# An Unclassifiable Unidimensional Theory without OTOP 

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#### Abstract

A countable unidimensional theory without the omitting types order property (OTOP) has prime models over pairs and is hence classifiable. We show that this is not true for uncountable unidimensional theories.


1 Introduction In [4], Shelah settled the Main Gap for countable theories. Roughly speaking, this result shows that for a countable first order theory $T$ either there is a reasonable set of cardinal invariants which describe the isomorphism type of models of $T$ or in some technical sense the class of models of $T$ is chaotic. For countable unidimensional theories this dichotomy is explained by two properties, prime models over pairs (PMOP) and the omitting types order property (OTOP) both of which we proceed to define.

A stable theory $T$ is said to have PMOP if whenever $M_{0} \prec M_{i} \prec N \models T$ for $i=1,2$ and $M_{1}$ is independent from $M_{2}$ over $M_{0}$ then there is a prime model over $M_{1} \cup M_{2}$. If a unidimensional theory $T$ of cardinality $\lambda$ has PMOP then the class of models of $T$ is well behaved. For instance, every model of $T$ is determined up to isomorphism by its cardinality and a choice of relatively $|T|^{+}$-saturated submodel. This implies that $I(\kappa, T) \leq 2^{2^{T T}}$ for all $\kappa>|T|$.

A theory $T$ has the OTOP if there is a type $p(\bar{x}, \bar{y}, \bar{z})$ so that for every $\lambda$ there is $M$, a model of $T$, and a sequence $\left\langle\bar{a}_{\alpha}: \alpha<\lambda\right\rangle$ in $M$ so that $p\left(\bar{x}, \bar{a}_{\alpha}, \bar{a}_{\beta}\right)$ is realized in $M$ if and only if $\alpha<\beta$. This is a particular instance of the more general concept of having an $L_{\infty, \omega}$-definable order. $T$ has an $L_{\infty, \omega}$-definable order if there is an $L_{\infty, \omega^{-}}$ formula $\varphi(x, y)$ so that for every $\lambda$ there is $M$, a model of $T$, and a sequence $\left\langle\bar{a}_{\alpha}\right.$ : $\alpha<\lambda)$ in $M$ so that $M \models \varphi\left(\bar{a}_{\alpha}, \bar{a}_{\beta}\right)$ if and only if $\alpha<\beta$.

Theories which have $L_{\infty, \omega}$-definable orders were first studied in Shelah [3]. A theory $T$ which has an $L_{\infty, \omega}$-definable order has a chaotic class of models and is unclassifiable. A theory $T$ is said to be unclassifiable if for any regular cardinal $\lambda>|T|$ there are two nonisomorphic models of $T$ which can be forced to be isomorphic by a forcing notion which preserves cardinals (and adds no new small subsets of $\lambda$ ). A $\lambda$ closed forcing is an example of such a forcing. The class of models of such a theory
cannot have a reasonable set of cardinal invariants which determine isomorphism. It should be remarked that any unidimensional theory with PMOP is not unclassifiable (is classifiable).

A countable unidimensional theory which does not have the OTOP has PMOP (see 4], ch. 12). In this paper we show that this is not true for uncountable unidimensional theories. Precisely we give an example of a unidimensional theory with no $L_{\infty, \omega}$-definable order which nevertheless is unclassifiable.

We begin in this section by describing a particular theory, $T$, which has cardinality $2^{\aleph_{0}}$ and is unidimensional with $R^{\infty}(x=x)=2$. In Section 2 we prove that $T$ has quantifier elimination and that one cannot code arbitrarily long $L_{\infty, \omega}$-definable orders in models of $T$. In the third section we prove that $T$ is unclassifiable. In the final section we consider a family of examples similar to $T$ and give a characterization of those which are classifiable and those which are not.

Let us first describe the particular theory $T$ somewhat informally. The language for $T$ will be multisorted. One sort, $V$, will contain a vector space over the two element field. There will be predicates to pick out a descending chain of subspaces $U_{n}$ where $\left[U_{n}: U_{n+1}\right]=2$. It will be convenient to introduce predicates for all the cosets of $U_{n}$. There will be continuum many other sorts $S_{\eta}$ for $\eta \in 2^{\omega}$. Each of these can be thought of as a cover of $V$ and the fiber in $S_{\eta}$ above any $v \in V$ is acted on regularly by $V$. For any fixed $\eta$, there is no interaction between the fibers in $S_{\eta}$ above different elements from $V$. However, for every $v \in V$, there is interaction between the fiber above $v$ in $S_{\eta}$ and $S_{\mu}$ for $\eta, \mu \in 2^{\omega}, \eta \neq \mu$. Suppose that $\left.\eta\right|_{n}=\left.\mu\right|_{n}$ but $\eta(n) \neq \mu(n)$. $U_{n}$, the subspace of $V$, induces an equivalence relation with $2^{n}$ many classes on the fiber above $v$ in both $S_{\eta}$ and $S_{\mu}$. We will introduce relations between $S_{\eta}$ and $S_{\mu}$ which will give a bijection between these sets of equivalence classes.

Now we give a formal description of $T$. We begin with a definition.
Definition 1.1 Given $\mu \neq \eta \in 2^{\omega}, i(\mu, \eta)=$ the greatest $n<\omega$ for which $\left.\mu\right|_{n}=\left.\eta\right|_{n}$.
The language $L$ is defined as follows.

1. $L$ has sorts $\{V\} \cup\left\{S_{\eta}: \eta \in 2^{\omega}\right\}$.
2. In the sort $V$ there is a constant 0 , a binary function + , and unary predicates $\left\{U_{n}^{j}: n<\omega, j<2^{n}\right\}$.
3. For each $\eta \in 2^{\omega}$, there are functions $f_{\eta}: S_{\eta} \rightarrow V, g_{\eta}: S_{\eta} \times S_{\eta} \rightarrow V$ and $h_{\eta}$ : $V \times S_{\eta} \rightarrow S_{\eta}$.
4. For each pair $\mu \neq \eta \in 2^{\omega}$, there are binary relations $R_{(\mu, \eta)}^{j} \subset S_{\mu} \times S_{\eta}$ for $j<$ $2^{i(\mu, \eta)}$.

