#### 7. DEGREES OF INTERPRETABILITY

Suppose PA $\dashv$  T. We shall use A, B, etc. for extensions of T. (Thus, T, A, B, etc. are essentially reflexive.) The relation  $\leq$  of interpretability is reflexive and transitive. Thus, the relation  $\equiv$  of mutual interpretability (restricted to extensions of T) is an equivalence relation; its equivalence classes will be called *degrees* (of interpretability) and will be written a, b, c, etc.  $D_T$  is the set of degrees of extensions of T. A is of degree a if  $A \in A$  and  $A \in A$  and  $A \in A$  and the relation  $A \in A$  and

§1. Algebraic properties. In this § we restrict ourselves to purely algebraic properties of  $D_T$ . First we define the theory  $A^T$  and the operations  $\downarrow$  and  $\uparrow$  on theories as follows.

$$A^{T} = T + \{Con_{A \mid k}: k \in N\},\$$
  
 $A \downarrow B = T + \{Con_{A \mid k} \lor Con_{B \mid k}: k \in N\},\$   
 $A \uparrow B = T + \{Con_{A \mid k} \land Con_{B \mid k}: k \in N\}.$ 

From Lemma 6.2 and Theorem 6.6, we get the following:

**Lemma 1.** (a)  $A \le B$  iff  $A^T \dashv B$ . Thus,  $A^T \equiv A$  and  $A \le B$  iff  $A^T \dashv B^T$ . (b)  $A \le B$ , C iff  $A \le B \downarrow C$ ,

(c) A, B  $\leq$  C iff A $\uparrow$ B  $\leq$  C.

The following lemma is little more than a restatement of Lemma 4.4.

**Lemma 2.** If  $\theta$  is  $\Pi_1$  and  $A \vdash \theta$ , there is a k such that  $PA \vdash Con_{A \mid k} \to \theta$ .

Instead of  $A \downarrow B$  it is sometimes convenient to use the theory  $A \lor B$  defined by  $A \lor B = \{ \phi \lor \psi \colon \phi \in A \& \psi \in B \}.$ 

Th(A $\lor$ B) = Th(A)  $\cap$  Th(B). Evidently, A $\downarrow$ B  $\dashv$  A $\lor$ B and, by Lemma 2, A $\lor$ B  $\dashv_{\Pi_1}$ A $\downarrow$ B. But then, by Theorem 6.6, that A $\lor$ B  $\leq$  A $\downarrow$ B and so A $\lor$ B  $\equiv$  A $\downarrow$ B. It follows that for every sentence  $\varphi$ , (A +  $\varphi$ ) $\downarrow$ (A +  $\neg \varphi$ )  $\leq$  A.

From Lemma 2 and Lemma 6.1 we get:

**Lemma 3.** For every  $\Pi_1$  sentence  $\pi$ ,  $T + \pi \le A \uparrow B$  iff  $A \uparrow B \vdash \pi$  iff there are  $\Pi_1$  sentences  $\varphi$ ,  $\psi$  such that  $A \vdash \varphi$ ,  $B \vdash \psi$ , and  $T + \varphi \land \psi \vdash \pi$ .

For A  $\in$  a and B  $\in$  b, let a  $\cap$  b = d(A $\downarrow$ B) and a  $\cup$  b = d(A $\uparrow$ B). By Lemma 1,  $\cap$  and  $\cup$ 

are well–defined,  $a \cap b$  is the g.l.b. of a and b and  $a \cup b$  is l.u.b. of a and b. Thus, we have proved part of the following:

### **Theorem 1.** $\mathbf{D}_{\mathsf{T}}$ is a distributive lattice.

To prove distributivity we need the following lemma whose proof is left to the reader.

**Lemma 4.** (a) For every k, there is an m such that

$$PA \vdash Con_{(A \lor B) \mid m} \rightarrow Con_{A \mid k} \lor Con_{B \mid k}$$

(b) For every k, there is an m such that

$$PA \vdash Con_{A \mid m} \lor Con_{B \mid m} \rightarrow Con_{(A \lor B) \mid k}$$

**Proof of Theorem 1.** Let  $D = A^T \lor (B \uparrow C)$  and  $E = (A \lor B) \uparrow (A \lor C)$ . To prove that  $D_T$  is distributive, it is, by Lemma 1, sufficient to show that  $D \dashv F$ .

Let k be arbitrary. By Lemma 4 (a), there is an m such that

$$PA \vdash Con_{(A \lor B) \mid m} \rightarrow Con_{A \mid k} \lor Con_{B \mid k}$$

$$PA \vdash Con_{(A \lor C) \mid m} \rightarrow Con_{A \mid k} \lor Con_{C \mid k}$$

But then

$$E \vdash Con_{A \mid k} \lor (Con_{B \mid k} \land Con_{C \mid k}).$$

It follows that D $\dashv$  E. The proof that D $\vdash$  E is similar.

 $\mathbf{D}_{\mathrm{T}}$  has a minimal element  $0_{\mathrm{T}} = \mathrm{d}(\mathrm{T})$  and a maximal element  $1_{\mathrm{T}}$ , the common degree of all inconsistent theories.

In our next result we answer a number of standard questions concerning  $D_T$ ; in particular, it follows that  $D_T$  is dense.

**Theorem 2.** Suppose  $a < b < 1_T$ ,  $d_0 \not \le a$ , and  $b \not \le d_1$ . There are then degrees  $c_0$ ,  $c_1$  such that  $a < c_i < b$ ,  $d_0 \not \le c_i \not \le d_1$ ,  $i = 0, 1, c_0 \cap c_1 = a$ , and  $c_0 \cup c_1 = b$ .

We derive this from:

**Lemma 5.** Suppose X is r.e. and monoconsistent with PA. Let  $\theta$  be any true  $\Pi_1$  sentence. There are then  $\Pi_1$  sentences  $\theta_0$ ,  $\theta_1$  such that

- (i) PA $\vdash \theta_0 \lor \theta_1$ ,
- (ii)  $PA \vdash \theta_0 \land \theta_1 \rightarrow \theta$ ,
- (iii)  $\theta_i^J \notin X$ , i, j = 0, 1.

**Proof.** We may assume that if  $\varphi \in X$  and PAF  $\varphi \to \psi$ , then  $\psi \in X$ . Let  $\theta := \forall x \gamma(x)$ , where  $\gamma(x)$  is PR. Let R(k,m) be a primitive recursive relation such that  $X = \{k: \exists mR(k,m)\}$  and let  $\rho(x,y)$  be a PR binumeration of R(k,m). Finally, let  $\theta_0$  and  $\theta_1$  be such that

$$(1) \qquad PA \vdash \theta_0 \leftrightarrow \forall y \big( (\rho(\theta_0, y) \vee \neg \gamma(y)) \to \exists z \leq y \rho(\theta_1, z) \big),$$

(2)  $PA \vdash \theta_1 \leftrightarrow \forall z (\rho(\theta_1, z) \to \exists y < z (\rho(\theta_0, y) \lor \neg \gamma(y))).$ 

Then (i) and (ii) follow directly (cf. Lemma 1.3).

Suppose  $\theta_0 \in X$  or  $\theta_1 \in X$  and let m be the least number such that  $R(\theta_0,m)$  or  $R(\theta_1,m)$ . If  $R(\theta_1,m)$ , then  $\theta_1 \in X$ . Also, by (2), and since  $\theta$  is true,  $PA \vdash \neg \theta_1$ , a contradiction. It follows that not  $R(\theta_1,m)$  and, therefore,  $R(\theta_0,m)$ . But then  $\theta_0 \in X$  and, by (1),  $PA \vdash \neg \theta_0$ , again a contradiction. Thus,  $\theta_0 \notin X$  and  $\theta_1 \notin X$ .

Finally, if  $\neg \theta_i \in X$ , then, by (i),  $\theta_{1-i} \in X$ . It follows that  $\neg \theta_0 \notin X$  and  $\neg \theta_1 \notin X$ .

**Proof of Theorem 2.** Let  $A \in a$ ,  $B \in b$ ,  $D_i \in d_i$ . By Orey's compactness theorem (Theorem 6.5) there are sentences  $\psi$ ,  $\chi$  such that  $B \vdash \psi$ ,  $\psi \not\subseteq A$ ,  $D_0 \vdash \chi$ ,  $\chi \not\subseteq A$ . By Theorem 6.6, there is a  $\Pi_1$  sentence  $\pi$  such that  $B \vdash \pi$ ,  $A \not\vdash \pi$ , and  $D_1 \not\vdash \pi$ . Let

$$\begin{split} X_0 &= \{\phi\colon A \vdash \phi \vee \pi\}, & X_1 &= \{\phi\colon \psi \leq A + \neg \phi\}, \\ X_2 &= \{\phi\colon \chi \leq A + \neg \phi\}, & X_3 &= \{\phi\colon D_1 \vdash \phi \vee \pi\}. \end{split}$$

Let  $X = X_0 \cup X_1 \cup X_2 \cup X_3$ . Then X is r.e. (cf. Lemma 6.5). It is also easy to verify that X is monoconsistent with PA. By Lemma 5, there are then  $\Pi_1$  sentences  $\theta_0$ ,  $\theta_1$  such that

- (1)  $PA \vdash \theta_0 \lor \theta_1$ ,
- (2)  $PA \vdash \theta_0 \land \theta_1 \rightarrow Con_B$ ,
- (3)  $\theta_i^J \notin X$ , i, j = 0, 1.

