We shall now say that P is <u>reducible</u> to Q if

$$P(\vec{\alpha}, \vec{x}) \longmapsto Q(\lambda y G_1(y, \vec{\alpha}, \vec{x}), \dots, \lambda y G_m(y, \vec{\alpha}, \vec{x}), F_1(\vec{\alpha}, \vec{x}), \dots, F_k(\vec{\alpha}, \vec{x}))$$

where  $G_1, ..., G_m, F_1, ..., F_k$  are total and recursive.

19.3. PROPOSITION. If P is  $\Pi_n^1$  and Q is reducible to P, then P is  $\Pi_n^1$ ; and similarly with  $\Sigma_n^1$  or  $\Delta_n^1$  in place of  $\Pi_n^1$ .  $\Box$ 

The analogue of the table in §12 is the following table.

P,Q	$\neg P$	$P \lor Q$	P & Q	$\forall \alpha P$	$\exists \alpha P$	QxP
$\Pi^1_n$	$\Sigma_n^1$	$\Pi_n^1$	$\Pi^1_n$	$\Pi_n^1$	$\Sigma_{n+1}^1$	$\Pi_n^1$
$\Sigma_n^1$	$\Pi_n^1$		$\Sigma_n^1$		$\Sigma_n^1$	
$\Delta_n^1$	$\Delta_{n}^{1}$	$\Delta_n^1$	$\Delta_n^1$	$\Pi^1_n$		$\Delta_n^1$

It is proved and used in the same way as the earlier table.

The classification of analytical relations into the  $\Pi_n^1$  and  $\Sigma_n^1$  relations is called the <u>analytical hierarchy</u>.

19.4. ANALYTICAL ENUMERATION THEOREM. For every n, m, and k, there is a  $\Pi_n^1(m,k+1)$ —ary function which enumerates the class of  $\Pi_n^1(m,k)$ —ary relations; and similarly with  $\Sigma_n^1$  for  $\Pi_n^1$ .

**Proof.** Suppose, for example, we want to enumerate the  $\Pi_2^1$ (1,1)-ary relations. Every such relation R is of the form  $\forall \alpha \exists \beta P$  where P is  $\Pi_1^0$ by the remarks after 19.1. Thus if Q is  $\Pi_1^0$  and enumerates the  $\Pi_1^0$  (3,1)-ary relations, then  $\forall \alpha \exists \beta Q(\alpha, \beta, \gamma, x, e)$  is the desired enumerating function.  $\Box$ 

19.5. ANALYTICAL HIERARCHY THEOREM. For each *n*, there is a  $\Pi_n^1$  set which is not  $\Sigma_n^1$ , hence not  $\Pi_k^1$  or  $\Sigma_k^1$  for any k < n. The same holds with  $\Pi_n^1$  and  $\Sigma_n^1$  interchanged.

*Proof.* As in the arithmetical case.  $\Box$ 

## 20. The Projective Hierarchy

The results of the last section can be relativized to a class  $\Phi$  of total functions of number variables. A particularly interesting case is that in which  $\Phi$ 

is the class of all such functions. Replacing the functions by their contractions, we see we are relativizing to the class  $\mathbb{R}$  of reals. Note that by 18.1, a function is recursive in  $\mathbb{R}$  iff it is obtained from a recursive function by replacing some of the unary function variables by names of particular reals. The same then holds with *recursive* replaced by  $\Pi_k^1$  or  $\Sigma_k^1$ .

A relation is <u>projective</u> if it is analytical in  $\mathbb{R}$ . The analytical hierarchy relativized to  $\mathbb{R}$  is called the <u>projective hierarchy</u>. (It is customary to write a boldface  $\Pi_n^1$  for  $\Pi_n^1$  in  $\mathbb{R}$  and similarly for  $\Sigma$  and  $\Delta$ . We avoid this notation, since boldface is sometimes hard to distinguish from lightface.) The theory of the projective hierarchy antedates that of the analylytical hierearchy; it was begun by Lusin, Suslin, and Sierpinski.

The Enumeration Theorem does not hold in its usual form for  $\Pi_n^1$  in **R**; but we shall prove a modified form. We say that a (m+1,k)-ary relation Q<u>**R**-enumerates</u> a class  $\Phi$  of (m,k)-ary relations if for every R in  $\Phi$ , there is a  $\beta$ . such that  $R(\vec{\alpha}, \vec{x}) \leftrightarrow Q(\vec{\alpha}, \vec{x}, \beta)$  for all  $\vec{\alpha}$  and  $\vec{x}$ .

20.1. PROJECTIVE ENUMERATION THEOREM. For every n, m, and k, there is a (m+1,k)-ary  $\Pi_n^1$  relation which **R**-enumerates the class of (m,k)-ary  $\Pi_n^1$  in **R** relations; and similarly with  $\Sigma_n^1$  for  $\Pi_n^1$ .