Now let $M$ be the following $L$-structure.

1. $V$ is a countable abelian group in which every element other than the identity has order $2 .+$ and 0 are interpreted by the group operation and the identity element of $V$, respectively. For each $n<\omega, U_{n}^{0}$ is a subgroup of $V$ such that $U_{n+1}^{0} \subset U_{n}^{0}$ and $\left[U_{n}^{0}: U_{n+1}^{0}\right]=2$. Finally, for each $n<\omega,\left\{U_{n}^{j}: j<2^{n}\right\}$ is a list of the cosets of $U_{n}^{0}$ in $V$.
2. For each $\eta \in 2^{\omega}, S_{\eta}=V \times V$ and the functions are interpreted as follows.
(a) $f_{\eta}(x, y)=y$.
(b) $g_{\eta}\left((x, y),\left(x^{\prime}, y^{\prime}\right)\right)=x-x^{\prime}$ provided that $y=y^{\prime}$ and is undefined otherwise.
(c) $h_{\eta}(u,(x, y))=(u+x, y)$.
3. Given $\mu \neq \eta \in 2^{\omega}, R_{(\mu, \eta)}^{j}\left((x, y),\left(x^{\prime}, y^{\prime}\right)\right)$ holds just if $y=y^{\prime}$ and $x-x^{\prime} \in$ $U_{i(\mu, \eta)}^{j}$.
We define $T$ to be the theory of $M$ in the language $L$.
Throughout the rest of this paper we will drop the subscripts in the functions $f_{\eta}, g_{\eta}, h_{\eta}$ whenever there is no confusion as to their domains. We end this section with an easy, if somewhat inelegant, lemma which will be of use later.

Lemma 1.2 Let $N$ be any model of $T$ and $\mu, \eta, \tau$ distinct members of $2^{\omega}$. Suppose that $a, a^{\prime} \in S_{\mu}, b, b^{\prime} \in S_{\tau}$, and $c, c^{\prime} \in S_{\eta}$.
(i) If $N \models R_{(\mu, \tau)}^{j}(a, b)$ for some $j<2^{i(\mu, \tau)}$ and if $u, u^{\prime} \in V$ with $u-u^{\prime} \in U_{i(\mu, \tau)}^{0}$, then $N \models R_{(\mu, \tau)}^{j}\left(h(u, a), h\left(u^{\prime}, b\right)\right)$.
(ii) If $N \models R_{(\mu, \tau)}^{j}(a, b) \wedge R_{(\mu, \tau)}^{j}\left(a^{\prime}, b^{\prime}\right)$ for some $j<2^{i(\mu, \tau)}$ and if $u, u^{\prime} \in V$ with $u-u^{\prime} \in U_{i(\mu, \tau)}^{0}$, then for any $k<2^{i(\mu, \tau)}$

$$
N \models R_{(\mu, \tau)}^{k}(a, h(u, b)) \longleftrightarrow R_{(\mu, \tau)}^{k}\left(a^{\prime}, h\left(u^{\prime}, b^{\prime}\right)\right) .
$$

(iii) Let $n_{0}=i(\mu, \eta), n_{1}=i(\mu, \tau)$, and $n_{2}=i(\tau, \eta)$ and suppose that $n_{2} \leq n_{0}$. Now choose $j_{l}<2^{n_{l}}$ for $l=0,1,2$ and suppose that

$$
N \models R_{(\mu, \eta)}^{j_{0}}(a, c) \wedge R_{(\mu, \tau)}^{j_{1}}(a, b) \wedge R_{(\tau, \eta)}^{j_{2}}(b, c) .
$$

Then

$$
N \models R_{(\mu, \eta)}^{j_{0}}\left(a^{\prime}, c^{\prime}\right) \wedge R_{(\mu, \tau)}^{j_{1}}\left(a^{\prime}, b^{\prime}\right) \quad \Longrightarrow \quad N \models R_{(\tau, \eta)}^{j_{2}}\left(b^{\prime}, c^{\prime}\right) .
$$

Proof: (i) and (ii) are immediate consequences of the definition of $T$. (iii): It suffices to show that (iii) holds when $N$ is our original model $M$. First note that $\left.\mu\right|_{n_{2}}=$ $\left.\eta\right|_{n_{2}}=\left.\tau\right|_{n_{2}}$ - since $i(\tau, \eta)=n_{2} \leq i(\mu, \eta)$ - and we conclude that $n_{2} \leq n_{1}$. Since we are working in $M$ we may assume that $a, a^{\prime}, b, b^{\prime}, c, c^{\prime} \in V(M)$ and our hypotheses imply that

$$
\begin{equation*}
(a-c),\left(a^{\prime}-c^{\prime}\right) \in U_{n_{0}}^{j_{0}},(a-b),\left(a^{\prime}-b^{\prime}\right) \in U_{n_{1}}^{j_{1}} \quad \text { and } \quad b-c \in U_{n_{2}}^{j_{2}} \tag{1}
\end{equation*}
$$

and it remains to prove that $b^{\prime}-c^{\prime} \in U_{n_{2}}^{j_{2}}$. Now

$$
(b-c)-\left(b^{\prime}-c^{\prime}\right)=\left(\left(a^{\prime}-b^{\prime}\right)-(a-b)\right)+\left((a-c)-\left(a^{\prime}-c^{\prime}\right)\right)
$$

and so the required result follows from (1) and the fact that $n_{2} \leq \min \left(n_{1}, n_{0}\right)$.

## 2 T does not have OTOP

Proposition 2.1 Let $N$ be any model of $T$. Then the set of partial isomorphisms $\{\rho: D \rightarrow N: D \subset N$ is finite $\}$ has the back and forth property.

