Characterizing the Join-Irreducible Medvedev Degrees

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Abstract We characterize the join-irreducible Medvedev degrees as the degrees of complements of Turing ideals, thereby solving a problem posed by Sorbi. We use this characterization to prove that there are Medvedev degrees above the second-least degree that do not bound any join-irreducible degrees above this second-least degree. This solves a problem posed by Sorbi and Terwijn. Finally, we prove that the filter generated by the degrees of closed sets is not prime. This solves a problem posed by Bianchini and Sorbi.

1 Introduction

We present solutions to three problems concerning the Medvedev degrees. A mass problem is a set $A \subseteq \omega^{\omega}$. For mass problems A and B, we say that A Medvedev reduces to \mathcal{B} ($A \leq_M \mathcal{B}$) if there is a Turing functional Φ such that $\Phi(\mathcal{B}) \subseteq A$. That is, $\Phi(f) \in A$ for all $f \in \mathcal{B}$. We say that A and \mathcal{B} are Medvedev equivalent ($A \equiv_M \mathcal{B}$) if $A \leq_M \mathcal{B}$ and $\mathcal{B} \leq_M A$. The equivalence class [A] is called the Medvedev degree of A, and the structure $\mathfrak{M} = (2^{\omega^{\omega}} / \equiv_M, \leq_M)$ is called the Medvedev degrees. See Sorbi [15] for a good introduction to the theory of the Medvedev degrees.

For $f, g \in \omega^{\omega}$, let $f \oplus g$ be the function $(f \oplus g)(2n) = f(n)$ and $(f \oplus g)(2n+1) = g(n)$. For $m \in \omega$ and $f \in \omega^{\omega}$, let $m \cap f$ be the function $(m \cap f)(0) = m$ and $(m \cap f)(n + 1) = f(n)$. In general, ' \cap ' denotes string concatenation. Functions $f \in \omega^{\omega}$ are interpreted as ω -length strings when appropriate. For a mass problem \mathcal{A} , let $m \cap \mathcal{A} = \{m \cap f \mid f \in \mathcal{A}\}$. Given mass problems \mathcal{A} and \mathcal{B} , let $\mathcal{A} + \mathcal{B} = \{f \oplus g \mid f \in \mathcal{A} \land g \in \mathcal{B}\}$ and let $\mathcal{A} \times \mathcal{B} = 0 \cap \mathcal{A} \cup 1 \cap \mathcal{B}$. Then $[\mathcal{A}] + [\mathcal{B}] = [\mathcal{A} + \mathcal{B}]$ is the *join* (i.e., \leq_{M} -least upper bound) of $[\mathcal{A}]$ and $[\mathcal{B}]$, while $[\mathcal{A}] \times [\mathcal{B}] = [\mathcal{A} \times \mathcal{B}]$ is the *meet* (i.e., \leq_{M} -greatest lower bound) of $[\mathcal{A}]$ and $[\mathcal{B}]$. Hence \mathfrak{M} is a lattice. In fact, \mathfrak{M} is a distributive lattice, meaning that join and meet distribute over each other: $\mathbf{a} + (\mathbf{b} \times \mathbf{c}) = (\mathbf{a} + \mathbf{b}) \times (\mathbf{a} + \mathbf{c})$ and

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 $\mathbf{a} \times (\mathbf{b} + \mathbf{c}) = (\mathbf{a} \times \mathbf{b}) + (\mathbf{a} \times \mathbf{c})$. Notation for join and meet appears in the literature variously as +, \times , as \vee , \wedge , and confusingly as \wedge , \vee . We choose the +, \times notation to avoid conflict with the logical notation and to match Sorbi and Terwijn [16].

 \mathfrak{M} has a least element $\mathbf{0} = [\omega^{\omega}]$ (and any \mathcal{A} containing a recursive function has this degree), a second-least element $\mathbf{0}' = [\{f \mid f >_{\mathrm{T}} 0\}]$, and a greatest element $\mathbf{1} = [\varnothing]$. (The Medvedev degree $\mathbf{0}'$ has little to do with 0', the Turing jump of the 0 function. Here $\mathbf{0}'$ always refers to the second-least Medvedev degree.)

In any lattice, an element **a** is called *join-reducible* if there are **x**, **y** < **a** such that $\mathbf{a} = \mathbf{x} + \mathbf{y}$. Otherwise, **a** is called *join-irreducible*. Dually, **a** is called *meet-reducible* if there are **x**, **y** > **a** such that $\mathbf{a} = \mathbf{x} \times \mathbf{y}$. Otherwise, **a** is called *meet-irreducible*. Dyment [3] characterized the meet-reducible Medvedev degrees in the following theorem. Its corollary helps identify meet-irreducible Medvedev degrees.

Theorem 1.1 ([3]) A Medvedev degree **a** is meet-reducible if and only if $\mathbf{a} = [\mathcal{A}]$ for a mass problem \mathcal{A} for which there are r.e. sets $V_0, V_1 \subseteq \omega^{<\omega}$ such that

(i) $(\forall f \in \mathcal{A})(\exists \sigma \in V_0 \cup V_1)(\sigma \subset f),$

(ii) the following mass problems are \leq_M -incomparable:

$$\{f \in \mathcal{A} \mid (\exists \sigma \in V_0)(\sigma \subset f)\}$$
 and $\{f \in \mathcal{A} \mid (\exists \sigma \in V_1)(\sigma \subset f)\}.$

Corollary 1.2 ([3]) If A is a mass problem such that $\sigma \cap A \subseteq A$ for all $\sigma \in \omega^{<\omega}$, then [A] is meet-irreducible.

In particular, **0'** is meet-irreducible because $\sigma \cap f >_T 0$ whenever $\sigma \in \omega^{<\omega}$ and $f >_T 0$.

The problem of characterizing the join-irreducible Medvedev degrees was posed in [15]. In Section 2, we prove that $\mathbf{a} \in \mathfrak{M}$ is join-irreducible if and only if $\mathbf{a} = [\omega^{\omega} - \mathfrak{L}]$ for some Turing ideal \mathfrak{L} .

We have seen that \mathfrak{M} is a distributive lattice with **0** and **1**. In fact, \mathfrak{M} is a Brouwer algebra. A *Brouwer algebra* is a distributive lattice with **0** and **1** such that for every **a** and **b** there is a least **c** such that $\mathbf{a} + \mathbf{c} \ge \mathbf{b}$. This least **c** is denoted by $\mathbf{a} \to \mathbf{b}$. For mass problems \mathcal{A} and \mathcal{B} , define $\mathcal{A} \to \mathcal{B} = \{e^{\frown}g \mid (\forall f \in \mathcal{A})(\Phi_e(f \oplus g) \in \mathcal{B})\}$. Then $[\mathcal{A}] \to [\mathcal{B}] = [\mathcal{A} \to \mathcal{B}]$. A Brouwer algebra is dual to a Heyting algebra, but \mathfrak{M} is proved not to be a Heyting algebra in Sorbi [12].

Brouwer algebras give semantics for propositional logic. For any Brouwer algebra \mathfrak{B} , a *valuation* is a function ν : propositional variables $\rightarrow \mathfrak{B}$. A valuation ν extends to all propositional formulas φ by defining

$$v(\varphi \land \psi) = v(\varphi) + v(\psi),$$

$$v(\varphi \lor \psi) = v(\varphi) \times v(\psi),$$

$$v(\varphi \to \psi) = v(\varphi) \to v(\psi), \text{ and }$$

$$v(\neg \varphi) = v(\varphi) \to 1.$$

A propositional formula φ is called *valid* in \mathfrak{B} if $v(\varphi) = \mathbf{0}$ for every valuation v. Let Th(\mathfrak{B}) denote the set of propositional formulas valid in \mathfrak{B} . The axioms of intuitionistic logic are valid in every Brouwer algebra \mathfrak{B} , so IPC \subseteq Th(\mathfrak{B}) \subseteq CPC for every Brouwer algebra \mathfrak{B} . Here IPC denotes intuitionistic logic and CPC denotes classical logic. Logics *L* for which IPC $\subseteq L \subseteq$ CPC are called *intermediate logics*.

Providing semantics for propositional logic was one of Medvedev's main motivations behind introducing \mathfrak{M} , and he proved Th(\mathfrak{M}) = JAN in Medvedev [8]. JAN denotes the logic IPC $+\neg p \lor \neg \neg p$ named after Jankov who studied it in Jankov [5]. In any Brouwer algebra \mathfrak{B} , the quotient of \mathfrak{B} by the principal filter generated by $\mathbf{a} \in \mathfrak{B}$ is denoted by \mathfrak{B}/\mathbf{a} . The quotient \mathfrak{B}/\mathbf{a} is isomorphic to the interval $[\mathbf{0}, \mathbf{a}]$ which is a Brouwer algebra under the operations inherited from \mathfrak{B} . Logics of the form Th(\mathfrak{M}/\mathbf{a}) have been studied in Skvortsova [10], Sorbi [14], and Sorbi and Terwijn [16]. (Skvortsova and Dyment are the same person. Dyment married and became Skvortsova.) The results in Section 3 and Section 4 are motivated by the following question which remains open.

Question 1.3 ([16]) Is $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \operatorname{JAN}$ for all $\mathbf{a} >_{\mathrm{M}} \mathbf{0}'$?

Sorbi and Terwijn's study of Question 1.3 in [16] lead them to ask whether every degree $>_M 0'$ bounds a join-irreducible degree $>_M 0'$ because a "yes" answer to this question implies a "yes" answer to Question 1.3. However, Sorbi and Terwijn conjectured that there is a degree $>_M 0'$ that bounds no join-irreducible degree $>_M 0'$, and we prove that this is correct in Section 3. In Section 4 we provide slight extensions to some of the results in [14], thereby widening the class of degrees **a** for which Th(\mathfrak{M}/\mathbf{a}) \subseteq JAN is known.

Lastly, in Section 5 we use techniques similar to those used to characterize the join-irreducible degrees to prove that the filter generated by the degrees of mass problems closed in ω^{ω} is not prime. This problem was posed in Bianchini and Sorbi [2] and in Sorbi [15].

2 Characterizing the Join-Irreducible Medvedev Degrees

A *Turing ideal* is a set $\mathcal{I} \subseteq \omega^{\omega}$ that is closed downward under \leq_{T} (i.e., $f \in \mathcal{I} \land g \leq_{\mathrm{T}} f \rightarrow g \in \mathcal{I}$) and closed under \oplus (i.e., $f, g \in \mathcal{I} \rightarrow f \oplus g \in \mathcal{I}$). We prove that $\mathbf{a} \in \mathfrak{M}$ is join-irreducible if and only if $\mathbf{a} = [\omega^{\omega} - \mathcal{I}]$ for some Turing ideal \mathcal{I} . We frequently use the following well-known lemma without mention.

Lemma 2.1 (see [1] Section III.2) In a distributive lattice, **a** is join-irreducible if and only if for all **x** and **y**, $\mathbf{a} \le \mathbf{x} + \mathbf{y}$ implies $\mathbf{a} \le \mathbf{x}$ or $\mathbf{a} \le \mathbf{y}$. Dually, **a** is meet-irreducible if and only if for all **x** and **y**, $\mathbf{a} \ge \mathbf{x} \times \mathbf{y}$ implies $\mathbf{a} \ge \mathbf{x}$ or $\mathbf{a} \ge \mathbf{y}$.