**Proof.** As in the proof of the analytical case, it is enough to do this for  $\Sigma_1^0$ , i.e., RE. If R is RE in **R**, it is RE in a finite sequence  $\vec{\alpha}$  of reals. If e is a  $\vec{\alpha}$ -index of R, then by (3) of §18,

$$\begin{aligned} R(\vec{\beta}, \vec{x}) & \longleftrightarrow \{e\}^{\vec{\alpha}}(\vec{\beta}, \vec{x}) \text{ is defined} \\ & \longleftrightarrow \{e\}(\vec{\beta}, \vec{\alpha}, \vec{x}) \text{ is defined} \\ & \longleftrightarrow W_e(\vec{\beta}, \vec{\alpha}, \vec{x}). \end{aligned}$$

Choose  $\gamma$  so that  $(\gamma)_0(0) = e$  and  $(\gamma)_i = \alpha_i$  for  $1 \le i \le n$ . Then the right side becomes  $W_{(\gamma)_0(0)}(\vec{\beta}, (\gamma)_1, ..., (\gamma)_m, \vec{x})$ . This is an RE relation P of  $\vec{\beta}, \vec{x}, \gamma$ ; and P is the desired enumerating relation.  $\Box$ 

We leave it to the reader to derive a Projective Hierarchy Theorem from this; the examples will now be (1,0)-ary. (Every (0,k)-ary relation is recursive in  $\mathbb{R}$ .)

20.2. PROPOSITION. Let P be defined by  $P(\vec{\alpha}, \vec{x}, y) \leftrightarrow P_y(\vec{\alpha}, \vec{x})$ . Then P is  $\Pi_n^1$  in R iff each  $P_y$  is  $\Pi_n^1$  in R; and similarly with  $\Sigma_n^1$  or  $\Delta_n^1$  in place of  $\Pi_n^1$ .

Proof. If P is  $\Pi_n^1$  in R, each  $P_y$  is clearly  $\Pi_n^1$  in R. Now suppose that each  $P_y$  is  $\Pi_n^1$  in R. By the Projective Enumeration Theorem, there is a  $\Pi_n^1$ relation Q and a  $\beta_y$  for each y such that  $P_y(\vec{\alpha}, \vec{x}) \leftrightarrow Q(\vec{\alpha}, \vec{x}, \beta_y)$ . Choose  $\beta$  so that  $(\beta)_y = \beta_y$  for all y. Then  $P(\vec{\alpha}, \vec{x}, y) \leftrightarrow Q(\vec{\alpha}, \vec{x}, (\beta)_y)$ . Thus P is  $\Pi_n^1$  in  $\beta$  and hence in R.  $\Box$ 

The further study of the analytical and projective hierarchies is known as Descriptive Set Theory, and is a hybrid of Recursion Theory and Set Theory. We shall prove only one result. We shall prove it for the projective hierarchy; the analogue for the analytical hierarchy is more difficult both to state and to prove.

We recall a definition from measure theory. Let X be a space and let  $\Lambda$  be a class of subsets of X. We say that  $\Lambda$  is a  $\sigma$ -ring if: (a) the complement of every set in  $\Lambda$  is in  $\Lambda$ ; (b) every countable union of sets in  $\Lambda$  is in  $\Lambda$ ; (c)  $X \in \Lambda$ . From (a) and (b) it follows that: (d) every countable intersection of sets in  $\Lambda$  is in  $\Lambda$ . If  $\Gamma$  is any collection of subsets of X, there is a smallest  $\sigma$ -ring including  $\Gamma$ ; it is the intersection of all of the  $\Sigma$ -rings which include  $\Gamma$ .

20.3 PROPOSITION. The class of  $\Delta_n^1$  in  $\mathbb{R}$  (m,k)-ary relations is a  $\sigma$ -ring in  $\mathbb{R}^{m,k}$ .

*Proof.* In view of the table, it is enough to show that the union Q of a sequence  $\{P_j\}$  of such relations is  $\Delta_n^1$  in  $\mathbb{R}$ . Defining  $P(\vec{\alpha}, \vec{x}, j) \leftrightarrow P_j(\vec{\alpha}, \vec{x})$ , P is  $\Delta_n^1$  by 20.2. Since  $Q(\vec{\alpha}, \vec{x}) \leftrightarrow \exists j P(\vec{\alpha}, \vec{x}, j)$ , Q is  $\Delta_n^1$  in  $\mathbb{R}$  by the table.  $\Box$ 

An (m,k)-ary relation is <u>Borel</u> if it belongs to the smallest  $\sigma$ -ring in  $\mathbb{R}^{m,k}$ which contains all the recursive (m,k)-ary relations. By 20.3, every Borel relation is  $\Delta_1^1$  in **R**. We shall prove that the converse also holds.