Proof: Suppose $\bar{a}, \bar{b} \subset N$ satisfy the same atomic formulas (we will write $\bar{a} \equiv_{0} \bar{b}$ ) and fix $c \in N$. We must find $d \in N$ so that $\bar{a} c \equiv_{0} \bar{b} d$. Denote by $A, B$ the substructures
of $N$ generated by $\bar{a}, \bar{b}$, respectively. Then $A$ is finite and $A \equiv_{0} B$ in the obvious way. We may clearly assume that $c \notin A$. Now there are various possibilities.
Case 1: $\quad c \in V(N)$. Since $V(A)$ is finite, we can find $d \in V(N) \backslash V(B)$ so that

$$
\begin{equation*}
\text { for all } n<\omega, \quad c-d \in U_{n}^{0}(N) . \tag{2}
\end{equation*}
$$

We claim that this $d$ works. It is enough to verify that the following holds.
(i) For any $u \in V(A)$, any $n<\omega$ and $j<2^{n}$,

$$
u+c \in U_{n}^{j}(N) \quad \text { if and only if } \quad u^{\prime}+d \in U_{n}^{j}(N)
$$

where $u^{\prime}$ corresponds to $u$ in $V(B)$.
(ii) For any $\eta, \mu \in 2^{\omega}, v \in S_{\eta}(A), w \in S_{\mu}(A)$ and $j<2^{i(\mu, \eta)}$,

$$
N \models R_{(\mu, \eta)}^{j}(w, h(c, v)) \quad \text { if and only if } \quad N \models R_{(\mu, \eta)}^{j}\left(w^{\prime}, h\left(d, v^{\prime}\right)\right)
$$

where $v^{\prime}$ corresponds to $v$ in $S_{\eta}(B)$ and $w^{\prime}$ corresponds to $w$ in $S_{\mu}(B)$.
This is an easy consequence of (2) above and Lemma 1.2 (ii).
Case 2: For some $\eta \in 2^{\omega}, c \in S_{\eta}(N)$.
Subcase 1: For all $e \in A \backslash V(A), f(e) \neq f(c)$. In this case, let $u=f(c)$ and choose $u^{\prime} \in V(N)$, as in Case 1, so that $A u \equiv_{0} B u^{\prime}$. It is now trivial to check that any $d \in$ $f_{\eta}^{-1}\left(u^{\prime}\right)$ works.
Subcase 2: There is $e \in S_{\eta}(A)$ with $f(e)=f(c)$. Let $e^{\prime} \in B$ correspond to $e$. Now let $u=g(c, e)$ and choose $u^{\prime} \in V(N)$, as in Case 1 , so that $A u \equiv_{0} B u^{\prime}$. Then since $c=h(u, e)$, taking $d$ to be $h\left(u^{\prime}, e^{\prime}\right)$ works. All that remains is
Subcase 3: The set $Q=\left\{\tau \in 2^{\omega}\right.$ : there is $e \in S_{\tau}(A)$ with $\left.f(e)=f(c)\right\}$ is nonempty and does not contain $\eta$. In this case, let $n_{0}=\max \{i(\tau, \eta): \tau \in Q\}$ and choose $\mu \in Q$ with $i(\mu, \eta)=n_{0}$. Further, let $e \in S_{\mu}(A)$ satisfy $f(e)=f(c)$. Now for some $j_{0}<$ $2^{n_{0}}$ we have $N \models R_{(\mu, \eta)}^{j_{0}}(e, c)$. Let $e^{\prime}$ correspond to $e$ in $B$ and choose $d \in S_{\eta}(N)$ so that $N \models R_{(\mu, \eta)}^{j_{0}}\left(e^{\prime}, d\right)$. We claim that this $d$ works.
By virtue of our hypotheses, it suffices to check that if $\tau \in 2^{\omega}, j_{2}<2^{i(\tau, \eta)}, z \in S_{\tau}(A)$ and $u \in V(A)$, then

$$
N \models R_{(\tau, \eta)}^{j_{2}}(z, h(u, c)) \quad \Longrightarrow \quad N \models R_{(\tau, \eta)}^{j_{2}}\left(z^{\prime}, h\left(u^{\prime}, d\right)\right)
$$

where $z^{\prime}, u^{\prime}$ correspond to $z, u$ respectively in $B$.
So suppose that $N \models R_{(\tau, n)}^{j_{2}}(z, h(u, c))$. By Lemma-1.2(ii) we may assume that $u=u^{\prime}=0$. Again by Lemman 1.2(ii), we may assume that $\tau \neq \mu$ (note that if $\tau$ were equal to $\mu$, we would have $g(e, z) \equiv_{0} g\left(e^{\prime}, z^{\prime}\right)$ ). Now observe that $f(z)=f(c)=$ $f(e)$ (hence $\tau \in Q$ ) and so for some $j_{1}<2^{i(\mu, \tau)}$,

$$
N \models R_{(\mu, \tau)}^{j_{1}}(e, z) \wedge R_{(\mu, \tau)}^{j_{1}}\left(e^{\prime}, z^{\prime}\right) .
$$

Now since $\tau \in Q$, it follows that $i(\tau, \eta) \leq n_{0}$ and thus Lemma 1.2 (iii) implies that $N \models R_{(\tau, \eta)}^{j_{2}}\left(z^{\prime}, d\right)$, as required. This completes the proof of the proposition.

The next corollary is an immediate consequence of Proposition 2.1.
Corollary 2.2 $\quad T$ admits quantifier elimination.
The following is now immediate from the quantifier elimination.
Corollary 2.3 $T$ is superstable, unidimensional, and $R^{\infty}\left(S_{\eta}\right)=2$ for all $\eta<2^{\omega}$.
Corollary 2.4 T does not have an $L_{\infty, \omega}$-definable order. In particular, $T$ does not have OTOP.