Proof Suppose **a** is join-irreducible and $\mathbf{a} \leq \mathbf{x} + \mathbf{y}$. Then

$$\mathbf{a} = \mathbf{a} \times (\mathbf{x} + \mathbf{y}) = (\mathbf{a} \times \mathbf{x}) + (\mathbf{a} \times \mathbf{y}).$$

Thus $\mathbf{a} = \mathbf{a} \times \mathbf{x}$ or $\mathbf{a} = \mathbf{a} \times \mathbf{y}$ which means $\mathbf{a} \le \mathbf{x}$ or $\mathbf{a} \le \mathbf{y}$. Conversely, if \mathbf{a} is join-reducible, then by definition there are $\mathbf{x}, \mathbf{y} < \mathbf{a}$ with $\mathbf{a} = \mathbf{x} + \mathbf{y}$. The proof for the meet-irreducible case is obtained by dualizing the proof for the join-irreducible case.

For a mass problem \mathcal{A} , let $C(\mathcal{A})$ denote the *Turing upward-closure* of \mathcal{A} : $C(\mathcal{A}) = \{f \mid (\exists g \in \mathcal{A})(f \geq_T g)\}$. A mass problem \mathcal{A} is called *Turing upward-closed* if $\mathcal{A} = C(\mathcal{A})$. The identity functional witnesses $C(\mathcal{A}) \leq_M \mathcal{A}$ for any mass problem \mathcal{A} , and if \mathcal{A} and \mathcal{B} are mass problems such that \mathcal{A} is Turing upward-closed, then $\mathcal{A} \leq_M \mathcal{B}$ if and only if $\mathcal{B} \subseteq \mathcal{A}$. Our starting point is the following observation.

Lemma 2.2 ([15]) If \mathcal{A} is a mass problem such that $[\mathcal{A}]$ is join-irreducible, then $\omega^{\omega} - C(\mathcal{A})$ is a Turing ideal.

Proof We prove the contrapositive. If $\omega^{\omega} - C(\mathcal{A})$ is not a Turing ideal, then there are $f, g \notin C(\mathcal{A})$ with $f \oplus g \in C(\mathcal{A})$. This means that $\{f\}, \{g\} \not\geq_{M} \mathcal{A}$ but $\{f\} + \{g\} \geq_{M} \mathcal{A}$. Thus $[\mathcal{A}]$ is join-reducible.

The next lemma is the main step in our characterization.

Lemma 2.3 If A is a mass problem such that [A] is join-irreducible, then $A \equiv_M C(A)$

Proof We prove the contrapositive. Suppose $A \not\equiv_M C(A)$. Then it must be that $A \not\leq_M C(A)$. We find mass problems X and Y such that $X, Y \not\geq_M A$ but $X + Y \geq_M A$. Thus [A] is join-reducible.

To find \mathcal{X} and \mathcal{Y} , first find a sequence $(h_n \mid n \in \omega)$ of functions and a sequence $(e_n \mid n \in \omega)$ of indices such that

- (i) $\Phi_{e_n}(h_n) \in \mathcal{A}$ for all $n \in \omega$,
- (ii) $\Phi_n(h_{2n}) \notin A$ and $\Phi_n(h_{2n+1}) \notin A$ for all $n \in \omega$, and
- (iii) $h_n(0) = \langle n, e_0, e_1, \dots, e_{n-1} \rangle$ for all $n \in \omega$.

We find the desired sequences by iterating the following claim.

Claim 2.4 If $A \not\leq_M C(A)$, then for every $e, m \in \omega$ there is an $h \in C(A)$ such that h(0) = m and $\Phi_e(h) \notin A$.

Proof of claim Suppose not. Then there are $e, m \in \omega$ such that h(0) = m implies $\Phi_e(h) \in \mathcal{A}$ for all $h \in C(\mathcal{A})$. Thus $h \mapsto \Phi_e(m \cap h)$ is a reduction witnessing $\mathcal{A} \leq_M C(\mathcal{A})$, a contradiction.

Suppose we have h_i and e_i for all i < n. To find h_n and e_n , let $e = \lfloor n/2 \rfloor$ and let $m = \langle n, e_0, e_1, \ldots, e_{n-1} \rangle$. By the claim, there is an $h_n \in C(\mathcal{A})$ such that $h_n(0) = m$ and $\Phi_e(h_n) \notin \mathcal{A}$. The fact that $h_n \in C(\mathcal{A})$ means that there is an e_n such that $\Phi_{e_n}(h_n) \in \mathcal{A}$.

Now set $\mathcal{X} = \{h_{2n} \mid n \in \omega\}$ and $\mathcal{Y} = \{h_{2n+1} \mid n \in \omega\}$. Then $\Phi_e(\mathcal{X}) \nsubseteq \mathcal{A}$ and $\Phi_e(\mathcal{Y}) \nsubseteq \mathcal{A}$ for each *e* by item (ii). Hence $\mathcal{X}, \mathcal{Y} \ngeq_M \mathcal{A}$. The following reduction witnesses $\mathcal{X} + \mathcal{Y} \ge_M \mathcal{A}$.

Given *h*, decompose *h* as $h = f \oplus g$ and decode f(0) and g(0) as $f(0) = \langle 2n, x_0, x_1, \ldots, x_{2n-1} \rangle$ and $g(0) = \langle 2m + 1, y_0, y_1, \ldots, y_{2m} \rangle$. If either f(0) or g(0) is not of the required form, then output the 0 function (as such an *h* cannot be in $\mathcal{X} + \mathcal{Y}$). Otherwise, output $\Phi_{x_{2m+1}}(g)$ if 2n > 2m + 1 and output $\Phi_{y_{2n}}(f)$ if 2m + 1 > 2n.

Suppose this reduction is applied to some $h = h_{2n} \oplus h_{2m+1} \in \mathcal{X} + \mathcal{Y}$. In this case, $f = h_{2n}, g = h_{2m+1}$, and f(0) and g(0) are of the required form by item (iii). So if 2n > 2m + 1 we output $\Phi_{e_{2m+1}}(h_{2m+1})$ and if 2m + 1 > 2n we output $\Phi_{e_{2n}}(h_{2n})$. Both alternatives are in \mathcal{A} by item (i). Thus $\mathcal{X} + \mathcal{Y} \ge_M \mathcal{A}$.

Theorem 2.5 A Medvedev degree **a** is join-irreducible if and only if $\mathbf{a} = [\omega^{\omega} - \mathbf{l}]$ for some Turing ideal \mathbf{l} .

Proof Suppose **a** is join-irreducible, and let \mathcal{A} be a mass problem such that $\mathbf{a} = [\mathcal{A}]$. Then $\mathcal{I} = \omega^{\omega} - C(\mathcal{A})$ is a Turing ideal by Lemma 2.2, $\mathcal{A} \equiv_{\mathrm{M}} C(\mathcal{A})$ by Lemma 2.3, and therefore $\mathcal{A} \equiv_{\mathrm{M}} C(\mathcal{A}) = \omega^{\omega} - \mathcal{I}$. Hence $\mathbf{a} = [\omega^{\omega} - \mathcal{I}]$ for the Turing ideal \mathcal{I} .

Conversely, suppose \mathcal{I} is a Turing ideal and let \mathcal{X} and \mathcal{Y} be mass problems such that $\mathcal{X}, \mathcal{Y} \not\geq_{\mathrm{M}} \omega^{\omega} - \mathcal{I}$. We show that $\mathcal{X} + \mathcal{Y} \not\geq_{\mathrm{M}} \omega^{\omega} - \mathcal{I}$. Observe $\mathcal{X}, \mathcal{Y} \not\subseteq \omega^{\omega} - \mathcal{I}$ for otherwise the identity functional would witness $\mathcal{X}, \mathcal{Y} \geq_{\mathrm{M}} \omega^{\omega} - \mathcal{I}$. Let $f \in \mathcal{X} \cap \mathcal{I}$ and let $g \in \mathcal{Y} \cap \mathcal{I}$, thereby making $f \oplus g \in (\mathcal{X} + \mathcal{Y}) \cap \mathcal{I}$. The function $f \oplus g$ is in $\mathcal{X} + \mathcal{Y}$, but it does not compute any member of $\omega^{\omega} - \mathcal{I}$. Therefore, $\mathcal{X} + \mathcal{Y} \not\geq_{\mathrm{M}} \omega^{\omega} - \mathcal{I}$. Hence, $[\omega^{\omega} - \mathcal{I}]$ is join-irreducible.

Theorem 2.5 is also valid for the *Muchnik degrees* \mathfrak{M}_{w} in place of \mathfrak{M} , a fact first noticed by Terwijn [17]. \mathfrak{M}_{w} is defined just as \mathfrak{M} , but with *Muchnik reducibility* (also called *weak reducibility*) \leq_{w} in place of \leq_{M} : $A \leq_{w} \mathcal{B}$ if for every $f \in \mathcal{B}$ there is a $g \in \mathcal{A}$ such that $f \geq_{T} g$. \mathfrak{M}_{w} is a Brouwer algebra with +, ×, and \rightarrow defined by $[\mathcal{A}]_{w} + [\mathcal{B}]_{w} = [\mathcal{A} + \mathcal{B}]_{w}, [\mathcal{A}]_{w} \times [\mathcal{B}]_{w} = [\mathcal{A} \times \mathcal{B}]_{w}$, and $[\mathcal{A}]_{w} \rightarrow [\mathcal{B}]_{w} = [\{g \mid (\forall f \in \mathcal{A})(\exists h \in \mathcal{B})(h \leq_{T} f \oplus g)\}]_{w}$. The proof of Lemma 2.2 also works for \mathfrak{M}_{w} , and it is easy to check that $\mathcal{A} \equiv_{w} C(\mathcal{A})$ for any mass problem \mathcal{A} (i.e., the \mathfrak{M}_{w} analogue of Lemma 2.3 is trivial). This gives the forward direction of Theorem 2.5 for \mathfrak{M}_{w} .

3 Degrees That Bound No Join-Irreducible Degrees $>_M 0'$

Recall that JAN is the intermediate logic IPC $+\neg p \lor \neg \neg p$. The results of this section and the next are motivated by Question 1.3: Is Th $(\mathfrak{M}/\mathbf{a}) \subseteq$ JAN for every $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$?

Th($\mathfrak{M} / \mathbf{0}'$) = CPC because $\mathfrak{M} / \mathbf{0}' \cong [\mathbf{0}, \mathbf{0}'] = \{\mathbf{0}, \mathbf{0}'\}$. In fact, $\mathbf{0}'$ is the only degree for which Th($\mathfrak{M} / \mathbf{0}'$) = CPC. This is because if $\mathbf{a} >_{\mathrm{M}} \mathbf{0}'$, then $\mathbf{0}' \to \mathbf{a} = \mathbf{a}$, hence $\mathbf{0}' \times (\mathbf{0}' \to \mathbf{a}) = \mathbf{0}'$. Thus, let $p = \mathbf{0}'$ to see that the formula $p \vee \neg p$ is not valid in Th($\mathfrak{M} / \mathbf{a}$).