Let A and B be subsets of a space X. We say that a subset C of X <u>separates</u> A and B if  $A \subseteq C$  and  $B \subseteq C^{\mathbb{C}}$ . This clearly implies that A and B are disjoint.

20.4. SEPARATION THEOREM. Any two disjoint  $\Sigma_1^1$  in  $\mathbb{R}$  (m,k)—ary relations can be separated by a Borel relation.

Proof. To make the notation simpler, let m = 1 and k = 0. Say that A is <u>inseparable</u> from B if no Borel relation separates A and B. We shall first prove the following lemma: If  $\bigcup_{i \in \omega} A_i$  is inseparable from  $\bigcup_{j \in \omega} B_j$  then there are i and j such that  $A_i$  is inseparable from  $B_j$ . Suppose, on the contrary, that for every i and j, there is a Borel relation  $C_{i,j}$  which separates  $A_i$  and  $B_j$ . If  $C = \bigcap_{j \in \omega} \bigcup_{i \in \omega} C_{i,j}$  then C is Borel and separates  $\bigcup_{i \in \omega} A_i$  and  $\bigcup_{j \in \omega} B_j$ 

Now assume that P and Q are inseparable  $\Sigma_1^1$  in  $\mathbb{R}$  relations; we shall show that P and Q are not disjoint. Using the remarks after 19.1, we can write

$$P(\alpha) \longmapsto \exists \beta \forall n R(\overline{\alpha}(n), \overline{\beta}(n)),$$
$$Q(\alpha) \longmapsto \exists \gamma \forall n R'(\overline{\alpha}(n), \overline{\gamma}(n)),$$

where R and R' are recursive in **R**. For  $z, w \in Seq$ , let

$$P_{z,w}(\alpha) \longmapsto z = \overline{\alpha}(lh(z)) \& \exists \beta(w = \overline{\beta}(lh(w)) \& \forall nR(\overline{\alpha}(n),\overline{\beta}(n))),$$

and define  $Q_{z,w}$  similarly but with R replaced by R'. It is clear that

$$P_{z,w} = \bigcup_{m \in \omega} \bigcup_{p \in \omega} P_{z_*} < m >, w_*$$

and similarly for  $Q_{z,w}$ 

We shall define  $\alpha(n)$ ,  $\beta(n)$ , and  $\gamma(n)$  by induction on n so that  $P_{\overline{\alpha}(n),\overline{\beta}(n)}$ and  $Q_{\overline{\alpha}(n),\overline{\gamma}(n)}$  are inseparable. Since  $P = P_{<>,<>}$  and  $Q = Q_{<>,<>}$ , this holds for n = 0. Suppose it holds for some n. By our lemma, there are i, j, k, and l so that  $P_{\overline{\alpha}(n),<i>,\overline{\beta}(n),<j>}$  is inseparable from  $Q_{\overline{\alpha}(n),<k>,\overline{\gamma}(n),<l>}$ . Then i = k; for otherwise,  $\{\delta: \overline{\delta}(n+1) = \overline{\alpha}(n)^* < i>\}$  is a recursive (and hence Borel) set which separates  $P_{\overline{\alpha}(n),<i>,\overline{\beta}(n),<j>}$  from  $Q_{\overline{\alpha}(n),<k>,\overline{\gamma}(n),<l>}$ . Thus we may take  $\alpha(n) = i$ ,  $\beta(n) = j$ , and  $\gamma(n) = l$ .

For each *n*,  $P_{\overline{\alpha}(n),\overline{\beta}(n)}$  is inseparable from  $Q_{\overline{\alpha}(n),\overline{\gamma}(n)}$ ; so, they are both non-empty. This implies that  $R(\overline{\alpha}(n),\overline{\beta}(n))$  and  $R'(\overline{\alpha}(n),\overline{\gamma}(n))$  for all *n*. Hence  $P(\alpha)$  and  $Q(\alpha)$ ; so *P* and *Q* are not disjoint.  $\Box$ 

20.5. Suslin's Theorem. A relation is Borel iff it is  $\Delta_1^1$  in R.

*Proof.* We have already seen that every Borel relation is  $\Delta_1^1$  in  $\mathbb{R}$ . Now let P be  $\Delta_1^1$  in  $\mathbb{R}$ . Then P and  $\neg P$  are  $\Sigma_1^1$  in  $\mathbb{R}$ ; so by the Separation Theorem, there is a  $\Delta_1^1$  in  $\mathbb{R}$  relation which separates P and  $\neg P$ . But the only relation which separates P and  $\neg P$  is P.  $\Box$