

## Co-contractions of graphs and right-angled Artin groups

SANG-HYUN KIM

We define an operation on finite graphs, called *co-contraction*. Then we show that for any co-contraction  $\hat{\Gamma}$  of a finite graph  $\Gamma$ , the right-angled Artin group on  $\Gamma$  contains a subgroup which is isomorphic to the right-angled Artin group on  $\hat{\Gamma}$ . As a corollary, we exhibit a family of graphs, without any induced cycle of length at least 5, such that the right-angled Artin groups on those graphs contain hyperbolic surface groups. This gives the negative answer to a question raised by Gordon, Long and Reid.

[20F65](#), [20F36](#); [05C25](#)

### 1 Introduction

In this paper, by a *graph* we mean a finite graph without loops and without multi-edges. A *right-angled Artin group* is a group defined by a presentation with a finite generating set, where the relators are certain commutators between the generators. Such a presentation naturally determines the *underlying graph*, where the vertices correspond to the generators and the edges to the pairs of commuting generators. It is known that the isomorphism type of a right-angled Artin group uniquely determines the isomorphism type of the underlying graph by Droms [6] and Kim, Makar-Limanov, Neggers and Roush [13]. Also, right-angled Artin groups possess various group theoretic properties. To name a few, right-angled Artin groups are linear by Humphries [12], Hsu and Wise [11] and Davis and Januszkiewicz [4], biorderable by Duchamp and Thibon [8], biautomatic by Van Wyk [20] and moreover, admitting free and cocompact actions on finite-dimensional CAT(0) cube complexes by Charney and Davis [1], Meier and Van Wyk [15] and Niblo and Reeves [17].

On the other hand, it is interesting to ask what we can say about the isomorphism type of the underlying graph, if a right-angled Artin group satisfies a given group theoretic property. Let  $\Gamma$  be a graph. We denote the vertex set and the edge set of  $\Gamma$  by  $V(\Gamma)$  and  $E(\Gamma)$ , respectively. The *complement graph* of  $\Gamma$  is the graph  $\bar{\Gamma}$  defined by  $V(\bar{\Gamma}) = V(\Gamma)$  and  $E(\bar{\Gamma}) = \{\{u, v\} : \{u, v\} \notin E(\Gamma)\}$ . For a subset  $S$  of  $V(\Gamma)$  the *induced subgraph* on  $S$ , denoted by  $\Gamma_S$ , is defined to be the maximal subgraph of  $\Gamma$  with the vertex set  $S$ . This implies that  $V(\Gamma_S) = S$  and  $E(\Gamma_S) = \{\{u, v\} : u, v \in S \text{ and } \{u, v\} \in E(\Gamma)\}$ . If  $\Lambda$  is another graph, an *induced*  $\Lambda$  in  $\Gamma$  means an induced

subgraph isomorphic to  $\Lambda$  in  $\Gamma$ . We denote by  $C_n$  the cycle of length  $n$ . That is,  $V(C_n)$  is a set of  $n$  vertices, say  $\{v_1, v_2, \dots, v_n\}$ , and  $E(C_n)$  consists of the edges  $\{v_i, v_j\}$  where  $|i - j| \equiv 1 \pmod{n}$ . Let  $A(\Gamma)$  be the right-angled Artin group with its underlying graph  $\Gamma$ . Then, the following are true.

- $A(\Gamma)$  is coherent, if and only if  $\Gamma$  is *chordal*, ie  $\Gamma$  does not contain an induced  $C_n$  for any  $n \geq 4$ ; see Droms [5]. This happens if and only if  $[A(\Gamma), A(\Gamma)]$  is free; see H Servatius, Droms and B Servatius [19].
- $A(\Gamma)$  is a virtually 3-manifold group, if and only if each connected component of  $\Gamma$  is a tree or a triangle; see Droms [5] and Gordon [9]
- $A(\Gamma)$  is subgroup separable, if and only if no induced subgraph of  $\Gamma$  is a square or a path of length 3 by Metaftsis and Raptis [16]. This happens if and only if every subgroup of  $A(\Gamma)$  is also a right-angled Artin group, again by Droms [7].
- $A(\Gamma)$  contains a *hyperbolic surface group*, ie the fundamental group of a closed, hyperbolic surface, if there exists an induced  $C_n$  for some  $n \geq 5$  in  $\Gamma$ ; see Crisp and Wiest [3] and again Servatius, Droms and Servatius [19].

In [10], Gordon, Long and Reid proved that a word-hyperbolic (not necessarily right-angled) Coxeter group either is virtually free or contains a hyperbolic surface group. They also showed that certain (again, not necessarily right-angled) Artin groups do not contain a hyperbolic surface group, raising the following question.

**Question 1.1** Does  $A(\Gamma)$  contain a hyperbolic surface group if and only if  $\Gamma$  contains an induced  $C_n$  for some  $n \geq 5$ ?

In this paper, we give the negative answer to the above question. Let  $\Gamma$  be a graph and  $B$  be a set of vertices of  $\Gamma$  such that  $\Gamma_B$  is connected. The *contraction* of  $\Gamma$  relative to  $B$  is the graph  $\text{CO}(\Gamma, B)$  obtained from  $\Gamma$  by collapsing  $\Gamma_B$  to a vertex, and deleting loops or multi-edges. We define the *co-contraction*  $\overline{\text{CO}}(\Gamma, B)$  of  $\Gamma$  relative  $B$ , such that

$$\overline{\text{CO}}(\Gamma, B) = \overline{\text{CO}(\overline{\Gamma}, B)}.$$

Then we prove the following theorem, which will imply that  $A(\overline{C_n})$  contains  $A(\overline{C_5}) = A(C_5)$  and hence a hyperbolic surface subgroup, for  $n \geq 5$  (see Figure 3). An easy combinatorial argument shows that  $\overline{C_n}$  does not contain an induced cycle of length at least 5, for  $n > 5$ .

**Theorem** Let  $\Gamma$  be a graph and  $B$  be a set of vertices in  $\Gamma$ , such that  $\overline{\Gamma_B}$  is connected. Then  $A(\Gamma)$  contains a subgroup isomorphic to  $A(\overline{\text{CO}}(\Gamma, B))$ .

In this paper, the above theorem is proved in the following steps.

In [Section 2](#), we recall basic facts on right-angled Artin groups and HNN extensions. A *dual van Kampen diagram* is described. We owe the notation to Crisp and Wiest [\[3\]](#) where a closely related concept, a *dissection*, was defined and used with great clarity.

In [Section 3](#), we define *co-contraction* of a graph, and examine its properties.

In [Section 4](#), we prove the theorem by exhibiting an embedding of  $A(\overline{\text{CO}}(\Gamma, B))$  into  $A(\Gamma)$ . The main tool for the proof is a dual van Kampen diagram.

In [Section 5](#), we compute intersections of certain subgroups of right-angled Artin groups. From this, we deduce a more detailed version of the theorem describing some other choices of the embeddings.

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## 2 Preliminaries on right-angled Artin groups

Let  $\Gamma$  be a graph. The *right-angled Artin group on  $\Gamma$*  is the group presented as

$$A(\Gamma) = \langle v \in V(\Gamma) \mid [a, b] = 1 \text{ if and only if } \{a, b\} \in E(\Gamma) \rangle.$$

Each element of  $A(\Gamma)$  can be expressed as  $w = \prod_{i=1}^k c_i^{e_i}$ , where  $c_i \in V(\Gamma)$  and  $e_i = \pm 1$ . Such an expression is called a *word (of length  $k$ )* and each  $c_i^{e_i}$  is called a *letter* of the word  $w$ . We say the word  $w$  is *reduced*, if the length is minimal among the words representing the same element. For each  $i_0 = 1, 2, \dots, k$ , the word  $w_1 = \prod_{i=i_0}^k c_i^{e_i} \cdot \prod_{i=1}^{i_0-1} c_i^{e_i}$  is called a *cyclic conjugation* of  $w = \prod_{i=1}^k c_i^{e_i}$ . By a *subword* of  $w$ , we mean a word  $w' = \prod_{i=i_0}^{i_1} c_i^{e_i}$  for some  $1 \leq i_0 < i_1 \leq k$ . A letter or a subword  $w'$  of  $w$  is on the *left* of a letter or a subword  $w''$  of  $w$ , if  $w' = \prod_{i=i_0}^{i_1} c_i^{e_i}$  and  $w'' = \prod_{i=j_0}^{j_1} c_i^{e_i}$  for some  $i_1 < j_0$ .