Furthermore, if $\mathbf{a} >_M \mathbf{0}'$, then we cannot have $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) \supseteq \operatorname{JAN}$. It is an easy check that in any Brouwer algebra \mathfrak{B} with meet-irreducible $\mathbf{0}$ (such as the algebras \mathfrak{M}/\mathbf{a}), $\neg p \lor \neg \neg p \in \operatorname{Th}(\mathfrak{B})$ if and only if $\mathbf{1}$ is join-irreducible. However, if $\mathbf{a} >_M \mathbf{0}'$ is join-irreducible, then $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) = \operatorname{JAN}[14]$. Thus, if $\mathbf{a} >_M \mathbf{0}'$ and $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) \supseteq \operatorname{JAN}$, then $\neg p \lor \neg \neg p \in \operatorname{Th}(\mathfrak{M}/\mathbf{a})$ which implies that \mathbf{a} is join-irreducible which implies that $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) = \operatorname{JAN}$. Thus a "no" answer to Question 1.3 must yield a degree \mathbf{a} such that $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) \not\subseteq \operatorname{JAN}$ and $\operatorname{JAN} \not\subseteq \operatorname{Th}(\mathfrak{M}/\mathbf{a})$.

The following theorem shows that to give a "yes" answer to Question 1.3 it suffices to show that every $\mathbf{a} >_M \mathbf{0}'$ bounds a finite meet of join-irreducible degrees $>_M \mathbf{0}'$.

Theorem 3.1 ([14]) If **a** is a degree such that $\mathbf{a} \ge_{\mathbf{M}} \prod_{i=0}^{n} \mathbf{d}_{i}$ for join-irreducible degrees $\mathbf{d}_{i} >_{\mathbf{M}} \mathbf{0}'$, $i \le n$, then $\operatorname{Th}(\mathfrak{M} / \mathbf{a}) \subseteq \operatorname{JAN}$.

(The above theorem is stated more generally in [14]. Each degree \mathbf{d}_i for $i \leq n$ is allowed to be either join-irreducible or $\mathfrak{D}e$ -irreducible. See the parenthetical discussion following Theorem 4.1 for the definition of $\mathfrak{D}e$ -irreducible and an explanation of why we do not consider such degrees here. Theorem 4.1 is a restatement of [14], Theorem 2.11, which is the main tool used to prove Theorem 3.1.)

The degrees of the mass problems $\mathcal{B}_f = \{g \mid g \not\leq_T f\}$ play an important role in the study of Question 1.3. The following lemma from Sorbi [13] encapsulates the properties of the $[\mathcal{B}_f]$ s that we need in this section and the next.

Lemma 3.2 ([13])

- (i) Every $[\mathcal{B}_f]$ is join-irreducible.
- (ii) Every $\sum_{i=1}^{n} [\mathcal{B}_{f_i}]$ is meet-irreducible.
- (iii) Let V and J be finite sets and let U_v and I_j be finite sets for each $v \in V$ and $j \in J$. Let \mathbf{x}_u^v and \mathbf{y}_i^j be degrees of the form $[\mathcal{B}_f]$ for every $v \in V$, $u \in U_v$, $j \in J$, and $i \in I_j$. Let $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v$ and $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} \mathbf{y}_i^j$. Then $\mathbf{a} \leq_M \mathbf{b}$ if and only if

$$(\forall v \in V) (\exists j \in J) (\forall i \in I_j) (\exists u \in U_v) (\mathbf{x}_u^v \leq_{\mathbf{M}} \mathbf{y}_i^j).$$

(iv) In the notation of item (iii),

$$\mathbf{a} \to \mathbf{b} = \sum \left\{ \prod_{i \in I_j} \mathbf{y}_i^j \mid \left(\forall v \in V \right) \left(\prod_{i \in I_j} \mathbf{y}_i^j \not\leq_{\mathsf{M}} \prod_{u \in U_v} \mathbf{x}_u^v \right) \right\}$$

(where the empty join is **0**).

Proof Item (i) is by Theorem 2.5 and item (ii) is by Corollary 1.2. Item (iv) is proved in [13]. Item (iii) follows from item (iv) because $\mathbf{a} \leq_M \mathbf{b}$ if and only if $\mathbf{b} \rightarrow \mathbf{a} = \mathbf{0}$.

In [16] it is asked if every degree $\mathbf{a} >_M \mathbf{0}'$ bounds a join-irreducible degree $>_M \mathbf{0}'$, and it is conjectured that this is not the case based on the evidence provided by the following theorem.

Theorem 3.3 ([16]) There is a degree $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$ such that $\mathbf{a} \not\geq_{\mathbf{M}} [\mathcal{B}_f]$ for every $f >_{\mathbf{T}} \mathbf{0}$.

Our characterization of the join-irreducible degrees implies that every join-irreducible degree $>_{\mathbf{M}} \mathbf{0}'$ bounds some degree $[\mathcal{B}_f]$ with $f >_{\mathbf{T}} 0$. Thus the conjecture is correct.

Corollary 3.4 (to Theorem 2.5) If $\mathbf{a} >_M \mathbf{0}'$ is join-irreducible, then $\mathbf{a} \ge_M [\mathcal{B}_f]$ for some $f >_T 0$.

Proof If **a** is join-irreducible, then, by Theorem 2.5, $\mathbf{a} = [\omega^{\omega} - \mathbf{l}]$ for some Turing ideal \mathbf{l} . If $[\omega^{\omega} - \mathbf{l}] >_{\mathbf{M}} \mathbf{0}'$, then \mathbf{l} contains some function $f >_{\mathbf{T}} \mathbf{0}$. Thus $\omega^{\omega} - \mathbf{l} \subseteq \mathcal{B}_f$. Hence $\mathbf{a} = [\omega^{\omega} - \mathbf{l}] \ge_{\mathbf{M}} [\mathcal{B}_f]$.

Theorem 3.5 There is a degree $\mathbf{a} >_M \mathbf{0}'$ such that every degree \mathbf{x} with $\mathbf{0}' <_M \mathbf{x} \leq_M \mathbf{a}$ is join-reducible.

Proof By Theorem 3.3, let $\mathbf{a} >_M \mathbf{0}'$ be such that $\mathbf{a} \ngeq_M [\mathcal{B}_f]$ for every $f >_T 0$. This **a** is the desired degree because, by Corollary 3.4, if $\mathbf{a} \ge_M \mathbf{x}$ for some join-irreducible $\mathbf{x} >_M \mathbf{0}'$, then $\mathbf{a} \ge_M [\mathcal{B}_f]$ for some $f >_T 0$.

The degree **a** satisfying Theorem 3.3 was constructed by diagonalization in [16]. We can give somewhat more concrete examples of degrees satisfying Theorem 3.3 and Theorem 3.5. Recall the following definitions. Functions $f, g >_T 0$ are a *Turing minimal pair* if, for all $h, h \leq_T f, g$ implies $h \leq_T 0$. A function f has *minimal Turing degree* if, for all $h, h <_T f$ implies $h \leq_T 0$. Minimal pairs and minimal degrees exist. In fact, there are continuum many distinct minimal Turing degrees. See Lerman [6], Section II.4 and Section V.2.

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Theorem 3.6 If f and g are a minimal pair, then the degree $\mathbf{a} = [\mathcal{B}_f] \times [\mathcal{B}_g]$ witnesses Theorem 3.5.

Proof Let f and g be a minimal pair. Then $[\mathcal{B}_f], [\mathcal{B}_g] >_M \mathbf{0}'$ because $f, g >_T 0$. Thus $[\mathcal{B}_f] \times [\mathcal{B}_g] >_M \mathbf{0}'$ because $\mathbf{0}'$ is meet-irreducible by Corollary 1.2. To show that $[\mathcal{B}_f] \times [\mathcal{B}_g]$ bounds no join-irreducible degree $>_M \mathbf{0}'$, it suffices by Corollary 3.4 to show that $[\mathcal{B}_f] \times [\mathcal{B}_g]$ bounds no $[\mathcal{B}_h]$ for $h >_T 0$. This is true because f, g is a minimal pair, so for any $h >_T 0$, either $h \not\leq_T f$ or $h \not\leq g$. Thus, either $h \in \mathcal{B}_f$ or $h \in \mathcal{B}_g$, which means $\mathcal{B}_f \times \mathcal{B}_g$ contains a function $\equiv_T h$. \mathcal{B}_h contains no function $\leq_T h$; therefore, $\mathcal{B}_f \times \mathcal{B}_g \not\geq_M \mathcal{B}_h$.

We can extend the idea behind Theorem 3.6 to find a degree $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$ that does not bound any finite meet of join-irreducible degrees $>_{\mathbf{M}} \mathbf{0}'$. Several of our examples in this section and the next are of the form $[\bigcup_{i \in \omega} i \cap \mathcal{D}_i]$ for mass problems $\mathcal{D}_i, i \in \omega$.

Lemma 3.7 Let $\mathbf{d} = \left[\bigcup_{i \in \omega} i \cap \mathcal{D}_i\right]$ where $[\mathcal{D}_i] >_M \mathbf{0}'$ for each $i \in \omega$. Then $\mathbf{d} >_M \mathbf{0}'$.

Proof Suppose for a contradiction that Φ is a reduction witnessing $\mathbf{d} \leq_{\mathbf{M}} \mathbf{0}'$ (i.e., $\Phi(f) \in \bigcup_{i \in \omega} i \cap \mathcal{D}_i$ for all $f >_{\mathbf{T}} 0$). Let $\sigma \in \omega^{<\omega}$ be such that $\Phi(\sigma)(0) \downarrow$ and let $i = \Phi(\sigma)(0)$. Then $f \mapsto \Phi(\sigma \cap f)$ is a reduction witnessing $\mathbf{0}' \geq_{\mathbf{M}} [\mathcal{D}_i]$, contradicting $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$.

Theorem 3.8 There is a degree $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$ such that no degree \mathbf{x} with $\mathbf{0}' <_{\mathbf{M}} \mathbf{x} \leq_{\mathbf{M}} \mathbf{a}$ is of the form $\prod_{i=0}^{n} \mathbf{d}_{i}$ for join-irreducible degrees $\mathbf{d}_{i} >_{\mathbf{M}} \mathbf{0}'$, $i \leq n$.

Proof By Corollary 3.4, it suffices to find a degree $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$ which is not above any degree of the form $\prod_{i=0}^{n} [\mathcal{B}_{f_i}]$ where $f_i >_{\mathbf{T}} 0$ for each $i \leq n$. Let $\{g_i \mid i \in \omega\}$ be a countable collection of functions all of distinct minimal Turing degree. Let $\mathcal{A} = \bigcup_{i \in \omega} i \cap \mathcal{B}_{g_i}$ and put $\mathbf{a} = [\mathcal{A}]$. Lemma 3.7 proves that $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$.