Let  $e_i = d(A + \theta_i)$ , i = 0, 1. Then  $a \le e_i$ .  $b \not\le e_i$ , since  $\neg \theta_i \not\in X_1$ .  $d_0 \not\le e_i$ , since  $\neg \theta_i \not\in X_2$ . Let  $c_i = e_i \cap b$ . Then  $c_i < b$  and  $d_0 \not\le c_i$ . If  $c_i \le a$ , then, since  $\theta_i$  is  $\Pi_1$ ,  $A \vdash \theta_i \lor \pi$ , contradicting the fact that  $\theta_i \not\in X_0$ . Thus,  $c_i \not\le a$  and so  $a < c_i$ . Similarly,  $c_i \not\le d_1$ , since  $\theta_i \not\in X_3$ . By (1),  $c_0 \cap c_1 = a$ . By (2), Theorem 6.4, and Lemma 3,  $e_0 \cup e_1 \ge b$ , whence, by distributivity,  $c_0 \cup c_1 = b \cap (e_0 \cup e_1) = b$ .

From Lemma 5 we can also derive the following:

**Corollary 1.** T is not  $\Sigma_1$ -sound iff there are degrees  $a_0$ ,  $a_1 < 1_T$  such that  $a_0 \cup a_1 = 1_T$  (and  $a_0 \cap a_1 = 0_T$ ).

**Proof.** Suppose T is  $\Sigma_1$ —sound. Let  $a,b < 1_T$ ,  $A \in a$ ,  $B \in b$ . Then  $A \uparrow B$  is consistent and so  $a \cup b < 1_T$ . Next, suppose T is not  $\Sigma_1$ —sound. There is then a true  $\Pi_1$  sentence  $\theta$  such that  $T \vdash \neg \theta$ . Let  $\theta_i$  be as in Lemma 5 with X = Th(T). Let  $a_i = d(T + \theta_i)$ . Then  $a_i < 1_T$ , by Lemma 5 (iii), and  $a_0 \cap a_1 = 0_T$ , by Lemma 5 (i). Finally, by Lemma 3,  $(T + \theta_0) \uparrow (T + \theta_1) \vdash \theta$ . Since  $T \vdash \neg \theta$ , it follows that  $(T + \theta_0) \uparrow (T + \theta_1)$  is inconsistent and so  $a_0 \cup a_1 = 1_T$ .

By Corollary 1, if PA $\dashv$  S and S is  $\Sigma_1$ -sound but T is not, then  $D_S$  and  $D_T$  are not isomorphic. But suppose S and T are both  $\Sigma_1$ -sound. It is an open problem if this implies that  $D_S$  and  $D_T$  are isomorphic.

Given that there are  $c_0$ ,  $c_1 > a$  such that  $c_0 \cap c_1 = a$ , we may ask if any b such that  $a < b < 1_T$  caps to a in the sense that there is a c > a such that  $b \cap c = a$ . (Dually, b cups to a if there is a c < a such that  $b \cup c = a$ .) In our next result this question and its dual are answered in the negative. We write  $a \ll_{\cap} b$  to mean that a < b and b does not cap to a. Dually,  $a \ll_{\circ} b$  means that a < b and a does not cup to b.

**Theorem 3.** (a) Suppose  $0_T < a \not\le c$ . There is a b such that  $0_T < b \ll_0 a$  and  $b \not\le c$ . (b) Suppose  $c \not\le a < 1_T$ . There is a b such that  $a \ll_0 b < 1_T$  and  $c \not\le b$ .

**Proof.** (a) Let  $A \in a$  and  $C \in c$ . There is a  $\Pi_1$  sentence  $\theta$  such that  $A \vdash \theta$  and  $C \dashv \theta$ . Let  $X = Th(C + \neg \theta)$ . X is r.e. and (mono)consistent with  $T + \neg \theta$ . By Theorem 5.2, there is a  $\Pi_1$  sentence  $\psi \not\in X$  such that  $\psi$  is  $\Sigma_1$ —conservative over  $T + \neg \theta$ . Let  $B = T + \psi \vee \theta$  and b = d(B). Then  $0_T < b \not\leq c$  and  $b \leq a$ . Suppose  $b \cup d = a$ . Let  $D \in d$ . Then, by Lemma 6.2, there is an m such that  $T + \psi + Con_{D \mid m} \vdash \theta$  and so  $T + \neg \theta + \psi \vdash \neg Con_{D \mid m}$ . Since  $\psi$  is  $\Sigma_1$ —conservative over  $T + \neg \theta$ , it follows that  $T + \neg \theta \vdash \neg Con_{D \mid m}$  and so  $D \vdash \theta$ . Thus,  $d \geq b$  and so  $d = b \cup d = a$ .  $\bullet$ 

The proof of the following lemma from Lemma 6.2, Theorem 6.6, and Lemma 2 is straightforward.

Lemma 6. The following conditions are equivalent:

- (i)  $A \downarrow B \leq C$ .
- (ii)  $A \le C + \neg Con_{B \mid k}$  for every k.
- (iii)  $A \le C + \neg \theta$  for every  $\Pi_1$  sentence  $\theta$  such that  $B \vdash \theta$ .

Let  $\sigma$  be any  $\Sigma_1$  sentence. By Corollary 6.3, the degree  $d(A+\sigma)$  is uniquely determined by  $\sigma$  and d(A). Thus, we may denote the former by  $d(A) + \sigma$ . A degree of the form  $a + \sigma$  will be called a  $\Sigma_1$ –extension of a. If X is an r.e. set of  $\Sigma_1$  sentences, then, by Theorem 6.11 (b), d(A+X) is a  $\Sigma_1$ –extension of d(A).

Lemma 7. The following conditions are equivalent:

- (i) a  $\ll$  b.
- (ii) a < b and for every  $\Sigma_1$ —extension c of a, if b  $\leq$  c, then c =  $1_T$ .

**Proof.** Suppose (i) holds. Let  $A \in a$  and  $B \in b$ . Let  $\sigma$  be  $\Sigma_1$  and such that  $b \le a + \sigma$ . Then  $B \downarrow (A + \neg \sigma) \le (A + \sigma) \downarrow (A + \neg \sigma) \le A$ . Hence, by assumption,  $A + \neg \sigma \le A$ , whence  $A \vdash \neg \sigma$  and so  $a + \sigma = 1_T$ . Thus, (ii) holds.

Next suppose (ii) holds. Let c be such that  $b \cap c = a$ . Let  $A \in a$ ,  $B \in b$ ,  $C \in c$ . Let  $\theta$  be any  $\Pi_1$  sentence provable in C. It suffices to show that  $A \vdash \theta$ . By Lemma 6,  $B \le A + \neg \theta$ . But then, by assumption,  $A \vdash \theta$ , as desired.

**Lemma 8.** If  $\pi$  is  $\Pi_1$ ,  $A \le B + \pi$ , and  $\neg \pi$  is  $\Pi_1$ -conservative over A, then  $d(A) \ll_{\cap} d(B + \pi)$ .

**Proof.** Suppose  $B + \pi \le A + \sigma$ . Then, by Lemma 6.1,  $A + \sigma \vdash \pi$ , whence  $A + \neg \pi \vdash \neg \sigma$  and so  $A \vdash \neg \sigma$ , in other words,  $A + \sigma$  is inconsistent. Now use Lemma 7. **Proof of Theorem 3 (b).** Let  $A \in a$ ,  $C \in c$ . By Theorem 6.5, there is a sentence  $\psi$  such that  $C \vdash \psi \not\subseteq A$ . Let  $X = \{\phi: \psi \le A + \neg \phi\}$ . Then, by Theorem 5.2, there is a  $\Sigma_1$  sentence  $\chi \not\in X$  such that  $\chi$  is  $\Pi_1$ –conservative over A. Let  $B = A + \neg \chi$  and b = d(B). Then  $c \not\subseteq A$ 

From Theorem 6.4 and Lemma 8, and Theorem 5.1, we get the following (compare Theorem 6.2):

Corollary 2. 
$$d(A) \ll_{\cap} d(T+Con_A)$$
.

Theorem 3 (a) leads to the problem if for any  $a < 1_T$ , there is a b such that  $a < \!\!\! < b < 1_T$ . (The dual of this is false: if  $0_T < b < a$  and not  $0_T < \!\!\! < c$  a, then not  $b < \!\!\! < c$  a.) We now show that the answer is negative.

a is a *cupping* degree if  $a < 1_T$  and a cups to every b such that  $a \le b < 1_T$ . Let  $CON_T = \{a < 1_T : a = d(T + Con_\tau) \text{ for some PR binumeration } \tau(x) \text{ of } T\}$ . By Corollary 2.4,  $CON_T \ne \emptyset$ .

**Theorem 4.** Every member of CON<sub>T</sub> is a cupping degree.

**Proof.** Suppose  $a = d(T+Con_{\tau}) < 1_T$ , where  $\tau(x)$  is a PR binumeration of T. Let b be any degree such that  $a \le b < 1_T$ . Let  $B \in b$ . We want to define a degree d such that d  $\not\ge a$  and  $a \cup d \ge b$ . The obvious way to try is to let  $d = d(T+\theta)$ , where

$$\theta := \forall u(Prf_B(\bot,u) \rightarrow \exists z < uPrf_{\tau}(\bot,z)).$$

But it seems difficult to prove, and may not even be true, that  $d \not\ge a$  so we have to proceed in a somewhat different way.

Let  $\varphi$  be such that

$$PA \vdash \phi \leftrightarrow \forall z (Prf_{\tau}(\phi, z) \rightarrow \exists u \leq z Prf_{R}(\bot, u)),$$

and let

$$\psi := \forall u(\Pr f_B(\bot, u) \rightarrow \exists z < u \Pr f_\tau(\varphi, z)).$$

Then

- (1) T⊬ φ,
- (2)  $PA \vdash \phi \lor \psi$ ,
- (3)  $PA \vdash \phi \land \psi \rightarrow Con_B$ .

Clearly, PA $\vdash \neg \phi \rightarrow \Pr_{\tau}(\phi)$ . Since  $\neg \phi$  is  $\Sigma_1$ , we also have, by provable  $\Sigma_1$ –completeness, PA $\vdash \neg \phi \rightarrow \Pr_{\tau}(\neg \phi)$ . Thus,

(4)  $PA \vdash Con_{\tau} \rightarrow \varphi$ .

