## THE LENGTH OF THE PERIOD OF THE SIMPLE CONTINUED FRACTION OF $d^{1/2}$ .

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Let p(d) denote the length of the period of the simple continued fraction for  $d^{1/2}$  and  $\varepsilon$  the fundamental unit in the ring  $Z[d^{1/2}]$ . We prove that as  $d \to \infty$ ,

THEOREM 1.  $p(d) \leq 7/2\pi^{-2}d^{1/2}\log d + O(d^{1/2})$ .

THEOREM 2.  $\log \varepsilon \leq 3\pi^{-2} d^{1/2} \log d + O(d^{1/2})$ .

**THEOREM 3.**  $p(d) \neq o(d^{1/2}/\log \log d)$ .

THEOREM 4. If  $\log \varepsilon \neq o(d^{1/2} \log d)$  then also

 $p(d) \neq o(d^{1/2} \log d)$  .

Recently Hickerson [1] has proved that  $p(d) = O(d^{1/2+\delta})$  for every  $\delta > 0$ , and in fact a result somewhat more precise than this. Lehmer [2] has suggested that for arbitrarily large d, p(d) might be as large as  $0.30d^{1/2} \log d$ , and if this is indeed the case then Theorem 1 is almost the best possible result. In fact it is easy to show that  $p(d) = O(d^{1/2} \log d)$  using known results regarding  $\log \varepsilon$ , but the constant in Theorem 1 improves the best obtainable in this way.

Let  $\varepsilon_0$  denote the fundamental unit in the field  $Q(d^{1/2})$ ,  $[a_0, \overline{a_1, a_2}, \overline{\cdots a_{p(d)-1}, 2a_0}]$  the continued fraction for  $d^{1/2}$  and  $P_r/Q_r$  its rth convergent. Then as is well known  $\varepsilon = \varepsilon_0$  or  $\varepsilon_0^3$ . Thus by the result of Stephens [3],

$$\logarepsilon \leq 3\logarepsilon_{_0} \leq rac{3}{2}(1-e^{_{-1/2}}+\delta)d^{_{1/2}}\log d \;.$$

Now  $Q_0 = 1$ ,  $Q_1 = a_1 \ge 1$  and  $Q_{r+2} = a_{r+2}Q_{r+1} + Q_r \ge Q_{r+1} + Q_r$  and so by induction  $Q_r \ge u_{r+1}$ , the Fibonacci number, for  $r \ge 0$ . Now

$$egin{aligned} arepsilon &= P_{p(d)-1} + Q_{p(d)-1} d^{1/2} \ &> 2 d^{1/2} Q_{p(d)-1} - 1 \ &\geq 2 d^{1/2} u_{p(d)} - 1 \ &> \left\{ rac{1 + \sqrt{5}}{2} 
ight\}^{p(d)} extbf{,} \end{aligned}$$

and so  $p(d) < Ad^{1/2} \log d$  where A is approximately 5/4.

In exactly the same way, using  $a_r < d^{1/2}$  for  $0 \leq r < p(d)$  it is possible to show that  $p(d) \gg \log \varepsilon / \log d$ . Since  $d = 2^{2k+1}$  gives  $\varepsilon = (1 + \sqrt{2})^{2^k}$ , we find that for arbitrarily large d it is possible for  $p(d) \gg d^{1/2} / \log d$ , and it will be shown that this can be improved at least by replacing the  $\log d$  by  $\log \log d$ . Theorems 1 and 3 together show that the scope for sharpening the results is somewhat limited; nevertheless the remaining problem is important and worthy of further study, for as we mention in the concluding remarks, if it could be proved that  $p(d) = o(d^{1/2} \log d)$  this would imply also that  $\log \varepsilon = o(d^{1/2} \log d)$  a result which has been sought in vain for many years.

Throughout we use  $\varepsilon_1$  to denote the fundamental unit in  $Z[d^{1/2}]$ with norm + 1; then  $\varepsilon_1 = \varepsilon$  or  $\varepsilon^2$ . In accordance with established practice, if for given integers d and N there exist integers X and Y with  $X^2 - dY^2 = N$ , then we say that  $X + Yd^{1/2}$  is a solution of the equation  $x^2 - dy^2 = N$ . Given one such solution, all the members of the set  $\pm (X + Yd^{1/2})\varepsilon_1^n$  are also solutions, and this set is called a *class* of solutions. A given equation may well have more than one such class of solutions, but it is well known that the number of such classes is finite.