The expression  $w_1 = w_2$  shall mean that  $w_1$  and  $w_2$  are equal as words (letter by letter). On the other hand,  $w_1 =_{A(\Gamma)} w_2$  means that the words  $w_1$  and  $w_2$  represent the same element in  $A(\Gamma)$ . For an element  $g \in A(\Gamma)$  and a word  $w$ ,  $w =_{A(\Gamma)} g$  means that the word  $w$  is representing the group element  $g$ .  $1$  denotes both the trivial element in  $A(\Gamma)$  and the empty word, depending on the context.

Let  $w$  be a word representing the trivial element in  $A(\Gamma)$ . A *dual van Kampen diagram*  $\Delta$  for  $w$  in  $A(\Gamma)$  is a pair  $(\mathcal{H}, \lambda)$  satisfying the following ([Figure 1 \(c\)](#)):

- (i)  $\mathcal{H}$  is a set of transversely oriented simple closed curves and transversely oriented properly embedded arcs in general position, in an oriented disk  $D \subseteq \mathbb{R}^2$ .
- (ii)  $\lambda$  is a map from  $\mathcal{H}$  to  $V(\Gamma)$  such that  $\gamma$  and  $\gamma'$  in  $\mathcal{H}$  are intersecting *only if*  $\lambda(\gamma)$  and  $\lambda(\gamma')$  are adjacent in  $\Gamma$ .
- (iii) Enumerate the boundary points of the arcs in  $\mathcal{H}$  as  $v_1, v_2, \dots, v_m$  so that  $v_i$  and  $v_j$  are adjacent on  $\partial D$  if and only if  $|i - j| \equiv 1 \pmod{m}$ . For each  $i$ , let  $a_i$  be the label of the arc that intersects with  $v_i$ . Put  $e_i = 1$  if, at  $v_i$ , the orientation of  $\partial D$  coincides with the transverse orientation of the arc that  $v_i$  is intersecting, and  $e_i = -1$  otherwise. Then  $w$  is a cyclic conjugation of  $v_1^{e_1} v_2^{e_2} \dots v_m^{e_m}$ .

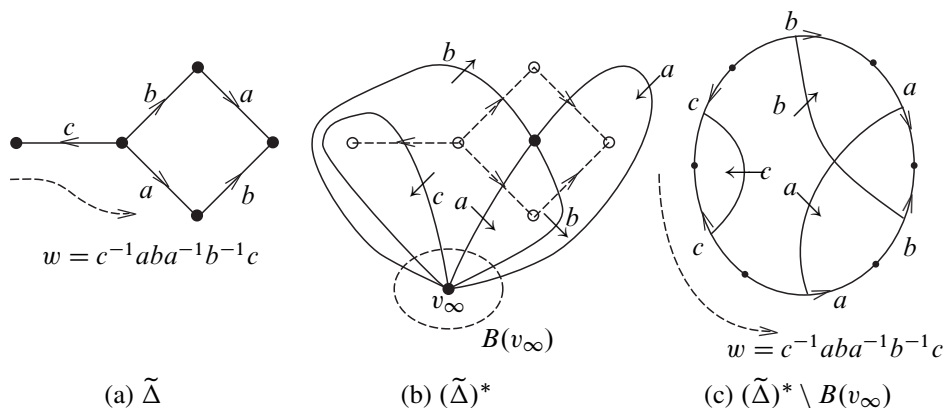


Figure 1: Constructing a dual van Kampen diagram from a van Kampen diagram, for  $w = c^{-1}aba^{-1}b^{-1}$  in  $\langle a, b, c \mid [a, b] = 1 \rangle$

Note that simple closed curves in a dual van Kampen diagram can always be assumed to be removed. Also, we may assume that two curves in  $\Delta$  are minimally intersecting, in the sense that there does not exist any bigon formed by arcs in  $\mathcal{H}$ . See Crisp and Wiest [3] for more details, as well as generalization of this definition to arbitrary compact surfaces, rather than a disk.

Let  $\tilde{\Delta} \subseteq S^2$  be a (standard) van Kampen diagram for  $w$ , with respect to a standard presentation  $A(\Gamma) = \langle V(\Gamma) \mid [u, v] = 1 \text{ if and only if } \{u, v\} \in E(\Gamma) \rangle$  (Figure 1). Consider  $\tilde{\Delta}^*$ , the dual of  $\tilde{\Delta}$  in  $S^2$ , and name the vertex which is dual to the face  $S^2 \setminus \tilde{\Delta}$  as  $v_\infty$ . Then for a sufficiently small ball  $B(v_\infty)$  around  $v_\infty$ ,  $\tilde{\Delta}^* \setminus B(v_\infty)$  can be considered as a dual van Kampen diagram with a suitable choice of the labeling map. Therefore a dual van Kampen diagram exists for any word  $w$  representing the trivial element in  $A(\Gamma)$ . Conversely, a van Kampen diagram  $\tilde{\Delta}$  for a word can be obtained from a dual van Kampen diagram  $\Delta$  by considering the dual complex again. So, the existence of a dual van Kampen diagram for a word  $w$  implies that  $w =_{A(\Gamma)} 1$ .

Given a dual van Kampen diagram  $\Delta$ , divide  $\partial D$  into segments so that each segment intersects with exactly one arc in  $\mathcal{H}$ . Let the label and the orientation of each segment be induced from those of the arc that intersects with the segment. The resulting labeled and directed graph structure on  $\partial D$  is called the *boundary* of  $\Delta$  and denoted by  $\partial\Delta$ .

We call each arc in  $\mathcal{H}$  labeled by  $q \in V(\Gamma)$  as a  $q$ -arc, and each segment in  $\partial\Delta$  labeled by  $q$  as a  $q$ -segment. Sometimes we identify the letter  $q^{\pm 1}$  of  $w$  with the corresponding  $q$ -segment. A connected union of segments on  $\partial\Delta$  is called an *interval*. By convention, a subword  $w_1$  of  $w$  shall also denote the corresponding interval (called  $w_1$ -interval) on  $\partial\Delta$ .

Now let  $\Delta = (\mathcal{H}, \lambda)$  be a dual van Kampen diagram on  $D \subseteq \mathbb{R}^2$ . Suppose  $\gamma$  is a properly embedded arc in  $D$ , which is either an element in  $\mathcal{H}$  or in general position with  $\mathcal{H}$ . Then one can cut  $\Delta$  along  $\gamma$  in the following sense. First, cut  $D$  along  $\gamma$  to get two disks  $D'$  and  $D''$ . Consider the intersections of the disks with the curves in  $\mathcal{H}$ . Then, let those curves in  $D'$  and  $D''$  inherit the transverse orientations and the labeling maps from  $\Delta$ . We obtain two dual van Kampen diagrams, one for each of  $D'$  and  $D''$ . Conversely, we can glue two dual van Kampen diagrams along identical words.  $\gamma$  is called an *innermost  $q$ -arc* if the interior of  $D'$  or  $D''$  does not intersect any  $q$ -arc.

**Definition 2.1** Let  $\Gamma$  be a graph. Let  $w$  be a word representing the trivial element in  $A(\Gamma)$ , and  $\Delta$  be a dual van Kampen diagram for  $w$ . Two segments on the boundary of  $\Delta$  are called a *canceling  $q$ -pair* if there exists a  $q$ -arc joining the segments. For any word  $w_1$ , two letters of  $w_1$  are called a *canceling  $q$ -pair* if there exist another word  $w'_1 =_{A(\Gamma)} w_1$  and a dual van Kampen diagram  $\Delta$  for  $w_1 w'^{-1}_1$ , such that the two letters are a canceling  $q$ -pair with respect to  $\Delta$ . A canceling  $q$ -pair is also called as a  $q$ -pair for abbreviation. A *canceling pair* is a canceling  $q$ -pair for some  $q \in V(\Gamma)$ .

For a group  $G$  and its subset  $P$ ,  $\langle P \rangle$  denotes the subgroup generated by  $P$ . For a subgroup  $H$  of  $A(\Gamma)$ ,  $w \in H$  shall mean that  $w$  represents an element in  $H$ .

**Lemma 2.2** Let  $\Gamma$  be a graph and  $q$  be a vertex of  $\Gamma$ . If a word  $w$  in  $A(\Gamma)$  has a  $q$ -pair, then  $w = w_1 q^{\pm 1} w_2 q^{\mp 1} w_3$  for some subwords  $w_1, w_2$  and  $w_3$  such that  $w_2 \in \langle \text{link}_\Gamma(q) \rangle$ . In this case,  $w$  is not reduced.

