

## Remarks on the truth definition

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Let  $S$  be a mathematical theory and suppose that there is given a sequence  $A_0(a), A_1(a), A_2(a), \dots$  of all the formulae with a variable  $a$ . We define  $g(A_i(a))$  by  $i$ . ( $g$  is a symbol outside of  $S$ .)

Tarski's truth theory ([4], [5]) shows that  $S$  is inconsistent, if  $S$  contains two formulae  $E(a)$  and  $Tr(a)$  satisfying the following conditions:

- |     |                                   |                      |
|-----|-----------------------------------|----------------------|
| (1) | $E(i), Tr(g(A_i(i))) \rightarrow$ | $i = 0, 1, 2, \dots$ |
|     | $\rightarrow E(i), Tr(g(A_i(i)))$ |                      |
| (2) | $A_i(a) \rightarrow Tr(i)$        | $i = 0, 1, 2, \dots$ |
| (3) | $Tr(i) \rightarrow A_i(a)$        | $i = 0, 1, 2, \dots$ |

(In this paper, we use Gentzen's sequence developed in [1].)

Contradiction follows even in the case that (2) and (3) are ascertained to be satisfied only when they do not actually contain the variable  $a$ . In fact, such sequences mean

$$(2)^* \quad A \rightarrow Tr(g(A))$$

$$(3)^* \quad Tr(g(A)) \rightarrow A$$

for all formulae  $A$  without free variable. Let  $E(a)$  be the  $m$ -th formula  $A_m(a)$ , and consider the special case

$$E(m), Tr(g(A_m(m))) \rightarrow$$

$$\rightarrow E(m), Tr(g(A_m(m)))$$

of (1). Applying (2)\* and (3)\* to the formula  $A_m(m)$ , we obtain

$$E(m), A_m(m) \rightarrow$$

and

$$\rightarrow E(m), A_m(m).$$

These two sequences imply a contradiction, since  $E(m)$  is  $A_m(m)$ .

Though (2) may be read "if  $A_i(a)$  holds, then  $Tr(i)$  also holds" and (3) may be read in an analogous manner, no contradiction may be derived if these

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colloquial expressions are formulated not in the form of axioms but in the schemata of inference :

$$(2)' \quad \frac{\rightarrow A_i(a)}{\rightarrow Tr(i)} \qquad (3)' \quad \frac{\rightarrow Tr(i)}{\rightarrow A_i(a)} .$$

In this paper we shall show that there are a great many consistent systems each of which contains two formulae  $E(a)$  and  $Tr(a)$ , satisfies (1) and admits (2)' and (3)'.

Let  $S_0$  be an arbitrary consistent system, which contains the theory of natural numbers ([1]) and does not contain predicate  $E(a)$  or  $Tr(a)$ .

The system  $S_1$  is called *E-Tr-extension* of  $S_0$ , if  $S_1$  is obtained from  $S_0$  by adding new predicates  $E(a)$  and  $Tr(a)$ , axioms (1) and inferences (2)', (3)' under the presupposition that a sequence  $A_0(a), A_1(a), A_2(a), \dots$  of all the formulae with a variable  $a$  (which may contain new predicates  $E$  and  $Tr$ ) is fixed and  $g(A_i(a))$  is defined by  $i$ .

**THEOREM.** *If  $S_0$  is consistent, then E-Tr-extension of  $S_0$  is also consistent.*

**PROOF.** First we shall prove the consistency of the system  $S_2$ , which is obtained from  $S_0$  by adding (1) and the following inferences (4);

$$(4) \quad \frac{Tr(i_1), \dots, Tr(i_n) \rightarrow Tr(i)}{Tr(i_1), \dots, Tr(i_n) \rightarrow A_i(a)} \quad \text{and} \quad \frac{Tr(i_1), \dots, Tr(i_n), Tr(i) \rightarrow}{Tr(i_1), \dots, Tr(i_n), A_i(a) \rightarrow}$$

where  $i, i_1, \dots, i_n$  are integers and  $i$  is different from  $i_1, \dots, i_n$ .

To prove this, we have only to prove that any sequence of the form

$$Tr(i_1), \dots, Tr(i_n) \rightarrow Tr(i) \quad (i \neq i_1 \text{ and } \dots \text{ and } i \neq i_n)$$

cannot be *provable from* (1) in  $S_0$ , which is easily proved by substitution of  $a \neq i$  for  $Tr(a)$ .

Now we shall prove the following lemma.

**LEMMA.** *Let  $S_3$  be a consistent extension of  $S_2$  and  $A_{i_0}(a)$  be provable in  $S_3$ . Then  $Tr(i_0)$  is consistent with  $S_3$ .*

**PROOF.** We have only to prove that  $Tr(i_0), \Gamma \rightarrow \Delta$  is provable in  $S_3$ , if  $\Gamma \rightarrow \Delta$  is provable from  $\rightarrow Tr(i_0)$  and  $S_3$ . If  $\Gamma \rightarrow \Delta$  is a beginning sequence of  $S_2$  or  $\rightarrow Tr(i_0)$ , then the lemma is clear. Here we have only to prove the lemma under the hypothesis that  $Tr(i_0), \Pi \rightarrow \Delta$  is provable in  $S_3$  for any upper sequence  $\Pi \rightarrow \Delta$  of  $\Gamma \rightarrow \Delta$ . We must treat many cases but every case is trivial and the lemma is proved.

In virtue of this lemma, we see that (2)' holds in the maximal consistent extension of  $S_2$  and that the theorem holds.

Using the results in [2] or [3], we see easily the following corollary.

**COROLLARY.** *If  $S$  is consistent, then there exists a consistent extension  $\tilde{S}$  of  $S$ , in which (1), (2)', (3)' hold.*

**References**

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