LEMMA 1. For each r,  $|P_r^2 - dQ_r^2| < 2d^{1/2}$ .

This is well known.

LEMMA 2. For a class K of solutions of  $x^2 - dy^2 = N$ , the g.c.d., (x, y) depends only upon K.

For if  $x_1 + y_1 d^{1/2}$  and  $x_2 + y_2 d^{1/2}$  belong to the same class, then for some integer n,

$$egin{aligned} &x_1+y_1d^{1/2}=\pm(x_2+y_2d^{1/2})arepsilon_1^n\ &=\pm(x_2+y_2d^{1/2})(a_n+b_nd^{1/2})\ , \end{aligned}$$

say. Thus  $(x_2, y_2)|(x_1, y_1)$  and similarly conversely.

A class K for which (x, y) = 1 is called a *primitive* class. The main result used in the proof of the theorems is

LEMMA 3. The number of primitive classes, f(N; d), of  $x^2 - dy^2 = N$  does not exceed  $2^{\omega(|N|)}$ . In the special case 2 || N,  $f(N; d) \leq 2^{\omega(|N|)-1}$ . Here  $\omega(N)$  denotes the number of distinct prime factors of N.

*Proof.* In the first place it suffices to consider the case in which (N, d) is square-free. For if  $(N, d) = k_1^2 k_2$  where  $k_2$  is square-free, (x, y) = 1 and  $x^2 - dy^2 = N$  then  $k_1 | x$  and so if  $x_1 = x/k_1$ ,  $N_1 = N/k_1^2$  and  $d_1 = d/k_1^2$  then  $x_1^2 - d_1y^2 = N_1$  with  $(x_1, y) = 1$ . For the latter equation we now have  $(N_1, d_1) = k_2$  which is square-free and so the total number of classes of primitive solutions of the given equation

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does not exceed  $2^{\omega(|N|)} \leq 2^{\omega(|N|)}$  in the general case, or  $2^{\omega(|N|)^{-1}} \leq 2^{\omega(|N|)^{-1}}$  in the special case 2||N| since in this case  $2||N_1|$  also. We suppose therefore from now on that (N, d) is square-free.

Let p denote any prime dividing N, and suppose that  $p^* || N$ ;

(i) if p | d then p | x, whence  $p^2 \not\mid dy^2$  otherwise we should find, since  $p \not\mid y$  that  $p^2 | d$  and  $p^2 | N$ . Hence s = 1 and so  $xy^{-1} \equiv 0 \pmod{p^s}$ .

(ii) if  $p \not\models d$  then p can divide neither x nor y, otherwise it would have to divide them both. Thus  $(xy^{-1})^2 \equiv d \pmod{p^s}$  and so if p is odd,  $xy^{-1} \equiv \pm a_p \pmod{p^s}$ .

(iii) if  $p \nmid d$ , p = 2 then  $(xy^{-1})^2 \equiv d \pmod{p^s}$  gives

(a) if  $s = 1, xy^{-1} \equiv d \pmod{2}$ , i.e.,  $xy^{-1} \equiv d \pmod{p^s}$ 

(b) if s = 2, since  $x^2 - dy^2 \equiv 0 \pmod{4}$  and both x and y are odd,  $d \equiv 1 \pmod{4}$  whence  $(xy^{-1})^2 \equiv 1 \pmod{4}$ , i.e.,  $xy^{-1} \equiv \pm 1 \pmod{4}$ , i.e.,  $xy^{-1} \equiv \pm 1 \pmod{p^s}$ 

(c) if  $s \ge 3$ , then  $d \equiv 1 \pmod{8}$  and now  $(xy^{-1})^2 \equiv d \pmod{2^s}$  gives  $xy^{-1} \equiv \pm a \pmod{2^{s-1}}$ .