Now consider any degree $\prod_{i=0}^{n} [\mathcal{B}_{f_i}]$, where $f_i >_T 0$ for each $i \leq n$. There is a $j \in \omega$ such that $g_j \not\geq_T f_i$ for each $i \leq n$. Thus, for this $j, [\mathcal{B}_{g_j}] \not\geq_M [\mathcal{B}_{f_i}]$ for each $i \leq n$. Therefore, $[\mathcal{B}_{g_j}] \not\geq_M \prod_{i=0}^{n} [\mathcal{B}_{f_i}]$ because $[\mathcal{B}_{g_j}]$ is meet-irreducible. Clearly, $[\mathcal{B}_{g_j}] \geq_M \mathbf{a}$, so $\mathbf{a} \not\geq_M \prod_{i=0}^{n} [\mathcal{B}_{f_i}]$ as well. \Box

For mass problems A_i , $i \in \omega$, the Medvedev degree $\left[\bigcup_{i \in \omega} i \cap A_i\right]$ is not in general the greatest lower bound of the degrees $[A_i]$, $i \in \omega$. Such greatest lower bounds need not even exist. For example, the degrees $[\mathcal{B}_{g_i}]$, $i \in \omega$ from Theorem 3.8 do not have a greatest lower bound. This follows from results in Dyment [4] which studies when countable collections of degrees have least upper bounds and greatest lower bounds.

If **a** is a degree such that $\mathbf{a} \not\geq_{\mathrm{M}} \mathbf{d}$ for all join-irreducible $\mathbf{d} >_{\mathrm{M}} \mathbf{0}'$, then $\mathbf{a} \to \mathbf{d} = \mathbf{d}$ for all join-irreducible $\mathbf{d} >_{\mathrm{M}} \mathbf{0}'$. The degree **a** constructed in Theorem 3.8 enjoys a similar property.

Theorem 3.9 There is a degree $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$ such that $\mathbf{a} \to \prod_{i=0}^{n} \mathbf{d}_{i} = \prod_{i=0}^{n} \mathbf{d}_{i}$ whenever $\mathbf{d}_{i} >_{\mathbf{M}} \mathbf{0}'$ and is join-irreducible for each $i \leq n$.

Proof As in Theorem 3.8, let $\{g_i \mid i \in \omega\}$ be a countable collection of functions all of distinct minimal Turing degree, let $\mathcal{A} = \bigcup_{i \in \omega} i \cap \mathcal{B}_{g_i}$, and put $\mathbf{a} = [\mathcal{A}]$. Suppose $\mathbf{d}_i >_{\mathbf{M}} \mathbf{0}'$ and is join-irreducible for each $i \leq n$. By Theorem 2.5, for

each $i \leq n$, let $\mathcal{I}_i \subseteq \omega^{\omega}$ be a Turing ideal such that $\mathbf{d}_i = [\omega^{\omega} - \mathcal{I}_i]$. Thus $\prod_{i=0}^{n} \mathbf{d}_i = [\bigcup_{i=0}^{n} i^{\frown} (\omega^{\omega} - \mathcal{I}_i)]$ and

$$\mathbf{a} \to \prod_{i=0}^{n} \mathbf{d}_{i} = \left[\left\{ e^{g} \mid (\forall f \in \mathcal{A}) \left(\Phi_{e}(f \oplus g) \in \bigcup_{i=0}^{n} i^{(\omega)}(\omega^{\omega} - \mathcal{I}_{i}) \right) \right\} \right]$$

We now describe a reduction witnessing $\mathbf{a} \to \prod_{i=0}^{n} \mathbf{d}_i \ge_{\mathrm{M}} \prod_{i=0}^{n} \mathbf{d}_i$.

Given $e^{\frown}g$, for each $i \le n + 1$ search for a string $i^{\frown}\sigma_i$ such that

$$\Phi_e((i \cap \sigma_i) \oplus g)(0) \downarrow$$
.

If there is a $k \le n$ such that

$$\Phi_e((i \cap \sigma_i) \oplus g)(0) = \Phi_e((j \cap \sigma_j) \oplus g)(0) = k$$

for two distinct $i, j \le n + 1$, choose the least such k and output $k^{\frown}g$. Otherwise, output 0.

Suppose we apply this reduction to $e^{\frown}g \in A \to \bigcup_{i=0}^{n} i^{\frown}(\omega^{\omega} - I_i)$. $\Phi_e(f \oplus g)$ must be total for each $f \in A$, and for each $i \in \omega$ there is an $f \in A$ with f(0) = i. Thus for each $i \leq n+1$ the search finds a string $i^{\frown}\sigma_i$ such that $\Phi_e((i^{\frown}\sigma_i) \oplus g)(0) \downarrow$. Moreover, each $i^{\frown}\sigma_i$ can be extended to a function in A, so $\Phi_e((i^{\frown}\sigma_i) \oplus g)(0) \leq n$ for each $i \leq n+1$. Therefore, there is a least $k \leq n$ for which there are distinct $i, j \leq n+1$ with $\Phi_e((i^{\frown}\sigma_i) \oplus g)(0) = \Phi_e((j^{\frown}\sigma_j) \oplus g)(0) = k$. The reduction outputs $k^{\frown}g$, so we must show that $k^{\frown}g \in \bigcup_{i=0}^{n} i^{\frown}(\omega^{\omega} - I_i)$ which means we must show that $g \notin I_k$. Suppose for a contradiction that $g \in I_k$. The functions g_i and g_j have distinct minimal degree, so either $g \nleq_T g_i$ or $g \nleq_T g_j$ $(g >_T 0$ because $\mathbf{a} \nvDash_M \prod_{i=0}^{n} \mathbf{d}_i$ by Theorem 3.8). For the sake of argument, suppose $g \nleq_T g_i$. Then $\sigma_i^{\frown}g \nleq_T g_i$ as well, so $\sigma_i^{\frown}g \in \mathcal{B}_{g_i}$ and $i^{\frown}\sigma_i^{\frown}g \in A$. However, $\Phi_e((i^{\frown}\sigma_i^{\frown}g) \oplus g) \in k^{\frown}(\omega^{\omega} - I_k)$ by the choice of $i^{\frown}\sigma_i$. This cannot be because $(i^{\frown}\sigma_i^{\frown}g) \oplus g \in I_k$; thus anything it computes is also in I_k .

By Corollary 4.6 below, the degree $\mathbf{a} = \left[\bigcup_{i \in \omega} i \cap \mathcal{B}_{g_i}\right]$ used to witness Theorem 3.8 and Theorem 3.9 satisfies $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \operatorname{JAN}$ and so does any degree that bounds it. There are, however, degrees $>_{\mathbf{M}} \mathbf{0}'$ that do not bound any degree of the form $\left[\bigcup_{i \in \omega} i \cap \mathcal{D}_i\right]$ where $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$ and is join-irreducible for each $i \in \omega$.

Theorem 3.10 There is a degree $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$ such that $\mathbf{a} \not\geq_{\mathbf{M}} [\bigcup_{i \in \omega} i \cap \mathcal{D}_i]$ whenever $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$ and is join-irreducible for each $i \in \omega$.

Proof Let \mathcal{D}_i be such that $[\mathcal{D}_i] >_M \mathbf{0}'$ and is join-irreducible for each $i \in \omega$. By Corollary 3.4, for every $i \in \omega$, there is an $f_i >_T \mathbf{0}$ such that $\mathcal{D}_i \ge_M \mathcal{B}_{f_i}$. The mass problem \mathcal{B}_{f_i} is Turing upward-closed for each $i \in \omega$, so $\mathcal{D}_i \subseteq \mathcal{B}_{f_i}$ for each $i \in \omega$. Thus $\bigcup_{i \in \omega} i \cap \mathcal{D}_i \subseteq \bigcup_{i \in \omega} i \cap \mathcal{B}_{f_i}$. Hence it suffices to find a degree $\mathbf{a} >_M \mathbf{0}'$ that does not bound any degree of the form $[\bigcup_{i \in \omega} i \cap \mathcal{B}_{f_i}]$, where $f_i >_T \mathbf{0}$ for each $i \in \omega$.

We use the same construction used in [16] to prove Theorem 3.3. Build mass problems $A_s \subseteq \{g \mid g >_T 0\}$ such that $\{g \mid g >_T 0\} - A_s$ is finite for each $s \in \omega$. Set $A_0 = \{g \mid g >_T 0\}$. At stage s + 1, choose $h_s >_T 0$ such that h_s does not compute any of the (finitely many) functions in $\{g \mid g >_T 0\} - A_s$. If $\Phi_s(h_s)$ is total and $>_T 0$, let $g_s = \Phi_s(h_s)$ and set $A_{s+1} = A_s - \{g_s\}$. Otherwise, set $A_{s+1} = A_s$. Put $A = \bigcap_{s \in \omega} A_s$ and put $\mathbf{a} = [A]$. To see $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$, observe that by construction $\Phi_s(h_s) \notin \mathcal{A}$ for each $s \in \omega$. Now let $f_i >_{\mathbf{T}} 0$ for each $i \in \omega$. We need to show that $\Phi_e(\mathcal{A}) \notin \bigcup_{i \in \omega} i \cap \mathcal{B}_{f_i}$ for every index *e*. To do this, we first show that the functions in $\{g \mid g >_{\mathbf{T}} 0\} - \mathcal{A}$ have distinct Turing degree. Suppose that g_i leaves \mathcal{A} at stage i + 1 and g_j leaves \mathcal{A} at stage j + 1for i+1 < j+1 (i.e., at stage i+1 we had $\Phi_i(h_i) = g_i >_{\mathbf{T}} 0$, and at stage j+1 we had $\Phi_j(h_j) = g_j >_{\mathbf{T}} 0$). Then $g_i \notin_{\mathbf{T}} g_j$ because otherwise $g_i \leq_{\mathbf{T}} g_j \leq_{\mathbf{T}} h_j$, contradicting that h_j was chosen $\not\geq_{\mathbf{T}} g_i$ at stage j + 1. Now suppose $\Phi_e(g)$ is total for all $g \in \mathcal{A}$. Fix any $\sigma \in \omega^{<\omega}$ such that $\Phi_e(\sigma)(0)\downarrow$, and let *n* be such that $\Phi_e(\sigma)(0) = n$. \mathcal{A} is missing at most one function $\equiv_{\mathbf{T}} f_n$, so let $g \in \mathcal{A}$ be such that $\sigma \subset g$ and $g \equiv_{\mathbf{T}} f_n$. Then $\Phi_e(g)(0) = n$, but $\Phi_e(g) \notin n \cap \mathcal{B}_{f_n}$. Hence $\Phi_e(\mathcal{A}) \notin \bigcup_{i \in \omega} i \cap \mathcal{B}_{f_i}$.

Question 3.11 Let **a** be the degree constructed in Theorem 3.10. Does

$$\mathbf{a} \rightarrow \left[\bigcup_{i \in \omega} i^{\frown} \mathcal{D}_i\right] = \left[\bigcup_{i \in \omega} i^{\frown} \mathcal{D}_i\right]$$

whenever $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$ and is join-irreducible for each $i \in \omega$? Is $\mathrm{Th}(\mathfrak{M} / \mathbf{a}) \subseteq \mathrm{JAN}$?