Proof: $\quad$ Suppose that there is an $L_{\infty, \omega}$-formula $\varphi(\bar{x}, \bar{y})$ so that for every $\lambda$ there is a model $M \models T$ and a sequence $\left\langle a_{\alpha}: \alpha<\lambda\right\rangle \subset M^{n}$ for some $n<\omega$ such that

$$
N \models \chi\left(\bar{a}_{\alpha}, \bar{a}_{\beta}\right) \quad \text { if and only if } \quad \alpha<\beta
$$

By standard arguments then there is a model $N \models T$, an indiscernible sequence $\left\langle\bar{a}_{\alpha}\right.$ : $\alpha<\omega\rangle \subset N^{n}$ for some $n<\omega$ such that

$$
N \models \chi\left(\bar{a}_{\alpha}, \bar{a}_{\beta}\right) \quad \text { if and only if } \quad \alpha<\beta
$$

By indiscernibility the pairs $\left(\bar{a}_{0}, \bar{a}_{1}\right)$ and $\left(\bar{a}_{1}, \bar{a}_{0}\right)$ have the same $L$-type in $N$. It then follows from Proposition 2.1 hat they have the same $L_{\infty, \omega}$-type in $N$, a contradiction.

## 3 T is unclassifiable

Definition 3.1 Let $L^{*}$ be any language and $N_{1}, N_{2} L^{*}$-structures. Then a set, $\Pi \neq$ $\varnothing$, of partial isomorphisms $\pi: N_{1} \rightarrow N_{2}$, is called a $\lambda$-back-and-forth system for $N_{1}, N_{2}$ if
(i) for all $\pi \in \Pi$, $|\operatorname{dom}(\pi)|<\lambda$, and
(ii) given $\pi \in \Pi$ and $C \subset N_{1}, D \subset N_{2}$ with $|C|,|D|<\lambda$, there are $\pi_{1}, \pi_{2} \in \Pi$ extending $\pi$ such that $C \subset \operatorname{dom}\left(\pi_{1}\right)$ and $D \subset \operatorname{ran}\left(\pi_{2}\right)$.

If there is a $\lambda$-back-and-forth system $\Pi$ between two models which in addition is closed under unions of chains of length less than $\lambda$, then we say that the models are strongly equivalent. Such a $\Pi$ is $\lambda$-closed as a forcing notion (see Kunen [2]) and hence satisfies the forcing condition listed above. Hyttinen and Shelah [1] show that for any unsuperstable theory $T$ and any regular cardinal $\lambda>|T|$ there are models $M$ and $N$ of cardinality $\lambda$ which are strongly equivalent and nonisomorphic.

We will prove that $T$ is unclassifiable by coding certain graphs into models of $T$. Given a graph $G$ let $I_{G}, E_{G}$ denote the vertex and edge sets, respectively, of $G$. We will be interested in those graphs which are symmetric and irreflexive, contain no triangles, and satisfy: $\left|I_{G}\right| \geq 2^{\aleph_{0}}$ and every vertex has valence at least 2 . Let $\mathbf{G}$ be the class of all such graphs, and let $L^{\prime}$ be the language of graphs.

We will show that for every $G \in \mathbf{G}$, there is a corresponding model $N_{G} \models T$ with $\left|N_{G}\right|=\left|I_{G}\right|$, so that if $G, H \in \mathbf{G}$, then
(i) $N_{G} \cong N_{H}$ implies $G \cong H$, and
(ii) if $\left|I_{G}\right|=\left|I_{H}\right|=\lambda>2^{\aleph_{0}}$ and $G, H$ are strongly equivalent, then $N_{G}$ and $N_{H}$ are strongly equivalent.

It easily follows that $T$ is unclassifiable.
We now fix a large saturated model $\mathbf{C} \vDash T$, and from here on we will work inside $\mathbf{C}$. Thus $V, S_{\eta}$, and so on, will mean $V(\mathbf{C}), S_{\eta}(\mathbf{C})$, and so on. We also denote $U_{n}^{0}$ simply by $U_{n}$ for all $n<\omega$ and we denote $\bigcap_{n<\omega} U_{n}$ by $U$. Clearly $U$ has $2^{\aleph_{0}}$ many cosets in $V$. Notice that by the definition of $T, V$ is a vector space over the field of two elements. With this in mind we will sometimes refer to the linear dependence (independence) of subsets of $V$, and given any $A \subset V$ we will denote by $\bar{A}$ the subspace generated by $A$.

We may assume that $M \prec \mathbf{C}$ where $M$ is our original model. Denote $V(M) \subset V$ by $V^{*}$. We will assume that $U \cap V^{*}$ is infinite. Now since $M$ is countable, the set of cosets of $U$ which intersect $V^{*}$ is countable; let us list this set as $\left\{W_{k}: k<\omega\right\}$.

Definition 3.2 The model $N \prec \mathbf{C}$ is said to be small, if for any coset, $W$ of $U$,

$$
W \cap V(N) \neq \varnothing \quad \text { if and only if } \quad W=W_{k} \text { for some } k<\omega
$$

We now describe a general method for constructing small models. First let us adopt the following notation: for $v \in V$, a sequence above $v$ is a sequence $\bar{a}_{v}=\left\langle a_{v, \eta}: \eta \in\right.$ $\left.2^{\omega}\right\rangle$ where for each $\eta \in 2^{\omega}, a_{v, \eta} \in S_{\eta}$ and $f\left(a_{v, \eta}\right)=v$.

Now let $Y$ be any subset of $U$ and let $V_{Y}=\overline{Y \cup V^{*}}$. For each $v \in V_{Y}$, choose a sequence over $v$, say $\bar{a}_{v}$. Now let $N$ be the substructure of $\mathbf{C}$ generated by $V_{Y} \cup\left\{\bar{a}_{v}\right.$ : $\left.v \in V_{Y}\right\}$. Then it is easily seen that $N$ is a model of $T$ (simply check that the set of partial isomorphisms, $\rho: N \rightarrow M$, has the back and forth property, precisely as in the proof of Proposition 2.1. It follows from Corollary 2.2 hat $N \prec \mathbf{C}$ and we also note that $|N|=|Y|+2^{\aleph_{0}}$. Moreover, $V(N)=V_{Y}$ and thus, since $Y \subset U$, it follows that $N$ is small.