Combining (i), (ii), and (iii) and using the Chinese Remainder Theorem, we see that  $xy^{-1}$  is congruent to one of at most

Next we prove that if  $x^2 - dy^2 = X^2 - dY^2 = N$  and if  $xy^{-1} \equiv XY^{-1} \pmod{N}$  then  $x + yd^{1/2}$  and  $X + Yd^{1/2}$  belong to the same class K. For

$$rac{x+yd^{1/2}}{X+Yd^{1/2}} = rac{(x+yd^{1/2})(X-Yd^{1/2})}{X^2-d\,Y^2} = rac{xX-dy\,Y}{N} + rac{-xY+Xy}{N}d^{1/2} = A+Bd^{1/2}\,, \;\; ext{say} \;.$$

Now B is an integer and A rational, and since  $A^2 - dB^2 = 1$  it follows that A too is an integer, and so that result of the lemma follows, except if 8|N.

Finally, if 8 | N then we find that if  $xy^{-1} \equiv XY^{-1} \pmod{1/2N}$  then  $x + yd^{1/2}$  and  $X + Yd^{1/2}$  belong to the same class; for if as above  $A + Bd^{1/2}$  denote their quotient, we find that B equals either an integer or else half an odd integer. In the former case the result follows as above. In the latter case we find  $(2A)^2 = d(2B)^2 + 4$  and since now 2B is an odd integer and  $4 \nmid d$ , 2A is also an odd integer, whence  $d \equiv 5 \pmod{8}$ . But this is inconsistent with  $x^2 - dy^2 \equiv 0 \pmod{8}$  where (x, y) = 1 and so this latter case does not arise. This concludes the proof.

LEMMA 4. If  $N(\varepsilon) = 1$ , then

$$f(N, d) + f(-N, d) \leq 2^{\omega(|N|)}$$

and

$$f(N, d) + f(-N, d) \leq 2^{\omega(|N|)-1} \quad if \quad 2||N.$$

*Proof.* After Lemma 3, it merely remains to prove that  $x^2 - dy^2 = N$  and  $X^2 - dY^2 = -N$  with  $xy^{-1} \equiv XY^{-1} \pmod{N}$ , or even  $(\mod 1/2N)$  if 8|N, is impossible. For we should obtain if

$$A + Bd^{1/2} = (x + yd^{1/2})(X + Yd^{1/2})^{-1}$$

that  $A^2 - dB^2 = -1$  with either A and B both integers, or else both half integers. Both cases are impossible if  $N(\varepsilon) = +1$ .

LEMMA 5. (1) If 
$$N(\varepsilon) = 1$$
 then  
 $p(d) \leq \sum_{0 < N < 2d^{1/2}} \{f(N; d) + f(-N; d)\}$ 

(2) If  $N(\varepsilon) = -1$  then

$$p(d) \leq \sum_{0 < N < 2d^{1/2}} f(N; d)$$
.

*Proof.* If  $0 \le m < n \le p(d) - 1$  then  $P_m + Q_m d^{1/2}$  and  $P_n + Q_n d^{1/2}$  are primitive solutions in distint classes; they are primitive since  $(P_r, Q_r) = 1$  and are in distinct classes since

$$1 < {P}_{\mathtt{m}} + {Q}_{\mathtt{m}} d^{\scriptscriptstyle 1/2} < {P}_{\mathtt{n}} + {Q}_{\mathtt{n}} d^{\scriptscriptstyle 1/2} \leqq arepsilon_{\mathtt{1}}$$
 .

Hence using Lemma 1,

$$p(d) \leq \text{the number of distinct primitive classes of all}$$
  
equations  $x^2 - dy^2 = N$  with  $-2d^{1/2} < N < 2d^{1/2}$   
 $= \sum_{-2d^{1/2} < N < 2d^{1/2}} f(N; d)$ , which gives (1).

If  $N(\varepsilon) = -1$  then the above reasoning applies if  $0 \le m < n \le 2p(d) - 1$  and so (2) follows, since if  $N(\varepsilon) = -1$ , f(N; d) = f(-N; d).

We remark that this result is best possible for example for the values d = 7, 13 respectively.

