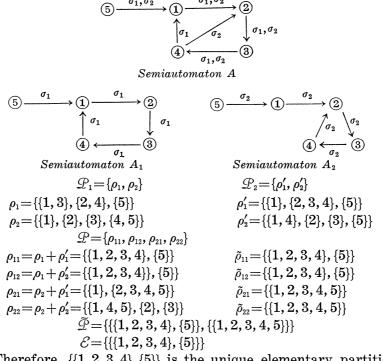
$$\widetilde{\mathcal{P}} = \{ \widetilde{\rho} ; \rho \in \mathcal{P} \}.$$

- (v) For each element  $\tilde{\rho}$  in  $\tilde{\mathcal{P}}$ , construct the set  $\tilde{\mathcal{P}}(\tilde{\rho}) = \{\tilde{\xi}; \tilde{\rho} < \tilde{\xi}, \tilde{\xi} \in \tilde{\mathcal{P}}\}.$
- (vi) Compute the following set:

$$\mathcal{E} = \tilde{\mathcal{P}} - \bigcup_{\tilde{\rho} \in \tilde{\mathcal{P}}} \tilde{\mathcal{P}}(\tilde{\rho}).$$

(vii)  $\mathcal{E}$  forms the set of all elementary partitions of S in A.

6. Example. Let  $A = (\{1, 2, 3, 4, 5\}, \{\sigma_1, \sigma_2\}, M)$  be a semiautomaton whose transition graph is the following:



Therefore,  $\{\{1, 2, 3, 4\}, \{5\}\}$  is the unique elementary partition of  $\{1, 2, 3, 4, 5\}$  in A.

## References

- [1] Ginzburg, A.: Algebraic Theory of Automata. Academic Press, New York—London (1968).
- [2] Yoeli, M., and A. Ginzburg: On homomorphic images of transition graphs. J. Franklin Inst., 278, 291-296 (1964).