Let  $d = d(T+\psi)$ . Then, since  $\psi$  and  $Con_{\tau}$  are  $\Pi_1$ , it follows from (3), (4), Lemma 3, and Theorem 6.4 that  $a \cup d \ge b$ . Suppose  $a \le d$ . Then  $T + \psi \vdash Con_{\tau}$ . But then, by (2) and (4),  $T \vdash \varphi$ , contradicting (1). Thus,  $a \not\le d$ . Let  $c = d \cap b$ . Then c < b. Finally,  $a \cup c = (a \cup d) \cap (a \cup b) = b$ . Thus, a is cupping.

Theorem 14', below, is an improvement of Theorem 4.

A set G of degrees is *cofinal* in  $\mathbf{D}_T$  if for every degree  $a < 1_T$ , there is a degree  $b \in G$  such that  $a \le b < 1_T$ .

# **Lemma 9.** $CON_T$ is cofinal in $D_T$ .

**Proof.** Suppose  $b < 1_T$ . By Corollary 2.4, even if T is not  $\Sigma_1$ –sound, there is a PR binumeration  $\beta(x)$  of a theory of degree b such that  $T + Con_{\beta}$  is consistent. By Theorem 6.4,  $b \le d(T + Con_{\beta})$ . By Theorem 2.8 (b), there is a PR binumeration  $\tau(x)$  of T such that  $T \vdash Con_{\tau} \leftrightarrow Con_{\beta}$ . Let  $a = d(T + Con_{\tau})$ . Then  $b \le a \in CON_T$ .

Let P be a property of degrees. We shall say that there are arbitrarily large degrees having property P if the set of degrees having P is cofinal in  $\mathbf{D}_T$ . Every sufficiently large degree has P if for every degree  $a < 1_T$ , there is a b such that  $a \le b < 1_T$  and every degree c such that  $b \le c < 1_T$  has P.

If a is cupping and a  $\leq$  b, b is cupping. Thus, from Theorem 4 and Lemma 9 we get:

Corollary 3. Every sufficiently large degree is a cupping degree.

Theorem 5. There is a degree which has no p.c.

The proof of this (and more) will be given in § 3 (Theorem 17).

In addition to the usual (finitary) distributive laws,  $D_T$  also satisfies the following infinitary distributive laws. Let G be a set of degrees.  $\bigcup G$  ( $\bigcap G$ ) is then the l.u.b. (g.l.b.) of G, if it exists.

**Theorem 6.** (a) If  $\bigcup G$  exists, then  $\bigcup G \cap b = \bigcup \{a \cap b : a \in G\}$ . (b) If  $\bigcap G$  exists, then  $\bigcap G \cup b = \bigcap \{a \cup b : a \in G\}$ .

By Theorem 6 (a), if a has no p.c., then {b:  $b \cap a = 0_T$ } has no l.u.b. In Lemma 23, below, we give a nontrivial example of a set G which has no g.l.b.

To prove Theorem 6 (b) we need the following:

Lemma 10. The following conditions are equivalent:

- (i)  $A \uparrow B \ge C$ .
- $\text{(ii) For all } (\Sigma_1) \text{ sentences } \chi \text{ and all } m \text{, if } A^T + \neg Con_{C \mid m} \dashv_{\Sigma_1} T + \chi \text{, then } B \vdash \neg \chi.$

**Proof.** Suppose (i) holds. Let  $\chi$  and m be such that  $A^T + \neg Con_{C \mid m} \dashv_{\Sigma_1} T + \chi$ . There is a k such that  $A^T + Con_{B \mid k} \vdash Con_{C \mid m}$ . It follows that  $T + \chi \vdash \neg Con_{B \mid k}$ , whence  $B \vdash \neg \chi$ . Thus, (ii) holds.

To prove that (ii) implies (i), suppose (i) fails, i.e.  $A^{\uparrow}B \ngeq C$ . There is then an m such that for every k,  $A^T + Con_{B \mid k} \nvdash Con_{C \mid m}$ . But then, by Theorem 4.3, there is a

 $\Sigma_1$  sentence  $\chi$  such that  $A^T + \neg Con_{C \mid m} \dashv_{\Sigma_1} T + \chi$  and  $T + \chi \not\vdash \neg Con_{B \mid k}$  for every k. Since  $\neg \chi$  is  $\Pi_1$ , it follows, by Lemma 2, that  $B \not\vdash \neg \chi$ . Thus, (ii) is false, as desired.  $\blacksquare$  **Proof of Theorem 6.** (a) Let  $c = \bigcup G$ . Clearly  $c \cap b$  is an upper bound of  $\{a \cap b: a \in G\}$ . Suppose d is any upper bound of  $\{a \cap b: a \in G\}$ . It is then sufficient to show that  $c \cap b \leq d$ . Let  $B \in b$ ,  $C \in c$ ,  $D \in d$ . Then  $A \downarrow B \leq D$  for every A such that  $d(A) \in G$ . But then, by Lemma  $A \cap A \subseteq A$  is  $A \cap A \subseteq A$  and every  $A \cap A \subseteq A$  and every  $A \cap A \subseteq A$  is  $A \cap A \subseteq A \subseteq A$ . And every  $A \cap A \subseteq A \subseteq A$  is  $A \cap A \subseteq A \subseteq A$ .

(b) Let  $c = \bigcap G$ . Clearly  $c \cup b$  is a lower bound of  $\{a \cup b : a \in G\}$ . Suppose d is any lower bound of  $\{a \cup b : a \in G\}$ . It is then sufficient to show that  $d \le c \cup b$ . Again let  $B \in b$  etc. Then  $D \le A \uparrow B$  for every A such that  $d(A) \in G$ . But then, by Lemma 10, for every such A, every m and every  $\Sigma_1$  sentence  $\chi$ , if  $B^T + \neg Con_{D \mid m} \dashv_{\Sigma_1} T + \chi$ , then  $A \vdash \neg \chi$ . It follows that for every m and every  $\Sigma_1$  sentence  $\chi$ , if  $B^T + \neg Con_{D \mid m} \dashv_{\Sigma_1} T + \chi$ , then  $C \vdash \neg \chi$ . Hence, again by Lemma 10,  $D \le C \uparrow B$  and so  $d \le c \cup b$ .

Suppose  $a \le b$ . Let [a,b] be the *interval*  $\{c: a \le c \le b\}$ . (We also write [a,b) for  $\{c: a \le c < b\}$  etc.) A natural (global) question concerning  $\mathbf{D}_T$  is if all intervals [a,b], where  $a < b < 1_T$ , are isomorphic (in the obvious sense). The answer is negative.

If c < d, let  $[d,c] = ([c,d], \ge)$ . Another natural question is, under what conditions [a,b] is isomorphic to [d,c], where a < b and c < d.

**Theorem 7.** (a) There are degrees a,  $b \in (0_T, 1_T)$  such that the intervals  $[0_T, a]$  and  $[0_T, b]$  are not isomorphic.

(b) Suppose a < b and c < d. Then [a,b] is not isomorphic to [d,c].

Theorem 7 (a) follows at once from our next two lemmas.

The interval  $[a_0,a_1]$ , where  $a_0 \le a_1$ , is said to satisfy the *reduction principle* if for any  $b_0$ ,  $b_1 \in [a_0,a_1]$ , if  $b_0 \cup b_1 = a_1$ , there are  $c_i \le b_i$ , i = 0, 1, such that  $c_0 \cap c_1 = a_0$  and  $c_0 \cup c_1 = a_1$ . A degree a is r.p. if  $[0_{T,a}]$  satisfies the reduction principle.

**Lemma 11.** If  $a = d(T+\theta)$ , where  $\theta$  is  $\Pi_1$ , then a is r.p.

**Proof.** Let  $b_0$ ,  $b_1$  be such that  $b_0 \cup b_1 = a$ . There are then  $\Pi_1$  sentences  $\psi_0$ ,  $\psi_1$  such that  $d(T+\psi_i) \le b_i$  and  $T+\psi_0 \wedge \psi_1 \vdash \theta$ . By Lemma 5.5, there are  $\Pi_1$  sentences  $\theta_0$ ,  $\theta_1$  such that  $T\vdash \theta_0 \vee \theta_1$ ,  $T\vdash \psi_i \to \theta_i$ , i=0,1,  $T\vdash \theta_0 \wedge \theta_1 \to \psi_0 \wedge \psi_1$ . Let  $c_i=d(T+\theta_i)$ , i=0,1. Then  $c_i \le b_i$ ,  $c_0 \cap c_1 = 0_T$ , and, by Lemma 3,  $c_0 \cup c_1 = b_0 \cup b_1 = a$ .

**Lemma 12.** There is a degree  $a < 1_T$  which is not r.p.

**Proof.** Let  $\pi$  be a  $\Pi_1$  sentence undecidable in T. In case T is not  $\Sigma_1$ -sound we also need to assume that  $\pi$  is  $\Sigma_1$ -conservative over T (cf. Theorem 5.2). We now effectively define r.e. sets  $X_k$  of  $\Pi_1$  sentences such that

(1)  $T + X_k + \pi^i$  is consistent, i = 0, 1,

- $(2) X_k \subseteq X_{k+1},$
- (3)  $T + X_k + \pi \not\vdash X_{k+1}$
- (4)  $T + X_k + \neg \pi \le T + X_{k+1}$ .

Let  $X_0 = \emptyset$ . Then (1) holds for k = 0. Now suppose (1) holds for k = n. By (the proof of) Lemma 2.1, we can effectively find a  $\Pi_1$  sentence  $\psi_n$  such that

(5)  $T + X_n + \pi^i + \neg \psi_n^i$  is consistent, i = 0, 1.

Let  $T_n =_{df} T + X_n + \neg \pi + \psi_n$ . It follows that

(6) there is no  $\Pi_1$  sentence  $\theta$  such that  $T_n \vdash \theta$  and  $T + \theta \vdash \neg \pi$ .