LEMMA 6. As  $x \to \infty$ 

(1) 
$$F(x) = \sum_{1 \le N \le x} 2^{\omega(N)} = cx \log x + O(x)$$
,

(2) 
$$A(x) = \sum_{\substack{1 \le N \le x \\ 2 \mid N}} 2^{\omega(N)} = \frac{2}{3} cx \log x + O(x)$$
,

(3) 
$$B(x) = \sum_{\substack{1 \le N \le x \\ 2 \neq N}} 2^{\omega(N)} = \frac{1}{3} cx \log x + O(x)$$
,

(4) 
$$C(x) = \sum_{\substack{1 \le N \le x \\ 4 \mid N}} 2^{\omega(N)} = \frac{1}{3} cx \log x + O(x) ,$$

(5) 
$$D(x) = \sum_{\substack{1 \le N \le x \\ 8|N}} 2^{\omega(N)} = \frac{1}{6} cx \log x + O(x) ,$$

(6) 
$$E(x) = \sum_{\substack{1 \le N \le x \\ 16 \mid N}} 2^{\omega(N)} = \frac{1}{12} cx \log x + O(x)$$
, where  $c = 6\pi^{-2}$ .

Proof. (1) The identity

$$2^{\omega(N)} = {\displaystyle \sum_{k^2 \mid N}} d \Big( {\displaystyle rac{N}{k^2}} \Big) \mu(k)$$

is easily proved by induction on the number of distinct prime factors of N. For if N is a prime or a prime power the result is immediate, and then the identity follows on observing that  $2^{\omega}$ , d and  $\mu$  are all multiplicative. Thus

$$\begin{split} F(x) &= \sum_{1 \leq N \leq x} \sum_{k^2 \mid N} d\Big(\frac{N}{k^2}\Big) \mu(k) \\ &= \sum_{1 \leq k \leq x^{1/2}} \sum_{1 \leq k_1 \leq xk^{-2}} d(k_1) \mu(k) \\ &= \sum_{1 \leq k \leq x^{1/2}} \mu(k) \sum_{1 \leq k_1 \leq xk^{-2}} d(k_1) \\ &= \sum_{1 \leq k \leq x^{1/2}} \mu(k) \left\{\frac{x}{k^2} \log \frac{x}{k^2} + O\Big(\frac{x}{k^2}\Big)\right\} \\ &= \sum_{1 \leq k \leq x^{1/2}} \frac{x \mu(k) \log x}{k^2} + O(x) \\ &= \frac{x \log x}{\zeta(2)} + O(x) \\ &= cx \log x + O(x) \;. \end{split}$$

(2) We have

$$egin{aligned} A(2x) &= \sum\limits_{\substack{1 < N \leq 2x \ 2 \mid N}} 2^{\omega(N)} \ &= \sum\limits_{\substack{1 \leq 1/2N < x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} + \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \ 2 \mid 1/2N}} 2^{\omega(2 \cdot 1/2N)} \ &= \sum\limits_{\substack{1 \leq 1/2N \leq x \$$

$$= \sum_{\substack{1 \le 1/2N \le x \\ 2|1/2N}} 2^{\omega(1/2N)} + \sum_{\substack{1 \le 1/2N \le x \\ 2|1/2N}} 2^{1+\omega(1/2N)}$$
  
=  $A(x) + 2B(x)$ .

Thus A(2x) + A(x) = 2A(x) + 2B(x) = 2F(x). We now prove by induction that

$$A(x) = 2 \sum_{r=1}^{\infty} (-1)^{r-1} F(x \cdot 2^{-r}) .$$

For, if x = 1, the result is clearly true since both sides vanish, and then if true for  $x \leq x_0$ , we have for  $x \leq 2x_0$ ,

$$A(x) = 2F\left(\frac{1}{2}x\right) - A\left(\frac{1}{2}x\right)$$

which is again of the required form, and this completes the induction. Now F(y) = 0 if y < 1 and so we have

$$A(x) = 2\sum_{r=1}^{k} (-1)^{r-1} F(x \cdot 2^{-r})$$
,

where

$$k = \left[\frac{\log x}{\log 2}\right].$$

Now by (1)

$$|F(y) - cy \log y| < Cy,$$

for some constant C and all y > 1. Thus

$$\left|A(x) - 2c\sum\limits_{r=1}^k (-1)^{r-1} rac{x}{2^r} {\cdot} \log rac{x}{2^r} 
ight| < 2C\sum\limits_{r=1}^k rac{x}{2^r} < 2Cx \; .$$