Finally, we note that the answer to Question 1.3 is "no" for \mathfrak{M}_w in place of \mathfrak{M} . Let $f >_T 0$ have minimal Turing degree, and let $\mathbf{a} = [\mathfrak{B}_f]_w$. Then, in \mathfrak{M}_w , $[\mathbf{0}, \mathbf{a}] = \{\mathbf{0}, \mathbf{0}', \mathbf{a}\}$ and JAN \subsetneq Th $(\mathfrak{M}_w/\mathbf{a}) \subsetneq$ CPC.

4 New Degrees Whose Corresponding Logic Is Contained in JAN

We extend Theorem 3.1 by proving that $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \operatorname{JAN}$ for degrees \mathbf{a} such that $\mathbf{a} \ge_{\mathrm{M}} [\bigcup_{i \in \omega} i \cap \mathcal{D}_i]$ for some collection of join-irreducible degrees $[\mathcal{D}_i] >_{\mathrm{M}} \mathbf{0}'$, $i \in \omega$.

A propositional formula is called *positive* if the connective '¬' does not appear in it. For a logic L, let L^+ denote the positive formulas in L, and for a Brouwer algebra \mathfrak{B} , let $\mathrm{Th}^+(\mathfrak{B})$ denote the set of positive formulas valid in \mathfrak{B} . JAN is the maximum intermediate logic L for which $L^+ = \mathrm{IPC}^+$ [5]. This means that $L^+ = \mathrm{IPC}^+$ implies $L \subseteq \mathrm{JAN}$ for any intermediate logic L. Thus $\mathrm{Th}^+(\mathfrak{B}) = \mathrm{IPC}^+$ implies $\mathrm{Th}(\mathfrak{B}) \subseteq \mathrm{JAN}$ for any Brouwer algebra \mathfrak{B} .

Let \mathfrak{B}_1 and \mathfrak{B}_2 be Brouwer algebras. An injection $f: \mathfrak{B}_1 \to \mathfrak{B}_2$ is called a *B*-embedding if it preserves $\mathbf{0}, \mathbf{1}, +, \times, \text{ and } \to (\text{and therefore also } \neg)$. An injection $f: \mathfrak{B}_1 \to \mathfrak{B}_2$ is called a B^+ -embedding if it preserves $\mathbf{0}, +, \times, \text{ and } \to (\text{but not necessarily } \mathbf{1} \text{ or } \neg)$. If \mathfrak{B}_1 *B*-embeds into \mathfrak{B}_2 , then $\text{Th}(\mathfrak{B}_2) \subseteq \text{Th}(\mathfrak{B}_1)$, and if \mathfrak{B}_1 B^+ -embeds into \mathfrak{B}_2 , then $\text{Th}^+(\mathfrak{B}_2) \subseteq \text{Th}^+(\mathfrak{B}_1)$. Both of these facts are easily checked in light of [9], Theorem VI.2.4. If $\mathbf{a} \leq \mathbf{b}$ are in a Brouwer algebra \mathfrak{B} , then $\mathfrak{B} / \mathbf{a} B^+$ -embeds into $\mathfrak{B} / \mathbf{b}$ by the identity. This implies that $\text{Th}^+(\mathfrak{B} / \mathbf{b}) \subseteq \text{Th}^+(\mathfrak{B} / \mathbf{a})$, and it follows that the \mathbf{a} for which $\text{Th}(\mathfrak{B} / \mathbf{a}) \subseteq \text{JAN}$ is upward-closed in any Brouwer algebra \mathfrak{B} .

Our goal is to B^+ -embed a certain class of Brouwer algebras into $\mathfrak{M} / \mathbf{a}$. For any set *X*, let $\operatorname{Fr}(X)$ denote the free distributive lattice generated by *X* and let $\mathbf{0} \oplus \operatorname{Fr}(X)$ denote $\operatorname{Fr}(X)$ with a new bottom element $\mathbf{0}$. The elements of $\operatorname{Fr}(X)$ are all of the form $\sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v$ where *V* and the U_v are finite sets of indices and the \mathbf{x}_u^v are all in *X* (see, for example, Balbes and Dwinger [1], Section V.3). For such representations, $\sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v \leq \sum_{j \in J} \prod_{i \in I_j} \mathbf{y}_i^j$ if and only if

$$(\forall v \in V)(\exists j \in J)(\forall i \in I_j)(\exists u \in U_v)(\mathbf{x}_u^v \leq \mathbf{y}_i^J).$$

If $\mathbf{a}, \mathbf{b} \in \operatorname{Fr}(X)$ are such that $\mathbf{a} \not\geq \mathbf{b}$, then $\mathbf{a} \to \mathbf{b}$ exists. To see this, let $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v$ and $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} \mathbf{y}_i^j$ be representations for \mathbf{a} and \mathbf{b} . Then check

$$\mathbf{a} \to \mathbf{b} = \sum \left\{ \prod_{i \in I_j} \mathbf{y}_i^j \mid \Big(\forall v \in V \Big) \Big(\prod_{i \in I_j} \mathbf{y}_i^j \nleq \prod_{u \in U_v} \mathbf{x}_u^v \Big) \right\}.$$

If $\mathbf{a} \ge \mathbf{b}$ are in Fr(X) for an infinite *X*, then $\mathbf{a} \to \mathbf{b}$ fails to exist because in this case Fr(X) has no least element. We see then that $\mathbf{a} \to \mathbf{b}$ exists for every $\mathbf{a}, \mathbf{b} \in \mathbf{0} \oplus Fr(X)$. If *X* is finite, then so are Fr(X) and $\mathbf{0} \oplus Fr(X)$. Hence, both are Brouwer algebras. Let Fr(n) denote the free distributive lattice with *n* generators. The logic $LM = \bigcap_{n \in \omega} Th(\mathbf{0} \oplus Fr(n))$ is called the *Medvedev logic of finite problems*. (LM is usually defined in terms of Brouwer algebras isomorphic to the $\mathbf{0} \oplus Fr(n)$. See [16] for details.) We take advantage of the fact that $LM^+ = IPC^+$ [8].

If X is infinite, then $\mathbf{0} \oplus \operatorname{Fr}(X)$ fails to be a Brouwer algebra only because it lacks a top element. Therefore, the notion of a B^+ -embedding makes sense when we allow \mathfrak{B}_1 or \mathfrak{B}_2 to be $\mathbf{0} \oplus \operatorname{Fr}(X)$. If we let $\mathbf{0} \oplus \operatorname{Fr}(X) \oplus \mathbf{1}$ denote $\operatorname{Fr}(X)$ with a new bottom element **0** and a new top element **1**, then $\mathbf{0} \oplus \operatorname{Fr}(X) \oplus \mathbf{1}$ is always a Brouwer algebra.

For any partial order (P, \leq_P) , let $Fr(P, \leq_P)$ denote the free distributive lattice generated by (P, \leq_P) . $Fr(P, \leq_P)$ is the quotient $Fr(P)/\equiv$ where, for

$$\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v$$
 and $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} \mathbf{y}_i^j$ in Fr(P),

 $\mathbf{a} \equiv \mathbf{b}$ if and only if $(\mathbf{a} \leq \mathbf{b}) \land (\mathbf{b} \leq \mathbf{a})$ and $\mathbf{a} \leq \mathbf{b}$ if and only if

$$(\forall v \in V) (\exists j \in J) (\forall i \in I_j) (\exists u \in U_v) (\mathbf{x}_u^v \leq_P \mathbf{y}_i^J).$$

 $Fr(P, \leq_P)$ is always a distributive lattice, and $\mathbf{0} \oplus Fr(P, \leq_P) \oplus \mathbf{1}$ is always a Brouwer algebra; also see [13].

The following lemmas facilitate our embeddings. Lemma 4.3 is a slight generalization of the claim in the proof of [13], Lemma 2.3 and of [10], Lemma 6. The embedding is done in Theorem 4.4 which is nearly identical to [14], Theorem 2.11. Part of the reason for reproducing the proof here is to (hopefully) correct the notational inconsistencies in the proof in [14]. We restate [14], Theorem 2.11, for reference.

Theorem 4.1 ([14], Theorem 2.11) Let $\mathbf{d} = \prod_{i=0}^{n} \mathbf{d}_{i}$ where $\mathbf{d}_{i} >_{\mathbf{M}} \mathbf{0}'$ and \mathbf{d}_{i} is joinirreducible for each $i \leq n$. Then $\mathbf{0} \oplus \operatorname{Fr}(P, \leq_{P}) \oplus \mathbf{1}$ B-embeds into $\mathfrak{M} / \mathbf{d}$ for every countable partial order (P, \leq_{P}) .

(The above theorem is stated more generally in [14]. Each degree \mathbf{d}_i , for $i \leq n$, is allowed to be either join-irreducible or $\mathfrak{D}e$ -irreducible. A degree \mathbf{a} is *dense* if it is of the form [\mathcal{A}] where \mathcal{A} is dense in ω^{ω} , and a degree \mathbf{d} is $\mathfrak{D}e$ -irreducible if $\mathbf{a} \rightarrow \mathbf{d} = \mathbf{d}$ for all dense degrees \mathbf{a} . We do not consider $\mathfrak{D}e$ -irreducible degrees in our version of [14], Theorem 2.11, which is Theorem 4.4 below, because in Theorem 4.4 we require that the mass problems \mathcal{D}_i (which play the role of the degrees \mathbf{d}_i in [14], Theorem 2.11) are Turing upward-closed. Mass problems that are Turing upward-closed are dense and hence their degrees are not $\mathfrak{D}e$ -irreducible.)

Lemma 4.2 ([3]) If $X \not\geq_M \mathcal{Y}$ are mass problems, then there is a $\mathcal{W} \subseteq X$ with $|\mathcal{W}| \leq \omega$ such that $\mathcal{W} \not\geq_M \mathcal{Y}$.

Proof $\mathfrak{X} \not\geq_{\mathrm{M}} \mathfrak{Y}$ means that there is no Turing functional Φ such that $\Phi(\mathfrak{X}) \subseteq \mathfrak{Y}$. Thus, for each functional Φ_e , there must be some function $f_e \in \mathfrak{X}$ such that $\Phi_e(f_e) \notin \mathcal{Y}$. Let \mathcal{W} consist of a choice of one such $f_e \in \mathcal{X}$ for each functional Φ_e .

Lemma 4.3 Let \mathcal{U} , \mathcal{V} , and \mathcal{F}_i , for $i \in \omega$, be mass problems such that $\bigcup_{i \in \omega} i \cap \mathcal{F}_i$ $\leq_{\mathrm{M}} \mathcal{U} + \mathcal{V}$ and $\sigma \cap \mathcal{U} \subseteq \mathcal{U}$ for all $\sigma \in \omega^{<\omega}$. Then there are mass problems \mathcal{V}_i , for $i \in \omega$, such that $\bigcup_{i \in \omega} i \cap \mathcal{V}_i \equiv_{\mathrm{M}} \mathcal{V}$ and $\mathcal{F}_i \leq_{\mathrm{M}} \mathcal{U} + \mathcal{V}_i$ for each $i \in \omega$.