For suppose  $T + \theta \vdash \neg \pi$ . Then  $T + \pi \vdash \neg \theta$  and so  $T \vdash \neg \theta$ , whence, by (5),  $T_n \not\vdash \theta$ .

Let  $X_{n+1} = \text{Th}(T_n) \cap \Pi_1$ . Let k = n+1. Then (1) is satisfied for i = 1 and, by (6), (1) is satisfied for i = 0. Moreover (2) and (4) hold for k = n. Finally,  $T + X_{n+1} \vdash \psi_n$  and so, by (5), (3) holds for k = n.

Let  $a_0 = d(T + \bigcup \{X_k : k \in N\})$ ,  $a_1 = d(T + \pi)$ , and  $a = a_0 \cup a_1$ . Since  $a_0 < 1_T$  and  $\pi$  is  $\Sigma_1$ —conservative over T, we have  $a < 1_T$ . We now show that a is not r.p. Let  $b_0$  and  $b_1$  be such that  $b_0 \le a_0$ ,  $b_1 \le a_1$ ,  $b_0 \cap b_1 = 0_T$ , and  $b_0 \cup b_1 \ge a_1$ . It is then sufficient to show that  $b_0 \cup b_1 \not\ge a_0$ .

Let  $\theta_{i,k}$  be  $\Pi_1$  sentences such that  $b_i = d(T + \{\theta_{i,k}: k \in N\})$ , i = 0, 1. We may assume that  $T + \theta_{i,k+1} \vdash \theta_{i,k}$  for i = 0, 1 and all k. By Lemma 3, there is then an m such that  $T + \theta_{0,m} \wedge \theta_{1,m} \vdash \pi$ .  $d(T + \theta_{0,m}) \leq b_0 \leq a_0$ . Thus, by (2), there is an n such that  $T + \theta_{0,m} \leq T + X_n$ . Since  $b_0 \cap b_1 = 0_T$ , for every k,  $T + \theta_{0,k} \vee \pi \leq T + \theta_{0,k} \vee (\theta_{0,m} \wedge \theta_{1,m}) \leq T + \theta_{0,m}$ . It follows that  $T + \theta_{0,k} \vee \pi \leq T + X_n$ , whence  $T + \theta_{0,k} \leq T + X_n + \neg \pi$  (cf. Corollary 6.3) and so, by (4),  $T + \theta_{0,k} \leq T + X_{n+1}$ . But this holds for all k, whence  $b_0 \leq d(T + X_{n+1})$ . Next, by (3),  $b_0 \cup b_1 \leq b_0 \cup a_1 \leq d(T + X_{n+1} + \pi) \not\geq a_0$ . It follows that  $b_0 \cup b_1 \not\geq a_0$  and so the proof is complete.

**Proof of Theorem 7 (b).** Let  $A \in a$  and let  $\pi$  be a  $\Pi_1$  sentence such that  $A \not\vdash \pi$ . Then  $[a,d(A^T+\pi)]$  satisfies the reduction principle (see the proof of Lemma 11). It follows that in [a,b] there is a degree e > a such that [a,e] satisfies the reduction principle. Thus, it is sufficient to show that the dual of the reduction principle is false in [c,d] whenever c < d.

Let  $C \in c$  and  $D \in d$ , and let  $\pi$  be such that  $C \not\vdash \pi$  and  $D \vdash \pi$ . Then, by Theorem 5.5 (b) with  $X = Th(C^T + \neg \pi)$ , there are  $\Sigma_1$  sentences  $\sigma_i$  such that  $C^T + \neg \sigma_i \equiv C^T + \sigma_{1-i}$ , i = 0, 1, and  $C^T + \neg \sigma_0 \land \neg \sigma_1 \not\vdash \pi$ . Let  $c_i = d(C^T + \sigma_i) = d(C^T + \neg \sigma_{1-i})$ . Then  $c_0 \cap c_1 = c$  and  $c_0 \cup c_1 \not\succeq d$ . Let  $d_i = c_i \cap d$ . Then  $d_0 \cap d_1 = c$  and  $d_0 \cup d_1 < d$ . Suppose now  $d_i \le e_i \le d$ , i = 0, 1, and  $e_0 \cap e_1 = c$ . We have to show that  $e_0 \cup e_1 < d$ . Let  $E_0 \in e_0$ .  $c_1 \cap e_0 = c_1 \cap d \cap e_0 = d_1 \cap e_0 \le e_1 \cap e_0 = c$ . It follows that  $(C^T + \neg \sigma_0) \not\downarrow E_0 \le C^T$ . But then, by Lemma 6, for every  $\Pi_1$  sentence  $\theta$ , if  $E_0 \vdash \theta$ , then  $C^T + \neg \sigma_0 \le C^T + \neg \theta$ , whence  $C^T + \sigma_0 \vdash \theta$ . It follows that  $e_0 \le c_0$  and so  $e_0 = d_0$ . Similarly,  $e_1 = d_1$ . Hence  $e_0 \cup e_1 = d_0 \cup d_1 < d$  and the proof is complete.

Theorem 7 (a) leads to the problem of determining the exact number of nonisomorphic intervals of  $D_T$ . This problem remains open.

We have actually proved more than is stated in Theorem 7. Let  $L = \{ \le, \cap, \cup, 0, 1 \}$  be the language of the theory of lattices with a bottom and a top element.

Formulated in L, the reduction principle is an  $\forall \exists$  sentence. Hence, by the proof of Theorem 7 (a), there are degrees a,  $b \in (0_T, 1_T)$  and an  $\forall \exists$  sentence of L which holds in  $[0_T, a]$  but not in  $[0_T, b]$ . (This is, so far, the only known way of proving that two intervals of  $\mathbf{D}_T$  are nontrivially nonisomorphic.) Similarly, the proof of Theorem 7 (b) shows that if a < b and c < d, there is an  $\exists \forall \exists$  sentence which is true in [a,b] and false in [d,c].

**§2.** A classification of degrees. When there is no risk of confusion we shall use  $\phi$  and X in place of  $T+\phi$  and T+X. Thus,  $d(\phi)$  is  $d(T+\phi)$ ,  $X<\phi$  means that  $T+X< T+\phi$ ,  $\phi\equiv\psi$  that  $T+\phi\equiv T+\psi$ , etc. We also write a  $\ll$ b to mean that a  $\ll$  $_0$ b. A  $\ll$ B means that  $d(A)\ll d(B)$ .  $\sigma$ ,  $\sigma_0$ , etc. will be used to denote  $\Sigma_1$  sentences and  $\pi$ ,  $\pi_0$ , etc. to denote  $\Pi_1$  sentences.

A degree a is  $\Phi$  if there is a  $\Phi$  sentence  $\phi$  such that  $a=d(\phi)$ . By the proof of Theorem 6.11 (a), it is clear that every degree is  $\Pi_2$  and  $\Sigma_2$ . This can be somewhat improved:

**Theorem 8.** Every degree is  $\Delta_2$ .

**Proof.** Let a be any degree. There is a primitive recursive set X of  $\Pi_1$  sentences such that a = d(X). Let  $\xi(x)$  be a PR binumeration of X and let  $\varphi$  be such that

$$\mathrm{PA} \vdash \phi \leftrightarrow \forall z \big( [\Pi_1](\phi,z) \to (\xi(z) \to \mathrm{Tr}_{\Pi_1}(z)) \big).$$

Then  $\varphi$  is  $\Pi_2$  and  $T + \varphi$  is a  $\Pi_1$ -conservative extension of T + X (cf. the proof of Theorem 5.4 (a)). It follows that  $a = d(\varphi)$ . Using Lemma 5.1 (i) and Lemma 1.3 (v) (applied to  $\neg \varphi$ ), we get:

$$PA \vdash \phi \leftrightarrow \forall z (\xi(z) \to \mathrm{Tr}_{\Pi_1}(z)) \vee \exists z \big( \neg [\Pi_1](\phi,z) \wedge \forall u < z (\xi(u) \to \mathrm{Tr}_{\Pi_1}(u)) \big).$$
 Thus,  $\phi$  is  $\Delta_2$ .

By Theorem 8, in terms of the arithmetical hierarchy, the only interesting (proper) subsets of  $D_T$  are the sets of  $B_1$  degrees,  $\Sigma_1$  degrees,  $\Pi_1$  degrees, and degrees which are both  $\Sigma_1$  and  $\Pi_1$ . (If T is not  $\Sigma_1$ –sound, there are also  $\Delta_1^T$  degrees other than  $0_T$  and  $1_T$ ; see e.g. the proof of Corollary 1.) The object of the rest of this  $\S$  is simply to show that these sets are different and that there is a non– $B_1$  degree. More detailed information about the  $\Sigma_1$  and the  $\Pi_1$  degrees will be given in the next  $\S$ .

Our next lemma is a restatement of Theorem 6.11 (b).

**Lemma 13.** If X is an r.e. set of  $\Sigma_1$  sentences, then d(X) is  $\Sigma_1$ .

The following lemma is occasionally useful.

**Lemma 14.** There exist a (primitive) recursive sequence  $\langle \sigma_k \rangle_{k < \omega}$  and a sentence  $\sigma$  such that (i)  $T + \sigma_{k+1} \vdash \sigma_{k'}$  for all k, (ii)  $\sigma_k < \sigma_{k+1'}$  for all k, (iii)  $\sigma \equiv \{\sigma_k : k \in \mathbb{N}\}$ .

This follows at once from (the proof of) Lemma 2.1 (applied to the sets  $\{\phi: Q + \phi \le T + \sigma_k\}$ ;  $\sigma_0 := 0 = 0$ ) and Lemma 13.