Hence

$$ig| A(x) - 2c \sum\limits_{r=1}^k (-1)^{r-1} rac{x}{2^r} \log x ig| < 2Cx + 2cx \log 2 \cdot \sum\limits_{r=1}^k r \cdot 2^{-r} < C_1 x \; .$$

Finally,

$$\begin{split} \sum_{r=1}^{k} (-1)^{r-1} \frac{x}{2^{r}} \log x &= \frac{1}{2} x \log x \cdot \frac{1 - \left(-\frac{1}{2}\right)^{k}}{1 - \left(-\frac{1}{2}\right)} \\ &= \frac{1}{3} x \log x \{1 + O(x^{-1})\} \\ &= \frac{1}{3} x \log x + O(\log x) , \end{split}$$

and so (2) follows.

- (3) now follows since B(x) = F(x) A(x).
- (4) follows since

$$C(x) = \sum_{\substack{1 < 1/2N \leq 1/2x \\ 21/12N \\ 21/12N}} 2^{\omega(2 \cdot 1/2N)} = A\left(\frac{1}{2}x\right).$$

(5) and (6) now follow similarly since D(x) = C(1/2x) and E(x) = D(1/2x).

*Proof of Theorem* 1. The idea of the proof is to combine the results of Lemmas 3-6. We have immediately that

$$p(d) \leq \sum_{1 \leq N \leq 2d^{1/2}} 2^{\omega(N)} = cd^{1/2} \log d + O(d^{1/2})$$

and the remainder of the proof deals with reducing the constant in the above. There are two ways of doing this; in the first place if 2||N, then the upper bound  $2^{\omega(N)}$  appearing above can immediately be halved in view of Lemmas 3 and 4; secondly depending upon the value of d, there are certain residue classes modulo 16 such that for any N belonging to one of them, the equation  $x^2 - dy^2 = N$  cannot have any primitive solutions at all. In each case, it is not possible to dispose of all the odd values of N in this way, and corresponding to these we always obtain a term

$$\sum_{\substack{1 \leq N \leq 2d^{1/2} \\ 2 
eq N}} 2^{\omega(N)} = B(2d^{1/2})$$
 .

There are various cases to consider.

(a)  $d \equiv 1 \pmod{8}$ . In this case, since x and y cannot both be even, we find that  $x^2 - dy^2 = N$  is either odd or divisible by 8. Thus we find that  $p(d) \leq B(2d^{1/2}) + D(2d^{1/2}) = 1/2cd^{1/2}\log d + O(d^{1/2})$ , as required.

(b)  $d \equiv 5 \pmod{8}$ . In this case, we find that if N is even, then  $2^2 || N$ , and accordingly

$$p(d) \leq B(2d^{_{1/2}}) + C(2d^{_{1/2}}) - D(2d^{_{1/2}}) = rac{1}{2}cd^{_{1/2}}\log d + O(d^{_{1/2}})$$
 .

(c) If  $d \equiv 2$  or  $3 \pmod{4}$  then N can be even only if 2 || N and we obtain

$$egin{aligned} p(d) &\leq B(2d^{1/2}) + \sum\limits_{\substack{1 < N \leq 2d^{1-2} \ 2 \mid |N|}} 2^{\omega(N)-1} \ &= B(2d^{1/2}) + rac{1}{2} \{A(2d^{1/2}) - C(2d^{1/2})\} \ &= rac{1}{2} c d^{1/2} \log d + O(d^{1/2}) \;. \end{aligned}$$

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It is to be noted for future reference that if  $4 \nmid d$ , then the 7c/12 of the theorem can be improved to 1/2c.