Given a graph $G \in \mathbf{G}$ we will construct the model $N_{G}$ as described above by making an appropriate choice for the set $Y$ and for the sequences $\bar{a}_{v}$. In order to code $G$, we will need to be able to recognize a given element of $V\left(N_{G}\right)$ as being one of two types according to which sequence was chosen over that element. The next lemmas demonstrate that this is possible.
Lemma 3.3 Let $v \in V$. Then there are sequences $\bar{a}_{v}=\left\langle a_{v, \eta}: \eta \in 2^{\omega}\right\rangle, \bar{b}_{v}=\left\langle b_{v, \eta}\right.$ : $\left.\eta \in 2^{\omega}\right\rangle$ over $v$ such that
(i) for all $\eta \neq \mu \in 2^{\omega}, \quad \models R_{(\mu, \eta)}^{0}\left(a_{v, \mu}, a_{v, \eta}\right)$;
(ii) for all $\eta \neq \mu \in 2^{\omega}, \quad \models \neg R_{(\mu, \eta)}^{0}\left(b_{v, \mu}, b_{v, \eta}\right)$.

Moreover, the sequences $\bar{a}_{v}, \bar{b}_{v}$ can be chosen for all $v \in V$ so that whenever $v_{1}, v_{2} \in$ $V$ satisfy $v_{1} \equiv{ }_{0} v_{2}$, then $\bar{a}_{v_{1}} \equiv{ }_{0} \bar{a}_{v_{2}}$ and $\bar{b}_{v_{1}} \equiv{ }_{0} \bar{b}_{v_{2}}$.

Proof: (i) is like (ii) except easier and the moreover clause is routine. (ii): Let $\eta_{0}, \ldots, \eta_{m} \in 2^{\omega}$. By compactness it suffices to show that in our original model $M$, there are $a_{0}, \ldots, a_{m} \in V(M)$ such that

$$
\text { for all } j \neq k \leq m, \quad a_{j}-a_{k} \in U_{i\left(\eta_{j}, \eta_{k}\right)-1} \backslash U_{i\left(\eta_{j}, \eta_{k}\right)}
$$

The proof is by induction on $m$, the case $m=1$ being trivial. So let us assume that $m>1$ and that $(\dagger)$ holds for all smaller $m$.

Let $n=\max \left\{i\left(\eta_{j}, \eta_{k}\right): j \neq k \leq m\right\}$ and assume without loss of generality that $n=i\left(\eta_{m-1}, \eta_{m}\right)$. By the inductive hypothesis, choose $a_{0}, \ldots, a_{m-1} \in V(M)$ so that $(\dagger)$ holds for these elements. Now choose $a_{m} \in V(M)$ so that $a_{m}-a_{m-1} \in U_{n-1} \backslash U_{n}$. We claim that $\left\{a_{0}, \ldots, a_{m}\right\}$ satisfies $(\dagger)$.

Fix $j<m-1$ and let $n_{1}=i\left(\eta_{j}, \eta_{m}\right), n_{2}=i\left(\eta_{j}, \eta_{m-1}\right)$. It suffices to show that $a_{j}-a_{m} \in U_{n_{1}-1} \backslash U_{n_{1}}$. Now since $a_{j}-a_{m-1} \in U_{n_{2}-1} \backslash U_{n_{2}}$ and $a_{m}-a_{m-1} \in$ $U_{n-1} \backslash U_{n}$, it is enough to show that $n>n_{1}$ and $n_{1}=n_{2}$. But this is an immediate consequence of the maximality of $n$ and Definition 1.1. Thus the lemma is proved.

We call any sequence satisfying (i) in the lemma a sequence of Type I over $v$, and any sequence satisfying (ii) a sequence of Type II over $v$.
Lemma 3.4 Suppose that $N \prec \mathbf{C}$ is small and $v \in V(N)$. Then $N$ does not contain both a Type I and a Type II sequence over v.

Proof: $\quad$ Suppose for a contradiction that $\bar{a}=\left\langle a_{\eta}: \eta \in 2^{\omega}\right\rangle \subset N$ is a Type I sequence over $v$ and $\bar{b}=\left\langle b_{\eta}: \eta \in 2^{\omega}\right\rangle \subset N$ is a Type II sequence over $v$. Since $N$ is small, there are $\mu \neq \eta \in 2^{\omega}$ and $k<\omega$ such that both $g\left(a_{\eta}, b_{\eta}\right)$ and $g\left(a_{\mu}, b_{\mu}\right)$ belong to $W_{k}$. Since $a_{\eta}=h\left(g\left(a_{\eta}, b_{\eta}\right), b_{\eta}\right)$ and $a_{\mu}=h\left(g\left(a_{\mu}, b_{\mu}\right), b_{\mu}\right)$, we obtain a contradiction from Lemma 1.2 i$)$.

We are now able to construct the models $N_{G}$ for $G \in \mathbf{G}$. Begin by fixing for each $v \in V$, a Type I sequence $\bar{a}_{v}$ and a Type II sequence $\bar{b}_{v}$ over $v$ as in the moreover clause of Lemma 3.3. Further fix a basis $X$ for $V^{*} \cap U$ and a subspace $Z \subset V^{*}$ such that $V^{*}=\left(V^{*} \cap U\right) \oplus Z(=\bar{X} \oplus Z)$.