Proof Let \mathcal{U} , \mathcal{V} , and \mathcal{F}_i , for $i \in \omega$, be as in the statement of the lemma. Let Φ be such that $\Phi(\mathcal{U} + \mathcal{V}) \subseteq \bigcup_{i \in \omega} i^{\widehat{}} \mathcal{F}_i$. For each $i \in \omega$, define $\mathcal{V}_i = \{g \in \mathcal{V} \mid (\exists \sigma \in \omega^{<\omega})(\Phi(\sigma \oplus g)(0) = i)\}$. $\mathcal{V} \leq_{\mathrm{M}} \bigcup_{i \in \omega} i^{\widehat{}} \mathcal{V}_i$ is clear. $\bigcup_{i \in \omega} i^{\widehat{}} \mathcal{V}_i \leq_{\mathrm{M}} \mathcal{V}$ by the reduction which, given g, searches for a $\sigma \in \omega^{<\omega}$ such that $\Phi(\sigma \oplus g)(0) \downarrow$ and outputs $\Phi(\sigma \oplus g)(0)^{\widehat{}} g$. To see $i^{\widehat{}} \mathcal{F}_i \leq_{\mathrm{M}} \mathcal{U} + \mathcal{V}_i$, consider the reduction which, given $f \oplus g$, searches for a $\sigma \in \omega^{<\omega}$ such that $\Phi(\sigma \oplus g)(0) = i$ and outputs $\Phi((\sigma^{\widehat{}} f) \oplus g)$. If $f \oplus g \in \mathcal{U} + \mathcal{V}_i$, then such a σ is found, $\sigma^{\widehat{}} f$ is in \mathcal{U} , and $\Phi((\sigma^{\widehat{}} f) \oplus g)$ is in $i^{\widehat{}} \mathcal{F}_i$.

Theorem 4.4 Let $\mathbf{d} = \left[\bigcup_{i \in \omega} i \cap \mathcal{D}_i\right]$ where $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$, $[\mathcal{D}_i]$ is join-irreducible, and \mathcal{D}_i is Turing upward-closed for each $i \in \omega$. Then $\mathbf{0} \oplus \operatorname{Fr}(2^{\omega}) B^+$ -embeds into $\mathfrak{M} / \mathbf{d}$.

Proof Let \mathcal{D}_i , for $i \in \omega$, be as in the statement of the theorem, let $\mathcal{D} = \bigcup_{i \in \omega} i \cap \mathcal{D}_i$, and let $\mathbf{d} = [\mathcal{D}]$. Lemma 3.7 proves that $\mathbf{d} >_{\mathbf{M}} \mathbf{0}'$. By Lemma 4.2, let $\mathcal{A} \subseteq \{f \mid f >_{\mathbf{T}} 0\}$ be a countable mass problem such that $\mathcal{A} \not\geq_{\mathbf{M}} \mathcal{D}$. Let $\{f_x \mid x \in 2^{\omega}\}$ be a collection of functions such that $f_x \mid_{\mathbf{T}} f_y$ for all $x, y \in 2^{\omega}$ with $x \neq y$ and that $f \not\leq_{\mathbf{T}} f_x$ for all $f \in \mathcal{A}$ and $x \in 2^{\omega}$. Such a sequence can be constructed via standard recursion-theoretic techniques: build a perfect tree whose paths are Turing incomparable and do not compute any functions in \mathcal{A} . See, for example, [6], Section II.4. Notice that $\mathcal{B}_{f_x} \leq_{\mathbf{M}} \mathcal{A}$ (because $\mathcal{A} \subseteq \mathcal{B}_{f_x}$) for each $x \in 2^{\omega}$.

Define $G: \mathbf{0} \oplus \operatorname{Fr}(2^{\omega}) \to \mathfrak{M}$ as follows. Let $G(\mathbf{0}) = \mathbf{0}$ and let $G(x) = [\mathscr{B}_{f_x}]$ on the generators $x \in 2^{\omega}$ of $\operatorname{Fr}(2^{\omega})$. Then extend G to all of $\mathbf{0} \oplus \operatorname{Fr}(2^{\omega})$ so that $G(\sum_{v \in V} \prod_{u \in U_v} x_u^v) = \sum_{v \in V} \prod_{u \in U_v} G(x_u^v)$. G preserves $\mathbf{0}$, +, and × by definition, and G is injective and preserves \to by Lemma 3.2, items (iii) and (iv). Hence G is a B^+ -embedding (this is essentially [13], Corollary 2.5). Now define $H: \mathbf{0} \oplus \operatorname{Fr}(2^{\omega}) \to \mathfrak{M}/\mathbf{d}$ by $H(\mathbf{a}) = G(\mathbf{a}) \times \mathbf{d}$ for all $\mathbf{a} \in \mathbf{0} \oplus \operatorname{Fr}(2^{\omega})$. This H is the desired B^+ -embedding. By definition, H preserves $\mathbf{0}$, +, and ×. We must show that H is injective and that H preserves \to .

Clearly, $H(\mathbf{a}) = \mathbf{0}$ if and only if $\mathbf{a} = \mathbf{0}$, so to show that H is injective we let $\mathbf{a}, \mathbf{b} \in \operatorname{Fr}(2^{\omega})$ be such that $H(\mathbf{a}) \leq_{\mathrm{M}} H(\mathbf{b})$ and show that $\mathbf{a} \leq \mathbf{b}$. Let $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} x_u^v$ be a representation for \mathbf{a} and let $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} y_i^j$ be a representation for \mathbf{b} . $H(\mathbf{a}) \leq_{\mathrm{M}} H(\mathbf{b})$ means that

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \times \mathbf{d} \leq_{\mathbf{M}} \sum_{j \in J} \prod_{i \in I_j} G(y_i^j) \times \mathbf{d}$$

Therefore,

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \times \mathbf{d} \leq_{\mathbf{M}} \sum_{j \in J} \prod_{i \in I_j} G(y_i^j) = \prod \left\{ \sum_{j \in J} G(y_{\alpha(j)}^j) \mid \alpha \in \prod_{j \in J} I_j \right\}$$

where the equality is by distributivity $(\prod_{j \in J} I_j \text{ denotes the Cartesian product of the } I_j s)$. Hence,

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \times \mathbf{d} \leq_{\mathbf{M}} \sum_{j \in J} G(y_{\alpha(j)}^j) \text{ for each } \alpha \in \prod_{j \in J} I_j.$$

Each $\sum_{j \in J} G(y_{\alpha(j)}^j)$ is meet-irreducible by Lemma 3.2, item (ii). Also, **d** $\leq_{\mathrm{M}} \sum_{j \in J} G(y_{\alpha(j)}^j)$ for each $\alpha \in \prod_{j \in J} I_j$ because $\sum_{j \in J} G(y_{\alpha(j)}^j) \leq_{\mathrm{M}} [\mathcal{A}]$ but $\mathbf{d} \leq_{\mathrm{M}} [\mathcal{A}]$. Thus,

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \leq_{\mathrm{M}} \sum_{j \in J} G(y_{\alpha(j)}^j) \text{ for each } \alpha \in \prod_{j \in J} I_j,$$

and this implies that

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \leq_{\mathbf{M}} \prod \left\{ \sum_{j \in J} G(y_{\alpha(j)}^j) \mid \alpha \in \prod_{j \in J} I_j \right\}.$$

The left-hand side of the above inequality is $G(\mathbf{a})$ and the right-hand side is $G(\mathbf{b})$. *G* is a B^+ -embedding, so we conclude $\mathbf{a} \leq \mathbf{b}$.

If either of $\mathbf{a}, \mathbf{b} \in \mathbf{0} \oplus \operatorname{Fr}(2^{\omega})$ is $\mathbf{0}$, then clearly $H(\mathbf{a} \to \mathbf{b}) = H(\mathbf{a}) \to H(\mathbf{b})$. So as before, let $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} x_u^v$ and let $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} y_i^j$ be in $\operatorname{Fr}(2^{\omega})$. We see $H(\mathbf{a} \to \mathbf{b}) \ge_M H(\mathbf{a}) \to H(\mathbf{b})$ because

$$H(\mathbf{a} \rightarrow \mathbf{b}) + H(\mathbf{a}) = H((\mathbf{a} \rightarrow \mathbf{b}) + \mathbf{a}) \ge_{\mathrm{M}} H(\mathbf{b})$$

To show that $H(\mathbf{a} \to \mathbf{b}) \leq_M H(\mathbf{a}) \to H(\mathbf{b})$, we show that if $\mathbf{z} \in \mathfrak{M}$ is such that $H(\mathbf{b}) \leq_M H(\mathbf{a}) + \mathbf{z}$, then $H(\mathbf{a} \to \mathbf{b}) \leq_M \mathbf{z}$. Suppose $H(\mathbf{b}) \leq_M H(\mathbf{a}) + \mathbf{z}$. That is,

$$\sum_{j \in J} \prod_{i \in I_j} G(y_i^j) \times \mathbf{d} \leq_{\mathbf{M}} \left(\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \times \mathbf{d} \right) + \mathbf{z}.$$
 (1)

Since $\mathbf{a} \to \mathbf{b} = \sum \{\prod_{i \in I_j} y_i^j \mid (\forall v \in V) (\prod_{i \in I_j} y_i^j \nleq \prod_{u \in U_v} x_u^v)\}$, we have $H(\mathbf{a} \to \mathbf{b}) = G(\mathbf{a} \to \mathbf{b}) \times \mathbf{d}$

$$= \sum \left\{ \prod_{i \in I_j} G(y_i^j) \mid \left(\forall v \in V \right) \left(\prod_{i \in I_j} G(y_i^j) \not\leq_{\mathsf{M}} \prod_{u \in U_v} G(x_u^v) \right) \right\} \times \mathbf{d}.$$

It suffices to show that, given $j \in J$, if $\prod_{i \in I_i} G(y_i^j)$ satisfies

$$\Big(\forall v \in V\Big)\Big(\prod_{i \in I_j} G(y_i^j) \not\leq_{\mathbf{M}} \prod_{u \in U_v} G(x_u^v)\Big)$$

then $\prod_{i \in I_j} G(y_i^j) \times \mathbf{d} \leq_M \mathbf{z}$. Suppose $\prod_{i \in I_j} G(y_i^j)$ is such a meet. Then we know

$$\left(\forall v \in V\right) \left(\exists u \in U_v\right) \left(\prod_{i \in I_j} G(y_i^j) \nleq_{\mathsf{M}} G(x_u^v)\right)$$

By choosing such a $u \in U_v$, for every $v \in V$, and by appealing to Lemma 3.2, items (i) and (ii), we see that there is an $\alpha \in \prod_{v \in V} U_v$ such that

$$\prod_{i \in I_j} G(y_i^j) \not\leq_{\mathsf{M}} \sum_{v \in V} G(x_{\alpha(v)}^v).$$
⁽²⁾