**Theorem 9.** (a) There is a  $\Pi_1$  degree which isn't  $\Sigma_1$ .

- (b) There is a  $\Sigma_1$  degree which isn't  $\Pi_1$ .
- (c) There is a degree other than  $0_T$  and  $1_T$  which is both  $\Sigma_1$  and  $\Pi_1$ .
- (d) There is a  $B_1$  degree which is neither  $\Sigma_1$  nor  $\Pi_1$ .
- (e) There is a degree which isn't B<sub>1</sub>.

**Proof.** (a) Let  $\pi$  be such that  $\neg \pi$  is  $\Pi_1$ —conservative over T and T $\not\vdash \neg \pi$ . Then, by Lemma 8,  $0_T << d(\pi)$  and so, by Lemma 7,  $d(\pi)$  is not  $\Sigma_1$ .  $\blacklozenge$ 

- (b) Let  $\langle \sigma_k \rangle_{k < \omega}$  and  $\sigma$  be as in Lemma 14. Suppose  $d(\sigma)$  is  $\Pi_1$  and let  $\pi$  be such that  $\sigma \equiv \pi$ . Then  $\pi \equiv \{\sigma_k : k \in N\}$  and so, by Lemma 14 (i), there is an m such that  $T + \sigma_m \vdash \pi$ . But then  $\{\sigma_k : k \in N\} \le \sigma_m$ , contradicting Lemma 14 (ii). Thus,  $d(\sigma)$  isn't  $\Pi_1$ .  $\blacklozenge$
- (d) The easiest way to prove this is to define  $\pi$  as in the proof of (a) and then  $\sigma$  as in the proof of (b), except that T is replaced by T +  $\pi$ . Then  $d(\pi \wedge \sigma)$  is neither  $\Sigma_1$  nor  $\Pi_1$ . Details are left to the reader.  $\blacklozenge$

Theorem 9 (c) will be derived from the following lemma, which will also be used later.

**Lemma 15.** There are  $\Pi_1$  sentences  $\theta_i$ , i = 0, 1, such that

- (i)  $T \not\vdash \theta_i$
- (ii)  $T \vdash \theta_0 \lor \theta_1$ ,
- (iii)  $T \vdash \theta_0 \land \theta_1 \rightarrow Con_{T}$
- (iv)  $T + Con_T \vdash \neg Pr_T(\theta_i)$
- (v)  $T + Con_T \vdash \theta_i$ ,
- $(vi) \quad \theta_i \equiv \neg \theta_{1-i}.$

**Proof.** Let  $\theta_i$ , i = 0, 1, be such that

$$PA \vdash \theta_i \longleftrightarrow \forall z (Prf_T(\theta_{i},z) \to \exists u {<} z {+} i Prf_T(\theta_{1-i},u)).$$

A standard argument proves (i). Formalizing this argument we get (iv). (ii) and (iii) are immediate. (v) follows from (iv). By (ii),

(1)  $PA \vdash Pr_T(\neg \theta_i) \rightarrow Pr_T(\theta_{1-i}).$ 

Also,

$$(2) \qquad \text{PA} \vdash \neg \theta_{1-i} \leftrightarrow \text{Pr}_{T}(\theta_{1-i}) \wedge \theta_{i}.$$

By Theorem 6.8,  $\theta_i \land \Pr_T(\neg \theta_i) \le \theta_i$ . By (1), it follows that  $\theta_i \land \Pr_T(\theta_{1-i}) \le \theta_i$  and so, by (2),  $\neg \theta_{1-i} \le \theta_i$ . But then, by (ii),  $\theta_i \equiv \neg \theta_{1-i}$ , i.e. (v) holds.

**Proof of Theorem 9 (c).** Let  $\theta_i$  be as in Lemma 15. Let  $a = d(\theta_0)$ . Then a is  $\Pi_1$  and, by Lemma 15 (vi), a is  $\Sigma_1$ . By Lemma 15 (i),  $a > 0_T$ . Finally, by Lemma 15 (i) and (ii),  $a < 1_T$ .

To prove Theorem 9 (e) we need the following:

**Lemma 16.** Suppose X is r.e. and for every k,  $X \mid k \ll X$ . Then if  $\varphi$  is  $B_1$  and  $X \leq \varphi$ , then  $X \ll \varphi$ . Thus, a fortiori d(X) is not  $B_1$ .

**Proof.**  $\varphi$  can be written in the form  $(\pi_0 \wedge \sigma_0) \vee ... \vee (\pi_n \wedge \sigma_n)$ . It is easily checked that for any degrees a, b, c, if  $a \ll b$  and  $a \ll c$ , then  $a \ll b \cap c$ . Thus, it is sufficient to show that if  $X \leq \pi \wedge \sigma$ , then  $X \ll \pi \wedge \sigma$ . Let  $\chi$  be a  $\Sigma_1$  sentence such that  $\pi \wedge \sigma \leq X + \chi$ . Then, by Lemma 7, it suffices to show that  $T + X \vdash \neg \chi$ . By assumption, there is a k such that  $T + X \vdash k + \chi \vdash \pi$ . Hence  $T + \pi \wedge \sigma \dashv T + X \vdash k + (\chi \wedge \sigma)$  and so  $X \leq X \vdash k + (\chi \wedge \sigma)$ . But then, since  $X \vdash k \ll X$ , by Lemma 7,  $T + X \vdash \neg (\chi \wedge \sigma)$ . But  $X \leq \pi \wedge \sigma$ . It follows that  $T + \pi \wedge \sigma \vdash \neg \chi$ , whence  $T + X + \chi \vdash \neg \chi$  and so  $T + X \vdash \neg \chi$ , as was to be shown.

**Proof of Theorem 9 (e).** By (the proof of) Theorem 5.2, we can effectively construct sentences  $\pi_n$  such that  $\neg \pi_n$  is  $\Pi_1$ -conservative over but not provable in  $T + \{\pi_k: k < n\}$ . Let  $X = \{\pi_k: k \in N\}$ . Then, by Lemma 8,  $X \mid k \ll X$  for all k. So, by Lemma 16, d(X) is not  $B_1$ .

§3.  $\Sigma_1$  and  $\Pi_1$  degrees. This § is devoted to a discussion of the  $\Sigma_1$  and  $\Pi_1$  degrees and the relations between them.

The l.u.b. of two  $\Pi_1$  degrees is  $\Pi_1$  and the g.l.b. of two  $\Pi_1$  ( $\Sigma_1$ ) degrees is  $\Pi_1$  ( $\Sigma_1$ ). Let us say that a is *high* if a  $\gg 0_T$ , *low* otherwise. Thus, by Lemma 7, a is low iff there is a  $\Sigma_1$  degree b such that a  $\leq$  b < 1 $_T$ . By Lemma 8, if  $\neg \pi$  is  $\Pi_1$ –conservative over T, d( $\pi$ ) is high. By Corollary 2, every member of CON $_T$  is high.

The following lemma is sometimes useful.

**Lemma 17.** Suppose a is high. Then for any b,  $[a \cap b,b)$  contains no  $\Sigma_1$  degree; in fact, if c is  $\Sigma_1$  and  $a \cap b \le c$ , then  $b \le c$ .

**Proof.** Let  $A \in a$ ,  $B \in b$ , and  $c = d(\sigma)$ . Suppose  $A \downarrow B \le T + \sigma$ . Then, by Lemma 6,  $A \le T + \sigma \land \neg Con_{B \mid k}$  for every k. Since a is high, it follows that  $T + \sigma \vdash Con_{B \mid k}$ , for every k, and so  $B \le T + \sigma$ .

**Theorem 10.** (a) The set of  $\Pi_1$  degrees is cofinal in  $\mathbf{D}_{\mathrm{T}}$ .

- (b) The set of  $\Sigma_1$  degrees is not cofinal in  $\mathbf{D}_T$ ; in fact, for every degree  $a > 0_T$ , there is a degree b < a such that [b,a) contains no  $\Sigma_1$  degree.
- (c) There is a low  $\Pi_1$  degree which is not  $\Sigma_1$ .
- **Proof.** (a) Since all members of CON<sub>T</sub> are  $\Pi_1$ , this follows from Lemma 9.  $\bullet$
- (b) By Theorem 3 (b), there is a high degree c such a  $\nleq$  c. Let b = c  $\cap$  a. Then, by Lemma 17, b is as desired.  $\blacklozenge$ 
  - (c) Let a be any low  $\Pi_1$  degree  $> 0_T$ . By Theorem 3 (b), there is a high  $\Pi_1$  degree

c  $\not\ge$  a. Let b = a ∩ c. Then b is low and and  $\Pi_1$ . Finally, by Lemma 17, b is not  $\Sigma_1$ . ■ Using Theorem 2 we can now prove the following corollary.

**Corollary 4.** (a) Suppose a is not  $\Pi_1$  and  $a \in (b,c)$ . There are then degrees b', c' such that  $a \in (b',c') \subseteq (b,c)$  and [b',c'] contains no  $\Pi_1$  degree.

- (b) Suppose a is not  $\Sigma_1$  and  $a \in (b,c)$ . There are then degrees b', c' such that  $a \in (b',c') \subseteq (b,c)$  and [b',c'] contains no  $\Sigma_1$  degree.
- **Proof.** (a) By Theorem 2, there are degrees  $b_0$ ,  $b_1$  such that  $b \le b_i < a$ , i = 0, 1, and  $b_0 \cup b_1 = a$ . Either  $[b_0,a]$  or  $[b_1,a]$  contains no  $\Pi_1$  degree. If not, then a would be the l.u.b. of two  $\Pi_1$  degrees and therefore  $\Pi_1$ . Suppose  $[b_i,a]$  contains no  $\Pi_1$  degree and let  $b' = b_i$ . By Theorem 2, there are degrees  $c_0$ ,  $c_1$  such that  $a < c_i \le c$ , i = 0, 1, and  $c_0 \cap c_1 = a$ . Either  $[b',c_0]$  or  $[b',c_1]$  contains no  $\Pi_1$  degree. For suppose  $d_i \in [b',c_i]$  and  $d_i$  is  $\Pi_1$ , i = 0, 1. Then  $d_0 \cap d_1 \in [b',a]$  and  $d_0 \cap d_1$  is  $\Pi_1$ , a contradiction. Suppose  $[b',c_i]$  contains no  $\Pi_1$  degree and let  $c' = c_i$ . Then b' and c' are as desired.  $\blacklozenge$
- (b) By a slight modification of the proof of Theorem 10 (b), which we leave to the reader, there is a degree b' such that  $b \le b' < a$  and [b',a] contains no  $\Sigma_1$  degree. The rest of the proof is the same as the proof of (a).