**Proof** There exists a word  $w' =_{A(\Gamma)} w$  and a dual van Kampen diagram  $\Delta$  for  $w w'^{-1}$ , such that a  $q$ -arc joins two segments of  $w$ .

Write  $w = w_1 q^{\pm 1} w_2 q^{\mp 1} w_3$ , where the letters  $q^{\pm 1}$  and  $q^{\mp 1}$  (identified with the corresponding segments on  $\partial\Delta$ ) are joined by a  $q$ -arc  $\gamma$  as in [Figure 2](#).

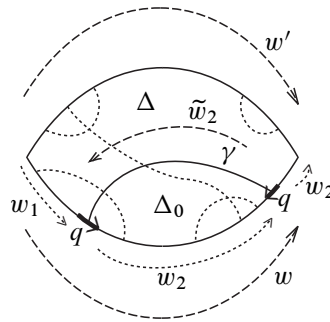


Figure 2: Cutting  $\Delta$  along  $\gamma$

Cut  $\Delta$  along  $\gamma$ , to get a dual van Kampen diagram  $\Delta_0$ , which contains  $w_2$  on its boundary. Give  $\Delta_0$  the orientation that coincides with the orientation of  $\Delta$  on  $w_2$ . Let  $\tilde{w}_2$  be the word, read off by following  $\gamma$  in the orientation of  $\Delta_0$ .  $\tilde{w}_2 \in \langle \text{link}_\Gamma(q) \rangle$ , for the arcs intersecting with  $\gamma$  are labeled by vertices in  $\text{link}_\Gamma(q)$ . Since  $\Delta_0$  is a dual van Kampen diagram for the word  $w_2 \tilde{w}_2$ , we have  $w_2 =_{A(\Gamma)} \tilde{w}_2^{-1} \in \langle \text{link}_\Gamma(q) \rangle$ .  $\square$

For  $S \subseteq V(\Gamma)$ , we let  $S^{-1} = \{q^{-1} : q \in S\}$  and  $S^{\pm 1} = S \cup S^{-1}$ . The following lemma is standard, and we briefly sketch the proof.

**Lemma 2.3** *Let  $\Gamma$  be a graph and  $S$  be a subset of  $V(\Gamma)$ . Then the following are true.*

- (1)  $\langle S \rangle$  is isomorphic to  $A(\Gamma_S)$ .
- (2) Each letter of any reduced word in  $\langle S \rangle$  is in  $S^{\pm 1}$ .

**Proof** (1) The inclusion  $V(\Gamma_S) \subseteq V(\Gamma)$  induces a map  $f: A(\Gamma_S) \rightarrow A(\Gamma)$ . Let  $w$  be a word representing an element in  $\ker f$ . Since  $w =_{A(\Gamma)} 1$ , there exists a dual van Kampen diagram  $\Delta$  for the word  $w$  in  $A(\Gamma)$ . Remove simple closed curves labeled by  $V(\Gamma) \setminus V(\Gamma_S)$ , if there is any. Since the boundary of  $\Delta$  is labeled by vertices in  $V(S)$ ,  $\Delta$  can be considered as a dual van Kampen diagram for the word  $w$  in  $A(\Gamma_S)$ . So we get  $w =_{A(\Gamma_S)} 1$ .

(2)  $w =_{A(\Gamma)} w'$  for some word  $w'$  such that the letters of  $w'$  are in  $S$ . Let  $\Delta$  be a dual van Kampen diagram for  $ww'^{-1}$ . If  $w$  contains a  $q$ -segment for some  $q \notin S$ , then a  $q$ -arc joins two segments in  $\Delta$ , and these segments must be in  $w$ . This is impossible by Lemma 2.2.  $\square$

From this point on,  $A(\Gamma_S)$  is considered as a subgroup of  $A(\Gamma)$ , for  $S \subseteq V(\Gamma)$ . Let  $H$  be a group and  $\phi: C \rightarrow D$  be an isomorphism between subgroups of  $H$ . Then

we define  $H*_\phi = \langle H, t \mid t^{-1}ct = \phi(c), \text{ for } c \in C \rangle$ , which is the HNN extension of  $H$  with the amalgamating map  $\phi$  and the stable letter  $t$ . Sometimes, we explicitly state what the stable letter is. If  $C = D$  and  $\phi$  is the identity map, then we let  $H*_C = \langle H, t \mid t^{-1}ct = t \text{ for } c \in C \rangle$ .

For a vertex  $v$  of a graph  $\Gamma$ , the *link of  $v$*  is the set

$$\text{link}_\Gamma(v) = \{u \in V(\Gamma) : u \text{ is adjacent to } v\}.$$

**Lemma 2.4** *Let  $\Gamma$  be a graph. Suppose  $\Gamma'$  is an induced subgraph of  $\Gamma$  such that  $V(\Gamma') = V(\Gamma) \setminus \{v\}$  for some  $v \in V(\Gamma)$ . Let  $C$  be the subgroup of  $A(\Gamma')$  generated by  $\text{link}_\Gamma(v)$ . Then the inclusion  $A(\Gamma') \hookrightarrow A(\Gamma)$  extends to the isomorphism  $f: A(\Gamma')*_C \rightarrow A(\Gamma)$  such that  $f(t) = v$ .*

**Proof** Immediate from the definition of right-angled Artin groups. □

We first note the following general lemma.

**Lemma 2.5** *Let  $H$  be a group and  $\phi: C \rightarrow D$  be an isomorphism between subgroups  $C$  and  $D$ . Let  $K$  be a subgroup of  $H$  and  $J = \langle K, t \rangle \leq H*_\phi$ . We let  $\psi: J \cap C \rightarrow J \cap D$  be the restriction of  $\phi$ . Then the inclusion  $J \cap H \hookrightarrow J$  extends to the isomorphism  $f: (J \cap H)*_\psi \rightarrow J$  such that  $f(\hat{t}) = t$ , where  $\hat{t}$  and  $t$  denote the stable letters of  $(J \cap H)*_\psi$  and  $H*_\phi$ , respectively.*

**Proof** Note that  $G = H*_\phi$  acts on a tree  $T$ , with a vertex  $v_0$  and an edge  $e_0 = \{v_0, t.v_0\}$  satisfying  $\text{Stab}(v_0) = H$  and  $\text{Stab}(e_0) = C$  [18]. Let  $T_0$  be the induced subgraph on  $\{j.v_0 : j \in J\}$ . For each vertex  $j.v_0$  of  $T_0$ , write  $j = k_1 t^{\epsilon_1} k_2 t^{\epsilon_2} \dots k_m t^{\epsilon_m}$ , where  $k_i \in K$  and  $\epsilon_i = \pm 1$  for each  $i$ . Then the following sequence in  $V(T_0)$

$$\begin{aligned} v_0 &= k_1.v_0, \\ k_1 t^{\epsilon_1}.v_0 &= k_1 t^{\epsilon_1} k_2.v_0, \\ k_1 t^{\epsilon_1} k_2 t^{\epsilon_2}.v_0 &= k_1 t^{\epsilon_1} k_2 t^{\epsilon_2} k_3.v_0, \\ &\dots \\ k_1 t^{\epsilon_1} k_2 t^{\epsilon_2} k_3 \dots t^{\epsilon_m}.v_0 &= j.v_0 \end{aligned}$$

gives rise to a path in  $T_0$  from  $v_0$  to  $j.v_0$ . Hence  $T_0$  is connected. Note that  $\psi: J \cap C = \text{Stab}_J(e_0) \rightarrow J \cap D = \text{Stab}_J(e_0)^t$ . Since  $J$  acts on a tree  $T_0$ , we have an isomorphism  $J \cong \text{Stab}_J(v_0)*_\psi = (J \cap H)*_\psi$ . □

### 3 Co-contraction of graphs

Let  $\Gamma$  be a graph and  $B \subseteq V(\Gamma)$ . We say  $B$  is *connected*, if  $\Gamma_B$  is connected.  $B$  is *anticonnected*, if  $\overline{\Gamma_B}$  is connected.