Now fix $G \in \mathbf{G}$ and let $Y_{G}$ be a linearly independent subset of $U$ which contains $X$ and such that $\left|Y_{G}\right|=\left|I_{G}\right|$. Index $Y_{G}$ as $\left\{y_{p}: p \in I_{G}\right\}$ and let $V_{G}=\overline{V^{*} Y_{G}}$. Observe that by our choice of $Y_{G}$ we have

$$
V_{G}=\overline{Y_{G}} \oplus Z
$$

Now for each $v \in V_{G}$ choose one of the sequences $\bar{a}_{v}$ or $\bar{b}_{v}$ over $v$ as follows.

1. If $v=y_{p}$ for some $p \in I_{G}$, then choose the sequence $\bar{a}_{v}$.
2. If $v=y_{p}+y_{q}$ for some $p, q \in I_{g}$ such that $(p, q) \in E_{G}$, then choose the sequence $\bar{a}_{v}$.
3. Otherwise choose the sequence $\bar{b}_{v}$.

For $v \in V_{G}$, let $\operatorname{seq}(v)$ denote the sequence chosen above $v$ (that is, seq $(v)$ is either $\bar{a}_{v}$ of $\bar{b}_{v}$ ) and now let $N_{G}$ be the substructure generated by $V_{G} \cup\left\{\operatorname{seq}(v): v \in V_{G}\right\}$. By our earlier discussion, $N_{G}$ is a small (elementary) submodel of $\mathbf{C}$ with $\left|N_{G}\right|=\left|I_{G}\right|$. Let us call $v \in V_{G}\left(=V\left(N_{G}\right)\right)$ of Type I (Type II) in $N_{G}$ just if seq $(v)$ is of Type I (Type II). Given distinct elements $v_{0}, v_{1}, v_{2}$ of $V_{G}$, call the set $\left\{v_{0}, v_{1}, v_{2}\right\}$ a notable triple if each $v_{i}$ is of Type I in $N_{G}$ and $v_{0}+v_{1}+v_{2}=0$. We leave to the reader to verify the following consequence of our assumptions on $\mathbf{G}$ and the linear independence of $Y_{G}$.

## Lemma 3.5

(i) For any $v \in V_{G}, v=y_{p}$ for some $p \in I_{G}$ if and only if $v$ belongs to at least two notable triples.
(ii) For any $p, q \in I_{G},(p, q) \in E_{G}$ if and only if $y_{p}+y_{q}$ is of Type I in $N_{G}$.

Combining Lemmas 3.4 and 3.5 yields
Corollary 3.6 If $G, H \in \mathbf{G}$ and $N_{G} \cong N_{H}$, then $G \cong H$.
Corollary 3.7 $I(\lambda, T)=2^{\lambda}$ for all $\lambda \geq 2^{\aleph_{0}}$.
It remains to prove that if $G, H \in \mathbf{G}$ have cardinality $\lambda>2^{\aleph_{0}}$ and are strongly equivalent then $N_{G}$ and $N_{H}$ are strongly equivalent.

It is convenient at this juncture to expand the languge $L$ to a new language $L^{1}$ by adding a new unary predicate $Q$ to the sort $V$. For any $G \in \mathbf{G}, N_{G}$ is made into an $L^{1}$-structure by interpreting $Q$ as follows:

$$
N_{G} \models Q(v) \text { just if } v \text { has Type I in } N_{G} \text {. }
$$

Now let us fix $G, H \in \mathbf{G}$ so that $\left|I_{G}\right|=\left|I_{H}\right|=\lambda>2^{\aleph_{0}}$ and $G, H$ are strongly equivalent. We also fix a $\lambda$-back-and-forth system $\Phi$ for $G, H$ which is closed under the union of chains of length less than $\lambda$. It certainly suffices to find a $\lambda$-back-and-forth system for $N_{G}, N_{H}$ with respect to the enriched language $L^{1}$ which is also closed under the union of short chains. Call a substructure $A \subset N_{G}$ full if for all $v \in V(A)$, $\operatorname{seq}(v) \subset A$, similarly for substructures of $N_{H}$. Given the definition of the sequences $\bar{a}_{v}, \bar{b}_{v}$, it is straightforward to verify the following.

Fact 3.8 Suppose that $A \subset N_{G}, B \subset N_{H}$ are full substructures. Then the following are equivalent.

1. There is an onto $L^{1}$-partial isomorphism $\theta: V(A) \rightarrow V(B)$.
2. There is an onto $L^{1}$-partial isomorphism $\bar{\theta}: A \rightarrow B$ such that $\theta \subset \bar{\theta}$ and for all $v \in V(A), \bar{\theta}(\operatorname{seq}(v))=\operatorname{seq}(\theta(v))$.
Let $L^{2}$ be the restriction of $L^{1}$ to the sort $V$. Then $V_{G}, V_{H}$ are naturally $L^{2}$-structures, and we conclude from the above fact that it suffices to find a $\lambda$-back-and-forth system for the $L^{2}$-structures $V_{G}, V_{H}$. We now proceed to define just such a system. Recall that we fixed a $\lambda$-back-and-forth system $\Phi$ for the graphs $G, H$. Given $\varphi \in \Phi$, we define a corresponding $L^{2}$-partial isomorphism, $\theta_{\varphi}: V_{G} \rightarrow V_{H}$, as follows. Suppose that $\operatorname{dom}(\varphi)=J \subset I_{G}$ and $\operatorname{ran}(\varphi)=K \subset I_{H}$. Then $\theta_{\varphi}$ is the unique linear map satisfying
3. $\operatorname{dom}\left(\theta_{\varphi}\right)=\overline{\left\{y_{p}: p \in J\right\}} \oplus Z$ and $\operatorname{ran}\left(\theta_{\varphi}\right)=\overline{\left\{y_{r}: r \in K\right\}} \oplus Z$,
4. $\theta_{\varphi}\left(y_{p}\right)=y_{\varphi(p)}$, and
5. $\left.\theta_{\varphi}\right|_{Z}=\left.\mathrm{id}\right|_{Z}$.