Theorem 10 (b) leads to the question if there are arbitrarily small  $\Sigma_1$  degrees. By our next result, the answer is affirmative; later we shall prove a stronger result (Theorem 15).

**Theorem 11.** If  $0_T < a$ , then there is a  $\Sigma_1$  and  $\Pi_1$  degree  $b \in (0_T, a)$ .

To prove this we need a lemma on partial conservativity.

**Lemma 18.** Let *X* be an r.e. set. There is then a PR formula  $\eta(y,x,z)$  such that for all k and  $\theta$ ,

- (i) if  $k \in X$ , then  $T + \theta \vdash \neg \exists z \eta(\theta, k, z)$ ,
- (ii) if  $k \notin X$ , then  $\exists z \eta(\theta, k, z)$  is  $\Pi_1$ -conservative over  $T + \theta$ .

The proof of Lemma 18 is similar to the proof of Lemma 5.3 (for  $\Gamma = \Pi_1$ ) and is left to the reader.

**Proof of Theorem 11.** Let  $\forall u \delta(u)$ , where  $\delta(u)$  is PR, be a  $\Pi_1$  sentence such that  $0_T < d(\forall u \delta(u)) < a$ . By Lemma 18, there is a PR formula  $\gamma(x,z)$  such that

- (1) if  $T \vdash \varphi$ , then  $T \vdash \neg \exists z \gamma(\varphi, z)$ ,
- (2) if  $T \not\vdash \varphi$ , then  $\exists z \gamma(\varphi, z)$  is  $\Pi_1$ —conservative over  $T + \varphi$ .

Let  $\theta$  be such that

(3) PAF  $\theta \leftrightarrow \forall u(\neg \delta(u) \rightarrow \exists z < u \gamma(\theta, z))$ , and let

$$\sigma := \exists z (\gamma(\theta, z) \land \forall u \leq z \delta(u)).$$

Then

- (4)  $PA \vdash \sigma \leftrightarrow \exists z \gamma(\theta, z) \land \theta$ ,
- (5)  $PA + \theta + \neg \exists z \gamma(\theta, z) \vdash \forall u \delta(u).$

It follows that

(6) T⊬ θ.

For suppose not. Then, by (1),  $T \vdash \neg \exists z \gamma(\theta, z)$  and so, by (5),  $T \vdash \forall u \delta(u)$ , contrary to the choice of  $\delta(u)$ .

By (3),  $\theta \le \forall u \delta(u)$  and so  $d(\theta) < a$ . By (6),  $0_T < d(\theta)$ . Finally, by (4), (6), (2),  $\sigma \equiv \theta$ . Thus,  $b = d(\sigma)$  is as claimed.

It is natural to ask if  $D_T$  is "generated" by some "small" set of degrees, for example, the set of  $\Sigma_1$  degrees. We prove two negative results, Theorems 12 and 13, and one partial positive result, Theorem 14 (and 14').

Let  $E_T$  be the set of l.u.b.s of (finite) sets of  $\Sigma_1$  degrees. Note that  $E_T$  is closed under  $\cap$ . By Lemma 15,  $CON_T \subseteq E_T$ .

**Theorem 12.** There is a  $\Pi_1$  degree not in  $E_T$ .

This is an immediate consequence of the following two lemmas.

**Lemma 19.** If  $a \in E_T$ , there is a smallest  $\Sigma_1$  degree  $\geq a$ .

**Proof.** Suppose  $a=d(\sigma_0)\cup...\cup d(\sigma_n)$ . Then  $d(\sigma_0\wedge...\wedge\sigma_n)$  is the smallest  $\Sigma_1$  degree  $\geq d(\sigma_0)\cup...\cup d(\sigma_n)$ . This can be seen as follows. Suppose  $d(\sigma_0)\cup...\cup d(\sigma_n)\leq d(\sigma)$ . Let  $\pi$  be such that  $T+\sigma_0\wedge...\wedge\sigma_n\vdash\pi$ . Then  $T+\sigma_0\vdash\sigma_1\wedge...\wedge\sigma_n\to\pi$ . Now,  $\sigma_1\wedge...\wedge\sigma_n\to\pi$  is a  $\Pi_1$  sentence. It follows that  $T+\sigma\vdash\sigma_1\wedge...\wedge\sigma_n\to\pi$ . But then  $T+\sigma_1\vdash\sigma\wedge\sigma_2\wedge...\wedge\sigma_n\to\pi$  and so  $T+\sigma\vdash\sigma_2\wedge...\wedge\sigma_n\to\pi$ . Continuing in this way we eventually get  $T+\sigma\vdash\pi$ , as desired.  $\blacksquare$ 

**Lemma 20.** There is a  $\Pi_1$  degree a for which there is no smallest  $\Sigma_1$  degree  $\geq$  a.

**Proof.** Let  $\langle \sigma_k \rangle_{k < \omega}$  and  $\sigma$  be as in Lemma 14. Let  $a = d(\neg \sigma)$ . Then a is  $\Pi_1$ . Now let  $\chi$  be any  $\Sigma_1$  sentence such that  $a \leq d(\chi)$ . Then  $T + \chi \vdash \neg \sigma$  and so  $T + \sigma \vdash \neg \chi$ . It follows that there is a k such that  $T + \sigma_k \vdash \neg \chi$  and so

(1)  $T + \chi \vdash \neg \sigma_k$ .

Since  $\sigma_k < \sigma$ , there is a sentence  $\pi$  such that  $T + \sigma \vdash \pi$  and  $T + \sigma_k \nvdash \pi$ . It follows that

- (2)  $T + \neg \pi \vdash \neg \sigma$ ,
- (3)  $T + \neg \pi \not\vdash \neg \sigma_k$ .

But then, by (2),  $a \le d(\neg \pi)$  and, by (1) and (3),  $\chi \not\leq \neg \pi$ . Thus,  $d(\chi)$  is not the smallest  $\Sigma_1$  degree  $\ge a$ .

A strengthening of Lemma 20 will be proved later (Lemma 23).

Let  $F_T$  be the set of l.u.b.s of (finite) sets of  $\Sigma_1$  and  $\Pi_1$  degrees. By Theorem 12,  $F_T \subsetneq E_T$ .

## Theorem 13. $F_T \neq D_T$ .

We need the following definition:  $A \iff B$  iff A < B and for every set X of  $\Sigma_1$  sentences, if  $B \dashv_{\Pi_1} A + X$ , then A + X is inconsistent. (Here X need not be r.e.) We write  $a \iff b$  to mean that  $A \iff B$  where  $A \in a$  and  $B \in b$ . (If  $A \equiv A'$  and  $B \equiv B'$ , then  $A \iff B$  iff  $A' \iff B'$ .) By Lemma 7,  $A \iff B$  implies  $A \iff B$ . As will become clear, the converse of this is not true. But if a is a and high, then a of a and a is a and high, then a of a and a is a and high, then a of a and a is a and a is a and high, then a of a and a is a and a if a is a and a is a and a if a is a and a is a and a if a is a if a is a and a if a is a if a if a is a if a if a if a is a if a if

**Lemma 21.** Suppose  $a \in F_T$  and for all  $\pi$ , if  $d(\pi) \le a$ , then  $d(\pi) \ll a$ . Then  $0_T \ll a$ .

**Proof.** By assumption there are  $\pi$ ,  $\sigma_0$ ,..., $\sigma_n$  such that  $a=d(\pi)\cup d(\sigma_0)\cup...\cup d(\sigma_n)$ . Also  $d(\pi)\ll a$ . Let  $A\in a$ . Then

(1)  $T + \sigma_i \le A \text{ for } i \le n.$ 

Moreover,  $\pi \ll \pi \land \sigma_0 \land ... \land \sigma_n$  and so, by Lemma 7,  $T + \pi \vdash \neg \sigma_0 \lor ... \lor \neg \sigma_n$ . But  $A \vdash \pi$  and so

(2)  $A \vdash \neg \sigma_0 \lor ... \lor \neg \sigma_n$ 

Let X be any set of  $\Sigma_1$  sentences such that

(3)  $A \dashv_{\Pi_1} T + X$ .

Then, by (2),  $T + X \vdash \neg \sigma_0 \lor ... \lor \neg \sigma_n$ , whence there is a  $k_0$  such that  $T + \sigma_0 \vdash \neg \land X \mid k_0 \lor \neg \sigma_1 \lor ... \lor \neg \sigma_n$ , and so, by (1) and (3),  $T + X \vdash \neg \sigma_1 \lor ... \lor \neg \sigma_n$ . Continuing in this way we eventually obtain the conclusion that T + X is inconsistent.

**Proof of Theorem 13.** We effectively construct sentences  $\psi_0$ ,  $\psi_1$ ,... such that if  $A_n = T + \{\psi_k : k < n\}$  and  $A = T + \{\psi_k : k \in N\}$ , then

- $(1) \qquad A_n \ll A_{n+1},$
- (2) not  $T \ll A$ .

Let a = d(A). Then for all  $\pi$ , if  $d(\pi) \le a$ , there is an n such that  $d(\pi) \le d(A_n)$ . Also  $d(A_n) \ll d(A_{n+1}) \le a$  and so  $d(\pi) \ll a$ . Thus, by (2) and Lemma 21,  $a \notin F_T$ .