(d) If  $d \equiv 0 \pmod{4}$ , then for a primitive solution of  $x^2 - dy^2 = N$ we must have either that x is odd, in which case N is also odd, or else x is even, y odd and 4|N. In the latter case we find that  $(1/2x)^2 - (1/4d)y^2 = 1/4N$  and so we obtain a primitive solution of the equation  $X^2 - (1/4d)Y^2 = 1/4N$ , in which moreover y is odd. Thus we have

either  $1/4d \equiv 0$  or  $1 \pmod{4}$  in which case 1/4N is odd or divisible by 4,

or  $1/4d \equiv 2$  or  $3 \pmod{4}$  in which case 1/4N is odd or 2 || 1/4N.

In the first case we obtain

$$egin{aligned} p(d) &\leq B(2d^{1/2}) + C(2d^{1/2}) - D(2d^{1/2}) + E(2d^{1/2}) \ &= rac{7}{12} c d^{1/2} \log d \, + \, O(d^{1/2}) \; , \end{aligned}$$

and in the second case we obtain similarly

$$egin{aligned} p(d) &\leq B(2d^{1/2}) + C(2d^{1/2}) - E(2d^{1/2}) \ &= rac{7}{12} c d^{1/2} \log d \, + \, O(d^{1/2}) \; , \end{aligned}$$

which concludes the proof.

LEMMA 7. As  $x \rightarrow \infty$ ,

$$F_1(x) = \sum_{1 \le N \le x} 2^{\omega(N)} \log \frac{x}{N} = cx \log x + O(x)$$
.

*Proof.* Let  $1 < \rho < x$ ; then

$$F_{1}(x) - F_{1}(x\rho^{-1}) = \sum_{1 \le N \le x} 2^{\omega(N)} \log \frac{x}{N} - \sum_{1 \le N \le x\rho^{-1}} 2^{\omega(N)} \log \frac{x}{\rho N}$$
$$= \sum_{1 \le N \le x\rho^{-1}} 2^{\omega(N)} \log \rho + \sum_{x\rho^{-1} < N \le x} 2^{\omega(N)} \log \frac{x}{N}$$

and so

$$\log \rho \cdot F(x\rho^{-1}) \leq F_1(x) - F_1(x\rho^{-1}) \leq \log \rho \cdot F(x) ,$$

since  $x/N < \rho$  for  $N > x \rho^{-1}$ .

Thus if  $1 < \rho^n \leq x < \rho^{n+1}$ , we find that

$$\log \rho \cdot \sum_{r=1}^{n} F(x \rho^{-r}) \leq F_1(x) - F_1(x \rho^{-n}) \leq \log \rho \cdot \sum_{r=0}^{n-1} F(x \rho^{-r}),$$

and so to complete the proof it suffices to show that

$$\log \rho \cdot \sum_{0}^{n-1} F(x \rho^{-r}) \longrightarrow cx \log x + O(x) \quad \text{as} \quad \rho \longrightarrow 1+ ,$$

where  $n = [(\log x / \log \rho)]$ .

Now for all y > 1, we have for some constant A,

$$cy \log y - Ay < F(y) < cy \log y + Ay$$
.

Thus

$$egin{aligned} \log 
ho \sum\limits_{0}^{n-1} F(x 
ho^{-r}) &< \log 
ho \sum\limits_{0}^{n-1} (cx \log x + Ax) 
ho^{-1} \ &< 
ho rac{\log 
ho}{
ho - 1} (cx \log x + Ax) \longrightarrow cx \log x + Ax \ & ext{ as } 
ho \longrightarrow 1+ \ . \end{aligned}$$

On the other hand

$$\begin{split} \log \rho \sum_{0}^{n-1} F(x \rho^{-r}) &> \log \rho \sum_{0}^{n-1} (cx \log x - cxr \log \rho - Ax) \rho^{-r} \\ &= \log \rho \cdot (cx \log x - Ax) \sum_{0}^{n-1} \rho^{-r} \\ &- cx (\log \rho)^2 \sum_{0}^{n-1} r \rho^{-r} \\ &= X - Y , \quad \text{say} . \end{split}$$

Now

$$X = \frac{\rho(cx\log x - Ax)\log\rho}{\rho - 1} \left\{ 1 - \frac{1}{\rho^n} \right\} \longrightarrow (cx\log x - Ax)(1 - x^{-1})$$

as  $\rho \to 1$ , since x lies between  $\rho^n$  and  $\rho^{n+1}$ . Also

$$Y < cx(\log \rho)^2 \sum_{0}^{\infty} r \rho^{-r} = \rho^2 cx \left\{ \frac{\log \rho}{\rho - 1} \right\}^2 \longrightarrow cx \quad \text{as} \quad \rho \longrightarrow 1 +$$

and so the result follows.