That $\theta_{\varphi}$ is a $L$-partial isomorphism follows from the fact that $Y_{G}, Y_{H} \subset U$. To see that $\theta_{\varphi}$ is in fact a $L^{2}$-partial isomorphism, we need only check that $\theta_{\varphi}$ preserves the predicate $Q$, but this is immediate since $\varphi$ is an $L^{\prime}$-partial isomorphism. Note also that since $|J|<\lambda$, we have $\left|\operatorname{dom}\left(\theta_{\varphi}\right)\right|<\lambda$ also. Now let

$$
\Theta=\left\{\theta_{\varphi}: \varphi \in \Phi\right\} .
$$

Lemma 3.9 $\Theta$ is a $\lambda$-back-and-forth system closed under the union of chains of length less than $\lambda$.

Proof: Let $\theta \in \Theta$ and suppose that $C \subset V_{G}$ satisfies $|C|<\lambda$. We must find $\theta^{\prime} \in \Theta$ so that $\theta \subset \theta^{\prime}$ and $C \subset \operatorname{dom}\left(\theta^{\prime}\right)$. By definition, there is $\varphi \in \Phi$ such that $\theta=\theta_{\varphi}$. Now by ( $\left\llcorner\right.$ ) above, there is $J^{\prime} \subset I_{G}$ such that $\left|J^{\prime}\right|<\lambda$ and

$$
C \subset \overline{\left\{y_{p}: p \in J^{\prime}\right\}} \oplus Z .
$$

Since $\Phi$ is a $\lambda$-back-and-forth system, there is $\varphi^{\prime} \in \Phi$ which extends $\varphi$ and satisfies $J^{\prime} \subset \operatorname{dom}\left(\varphi^{\prime}\right)$. Clearly taking $\theta^{\prime}=\theta_{\varphi^{\prime}}$ satisfies our requirements. The closure under unions of short chains is straightforward since $\Phi$ is closed under such unions, and hence the lemma is proved.
We conclude from our discussions above that $N_{G}$ and $N_{H}$ are indeed strongly equivalent and therefore that $T$ is unclassifiable.

4 Other examples In this section we will introduce a family of examples of unidimensional theories that explores the ways that PMOP can fail. For each limit ordinal $\delta<\omega_{1}$, let $\eta_{\delta}=\left\langle\delta_{n}: n \in \omega\right\rangle$ be an increasing sequence whose limit is $\delta$. Let $S=\omega_{1}^{<\omega} \cup\left\{\eta_{\delta}: \delta\right.$ a limit ordinal, $\left.\delta<\omega_{1}\right\}$. Fix $X \subseteq S$ and define $T_{X}$ as in Definition 1.1 except that $\eta$ now ranges over $X$ and not just $2^{\omega}$. So in particular there is a sort $S_{\eta}$ for every $\eta \in X$, and all corresponding functions and relations are also defined. Let $\bar{X}=\left\{\delta: \eta_{\delta} \in X\right\}$. We can repeat the proofs in Section 2 and get

Proposition 4.1 For $X$ as above, $T_{X}$ is unidimensional and does not have an $L_{\infty, \omega^{-}}$ definable order.
One way that prime models can be built is by inevitable construction.
Definition 4.2 Fix a complete theory $T$.

1. A type $p \in S(A)$ is called inevitable if $p$ is realized in any model $M$ of $T$ which contains $A$.
2. Suppose $M$ is a model of $T$ and $A \subseteq M$. We say that $M$ is inevitably constructed over $A$ if $M$ can be enumerated as $\left\langle a_{\alpha}: \alpha<\beta\right\rangle$ for some $\beta$ so that for all $\alpha<\beta$, $\operatorname{tp}\left(a_{\alpha} /\left\{a_{\gamma}: \gamma<\alpha\right\} \cup A\right)$ is inevitable.
If $M$ is inevitably constructed over $A$ then it is prime over $A$. A strong way for a theory $T$ to have PMOP would be to have inevitably constructed models over pairs of models independent over a common submodel, ICMOP.

We remind the reader that $\omega_{1}^{\omega}$ is topologized by the initial segment topology: that is, basic open sets are of the form $\left\{\eta \in \omega_{1}^{\omega}: \nu \subseteq \eta\right\}$ where $v \in \omega_{1}^{n}$ for some $n \in \omega$. We will prove

Proposition 4.3 For $X$ as above, $T_{X}$ has ICMOP if and only if

1. there is no uncountable $C \subseteq X$ which can be homeomorphically embedded into $2^{\omega}$, and
2. $\bar{X}$ is not stationary.

Proof: Fix a monster model of $T_{X}$. We leave it to the reader to verify that $T_{X}$ has ICMOP if and only if there is an enumeration of all the sorts of $T_{X},\left\langle S_{\eta_{i}}: i<\alpha\right\rangle$ so that for some (equivalently any) $a \in V$ and sequence $\left\langle a_{i}: i<\alpha\right\rangle$ with $a_{i} \in S_{\eta_{i}}$ and $f_{\eta_{i}}\left(a_{i}\right)=a$, for each $j<\alpha, \operatorname{tp}\left(a_{j} /\left\{a_{i}: i<j\right\}\right)$ is isolated. In fact, this type will be isolated if and only if $\eta_{j}$ is not in the closure of $\left\{\eta_{i}: i<j\right\}$ (in the topology induced by $\omega_{1}^{<\omega}$ on $X$ ). Hence $T_{X}$ has ICMOP if and only if $X$ can be enumerated so that every initial segment of the enumeration is closed; call such an enumeration a good enumeration.

