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ONE DIMENSION IN PS AND PSI

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It is with regret that the Editor learned of the death of Professor Ivo Thomas on February 2, 1976. This paper was under composition at that time.

PS is $\{CCCpqpp, CCpqCCqrCpr\}$ with substitution and detachment; **PSI** adds *Cpp*. We find characteristic matrices for their one-dimensional expressions, formed only from *C* and *p*. The matrix for **PS** is:

				2	
	0 *1 2 3	1	1	1	1
M1	*1	0	1	1	1
	2	0	0	1	0
	3	0	0	0	0

that for PSI, m2, is m1 without the value 3. We deal explicitly only with PS, leaving the reader to adapt the argument to PSI.

Equivalence **M1** is such that if α , β are equivalent they have the same value. Since **PS** contains both syllogisms, if α , β are equivalent, $C\varphi\alpha\varphi\beta$ holds. The latter point pervades proof of Theorem I, the former that of Theorem II.

It is obvious that **M1** assigns to p the quadruple (0123) and to each compound α a unique quadruple of 1's and 0's. Such quadruples will be called "the value" of the expression; for compound α "vol' α ". Six expressions are now listed with their values and a numeral which names the value. If the numeral is prefixed to p, the result names the corresponding expression.

Þ	(0123)	1
Срр	(1110)	2
ССррр	(0111)	3
СрСрр	(1100)	4
СССрррр	(1111)	5
<i>CCppCpCpp</i>	(1101)	6

The following Table records that Cipjp is equivalent to kp $(1 \le i, j \le 6)$, $(2 \le k \le 6)$. Justification of its entries requires the proving of seventy-two implications, of which each turns out to be a substitution in one of twenty-four. Suffice it to say that the shortest is CPP (using capitalized variables for arbitrary implications), the most often used is CPCQP, the longest and most specific is CCCcppCppCppCppCpp (which follows from CCCpqpp, CCpCpqCpq, and CCCprrCcqrCcpqr).

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		j ,				_		
	С	1	2	3	4	5	6	
	11	2	4	4	4	4	4	١
	2	3	5	3	6	5	6	/
i) 3	5	2	5	4	5	6	L L
ı) 4	3	5	3	5	5	5 ($\rangle k$
	5	3	2	3	4	5	6	
	6	3	2	3	2	5	5	

Four theorems follow:

Theorem I Every compound α is equivalent to some kp ($2 \le k \le 6$).

Proof is by induction on the number of C's in α . (i) α contains one C: then α is 2p, equivalent to itself since *CPP* is provable. (ii) α is $C\beta_{\gamma}$ with n + 1 C's and the theorem holds for all compound expressions with $\leq n$ C's. Three subcases arise.

a) β is elementary, γ compound: then β is 1p, and by the induction hypothesis γ is equivalent to some $jp (2 \le j \le 6)$ so $C\beta\gamma$ is equivalent to C1pjp which is equivalent to 4p (Table).

b) β is compound, γ elementary: $C\beta\gamma$ is equivalent to 3p or 5p. Proof is similar to that of a).

c) β , γ are compound: by the induction hypothesis β is equivalent to some ip, γ to some jp ($2 \le i, j \le 6$) so $C\beta\gamma$ is equivalent to some kp ($2 \le k \le 6$) (Table). (i), (ii) prove the theorem.

Theorem II If val'a is 5, a is equivalent to 5p.

Proof: By Theorem I α is equivalent to some $kp \ (2 \le k \le 6)$ and therefore has the value k. But only for k = 5 does kp have the value 5.

Theorem III If val' α is 5, α is provable.

Proof: If val' α is 5, by Theorem II, C5p α is provable, but 5p is provable being a substitution in an axiom, so α is provable.

Theorem IV If α is provable, val' α is 5.

Proof: $\mathfrak{M1}$ is a **PS** matrix. Theorems III and IV show that $\mathfrak{M1}$ is characteristic for the class of expressions investigated.

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Historical Note In [1] Łukasiewicz gave

С	1	2	3	4
*1	1	2	4	4
23	1 1 2	1	4	4
3	2	2	x	2
4	1	2	1	1

with x = 2 as a PS matrix, with x = 1 as a PSI matrix. The former is not characteristic for one-dimensional expressions since it satisfies 6p; however the latter can be used to do the work of $\mathfrak{M}1$. It has come to mind that C. A. Meredith in a letter of 8/23/58 gave as PS and PSI matrices $\mathfrak{M}1$ without the value 2, and $\mathfrak{M}2$. The former suffers from the same disability as Łukasiewicz's, but Meredith's mapping into two values may have sparked the idea underlying $\mathfrak{M}1$.

We remark finally that no finite matrix can be characteristic for either general system.

REFERENCE

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