**Definition 3.1** Let  $\Gamma$  be a graph and  $B \subseteq V(\Gamma)$ .

- (i) If  $B$  is connected, the *contraction of  $\Gamma$  relative to  $B$*  is the graph  $\text{CO}(\Gamma, B)$  defined by:

$$V(\text{CO}(\Gamma, B)) = (V(\Gamma) \setminus B) \cup \{v_B\}$$

$$E(\text{CO}(\Gamma, B)) = E(\Gamma_{V(\Gamma) \setminus B}) \cup \{\{v_B, q\} : q \in V(\Gamma) \setminus B \text{ and } \text{link}_\Gamma(q) \cap B \neq \emptyset\}$$

- (ii) If  $B$  is anticonnected, the *co-contraction of  $\Gamma$  relative to  $B$*  is the graph  $\overline{\text{CO}}(\Gamma, B)$  defined by:

$$V(\overline{\text{CO}}(\Gamma, B)) = (V(\Gamma) \setminus B) \cup \{v_B\}$$

$$E(\overline{\text{CO}}(\Gamma, B)) = E(\Gamma_{V(\Gamma) \setminus B}) \cup \{\{v_B, q\} : q \in V(\Gamma) \setminus B \text{ and } \text{link}_\Gamma(q) \supseteq B\}$$

- (iii) More generally, if  $B_1, B_2, \dots, B_m$  are disjoint connected subsets of  $V(\Gamma)$ , then inductively define

$$\text{CO}(\Gamma, (B_1, B_2, \dots, B_m)) = \text{CO}(\text{CO}(\Gamma, (B_1, B_2, \dots, B_{m-1})), B_m)$$

and if  $B_1, B_2, \dots, B_m$  are disjoint anticonnected subsets, then similarly,

$$\overline{\text{CO}}(\Gamma, (B_1, B_2, \dots, B_m)) = \overline{\text{CO}}(\overline{\text{CO}}(\Gamma, (B_1, B_2, \dots, B_{m-1})), B_m).$$

In a graph  $\Gamma$ , if  $B$  is connected, then  $\text{CO}(\Gamma, B)$  is obtained by (homotopically) collapsing  $\Gamma_B$  onto one vertex and removing any loops or multi-edges. If  $B$  is anticonnected, one has (see [Figure 3](#))

$$\overline{\text{CO}}(\Gamma, B) = \overline{\text{CO}(\overline{\Gamma}, B)}.$$

If  $B \subseteq V(\Gamma)$  and  $\text{link}_\Gamma(q) \supseteq B$ , then we say that  $q$  is a *common neighbor of  $B$* .

The following lemma states that the co-contraction of a set of anticonnected vertices can be obtained by considering a sequence of co-contractions of two nonadjacent vertices. The proof is immediate by considering the complement graphs.

**Lemma 3.2** Let  $\Gamma$  be a graph and  $B \subseteq V(\Gamma)$  be anticonnected. Then there exists a sequence of graphs

$$\Gamma_0 = \Gamma, \Gamma_1, \Gamma_2, \dots, \Gamma_p = \overline{\text{CO}}(\Gamma, B)$$

such that for each  $i = 0, 1, \dots, p-1$ ,  $\Gamma_{i+1}$  is a co-contraction of  $\Gamma_i$  relative to a pair of nonadjacent vertices of  $\Gamma_i$ .  $\square$



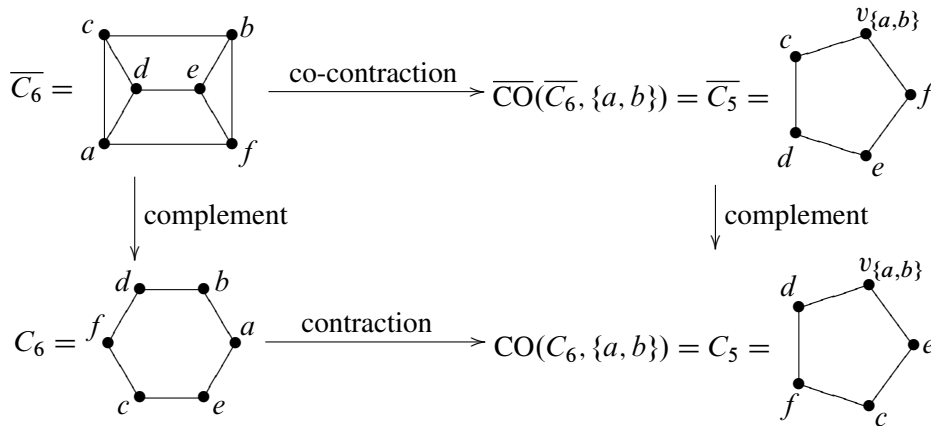


Figure 3: Note that  $\{q : \text{link}_{\overline{C_6}}(q) \supseteq \{a, b\}\} = \{c, f\}$ , ie  $c$  and  $f$  are common neighbors of  $\{a, b\}$ . Hence in  $\overline{\text{CO}}(\overline{C_6}, \{a, b\})$ ,  $v_{\{a,b\}}$  is adjacent to  $c$  and  $f$ . This can be also viewed by looking at the complement graph of  $\overline{C_6}$ , namely  $C_6$ , and collapsing the edge  $\{a, b\}$ .

**Lemma 3.3** (i) If  $B$  is a connected subset of  $p$  vertices of  $C_n$ , then  $\text{CO}(C_n, B) \cong C_{n-p+1}$ .

(ii) If  $B$  is an anticonnected subset of  $p$  vertices of  $\overline{C_n}$ , then  $\overline{\text{CO}}(\overline{C_n}, B) \cong \overline{C_{n-p+1}}$ .

**Proof** (1) is obvious. Considering the complement graphs, (2) follows from (1).  $\square$

### 4 Co-contraction of graphs and right-angled Artin groups

Let  $\Gamma$  be a graph and  $B$  be an anticonnected subset of  $V(\Gamma)$ . Fix a word  $\tilde{w} \in \langle B \rangle$  in  $A(\Gamma)$ . If a vertex  $x$  of  $\overline{\text{CO}}(\Gamma, B)$  is adjacent to  $v_B$ , then  $x$  is a common neighbor of  $B$  in  $\Gamma$ , and so,  $[x, \tilde{w}] =_{A(\Gamma)} 1$ . This implies that there exists a map  $\phi: A(\overline{\text{CO}}(\Gamma, B)) \rightarrow A(\Gamma)$  satisfying:

$$\phi(x) = \begin{cases} \tilde{w} & \text{if } x = v_B \\ x & \text{if } x \in V(\overline{\text{CO}}(\Gamma, B)) \setminus \{v_B\} = V(\Gamma) \setminus B \end{cases}$$

In this section, we show that this map  $\phi$  is injective for a suitable choice of the word  $\tilde{w}$ . First, we prove the injectivity for the case when  $B = \{a, b\}$  and  $\tilde{w} = b^{-1}ab$ .

**Lemma 4.1** Let  $\Gamma$  be a graph. Suppose  $a$  and  $b$  are nonadjacent vertices of  $\Gamma$ . Then there exists an injective map  $\phi: A(\overline{\text{CO}}(\Gamma, \{a, b\})) \rightarrow A(\Gamma)$  satisfying:

$$\phi(x) = \begin{cases} b^{-1}ab & \text{if } x = v_{\{a,b\}} \\ x & \text{if } x \in V(\Gamma) \setminus \{a, b\} \end{cases}$$

**Proof** Let  $\hat{\Gamma} = \overline{\text{CO}}(\Gamma, \{a, b\})$ ,  $\hat{v} = v_{\{a,b\}}$  and  $A = \{q : q \in V(\Gamma) \setminus \{a, b\}\}$ . For  $q \in A$ , let  $\hat{q}$  denote the corresponding vertex in  $\hat{\Gamma}$ , and  $\hat{A} = \{\hat{q} : q \in A\}$ .