There is an r.e. relation S(n,k,p,q) such that

(not T + 
$$\psi \ll$$
 T +  $\psi$  +  $\phi$ ) iff  $\exists p \forall q S(\psi, \phi, p, q)$ .

By Lemma 3.2 (b), there are a  $\Pi_1$  formula  $\sigma(x,y,z,u)$  and a  $\Sigma_1$  formula  $\sigma'(x,y,z,u)$  such that

- (3) if S(n,k,p,q), then  $T \vdash \sigma'(n,k,p,q)$ ,
- (4) Thus  $\sigma'(n,k,p,q) \to \sigma(n,k,p,q)$ ,
- (5) T + Y is consistent where Y =  $\{\neg \sigma(n,k,p,q): \text{not } S(n,k,p,q)\}$ .

Let  $A_0 = T$ . Suppose  $A_n$  has been defined and set  $\theta_n := \Lambda \{\psi_k : k < n\}$ . Then

- (6) not  $A_n \ll A_n + \varphi \text{ iff } \exists p \forall q S(\theta_n, \varphi, p, q).$
- By (3) and Lemma 5.2, there is a  $\Sigma_1$  formula  $\rho_n(x,y)$  such that
- (7)  $A_n \vdash \rho_n(\varphi, p) \rightarrow \sigma'(\theta_n, \varphi, p, q),$
- $\text{(8)} \qquad \text{if } \forall q S(\theta_n, \phi, p, q) \text{, then } \rho_n(\phi, p) \text{ is } \Pi_1\text{-conservative over } A_n.$
- By Theorem 5.4 (b), there is a formula  $\eta_n(x)$  such that
- (9)  $A_n + \eta_n(\phi)$  is a  $\Pi_1$ -conservative extension of  $A_n + \{\neg \rho_n(\phi, p): p \in N\}$ . Finally, let  $\psi_n$  be such that

(10)  $T \vdash \psi_n \leftrightarrow \eta_n(\psi_n)$ .

The formulas  $\rho_n(x,y)$ ,  $\eta_n(x)$  and the sentences  $\psi_n$  can be found effectively in n.

To prove (1) assume it is false. Then, by (6), there is a p such that  $\forall q S(\theta_n, \psi_n, p, q)$ . But then, by (8),  $\rho_n(\psi_n, p)$  is  $\Pi_1$ -conservative over  $A_n$ . By (9) and (10),  $A_{n+1} \vdash \neg \rho_n(\psi_n, p)$ . But, by Lemma 8, this implies that  $A_n \ll A_{n+1}$ , a contradiction. This proves (1).

Next we prove (2). Let Y be as in (5). Then T+Y is consistent. To prove that  $A\dashv_{\Pi_1} T+Y$  we first show that

(11)  $A_{n+1} + Y \dashv_{\Pi_1} A_n + Y$ .

Suppose  $A_{n+1} + Y \vdash \pi$ . Then there is a k such that

(12)  $A_{n+1} \vdash \neg \land Y \mid k \lor \pi$ .

By (1) and (6), for each p there is a  $q_p$  such that

(13) not  $S(\theta_n, \psi_n, p, q_p)$ .

By (12), (9), and (10),

 $A_n + \{\neg \rho_n(\psi_{n'}p): p \in N\} \vdash \neg \land Y \mid k \lor \pi.$ 

By (13), (4), (7),  $A_n + Y \vdash \neg \rho_n(\psi_{n'}p)$  for every p. It follows that  $A_n + Y \vdash \pi$ . This proves (11).

Since (11) holds for all n, it follows that  $A \dashv_{\Pi_1} T + Y$ . This proves (2) and so the proof is complete.

Let  $G_T$  be the set of degrees obtained from the  $\Sigma_1$  and the  $\Pi_1$  degrees by closing under  $\cap$  and  $\cup$ . It is an open problem if  $G_T \neq D_T$ .

The degree mentioned in Lemma 20 cannot be arbitrarily large: if a is high, there is a smallest  $\Sigma_1$  degree  $\geq$  a, namely  $1_T$ . Similarly, the degree a defined in the proof of Theorem 13 cannot be arbitrarily large; it is not  $>>> 0_T$ . This is explained, at least partially, by the following surprising:

**Theorem 14.** (a) Every sufficiently large degree is the l.u.b. of a  $\Sigma_1$  degree and a  $\Pi_1$  degree.

(b) Every sufficiently large degree is the l.u.b. of two  $\Sigma_1$  degrees.

**Proof.** We may assume that  $d(Con_T) < 1_T$ . By Lemma 9, it is sufficient to consider degrees a such that  $d(Con_T) \le a < 1_T$ . Let  $\pi_n := \forall u \delta_n(u)$ , where  $\delta_n(u)$  is PR, be  $\Pi_1$  sentences such that  $a = d(\{\pi_n : n \in N\})$ . We may assume that for all n,

- $(1) \qquad \mathsf{T} \vdash \pi_0 \to \mathsf{Con}_{\mathsf{T}},$
- $(2) \qquad T \vdash \pi_{n+1} \to \pi_n.$ 
  - (a) We define  $\Pi_1$  sentences  $\phi_n$  and  $\psi_n$  in the following way:
- (3)  $T \vdash \varphi_n \leftrightarrow \forall z (\Prf_T(\bigvee \{\varphi_k : k \le n\}, z) \to \exists u \le z \neg \delta_{n+1}(u)),$  $\psi_n := \forall u (\neg \delta_{n+1}(u) \to \exists z < u \Prf_T(\bigvee \{\varphi_k : k \le n\}, z)).$

It follows that

- (4)  $T \vdash \varphi_n \lor \psi_n$ ,
- (5)  $T + \varphi_n \wedge \psi_n \vdash \neg Pr_T(\bigvee \{\varphi_k : k \le n\}) \wedge \pi_{n+1}$
- $(6) T + \pi_{n+1} \vdash \psi_{n'}$

- (7)  $T + \neg \varphi_n \vdash \Pr_T(\bigvee \{\varphi_k : k < n\}),$
- (8)  $\neg \bigvee \{ \varphi_k : k < n \} \leq \pi_n$ .
- $(\bigvee \{\phi_k: k < 0\} := \bot.)$  (4), (5), (6) are standard.

Since  $\neg \phi_n$  is  $\Sigma_1$ , we have  $T + \neg \phi_n \vdash \Pr_T(\neg \phi_n)$ . Also, by (3),

$$T + \neg \phi_n \vdash \Pr_T(\bigvee \{\phi_k: k \le n\}).$$

But then (7) follows.

By Theorem 6.4, (8) follows from

(9)  $T + \pi_n \vdash \neg Pr_T(\vee \{\phi_k: k < n\}).$ 

By (1), (9) holds for n = 0. Suppose (9) holds for n = m. To show that it holds for n = m+1, we argue in T as follows: "Suppose  $\pi_{m+1}$ . Then, by (6),  $\psi_m$ . Also, by (2) and the inductive assumption,  $\neg Pr_T(\bigvee \{\phi_k: k < m\})$  and so, by (7),  $\phi_m$ . Finally, by (5),  $\neg Pr_T(\bigvee \{\phi_k: k < m+1\})$ , as desired." Thus, (9) holds for n = m+1. This proves (9) and so we have proved (8).

Next we show that for all n,

(10)  $T + \Lambda \{\psi_k : k < n\} + Con_T \vdash \varphi_n$ .

We first show that

- (11)  $T + Con_T \vdash \varphi_{0}$
- (12)  $T + \psi_n + \varphi_n \vdash \varphi_{n+1}$ .
- (11) follows from (7) with n = 0. (12) follows from (5) and (7).

Now (10) follows from (11) and (12).

Let  $a_0 = d(\{\neg \phi_k : k \in N\})$ ,  $a_1 = d(Con_T)$ ,  $a_2 = d(\{\phi_k : k \in N\})$ . Then, by Lemma 13,  $a_0$  is  $\Sigma_1$ .  $a_1$  is  $\Pi_1$ . By (8) and Orey's compactness theorem,  $a_0 \le a$  and, by hypothesis,  $a_1 \le a$ . But then  $a_0 \cup a_1 \le a$ . By (4) and (5),  $a_0 \cup a_2 \ge a$ . By (4) and (10),  $a_0 \cup a_1 \ge a_2$ . It follows that  $a_0 \cup a_1 \ge a$  and so  $a_0 \cup a_1 = a$ , as desired.  $\blacklozenge$ 

- (b) Let  $\theta_i$ , i=0,1, be as in Lemma 15. We define  $\Pi_1$  sentences  $\phi_n$  and  $\psi_n$  in the following way:
- (13)  $T \vdash \phi_n \leftrightarrow \forall z \big( \Prf_T(\theta_0 \lor \lor \{\phi_k : k \le n\}, z) \to \exists u \le z \neg \delta_{n+1}(u) \big),$   $\psi_n := \forall u \big( \neg \delta_{n+1}(u) \to \exists z < u \Prf_T(\theta_0 \lor \lor \{\phi_k : k \le n\}, z) \big).$

It follows that

- (14)  $T \vdash \varphi_n \lor \psi_{n'}$
- (15)  $T + \varphi_n \wedge \psi_n \vdash \neg Pr_T(\theta_0 \vee \bigvee \{\varphi_k : k \leq n\}) \wedge \pi_{n+1}$
- (16)  $T + \pi_{n+1} \vdash \psi_n$
- (17)  $T + \neg \varphi_n \vdash \Pr_T(\theta_0 \lor \bigvee \{\varphi_k : k < n\}),$
- $(18) \quad \neg \theta_0 \land \neg \bigvee \{ \phi_k : k < n \} \leq \pi_n.$

The proofs of (14) – (18) are almost the same as the proofs of (4) – (8).