Let us begin the proof of the proposition by proving that if (1) and (2) hold then $T_{X}$ has ICMOP. Let $X_{\alpha}=X \cap \alpha^{\omega}$; note that $X_{\alpha}$ is closed. For any $\alpha<\omega_{1}, \alpha^{\omega}$ can be homeomorphically embedded into $2^{\omega}$. It follows that $X_{\alpha}$ can also be so embedded. If for any $\alpha, X_{\alpha}$ is uncountable this would contradict condition 1. So each $X_{\alpha}$ is countable. Since $\bar{X}$ is not stationary, there is a cub $C \subseteq \omega_{1}$ so that $C \cap \bar{X}=\varnothing$. Suppose that $C=\left\langle\alpha_{i}: i<\omega_{1}\right\rangle$ and this enumeration is in increasing order. Since $C$ and $\bar{X}$ have empty intersection, for any limit ordinal $\delta, X_{\alpha_{\delta}}=\cup_{i<\delta} X_{\alpha_{i}}$. Hence, if one enumerates $X_{\alpha_{i+1}} \backslash X_{\alpha_{i}}$ in order type $\omega$ and places these enumerations in increasing order, one obtains an enumeration of $X$ in order type $\omega_{1}$ which is good.

For the other direction we will show first that if condition 1 fails then $X$ does not have a good enumeration. So suppose that $C \subseteq X$ is an uncountable set which can be homeomorphically embedded into $2^{\omega}$. Take any enumeration and concentrate on the first $\omega_{1}$ many elements of $C$ which are listed. There must be a countable subset of these which is dense. Hence this enumeration cannot be good.

It follows then that if condition 1 does not fail then all the sets $X_{\alpha}$ introduced above are countable. So $X$ is of size $\omega_{1}$. We will prove by induction on order type that for all $\alpha$ and all such $X$, if $\bar{X}$ is stationary then $X$ cannot have a good enumeration of order type $\alpha$.

The case where $\alpha=\beta+1$ is straightforward. If $\operatorname{cf}(\alpha)=\omega$ then suppose we have an enumeration of the required type of order type $\alpha$. Suppose $\alpha=\cup_{n \in \omega} \alpha_{n}$ and $X$ is enumerated $\left\langle\eta_{i}: i<\alpha\right\rangle$. Let $X_{n}=\left\{\eta_{i}: i<\alpha_{n}\right\}$. Since $\bar{X}$ is stationary, for some $n, \bar{X}_{n}$ is stationary, and since it can be enumerated in order type $\alpha_{n}$ this is a contradiction.

Finally, suppose that $\mathrm{cf}(\alpha)=\omega_{1}$. Let $\left\langle\alpha_{i}: i<\omega_{1}\right\rangle$ be a continuous increasing chain of ordinals whose union is $\alpha$. Define a function $g: \omega_{1} \rightarrow \omega_{1}$ so that $g(\beta)=$ the least $i$ such that $X_{\gamma}$ is enumerated before $\alpha_{i}$ for all $\gamma<\beta . g$ is continuous and cofinal so there is a cub $C$ of limit ordinals so that for every $\beta \in C, g(\beta)=\beta$. (Note that $g(\beta)$ is not necessarily greater than or equal to $\beta$ but by Fodor's lemma, the set of $\beta$ for which $g(\beta)<\beta$ is not stationary.)

Claim 4.4 Let $\bar{Y}$ be the set of $\delta \in \bar{X}$ for which $\eta_{\delta}$ is in the closure of $\cup_{\beta<\delta} X_{\beta}$. Then $\bar{Y}$ is stationary.

Proof: $\quad$ Suppose not. Then since $\bar{X}$ is stationary, there is an $n \in \omega$ with the following property: there is a stationary set $\bar{W} \subseteq \bar{X}$ such that for any $\delta \in \bar{W}, \eta_{\delta} \upharpoonright n \neq \zeta \upharpoonright n$ for any $\zeta \in \cup_{\beta<\delta} X_{\beta}$. But now by $n$ applications of Fodor's lemma, we can assume that $\eta_{\delta} \upharpoonright n$ is constant for all $\delta \in \bar{W}$ which is clearly a contradiction to the choice of $n$.

Define a function $f: C \cap \bar{Y} \rightarrow \omega_{1}$ as follows: for $\delta \in C \cap \bar{Y}, \eta_{\delta}$ is in the closure of $\cup_{\beta<\delta} X_{\beta}$. So since we are assuming we have a good enumeration, $\eta_{\delta}$ must be enumerated before $\alpha_{\delta}$. Let $f(\alpha)=i$ where $i$ is the least such that $\eta_{\delta}$ is enumerated
before $\alpha_{i}$. $f$ is regressive so there is a set $Z \subseteq X$ so that $\bar{Z} \subseteq \bar{Y} \cap C$ is stationary and $f(\bar{Z})$ is constant, say $\zeta$. But then all $\eta_{\alpha}$ for $\alpha \in \bar{Z}$ are enumerated before $\alpha_{\zeta}$ which contradicts the inductive assumption.

Proposition 4.5 If $X \subseteq S$ and $C \subseteq X$ is uncountable and can be homeomorphically embedded into $2^{\omega}$, then $T_{X}$ is unclassifiable.

Proof: The proof from Section 3 works here as well.
Remark 4.6 One thing which is not at issue here is whether or not $T_{X}$ is classifiable or unclassifiable. The definitions of these terms that we have given in this paper seem to leave the possibility that there may be unidimensional theories which are neither. Although that cannot be ruled out now, it can for these theories. To see this, note that what is really at issue in attempting to build a prime model over a pair of small models is what the structure of the collection of fibers over a point in the vector space is. If there is only one way up to isomorphism to define the fibers over a point in the vector space then the theory has PMOP and is classifiable. On the other hand, if there are two ways which are not isomorphic then the proof outlined in Section 3 will show that that theory is unclassifiable.

The point of this investigation was not to show that the dichotomy was satisfied for this class of examples but to isolate the exact properties that gave rise to unclassifiability. We still believe that all examples $T_{X}$ where $\bar{X}$ is stationary are unclassifiable but at this moment we only know that they do not have inevitably constructed models over pairs.

Acknowledgments We are grateful to Alan Dow, Claude Laflamme, and Chris Laskowski for several helpful discussions during the preparation of this paper.

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