Define  $\phi: A(\hat{\Gamma}) \rightarrow A(\Gamma)$  by:

$$\phi(x) = \begin{cases} b^{-1}ab & \text{if } x = \hat{v} \\ q & \text{if } x = \hat{q} \in \hat{A} \end{cases}$$

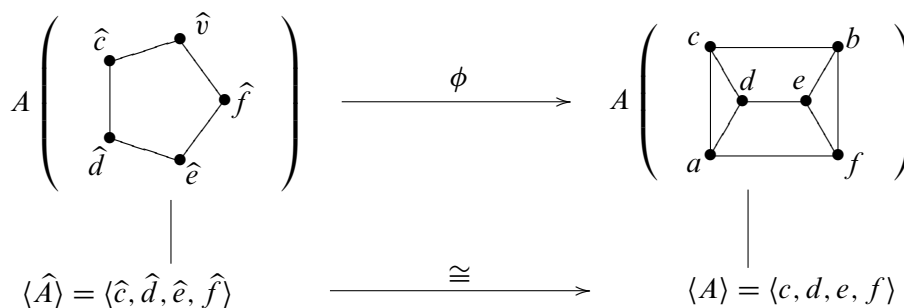
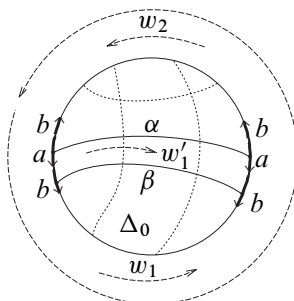


Figure 4:

Suppose  $\phi$  is not injective. Choose a word  $\hat{w}$  of the minimal length in  $\ker \phi \setminus \{1\}$ . Write  $\hat{w} = \prod_{i=1}^k \hat{c}_i^{e_i}$ , where  $\hat{c}_i \in \hat{A} \cup \{\hat{v}\}$  and  $e_i = \pm 1$ . As  $\hat{\Gamma}_{\hat{A}}$  is isomorphic to  $\Gamma_A$ ,  $\phi$  maps  $\langle \hat{A} \rangle$  isomorphically onto  $\langle A \rangle$  (Figure 4). So  $\hat{c}_i = \hat{v}$  for some  $i$ .

Let  $w = \prod_{i=1}^k \phi(\hat{c}_i)^{e_i}$ . Since  $w =_{A(\Gamma)} 1$ , there exists a dual van Kampen diagram  $\Delta = (\mathcal{H}, \lambda)$  for  $w$  in  $A(\Gamma)$ . In  $\Delta$ , choose an innermost  $a$ -arc  $\alpha$ . By considering a cyclic conjugation of  $\hat{w}$  if necessary, one may write  $\hat{w} = \hat{v}^{\pm 1} \cdot \hat{w}_1 \cdot \hat{v}^{\mp 1} \cdot \hat{w}_2$  and  $w = b^{-1}a^{\pm 1}b \cdot w_1 \cdot b^{-1}a^{\mp 1}b \cdot w_2$ , so that  $w_1 = \phi(\hat{w}_1)$ ,  $w_2 = \phi(\hat{w}_2)$  and  $\alpha$  joins the leftmost  $a^{\pm 1}$  of  $w$  and the  $a^{\mp 1}$  between  $w_1$  and  $w_2$  (Figure 5). Then the interval  $w_1$  does not contain any  $a$ -segment. Since each  $b$ -segment in  $w$  is adjacent to some  $a$ -segment, one sees that there does not exist any  $b$ -segment in  $w_1$ , either. Hence,  $w_1 \in \langle A \rangle = A(\Gamma_A)$  and  $\hat{w}_1 \in \langle \hat{A} \rangle = A(\hat{\Gamma}_{\hat{A}})$ . Note that  $\Gamma_A \cong \hat{\Gamma}_{\hat{A}}$ . Since  $\hat{w}_1$  is reduced, so is  $w_1$ .



$$w = b^{-1}a^{\pm 1}b \cdot w_1 \cdot b^{-1}a^{\mp 1}b \cdot w_2$$

Figure 5:  $\Delta$  in the proof of Lemma 4.1

Let  $\beta$  be the  $b$ -arc that meets the letter  $b$ , following  $a^{\pm 1}$  on the left of  $w_1$  in  $w$ .  $\beta$  does not intersect  $\alpha$ , for  $[a, b] \neq 1$ . Since  $w_1$  does not contain any  $b$ -segment,  $\beta$  intersects with the letter  $b^{-1}$  between  $w_1$  and  $w_2$ .

$w_1$  does not contain any canceling pair, for  $w_1$  is reduced. So each segment of  $w_1$  is joined to a segment in  $w_2$  by an arc in  $\mathcal{H}$ . Such an arc must intersect both  $\alpha$  and  $\beta$ . This implies that the segments in  $w_1$  are labeled by vertices in  $\text{link}_{\Gamma}(a) \cap \text{link}_{\Gamma}(b) = \phi(\text{link}_{\hat{\Gamma}}(\hat{v}))$ . It follows that  $\hat{w}_1 \in \langle \text{link}_{\hat{\Gamma}}(\hat{v}) \rangle$ , from the following diagram.

$$\begin{array}{ccccc} \hat{w}_1 \in \langle \hat{A} \rangle & \langle \text{link}_{\hat{\Gamma}}(\hat{v}) \rangle & \leq & \langle \hat{A} \rangle & \\ \downarrow & \downarrow \cong & & \phi \downarrow \cong & \\ w_1 \in \langle \text{link}_{\Gamma}(a) \cap \text{link}_{\Gamma}(b) \rangle & \leq & \langle A \rangle & & \end{array}$$

But then,  $\hat{w} = \hat{v}^{\pm 1} \hat{w}_1 \hat{v}^{\mp 1} \hat{w}_2 =_{A(\hat{\Gamma})} \hat{w}_1 \hat{w}_2$ , which contradicts to the minimality of the length of  $\hat{w}$ . □

**Theorem 4.2** *Let  $\Gamma$  be a graph and  $B$  be an anticonnected subset of  $V(\Gamma)$ . Then  $A(\Gamma)$  contains a subgroup isomorphic to  $A(\overline{\text{CO}(\Gamma, B)})$ .*

**Proof** Proof is immediate from Lemma 3.2 and Lemma 4.1. □

Figure 4 and Lemma 4.1 show the existence of an isomorphism:

$$\phi: A(C_5) \rightarrow \langle b^{-1}ab, c, d, e, f \rangle \leq A(\overline{C_6})$$

More generally, we have the following corollary.

**Corollary 4.3** (1)  $A(\overline{C_n})$  contains a subgroup isomorphic to  $A(\overline{C_{n-p+1}})$  for each  $1 \leq p \leq n$ .

- (2) If  $\Gamma$  contains an induced  $C_n$  or  $\overline{C_n}$  for some  $n \geq 5$ , then  $A(\Gamma)$  contains a hyperbolic surface group.

**Proof** (1) Immediate from [Lemma 3.3](#) and [Theorem 4.2](#).

- (2)  $A(C_n)$  contains a hyperbolic surface group for  $n \geq 5$  [[19](#)]. One has an embedding  $\phi: A(C_5) = A(\overline{C_5}) \hookrightarrow A(\overline{C_n})$ , for  $n \geq 5$ .  $\square$

A simple combinatorial argument shows that for  $n > 5$ , the induced subgraph of  $\overline{C_n}$  on any five vertices contains a triangle. So  $\overline{C_n}$  does not contain an induced  $C_m$  for any  $m \geq 5$ . From the [Corollary 4.3](#) (2), we deduce the negative answer to [Question 1.1](#) as follows.

**Corollary 4.4** *There exists an infinite family  $\mathcal{F}$  of graphs satisfying the following.*

- (i) *Each element in  $\mathcal{F}$  does not contain an induced  $C_n$  for  $n \geq 5$ .*
- (ii) *Each element in  $\mathcal{F}$  is not an induced subgraph of another element in  $\mathcal{F}$ .*
- (iii) *For each  $\Gamma \in \mathcal{F}$ ,  $A(\Gamma)$  contains a hyperbolic surface group.*

**Proof** Set  $\mathcal{F} = \{\overline{C_n} : n > 5\}$ .  $\square$

## 5 Contraction words

In [Lemma 4.1](#), the word  $b^{-1}ab$  was used to construct an injective map from the group  $A(\overline{CO}(\Gamma, \{a, b\}))$  into  $A(\Gamma)$ . This can be generalized by considering a *contraction word*, defined as follows.

- Definition 5.1** (1) Let  $\Gamma_0$  be an anticonnected graph. A sequence  $b_1, b_2, \dots, b_p$  of vertices of  $\Gamma_0$  is a *contraction sequence of  $\Gamma_0$* , if the following holds: for any  $(b, b') \in V(\Gamma_0) \times V(\Gamma_0)$ , there exists  $l \geq 1$  and  $1 \leq k_1 < k_2 < \dots < k_l \leq p$  such that,  $b_{k_1}, b_{k_2}, \dots, b_{k_l}$  is a path from  $b$  to  $b'$  in  $\overline{\Gamma}$ .
- (2) Let  $\Gamma$  be a graph and  $B$  be an anticonnected set of vertices of  $\Gamma$ . A reduced word  $w = \prod_{i=1}^p b_i^{e_i}$  is called a *contraction word of  $B$*  if  $b_i \in B$ ,  $e_i = \pm 1$  for each  $i$ , and  $b_1, b_2, \dots, b_p$  is a contraction sequence of  $\Gamma_B$ . An element of  $A(\Gamma)$  is called a *contraction element*, if it can be represented by a contraction word.