LEMMA 8. Let

$$A_{1}(x) = \sum_{\substack{1 \leq N \leq x \\ 2 \mid N}} 2^{\omega(N)} \log rac{x}{N}$$

with analogous definitions for  $B_1$ ,  $C_1$ , and  $D_1$ . Then the results of Lemma 6, (2)-(5) hold also for the functions  $A_1$  etc.

*Proof.* These results follow from Lemma 7 in exactly the same way as the corresponding results follow from Lemma 6(1).

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Proof of Theorem 2. We have for each convergent

$$\left| \, d^{\scriptscriptstyle 1/2} - \, rac{P_r}{Q_r} \, 
ight| < rac{1}{Q_r Q_{r+1}}$$
 ,

whence

$$rac{Q_{r+1}}{Q_r}\!<\!rac{1}{Q_r|P_r\!-\!Q_rd^{1/2}|}=\!rac{d^{1/2}+rac{P_r}{Q_r}}{|P_r^2-dQ_r^2|}\!<\!rac{2d^{1/2}+1}{N_r}$$
 ,

where

$$|P_r^2 - dQ_r^2| = N_r$$

Consider first the case  $N(\varepsilon) = -1$ . Then

$$egin{aligned} arepsilon_1 &= arepsilon^2 &= P_{2p(d)-1} + Q_{2p(d)-1} d^{1/2} \ &< (2d^{1/2}+1)Q_{2p(d)-1} \ &= (2d^{1/2}+1) \prod_0^{2p(d)-2} rac{Q_{r+1}}{Q_r} \ &< (2d^{1/2}+1) \prod_0^{2p(d)-2} rac{2d^{1/2}+1}{N_r} \ &= \prod_0^{2p(d)-1} rac{2d^{1/2}+1}{N_r} \ . \end{aligned}$$

Thus

$$egin{aligned} 2\logarepsilon &< \sum\limits_{0}^{2p(d)-1}\lograc{2d^{1/2}\,+\,1}{N_r} \ &\leq \sum\limits_{0< N< 2d^{1/2}}\{f(N;\,d)\,+\,f(-N;\,d)\}\lograc{2d^{1/2}\,+\,1}{N} \ &= \sum\limits_{0< N< 2d^{1/2}}\{f(N;\,d)\,+\,f(-N;\,d)\}\lograc{2d^{1/2}}{N}\,+\,O\{d^{-1/2}F(2d^{1/2})\} \ &= 2\sum\limits_{0< N< 2d^{1/2}}f(N;\,d)\lograc{2d^{1/2}}{N}\,+\,O(\log d) \;, \end{aligned}$$

since in this case f(N; d) = f(-N; d). Thus

$$egin{aligned} \log arepsilon &< \sum\limits_{0 < N < 2d^{1/2}} f(N;\,d) \log rac{2d^{1/2}}{N} + O(\log d) \ &< rac{1}{2} c d^{1/2} \log d \, + \, O(d^{1/2}) \; , \end{aligned}$$

as before, using Lemmas 7 and 8 in place of Lemma 6, since in this case  $4 \nmid d$ . In the case  $N(\varepsilon) = +1$ , we have

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$$egin{aligned} arepsilon &= P_{p(d)-1} + Q_{p(d)-1} d^{1/2} \ &< (2d^{1/2}+1)Q_{p(d)-1} \ &< \prod_{0}^{p(d)-1} rac{2d^{1/2}+1}{N_r} \ , \end{aligned}$$

as before.