Distributing $\sum_{v \in V} \prod_{u \in U_v} G(x_u^v)$ in the right-hand side of (1) yields

$$\prod_{i \in I_j} G(\mathbf{y}_i^j) \times \mathbf{d} \leq_{\mathbf{M}} \sum_{v \in V} G(\mathbf{x}_{a(v)}^v) + \mathbf{z}.$$

The degree $\sum_{v \in V} G(x_{\alpha(v)}^v)$ is a finite join of degrees of the form $[\mathcal{B}_f]$ and thus has a representative \mathcal{U} such that $\sigma \cap \mathcal{U} \subseteq \mathcal{U}$ for all $\sigma \in \omega^{<\omega}$. So, by Lemma 4.3, there are mass problems \mathcal{Z}_i for $i \in I_j$ and $\widehat{\mathcal{Z}}_i$ for $i \in \omega$ such that

$$\mathbf{z} = \left(\prod_{i \in I_j} [\mathcal{Z}_i]\right) \times \left[\bigcup_{i \in \omega} i \widehat{\mathcal{Z}}_i\right],$$

$$G(y_i^j) \leq_{\mathbf{M}} \sum_{v \in V} G(x_{\alpha(v)}^v) + [\mathcal{Z}_i] \text{ for each } i \in I_j, \text{ and}$$

$$[\mathcal{D}_i] \leq_{\mathbf{M}} \sum_{v \in V} G(x_{\alpha(v)}^v) + [\widehat{\mathcal{Z}}_i] \text{ for each } i \in \omega.$$

Each $G(y_i^j)$ is join-irreducible, and $G(y_i^j) \not\leq_M \sum_{v \in V} G(x_{\alpha(v)}^v)$ by (2). Thus $G(y_i^j) \leq_M [\mathbb{Z}_i]$ for each $i \in \omega$, so $\prod_{i \in I_j} G(y_i^j) \leq_M \prod_{i \in I_j} [\mathbb{Z}_i]$. Each $[\mathcal{D}_i]$ is join-irreducible by assumption, and $[\mathcal{D}_i] \not\leq_M \sum_{v \in V} G(x_{\alpha(v)}^v)$ because the right-hand side is $\leq_M [\mathbb{A}]$ but the left-hand side is not. It follows that $[\mathcal{D}_i] \leq_M [\widehat{\mathbb{Z}}_i]$ for each $i \in \omega$, and so $\widehat{\mathbb{Z}}_i \subseteq \mathcal{D}_i$ for each $i \in \omega$ because each \mathcal{D}_i is Turing upward-closed. Thus, $\bigcup_{i \in \omega} i \widehat{\mathbb{Z}}_i \subseteq \mathcal{D}$, so $\mathbf{d} \leq_M [\bigcup_{i \in \omega} i \widehat{\mathbb{Z}}_i]$. Therefore, $\prod_{i \in I_j} G(y_i^j) \times \mathbf{d} \leq_M (\prod_{i \in I_j} [\mathbb{Z}_i]) \times [\bigcup_{i \in \omega} i \widehat{\mathbb{Z}}_i] = \mathbf{z}$ as desired. \Box

Corollary 4.5 If $\mathbf{a} \ge_{\mathbf{M}} \mathbf{d}$ are degrees such that $\mathbf{d} = \left[\bigcup_{i \in \omega} i \cap \mathcal{D}_i\right]$ where $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$ and is join-irreducible for each $i \in \omega$, then $\mathbf{0} \oplus \operatorname{Fr}(2^{\omega}) B^+$ -embeds into $\mathfrak{M} / \mathbf{a}$.

Proof Let **a**, **d**, and \mathcal{D}_i , for $i \in \omega$, be as in the statement of the corollary. Let $\mathbf{d}_0 = \left[\bigcup_{i \in \omega} i \cap C(\mathcal{D}_i)\right]$ and notice that $\mathbf{d} \ge_M \mathbf{d}_0$. $\mathcal{D}_i \equiv_M C(\mathcal{D}_i)$ for each $i \in \omega$ by Lemma 2.3, so \mathbf{d}_0 satisfies the hypotheses of Theorem 4.4. Thus, $\mathbf{0} \oplus \operatorname{Fr}(2^{\omega})$ B^+ -embeds into $\mathfrak{M} / \mathbf{d}_0$ which B^+ -embeds into $\mathfrak{M} / \mathbf{a}$.

Corollary 4.6 If $\mathbf{a} \ge_{\mathbf{M}} \mathbf{d}$ are degrees such that $\mathbf{d} = \left[\bigcup_{i \in \omega} i \cap \mathcal{D}_i\right]$ where $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$ and is join-irreducible for each $i \in \omega$, then $\operatorname{Th}(\mathfrak{M} / \mathbf{a}) \subseteq \operatorname{JAN}$.

Proof The Brouwer algebra $\mathbf{0} \oplus \operatorname{Fr}(n) B^+$ -embeds into $\mathbf{0} \oplus \operatorname{Fr}(2^{\omega})$ for each *n*, and $\mathbf{0} \oplus \operatorname{Fr}(2^{\omega}) B^+$ -embeds into \mathfrak{M}/\mathbf{a} by Corollary 4.5. Thus, $\operatorname{Th}^+(\mathfrak{M}/\mathbf{a}) \subseteq \bigcap_{n \in \omega} \operatorname{Th}^+(\mathbf{0} \oplus \operatorname{Fr}(n)) = \operatorname{LM}^+ = \operatorname{IPC}^+$. So $\operatorname{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \operatorname{JAN}$.

Theorem 4.4 can be modified to *B*-embed $\mathbf{0} \oplus \operatorname{Fr}(2^{\omega}) \oplus \mathbf{1}$ into \mathfrak{M}/\mathbf{d} for degrees \mathbf{d} as in the statement of Theorem 4.4. However, if $\mathbf{a} \leq \mathbf{b}$ in a Brouwer algebra \mathfrak{B} , it is not in general the case that \mathfrak{B}/\mathbf{a} *B*-embeds into \mathfrak{B}/\mathbf{b} . So the proof of Corollary 4.5 fails for *B*-embedding $\mathbf{0} \oplus \operatorname{Fr}(2^{\omega}) \oplus \mathbf{1}$. Theorem 4.4 can also be modified to prove a more precise analogue of [14], Theorem 2.11 (restated as Theorem 4.1 above). Let $\mathbf{d} = [\bigcup_{i \in \omega} i \cap \mathfrak{D}_i]$ where $[\mathfrak{D}_i] >_{\mathbf{M}} \mathbf{0}', [\mathfrak{D}_i]$ is join-irreducible, and \mathfrak{D}_i is Turing upward-closed for each $i \in \omega$. Then $\mathbf{0} \oplus \operatorname{Fr}(P, \leq_P) \oplus \mathbf{1}$ *B*-embeds into \mathfrak{M}/\mathbf{d} for every countable partial order (P, \leq_P) .

5 \mathfrak{F}_{cl} Is Not Prime

Recall that a filter \mathfrak{F} in a lattice is called *prime* if $\mathbf{a} + \mathbf{b} \in \mathfrak{F} \to \mathbf{a} \in \mathfrak{F} \lor \mathbf{b} \in \mathfrak{F}$ for all \mathbf{a} and \mathbf{b} in the lattice. Theorem 2.5 can be phrased as a characterization of the prime principal filters in \mathfrak{M} : a degree \mathbf{a} generates a prime filter if and only if $\mathbf{a} = [\omega^{\omega} - \mathfrak{l}]$ for some Turing ideal \mathfrak{l} . There is little general theory of the nonprincipal filters in \mathfrak{M} , but several specific cases have been studied in Dyment [3], Sorbi [11], Bianchini and Sorbi [2], and Lewis, Shore, and Sorbi [7]. See also [15] for a summary of many of the results appearing in these works. We consider the filters \mathfrak{F} and \mathfrak{F}_{cl} .

Definition 5.1

- (i) A degree a is called *dense* (*closed*) if a = [A] for an A that is dense (closed) in ω^ω.
- (ii) \Im denotes the ideal generated by $\{a \mid a \text{ is dense}\}$.
- (iii) \mathfrak{F} denotes $\mathfrak{M} \mathfrak{I}$.
- (iv) \mathfrak{F}_{cl} denotes the filter generated by $\{a \mid a >_M 0 \text{ and is closed}\}$.

The join and meet of two dense degrees is dense [3], and the join and meet of two closed degrees is closed [2]. It follows that $\mathfrak{I} = \{\mathbf{b} \mid (\exists \mathbf{a} \ge_M \mathbf{b})(\mathbf{a} \text{ is dense})\}$ and $\mathfrak{F}_{cl} = \{\mathbf{b} \mid (\exists \mathbf{a} \le_M \mathbf{b})(\mathbf{a} >_M \mathbf{0} \text{ and is closed})\}$. The basic properties of $\mathfrak{I}, \mathfrak{F}, \text{ and } \mathfrak{F}_{cl}$ are as follows: \mathfrak{I} is a prime ideal [11], \mathfrak{F} is a prime filter [2], \mathfrak{I} is not principal [3], \mathfrak{F} and \mathfrak{F}_{cl} are not principal [2], and $\mathfrak{F}_{cl} \subsetneq \mathfrak{F}$ [2]. Both [2] and [15] ask for a proof that \mathfrak{F}_{cl} is not prime. We provide a proof of this fact now.

Lemma 5.2 For any $f \in \omega^{\omega}$ there are $\mathcal{A}, \mathcal{B} \subseteq \omega^{\omega}$ such that $\mathcal{A} + \mathcal{B} \ge_{M} \{f\}$ and, for any closed $\mathcal{C} \subseteq \omega^{\omega}$, if $\mathcal{A} \ge_{M} \mathcal{C}$ or $\mathcal{B} \ge_{M} \mathcal{C}$, then \mathcal{C} contains a recursive function.

Proof Fix a recursive bijection $\omega \leftrightarrow \omega^{<\omega}$. For $e, n \in \omega$, if

$$\forall m \forall \sigma (\exists \tau \supseteq \sigma) (\Phi_e(n^{\frown} \tau)(m) \downarrow),$$

then define $\eta(e, n, i) \in \omega^{<\omega}$ by induction on $i \in \omega$ as follows. Let $\eta(e, n, 0) = n^{\sigma}\sigma$, where σ is the least string such that $\Phi_e(n^{\sigma}\sigma)(0)\downarrow$. Given $\eta(e, n, i)$, let $\eta(e, n, i + 1) = \eta(e, n, i)^{\sigma}\sigma$, where σ is the least string such that

$$\Phi_e(\eta(e, n, i) \cap \sigma)(i+1) \downarrow$$

Let $f \in \omega^{\omega}$. We construct \mathcal{A} and \mathcal{B} such that

(i) if $g \in A$, then g(0) has the form

$$g(0) = \langle \ell, \langle n_0, x_0, y_0 \rangle, \dots, \langle n_{\ell-1}, x_{\ell-1}, y_{\ell-1} \rangle \rangle,$$

where $\ell \in \omega$ and $n_i \in \omega$, $x_i \in \{0, 1\}$, and $y_i \in \omega$ for each $i < \ell$;

- (ii) if $g \in A$ and $\langle n_e, 0, y_e \rangle$ is in the *e*th position of g(0), then
 - (a) $\exists m \exists \sigma (\forall \tau \supseteq \sigma) (\Phi_e(n_e \frown \tau)(m) \uparrow),$
 - (b) any $h \in \mathcal{B}$ with $h(0) = n_e$ is of the form $h = n_e \cap \sigma \cap f$, where $|\sigma| = y_e$;
- (iii) if $g \in A$ and $\langle n_e, 1, y_e \rangle$ is the *e*th position of g(0), then
 - (a) $\forall m \forall \sigma (\exists \tau \supseteq \sigma) (\Phi_e(n_e \cap \tau)(m) \downarrow),$
 - (b) any $h \in \mathcal{B}$ with $h(0) = n_e$ is of the form $h = \eta(e, n_e, i)^{-1} f$ for some $i \in \omega$;

(iv) the above properties hold with the roles of \mathcal{A} and \mathcal{B} reversed.