**Remark 5.2** If  $a$  and  $b$  are nonadjacent vertices in  $\Gamma$ , then any word in  $\langle a, b \rangle \setminus \{a^m b^n : m, n \in \mathbb{Z}\}^{\pm 1}$  is a contraction word of  $\{a, b\}$ .

We first note the following general lemma.

**Lemma 5.3** *Let  $\Gamma$  be a graph and  $g \in A(\Gamma)$ . Then  $g =_{A(\Gamma)} u^{-1}vu$  for some words  $u, v$  such that  $u^{-1}v^m u$  is reduced for each  $m \neq 0$ .*

**Proof** Choose words  $u, v$  such that  $u^{-1}vu$  is a reduced word representing  $g$  and the length of  $u$  is maximal. We will show that  $u^{-1}v^m u$  is reduced for any  $m \neq 0$ .

Assume that  $u^{-1}v^m u$  is not reduced for some  $m \neq 0$ . We may assume that  $m > 0$ . Let  $w$  be a reduced word for  $u^{-1}v^m u$ . Draw a dual van Kampen diagram  $\Delta$  for  $u^{-1}v^m u w^{-1}$ . Let  $v_i$  denote the  $v$ -interval on  $\partial\Delta$  corresponding to the  $i$ -th occurrence of  $v$  from the left in  $u^{-1}v^m u$  (Figure 6 (a)).

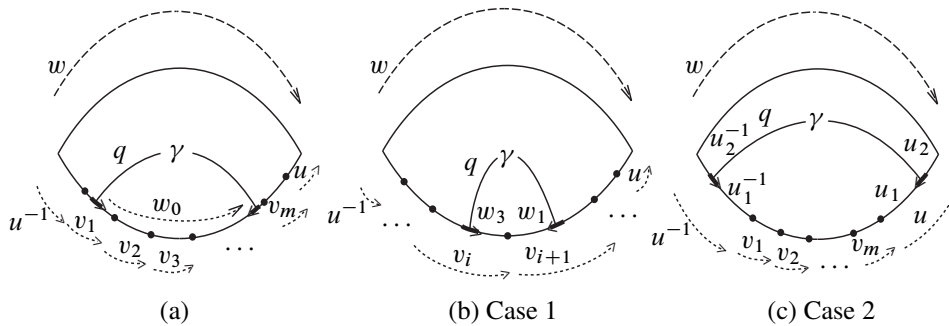


Figure 6: Proof of Lemma 5.3

By Lemma 2.2, there exists a  $q$ -arc  $\gamma$  joining two  $q$ -segments of  $u^{-1}v^m u$  for some  $q \in V(\Gamma)$ . Let  $w_0$  denote the interval between those two  $q$ -segments. We may choose  $q$  and  $\gamma$  so that the number of the segments in  $w_0$  is minimal. Then any arc intersecting with a segment in  $w_0$  must intersect  $\gamma$ . It follows that any letter in  $w_0$  should commute with  $q$ . Moreover,  $w_0$  does not contain any  $q$ -segment.

**Case 1** *The intervals  $u^{-1}$  and  $u$  do not intersect with  $\gamma$ .*

Because  $w_0$  does not contain any  $q$ -segment,  $\gamma$  joins  $v_i$  and  $v_{i+1}$  for some  $i$  (Figure 6 (b)). So one can write  $v = w_1 q^{\pm 1} w_2 q^{\mp 1} w_3$  for some subwords  $w_1, w_2, w_3$  of  $v$  such that and  $w_0 = w_3 w_1$ .  $[w_3, q] =_{A(\Gamma)} 1 =_{A(\Gamma)} [w_1, q]$ . So  $u^{-1}vu =_{A(\Gamma)} u^{-1}q^{\pm 1}w_1w_2w_3q^{\mp 1}u$ , which contradicts to the maximality of  $u$ .

**Case 2**  *$\gamma$  intersects  $u$ - or  $u^{-1}$ -interval.*

Suppose  $u^{-1}$  intersects  $\gamma$ . Since  $u^{-1}v$  is reduced,  $\gamma$  cannot intersect  $v_1$ . So,  $w_0$  contains  $v_1$ . Since  $w_0$  does not contain any  $q$ -segment,  $v$  does not contain the letters

$q$  or  $q^{-1}$  and so,  $\gamma$  cannot intersect any  $v_i$  for  $i = 1, \dots, m$ .  $\gamma$  should intersect with the  $u$ -interval of  $u^{-1}v^m u$  (Figure 6 (c)). This implies that  $\gamma$  intersects with the leftmost  $q$ -segment in the  $u$ -interval of  $u^{-1}v^m u$ . One can write  $u^{-1}v^m u = u_2^{-1}q^{\pm 1}u_1^{-1}v^m u_1 q^{\mp 1}u_2$  such that any letter in  $w_0 = u_1^{-1}v^m u_1$  commutes with  $q$ , ie  $[q, u_1] =_{A(\Gamma)} 1 =_{A(\Gamma)} [q, v]$ . But then  $u^{-1}v u =_{A(\Gamma)} u_2^{-1}u_1^{-1}v u_1 u_2$ , which is a contradiction to the assumption that  $u^{-1}v u$  is reduced.  $\square$

**Lemma 5.4** (1) Any reduced word for a contraction element is a contraction word.

(2) Any nontrivial power of a contraction element is a contraction element.

**Proof** (1) Let  $w = \prod_{i=1}^p b_p^{e_i}$  be a contraction word of an anticonnected set  $B$  in  $V(\Gamma)$ . Here,  $b_i \in B$  and  $e_i = \pm 1$  for each  $i$ . Suppose  $w'$  is a reduced word, such that  $w' =_{A(\Gamma)} w$ . There exists a dual van Kampen diagram  $\Delta$  for  $ww'^{-1}$ . Note that any properly embedded arc of  $\Delta$  meets both of the intervals  $w$  and  $w'$ , since  $w$  and  $w'$  are reduced (Lemma 2.2). Now let  $b, b' \in B$ .  $w$  is a contraction word, so one can find  $l \geq 1$  and  $1 \leq k_1 < k_2 < \dots < k_l \leq p$  such that,  $b_{k_i}$  and  $b_{k_{i+1}}$  are nonadjacent for each  $i = 1, \dots, l-1$ , and  $b = b_{k_1}, b' = b_{k_l}$ . Let  $\gamma_i$  be the arc that intersects with the segment  $b_{k_i}$  of  $w$ . Since  $\gamma_1, \gamma_2, \dots, \gamma_l$  are all disjoint, the boundary points of those arcs on  $w'$  will yield the desired subsequence of the letters of  $w'$ .

(2) Let  $u^{-1}v u$  be a reduced word for  $g$  as in Lemma 5.3. Note that a sequence, containing a contraction sequence as a monotonic subsequence, is again a contraction sequence. So the reduced word  $u^{-1}v^m u$  is a contraction word of  $B$ , for each  $m \neq 0$ .  $\square$

**Definition 5.5** Let  $\Gamma$  be a graph, and  $P$  and  $Q$  be disjoint subsets of  $V(\Gamma)$ . Suppose  $P_1$  is a set of words in  $\langle P \rangle \leq A(\Gamma)$ . A canonical expression for  $g \in \langle P_1, Q \rangle$  with respect to  $\{P_1, Q\}$  is a word  $\prod_{i=1}^k c_i^{e_i}$ , where

- (i)  $c_i \in P_1 \cup Q$
- (ii)  $e_i = 1$  or  $-1$
- (iii)  $\prod_{i=1}^k c_i^{e_i} =_{A(\Gamma)} g$

such that  $k$  is minimal.  $k$  is called the length of the canonical expression.

**Remark 5.6** In the above definition, a canonical expression exists for any element in  $\langle P_1, Q \rangle$ . In the case when  $P_1 \subseteq P$ , a word is a canonical expression with respect to  $\{P_1, Q\}$ , if and only if it is reduced in  $A(\Gamma)$ .