Thus

$$\begin{split} \log \varepsilon &< \sum_{0}^{p(d)-1} \log \frac{2d^{1/2}+1}{N_r} \\ &\leq \sum_{0 < N < 2d^{1/2}} \{f(N;\,d) + f(-N;\,d)\} \log \frac{2d^{1/2}+1}{N} \\ &= \sum_{0 < N < 2d^{1/2}} \{f(N;\,d) + f(-N;\,d)\} \log \frac{2d^{1/2}}{N} + O(\log d) \\ &\leq \frac{1}{2} c d^{1/2} \log d + O(d^{1/2}) \;, \end{split}$$

as before, provided  $4 \nmid d$ .

Finally if 4|d we observe that  $\varepsilon = \eta$  or  $\eta^2$  where  $\eta$  is the fundamental unit of the ring  $Z[((1/4)d)^{1/2}]$ . Then the result for this case follows by descent since now  $\log \varepsilon \leq 2 \log \eta$ .

This concludes the proof of Theorem 2.

Proof of Theorem 3. We have as before

$$\logarepsilon < \sum_{r=0}^{p(d)-1} \log rac{2d^{1/2}+1}{N_r}$$

and so for any K satisfying  $1 < K < 2d^{1/2}$ 

$$\begin{split} \log \varepsilon &< \sum_{r=0}^{p(d)-1} \log \frac{2d^{1/2}}{N_r} + O(\log d) \\ &= \sum_{N_r \leq K \atop 0 \leq r < p(d)} \log \frac{2d^{1/2}}{N_r} + \sum_{0 \leq r < p(d)} \log \frac{2d^{1/2}}{N_r} + O(\log d) \\ &< \sum_{1 \leq N \leq K} \{f(N; d) + f(-N; d)\} \log 2d^{1/2} \\ &+ p(d) \log \frac{2d^{1/2}}{K} + O(\log d) \\ &< A \log d \cdot K \log K + \frac{1}{2} p(d) \log (4dK^{-2}) + O(K \log d) \:. \end{split}$$

In particular taking  $K = 2d^{1/2}(\log d)^{-3}$  we obtain

$$\logarepsilon < 3p(d)\log\log d + o(d^{1/2})$$
 .

Now for  $d = 2^{2k+1}$  we have  $\varepsilon = (1 + \sqrt{2})^{2^k}$ , i.e.,  $\log \varepsilon > Ad^{1/2}$  where A > 0 and so  $p(d) \neq o(d^{1/2}/\log \log d)$ , as required.

Proof of Theorem 4. If  $\log \varepsilon \neq o(d^{1/2} \log d)$ , then there exists a positive constant  $c_1 < c$  so that for infinitely many values of d,  $\log \varepsilon > c_1 d^{1/2} \log d$ . Let g(N; d) denote the number of distinct primitive classes of solutions of  $x^2 - dy^2 = N$  for which x/y occurs as a convergent to the continued fraction for  $d^{1/2}$ . Then

$$2p(d) \ge \sum_{-2d^{1/2} < N < 2d^{1/2}} g(N; d)$$

and

$$\log \varepsilon < \sum_{-2d^{1/2} < N < 2d^{1/2}} g(N; d) \log \frac{2d^{1/2}}{|N|} + O(\log d) \;.$$

Thus if  $k \geq 1$ ,

$$egin{aligned} \logarepsilon & -2p(d)\log k <_{2d^{1/2} < N < 2d^{1/2}} g(N;\,d)\lograc{2d^{1/2}}{k\,|\,N|} + O(\log d) \ & \leq \sum\limits_{0 < |N| < 2d^{1/2}k^{-1}} g(N;\,d)\lograc{2d^{1/2}}{k|\,N|} + O(\log d) \ & \leq \sum\limits_{0 < N < 2d^{1/2}k^{-1}} 2^{\omega(N)}\lograc{2d^{1/2}k^{-1}}{N} + O(\log d) \end{aligned}$$

since  $g(N; d) \leq f(N; d)$ . Thus

$$egin{aligned} \log arepsilon &- 2p(d) \log k < F_1(2d^{1/2}k^{-1}) + O(\log d) \ &< cd^{1/2}k^{-1}\log d + O(d^{1/2}) \,. \end{aligned}$$

Thus if  $k > c/c_1$ , we have for infinitely many values of d,

$$p(d) > rac{kc_1-c}{2k\log k} d^{1/2}\log d + O(d^{1/2})$$
 ,

as required.

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