We construct \mathcal{A} and \mathcal{B} in stages. The construction is similar to the construction in Lemma 2.3 in that if g goes into \mathcal{A} before h goes into \mathcal{B} , then h(0) codes how to

recover *f* from *g*, and similarly with the roles of *A* and *B* reversed. Start at stage 0 with $A = \emptyset$, $B = \emptyset$, $s = \langle \rangle$, and $t = \langle \rangle$.

Stage e + 1 Set $n_e = e^{-t}$.

Case 1 $\exists m \exists \sigma (\forall \tau \supseteq \sigma) (\Phi_e(n_e \cap \tau)(m) \uparrow)$. Choose such a σ and put $n_e \cap \sigma \cap f$ in \mathcal{A} . Update $s = s \cap \langle n_e, 0, |\sigma| \rangle$.

Case 2 $\forall m \forall \sigma (\exists \tau \supseteq \sigma) (\Phi_e(n_e \cap \tau)(m) \downarrow)$. Put the functions $\eta(e, n_e, i) \cap 1 \cap f$ in \mathcal{A} for each $i \in \omega$. Update $s = s \cap \langle n_e, 1, 0 \rangle$.

Repeat the above procedure with the roles of A and B reversed and the roles of s and t reversed. This completes stage e + 1. Then go on to stage e + 2. This completes the construction.

Suppose $A \ge_M C$ where C is closed. We show that C contains a recursive function. The proof with \mathcal{B} in place of A is the same. Let $\Phi_e(A) \subseteq C$. Consider stage e + 1 of the above construction. Case 1 must not have occurred because otherwise A would contain a function $n_e^{-\sigma} \cap f$ such that $\Phi_e(n_e^{-\sigma} \cap f)$ is not total. Thus case 2 occurred, and so A contains the function $\eta(e, n_e, i) \cap 1^{-f} f$ for each $i \in \omega$. Let k be the recursive function $k = n_e^{-\sigma_0} \cap 0^{-\sigma_1} \cap 0^{-\sigma_2} \cap 0^{-\cdots}$, where $\eta(e, n_e, i) = n_e^{-\sigma_0} \cap 0^{-\cdots} \cap 0^{-\sigma_i}$ for each $i \in \omega$ (think of k as the "limit" of the strings $\eta(e, n_e, i)$ as $i \to \infty$). Then $\Phi_e(\eta(e, n_e, i) \cap 1^{-f}) \in C$ and $\Phi_e(\eta(e, n_e, i) \cap 1^{-f}) \upharpoonright i = \Phi_e(k) \upharpoonright i$ for each $i \in \omega$. Thus C contains the recursive function $\Phi_e(k)$ because C is closed.

We now describe a uniform procedure for producing f from $g \oplus h \in A + B$. First decode h(0) as $h(0) = \langle \ell, \langle n_0, x_0, y_0 \rangle, \dots, \langle n_{\ell-1}, x_{\ell-1}, y_{\ell-1} \rangle \rangle$ and look for g(0) among the n_e . If $\langle g(0), 0, y_e \rangle$ appears in h(0) at position e, then output g from position $y_e + 1$ onward as in this case $g = \sigma^{-1} f$ for some string σ of length $y_e + 1$. If $\langle g(0), 1, 0 \rangle$ appears in h(0) at position e, then $g = \eta(e, g(0), i)^{-1} f$ for some $i \in \omega$. Compute which i by successively computing the $\eta(e, g(0), j)$, matching them against g, and checking if the next bit of g is 0 (in which case compute $\eta(e, g(0), j + 1)$) or 1 (in which case j = i). Output f once i is found.

The number g(0) appears among the n_e coded into h(0) if g went into \mathcal{A} before h went into \mathcal{B} . Otherwise, h went into \mathcal{B} before g went into \mathcal{A} , so h(0) appears among the n_e coded in g(0). In this case, switch the roles of g and h and apply the above procedure to compute f.

Theorem 5.3 \mathfrak{F}_{cl} is not prime. In fact, if $\mathfrak{G} \subseteq \mathfrak{F}_{cl}$, $\mathfrak{G} \neq \{1\}$ is a filter, then \mathfrak{G} is not prime.

Proof Suppose $\mathfrak{G} \subseteq \mathfrak{F}_{cl}$ is a filter such that $\mathfrak{G} \neq \{1\}$. Let $f >_T 0$ be such that $[\{f\}] \in \mathfrak{G}$. Let $\mathcal{A}, \mathcal{B} \subseteq \omega^{\omega}$ be as in Lemma 5.2 for this f. Let $\mathbf{a} = [\mathcal{A}]$ and $\mathbf{b} = [\mathcal{B}]$. Then $\mathbf{a}, \mathbf{b} \notin \mathfrak{G}$ because $\mathbf{a}, \mathbf{b} \notin \mathfrak{F}_{cl}$, but $\mathbf{a} + \mathbf{b} \in \mathfrak{G}$ because $\mathbf{a} + \mathbf{b} \ge_M [\{f\}]$.

If **x** and **y** are degrees such that **y** is closed and $\mathbf{y} \leq_M \mathbf{x}$, then there is no dense degree **z** such that $\mathbf{y} \leq_M \mathbf{x} + \mathbf{z}$ [7]. Therefore, if $\mathfrak{G} \subseteq \mathfrak{F}_{cl}$, $\mathfrak{G} \neq \{\mathbf{1}\}$ is a filter, then any degrees **a** and **b** witnessing that \mathfrak{G} is not prime must both be in $\mathfrak{F} - \mathfrak{G}$.

The results of Section 3 suggest two new filters to study.

Definition 5.4

(i) (b) denotes the filter generated by

 $\{\mathbf{d} \mid \mathbf{d} >_{\mathbf{M}} \mathbf{0}' \text{ and is join-irreducible} \}.$

(ii) \mathfrak{H} denotes the filter generated by

$$\left\{ \left[\bigcup_{i \in \omega} i^{\frown} \mathcal{D}_i \right] \mid (\forall i \in \omega) ([\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}' \text{ and is join-irreducible}) \right\}.$$

(§) is exactly the set of all degrees **b** for which $\mathbf{b} \ge_{\mathrm{M}} \prod_{i=0}^{n} \mathbf{d}_{i}$ for some joinirreducible degrees $\mathbf{d}_{i} >_{\mathrm{M}} \mathbf{0}'$, $i \le n$, and \mathfrak{H} is exactly the set of all degrees **b** for which $\mathbf{b} \ge_{\mathrm{M}} [\bigcup_{i \in \omega} i \cap \mathcal{D}_{i}]$ for some join-irreducible degrees $[\mathcal{D}_{i}] >_{\mathrm{M}} \mathbf{0}'$, $i \in \omega$.

Theorem 5.5 $\mathfrak{F}_{cl} \subsetneq \mathfrak{G} \subsetneq \mathfrak{F} \subsetneq \{\mathbf{a} \mid \mathbf{a} >_M \mathbf{0}'\}$. $\mathfrak{G} \nsubseteq \mathfrak{F}$ (hence also $\mathfrak{H} \nsubseteq \mathfrak{F}$). Neither \mathfrak{G} nor \mathfrak{H} is principal.

Proof Every closed degree $>_{M} 0$ bounds a join-irreducible degree $>_{M} 0'$ [16]. Hence $\mathfrak{F}_{cl} \subseteq \mathfrak{G}$. $\mathfrak{G} \subseteq \mathfrak{G}$ is clear. To see $\mathfrak{G} \not\subseteq \mathfrak{F}$, observe that every \mathcal{B}_{f} is dense, so if $f >_{T} 0$, then $[\mathcal{B}_{f}] \in \mathfrak{G} - \mathfrak{F}$. This also shows $\mathfrak{G} \not\subseteq \mathfrak{F}_{cl}$. The degree constructed in Theorem 3.8 witnesses $\mathfrak{F} \not\subseteq \mathfrak{G}$. The degree constructed in Theorem 3.10 witnesses $\{\mathbf{a} \mid \mathbf{a} >_{M} \mathbf{0}'\} \not\subseteq \mathfrak{F}$. We show that \mathfrak{G} is not principal. The proof for \mathfrak{F} is the same. First, if \mathcal{A} is countable and contains no recursive functions, then there is a function $f >_{T} 0$ such that $g \not\leq_{T} f$ for all $g \in \mathcal{A}$. Thus $\mathcal{B}_{f} \leq_{M} \mathcal{A}$ (as $\mathcal{A} \subseteq \mathcal{B}_{f}$) for this f. Every $[\mathcal{B}_{f}]$ for $f >_{T} 0$ is in \mathfrak{G} , so every $[\mathcal{A}]$ where \mathcal{A} is countable and contains no recursive function is in \mathfrak{G} . If \mathfrak{G} were principal, it would be generated by a degree \mathbf{x} such that $\mathbf{x} \leq_{M} [\mathcal{A}]$ for all countable \mathcal{A} not containing a recursive function. By Lemma 4.2, the only such \mathbf{x} are $\mathbf{0}$ and $\mathbf{0}'$. We know $\mathbf{0}$ and $\mathbf{0}'$ are not in \mathfrak{G} , so \mathfrak{G} cannot be principal.

We end with a question.

Question 5.6

- (i) Is $\mathfrak{F} \subseteq \mathfrak{G}$? Is $\mathfrak{F} \subseteq \mathfrak{H}$?
- (ii) Is (S prime? Is S prime?
- (iii) Is $\{\mathbf{a} \mid \operatorname{Th}(\mathfrak{M} / \mathbf{a}) \subseteq \operatorname{JAN}\}$ a filter?

To prove that $\{\mathbf{a} \mid \operatorname{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \operatorname{JAN}\}\$ is a filter, it suffices to prove that $\operatorname{Th}(\mathfrak{M}/(\mathbf{a} \times \mathbf{b})) \subseteq \operatorname{JAN}\$ whenever both $\operatorname{Th}(\mathfrak{M}/\mathbf{a})\$ and $\operatorname{Th}(\mathfrak{M}/\mathbf{b})\$ are \subseteq JAN because $\{\mathbf{a} \mid \operatorname{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \operatorname{JAN}\}\$ is upward-closed in \mathfrak{M} .

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