Now we compute intersections of certain subgroups of  $A(\Gamma)$ .

**Lemma 5.7** *Let  $\Gamma$  be a graph,  $P, Q$  be disjoint subsets of  $V(\Gamma)$  and  $P_1$  be a set of words in  $\langle P \rangle \leq A(\Gamma)$ . Let  $R$  be any subset of  $V(\Gamma)$ .*

- (1) *If  $w$  is a canonical expression with respect to  $\{P_1, Q\}$ , then there does not exist a  $q$ -pair of  $w$  for any  $q \in Q$ .*
- (2)  *$\langle P_1, Q \rangle \cap \langle R \rangle \subseteq \langle P_1, Q \cap R \rangle$ . Moreover, the equality holds if  $P \subseteq R$ .*
- (3) *Let  $\tilde{w}$  be a contraction word of  $P$ , and  $P_1 = \{\tilde{w}\}$ . Assume  $P \not\subseteq R$ . Then  $\langle P_1, Q \rangle \cap \langle R \rangle = \langle Q \cap R \rangle$ .*

**Proof** (1) Let  $w$  be a canonical expression, Suppose there exists a  $q$ -pair of  $w$  for some  $q \in Q$ . Then by Lemma 2.2, one can write  $w = w_1 q^{\pm 1} w_2 q^{\mp 1} w_3$  for some subwords  $w_1, w_2$  and  $w_3$  such that  $w_2 \in \langle \text{link}_\Gamma(q) \rangle$ . It follows that  $w =_{A(\Gamma)} w'' = w_1 w_2 w_3$ . Since  $P \cap Q = \emptyset$ ,  $w_1, w_2$  and  $w_3$  are also canonical expressions with respect to  $\{P_1, Q\}$ . This contradicts to the minimality of  $k$ .

(2) Let  $w$  be a canonical expression of an element in  $\langle P_1, Q \rangle \cap \langle R \rangle$ , and  $w' =_{A(\Gamma)} w$  be a reduced word. Consider a dual van Kampen diagram  $\Delta$  for  $ww'^{-1}$ .

Suppose that there exists a  $q$ -segment in  $w$ , for some  $q \in Q$ . Then by (1), the  $q$ -segment should be joined, by a  $q$ -arc, to another  $q$ -segment of  $w'$ . Since  $w'$  is a reduced word representing an element in  $\langle R \rangle$ , each segment of  $w'$  is labeled by  $R^{\pm 1}$  (Lemma 2.3 (2)). Therefore,  $q \in Q \cap R$ .

If  $P \subseteq R$ , then  $\langle P_1, Q \cap R \rangle \subseteq \langle P_1, Q \rangle \cap \langle R \rangle$  is obvious.

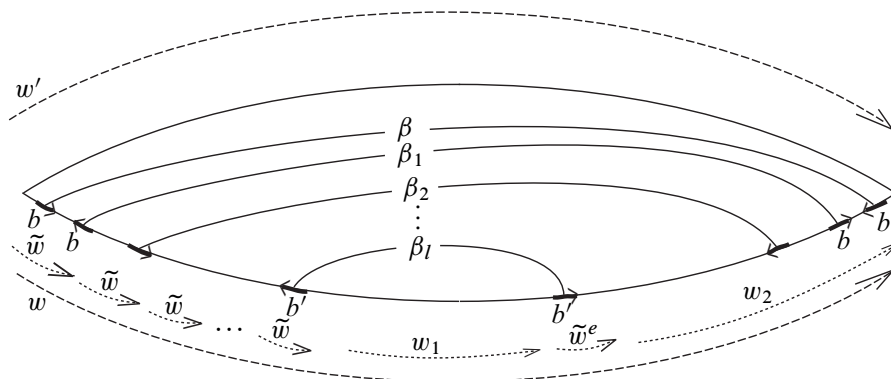
(3)  $\langle Q \cap R \rangle \subseteq \langle P_1, Q \rangle \cap \langle R \rangle$  is obvious.

To prove the converse, suppose  $w \in (\langle P_1, Q \rangle \cap \langle R \rangle) \setminus \langle Q \cap R \rangle$ .  $w$  is chosen so that  $w$  is a canonical expression with respect to  $\{P_1, Q\}$ , and the length (as a canonical expression) is minimal.

Let  $w = \prod_{i=1}^k c_i^{e_i}$  ( $c_i \in \{P_1, Q\}$ ,  $e_i = \pm 1$ ),  $w'$  be a reduced word satisfying  $w' =_{A(\Gamma)} w$ , and  $\Delta = (\mathcal{H}, \lambda)$  be a dual van Kampen diagram for  $ww'^{-1}$  (Figure 7). From the proof of (2),  $c_i \in P_1 \cup (Q \cap R) = \{\tilde{w}\} \cup (Q \cap R)$  for each  $i$ . Also, any shorter canonical expression than  $w$ , for an element in  $\langle P_1, Q \rangle \cap \langle R \rangle$ , is in  $\langle Q \cap R \rangle$ . This implies that  $c_1 = \tilde{w} = c_k$ . Note that each segment of  $w'$  is labeled by  $R^{\pm 1}$ .

Now suppose  $c_i = \tilde{w}$  for some  $i$ . Fix  $b \in P \setminus R$ . Choose the  $b$ -arc  $\beta$  that intersects with the leftmost  $b$ -segment in  $w$  on  $\partial\Delta$ . Note that this  $b$ -segment is contained in the leftmost  $\tilde{w}$ -interval in  $w$ .

Write  $w = \tilde{w}^m w_1 \tilde{w}^e w_2$  for some subwords  $w_1, w_2$  of  $w$ ,  $m \in \mathbb{Z} \setminus \{0\}$  and  $e \in \{1, -1\}$ . Here,  $w_1$  and  $w_2$  are chosen so that the letters of  $w_1$  are in  $(Q \cap R)^{\pm 1}$  and  $\beta$  intersects with a segment in the interval  $\tilde{w}^e w_2$ . Without loss of generality, we may assume  $m > 0$ .

Figure 7:  $\Delta$  in the proof of Lemma 5.7

Let  $b'$  be any element in  $P$ . By Lemma 5.4, any reduced word for  $\tilde{w}^m$  is a contraction word of  $P$ . So, one can find a sequence of arcs  $\beta_1, \beta_2, \dots, \beta_l \in \mathcal{H}$  such that

- (i)  $\lambda(\beta_1) = b, \lambda(\beta_l) = b'$ ,
- (ii)  $\lambda(\beta_i)$  and  $\lambda(\beta_{i+1})$  are nonadjacent in  $\Gamma$ , for each  $i = 1, 2, \dots, l-1$ , and
- (iii) each  $\beta_i$  intersects with a segment in the interval  $\tilde{w}^e w_2$ .

Note that (iii) comes from the assumptions that  $\beta_i$  does not join two segments from  $\tilde{w}^m$  (by reducing  $\tilde{w}^m$  first), and that the letters of  $w_1$  are in  $(Q \cap R)^{\pm 1}$ , which is disjoint from  $P$ .

As in the proof of (2), each segment of  $w_1$  is joined to a segment in  $w'$ . In particular,  $[b', w_1] = [\lambda(\beta_l), w_1] =_{A(\Gamma)} 1$ . Since this is true for any  $b' \in P$ ,  $w =_{A(\Gamma)} w_1 \tilde{w}^{m+e} w_2$ . One has  $\tilde{w}^{m+e} w_2 \in ((P_1, Q) \cap \langle R \rangle) \setminus \langle Q \cap R \rangle$ , for  $w \notin \langle Q \cap R \rangle$  and  $w_1 \in \langle Q \cap R \rangle$ . By the minimality of  $w$ , we have  $w_1 = 1$ . This argument continues, and finally one can write  $w = \tilde{w}^{m'}$  for some  $m' \neq 0$ . In particular, any reduced word for  $w$  is a contraction word of  $P$  (Lemma 5.4). This is impossible since  $w \in \langle R \rangle$  and  $P \not\subseteq R$ .  $\square$

**Lemma 5.8** Let  $\Gamma$  be a graph,  $B$  be an anticonnected set of vertices of  $\Gamma$  and  $g$  be a contraction element of  $B$ . Then there exists an injective map  $\phi: A(\overline{\text{CO}}(\Gamma, B)) \rightarrow A(\Gamma)$  satisfying:

$$\phi(x) = \begin{cases} g & \text{if } x = v_B \\ x & \text{if } x \in V(\Gamma) \setminus B \end{cases}$$

**Proof** As in the proof of Lemma 4.1, let  $\hat{\Gamma} = \overline{\text{CO}}(\Gamma, B)$ ,  $\hat{v} = v_B$  and  $A = \{q : q \in V(\Gamma) \setminus B\}$ . For  $q \in A$ , let  $\hat{q}$  denote the corresponding vertex in  $\hat{\Gamma}$ , and  $\hat{A} = \{\hat{q} : q \in A\}$ .



There exists a map  $\phi: A(\widehat{\Gamma}) \rightarrow A(\Gamma)$  satisfying:

$$\phi(x) = \begin{cases} g & \text{if } x = \widehat{v} \\ q & \text{if } x = \widehat{q} \in \widehat{A} \end{cases}$$

To prove that  $\phi$  is injective, we use an induction on  $|A|$ .

If  $A = \emptyset$ , then  $V(\Gamma) = B$  and  $\widehat{\Gamma}$  is the graph with one vertex  $\widehat{v}$ . So,  $\phi$  maps  $\langle \widehat{v} \rangle = A(\widehat{\Gamma}) \cong \mathbb{Z}$  isomorphically onto  $\mathbb{Z} \cong \langle g \rangle \leq A(\Gamma)$ .

Assume the injectivity of  $\phi$  for the case when  $|A| = k$ , and now let  $|A| = k + 1$ .

Choose any  $t \in A$ . Let  $A_0 = A \setminus \{t\}$  and  $\widehat{A}_0 = \{\widehat{q} : q \in A_0\}$ . Let  $\Gamma_0$  be the induced subgraph on  $A_0 \cup B$  in  $\Gamma$ , and  $\widehat{\Gamma}_0$  be the induced subgraph on  $\widehat{A}_0 \cup \{\widehat{v}\}$  in  $\widehat{\Gamma}$ . We consider  $A(\Gamma_0)$  and  $A(\widehat{\Gamma}_0)$  as subgroups of  $A(\Gamma)$  and  $A(\widehat{\Gamma})$ , respectively, so that  $A(\Gamma_0) = \langle A_0, B \rangle$  and  $A(\widehat{\Gamma}_0) = \langle \widehat{A}_0, \widehat{v} \rangle$ . Let  $K = \langle A_0, g \rangle = \phi(A(\widehat{\Gamma}_0))$  and  $J = \langle A, g \rangle = \phi(A(\widehat{\Gamma}))$ . By the inductive hypothesis,  $\phi$  maps  $A(\widehat{\Gamma}_0)$  isomorphically onto  $K$  (Figure 8).

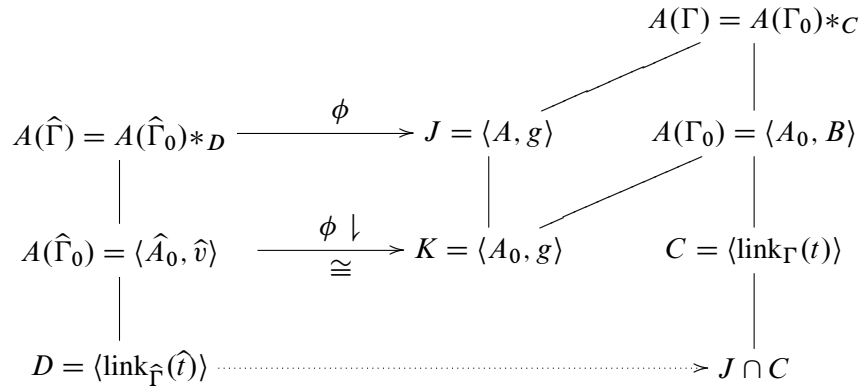


Figure 8: Proof of Lemma 5.8. Note that  $V(\Gamma) = A \sqcup B = A_0 \cup \{t\} \cup B$  and  $V(\widehat{\Gamma}) = \widehat{A} \sqcup \{\widehat{v}\} = \widehat{A}_0 \cup \{\widehat{t}\} \cup \{\widehat{v}\}$ .

From Lemma 2.4, we can identify  $A(\Gamma) = A(\Gamma_0) *_{C}$ , where  $C = \langle \text{link}_{\Gamma}(t) \rangle$  and  $t$  is the stable letter. Since  $J = \langle A_0, g, t \rangle = \langle K, t \rangle$ , Lemma 2.5 implies that we can also identify  $J = (J \cap A(\Gamma_0)) *_{J \cap C}$ , where  $t$  is the stable letter again. Also, we identify  $A(\widehat{\Gamma}) = A(\widehat{\Gamma}_0) *_{D}$ , where  $D = \langle \text{link}_{\widehat{\Gamma}}(\widehat{t}) \rangle$  and  $\widehat{t}$  is the stable letter.

By Lemma 5.7 (2),  $J \cap A(\Gamma_0) = \langle g, A \rangle \cap \langle A_0, B \rangle = \langle g, A \cap (A_0 \cup B) \rangle = \langle g, A_0 \rangle = \phi(A(\widehat{\Gamma}_0))$ .

Applying [Lemma 5.7](#) (2) and (3) for the case when  $R = \text{link}_\Gamma(t)$ , we have:

$$\begin{aligned} J \cap C &= \langle g, A \rangle \cap \langle \text{link}_\Gamma(t) \rangle \\ &= \begin{cases} \langle \text{link}_\Gamma(t) \cap A, g \rangle & \text{if } B \subseteq \text{link}_\Gamma(t) \\ \langle \text{link}_\Gamma(t) \cap A \rangle & \text{otherwise} \end{cases} \end{aligned}$$

From the definition of a co-contraction, we note that:

$$D = \text{link}_{\widehat{\Gamma}}(\widehat{t}) = \begin{cases} \{\widehat{q} : q \in \text{link}_\Gamma(t) \cap A\} \cup \{\widehat{v}\} & \text{if } B \subseteq \text{link}_\Gamma(t) \\ \{\widehat{q} : q \in \text{link}_\Gamma(t) \cap A\} & \text{otherwise} \end{cases}$$

Hence,  $J \cap C = \phi(D)$ . This implies that  $\phi: A(\widehat{\Gamma}) \rightarrow J$  is an isomorphism, as follows.

$$\begin{array}{ccccccc} D & \leq & A(\widehat{\Gamma}_0) & \leq & A(\widehat{\Gamma}_0)*_D & = & A(\widehat{\Gamma}) \\ \downarrow \cong & & \downarrow \cong & & & & \phi \downarrow \\ J \cap C & \leq & K = J \cap A(\Gamma_0) & \leq & (J \cap A(\Gamma_0))*_{J \cap C} & = & J \quad \square \end{array}$$

Now the following theorem is immediate by an induction on  $m$ .

**Theorem 5.9** *Let  $\Gamma$  be a graph and  $B_1, B_2, \dots, B_m$  be disjoint subsets of  $V(\Gamma)$  such that each  $B_i$  is anticonnected. For each  $i$ , let  $v_{B_i}$  denote the vertex corresponding to  $B_i$  in  $\overline{\text{CO}}(\Gamma, (B_1, B_2, \dots, B_m))$ , and  $g_i$  be a contraction element of  $B_i$ . Then there exists an injective map  $\phi: A(\overline{\text{CO}}(\Gamma, (B_1, B_2, \dots, B_m))) \rightarrow A(\Gamma)$  satisfying:*

$$\phi(x) = \begin{cases} g_i & \text{if } x = v_{B_i}, \text{ for some } i \\ x & \text{if } x \in V(\Gamma) \setminus \bigcup_{i=1}^m B_i \end{cases}$$

We conclude this article by noting that there is another partial answer to the question of which right-angled Artin groups contain hyperbolic surface groups. Namely, if  $\Gamma$  does not contain an induced cycle of length  $\geq 5$ , and either  $\Gamma$  does not contain an induced  $C_4$  (hence chordal), or  $\Gamma$  is triangle-free (hence bipartite), then  $A(\Gamma)$  does not contain a hyperbolic surface group [14]. In [2], an independent study by Crisp, Sapir and Sageev proves a similar result, as well as the complete classification of graphs with up to eight vertices, on which the corresponding right-angled Artin groups contain hyperbolic surface subgroups.

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*Department of Mathematics, University of Texas  
1 University Station C1200, Austin, TX 78712-0257, USA*

[shkim@math.utexas.edu](mailto:shkim@math.utexas.edu)

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