Next we show that for all n,

(19)  $T + \Lambda \{\psi_k : k < n\} + \theta_0 \wedge \theta_1 \vdash \phi_n$ 

We first show that

- (20)  $T + \theta_0 \wedge \theta_1 \vdash \varphi_0$ ,
- $(21) \quad T + \psi_n + \varphi_n \vdash \varphi_{n+1}.$
- (20) follows from Lemma 15 (iii) and (iv) and (17) with n=0. (21) follows from (15) and (17).

Now (19) follows from (20) and (21).

Let  $a_0 = d(\neg \theta_0 + \{\neg \phi_k : k \in N\})$ ,  $a_1 = d(\theta_0)$ . Then  $a_0$  is  $\Sigma_1$ . By Lemma 15 (vi),  $a_1$  is  $\Sigma_1$ . By (18),  $a_0 \le a$  and, by Lemma 15 (v),  $a_1 \le a$ . But then  $a_0 \cup a_1 \le a$ . By Lemma 15 (ii), (14), (19), and (15),  $a_0 \cup a_1 \ge a$ . Thus,  $a_0 \cup a_1 = a$ , as desired.

The proof of Theorem 14 actually yields the following stronger result; Theorem 14' is also an improvement of Theorem 4.

**Theorem 14'.** (a) Suppose  $a \in CON_T$  and  $a \le b < 1_T$ . There is then a  $\Sigma_1$  degree c such that  $a \cup c = b$ .

(b) Suppose  $a \in CON_T$ . There are then degrees  $a_0$ ,  $a_1$  such that (i)  $a_0$  and  $a_1$  are both  $\Sigma_1$  and  $\Pi_1$ , (ii)  $a_0 \cap a_1 = 0_T$ , (iii)  $a_0 \cup a_1 = a$ , (iv) for every degree  $b \ge a$ , there is a  $\Sigma_1$  degree  $b_i$  such that  $a_i \cup b_i = b$ , i = 0, 1.

One way to strengthen Theorem 12 would be to show that there is a  $\Pi_1$  degree  $a > 0_T$  such that no  $\Sigma_1$  degree cups to a. This, however, is not the case:

**Theorem 15.** For every  $\Pi_1$  degree  $a > 0_T$ , there is a  $\Sigma_1$  (and  $\Pi_1$ ) degree which cups to a.

**Proof.** The following proof is similar to that of Theorem 11. Let  $\pi$  be such that  $a = d(\pi)$  and let  $\delta(u)$  be a PR formula such that  $\pi := \forall u \delta(u)$ . By Lemma 18, there is a PR formula  $\eta(x,y,z)$  such that for all  $\varphi$ ,  $\psi$ ,

- (1) if  $T + \varphi \vdash \pi$ , then  $T + \psi \vdash \neg \exists z \eta(\varphi, \psi, z)$ ,
- (2) if  $T + \varphi \not\vdash \pi$ , then  $\exists z \eta(\varphi, \psi, z)$  is  $\Pi_1$ -conservative over  $T + \psi$ .

Next let  $\theta$  and  $\chi$  be such that

(3)  $T \vdash \theta \leftrightarrow \forall u (\neg \delta(u) \to \exists z < u \eta(\chi, \theta, z)),$  $T \vdash \chi \leftrightarrow \forall z (\eta(\chi, \theta, z) \to \exists u \le z \neg \delta(u)).$ 

Then

- (4)  $T \vdash \theta \lor \chi$
- (5)  $T \vdash (\theta \land \chi) \rightarrow \pi$ .

We now show that

(6)  $T + \chi \not\vdash \pi$ .

Suppose not. Then, by (1) and (3),  $T + \theta \vdash \pi$ . But then, by (4),  $T \vdash \pi$ , contrary to assumption. This proves (6).

Now let

$$\sigma := \exists z (\eta(\chi, \theta, z) \land \forall u \leq z \delta(u)).$$

Then

$$T \vdash \sigma \leftrightarrow \exists z \eta(\chi, \theta, z) \land \theta.$$

By (3),  $d(\theta) \le a$ . By (2) and (6),  $\sigma = \theta$ . Thus,  $d(\sigma)$  is  $\Sigma_1$  and  $\Pi_1$ . Let  $b = a \cap d(\chi)$ . By (6),  $a \not\leq d(\chi)$  and so b < a. By (5),  $d(\sigma) \cup b = a$ . Thus,  $d(\sigma)$  cups to a. Also note that b (is  $\Pi_1$  and)  $d(\sigma) \cap b = 0_T$ .

The problem if for every degree  $a > 0_T$ , there is a  $\Sigma_1$  degree which cups to a

remains open. By Theorem 14, this is true of every sufficiently large degree.

Our next task is to show that the result of interchanging  $\Sigma_1$  and  $\Pi_1$  in Theorem 15 is false.

**Theorem 16.** There is a  $\Sigma_1$  degree  $a > 0_T$  such that no  $\Pi_1$  degree cups to a.

Let  $\xi(x)$  be as in Lemma 5.8 with n=1 and let  $a=d(\{\xi(k):k\in N\})$ . Then  $a>0_T$  and no  $\Pi_1$  degree cups to a (see the proof of Theorem 3 (a)). To obtain a  $\Sigma_1$  degree satisfying these conditions we first prove the following refinement of Lemma 5.8 (for n = 1).

**Lemma 22.** There are  $\Pi_1$  formulas  $\xi(x)$ ,  $\eta(x)$  and  $\Sigma_1$  sentences  $\chi_k$  such that

- (i) T⊬ ξ(k)
- (ii)  $T\vdash \eta(k) \rightarrow \xi(k)$ ,
- $T \vdash \xi(k+1) \rightarrow \eta(k)$ (iii)
- $\xi(k)$  is  $\Sigma_1$ -conservative over  $T + \neg \eta(k)$ , (iv)
- (v)  $\{\xi(\mathbf{k})\colon \mathbf{k}\in\mathbf{N}\}\equiv\{\chi_{\mathbf{k}}\colon \mathbf{k}\in\mathbf{N}\}.$

**Proof.** We combine the ideas of the proofs of Lemma 5.8 and Theorem 11. By Lemma 18, there is a PR formula  $\gamma(x,z)$  such that for all  $\varphi$ ,

- if T $\vdash \varphi$ , then T $\vdash \neg \exists z \gamma(\varphi, z)$ , (1)
- (2) if  $T \not\vdash \varphi$ , then  $\exists z \gamma(\varphi, z) \Pi_1$ –conservative over  $T + \varphi$ .

Let  $\delta(u)$  be an arbitrary PR formula. Let  $\kappa(z,u,x,y)$  and  $\nu(z,u,x,y)$  be  $\Pi_1$  formulas and  $\mu(z,u,x,y,v)$  a PR formula such that

- $T \vdash \kappa(z,u,x,y) \leftrightarrow \forall v \mu(z,u,x,y,v),$ (3)
- (4) $T \vdash \neg v(z,u,x,0),$
- $T \vdash \kappa(\delta, u, k, y) \leftrightarrow \nu(\delta, u, k, y) \lor \forall \nu([\Sigma_1](\neg \eta_\delta(k) \land \xi_\delta(k), v) \to \neg \mathrm{Prf}_T(\xi_\delta(k), v)),$ (5)
- (6)  $T \vdash \nu(\delta, u, k, y+1) \leftrightarrow \forall \nu(\neg \mu(\delta, u, k+1, y, v) \rightarrow \exists z < \max\{u, v\} \gamma(\eta_{\delta}(k), z)),$

$$\begin{split} \xi_\delta(x) &:= \forall u (\delta(u) \to \kappa(\delta, u, x, u \stackrel{\bullet}{-} x)), \\ \eta_\delta(x) &:= \forall u (\delta(u) \to v(\delta, u, x, u \stackrel{\bullet}{-} x)). \end{split}$$

As in the proof of Lemma 5.8, (5) implies that

$$(7) \qquad T \vdash \, \xi_{\delta}(k) \leftrightarrow \eta_{\delta}(k) \vee \, \forall v([\Sigma_{1}](\neg \eta_{\delta}(k) \wedge \xi_{\delta}(k), v) \rightarrow \neg Prf_{T}(\xi_{\delta}(k), v)).$$

Let

$$\eta_{s}'(x):=\forall u(\delta(u)\to \nu(\delta,u,x,(u\mathrel{\ \stackrel{\scriptscriptstyle \bullet}{\;}\ }(x+1))+1)).$$

Then, by (6),

$$(8) \qquad T\vdash \eta_{\delta}'(k) \leftrightarrow \forall uv \big(\delta(u) \wedge \neg \mu(\delta,u,k+1,u \mathrel{\dot{-}}(k+1),v) \to \exists z < max\{u,v\} \gamma(\eta_{\delta}(k),z) \big).$$

Let

$$\chi_{\delta,k} := \exists z \big( \gamma(\eta_{\delta}(k),z) \, \wedge \, \forall uv \leq z \neg (\delta(u) \, \wedge \, \neg \mu(\delta,u,k+1,u \, \mathring{\boldsymbol{-}} \, (k+1),v)) \big).$$

Then  $\chi_{\delta k}$  is  $\Sigma_1$  and (cf. Lemma 1.3)

$$(9) \text{ Th } \chi_{\delta,k} \longleftrightarrow \exists z \gamma(\eta_{\delta}(k),z) \land \forall uv \big(\delta(u) \land \neg \mu(\delta,u,k+1,u \div (k+1),v) \to \\ \exists z < \max\{u,v\}\gamma(\eta_{\delta}(k),z)\big)$$

and so, by (8),

 $(10) \quad T \vdash \chi_{\delta,k} \leftrightarrow \exists z \gamma (\eta_{\delta}(k),z) \land \eta_{\delta}'(k).$ 

We now show that

- (11)  $T \vdash \eta_{\delta}(k) \rightarrow \xi_{\delta}(k)$ ,
- (12) if  $T \vdash \xi_{\delta}(k)$ , then  $T \vdash \eta_{\delta}(k)$ ,
- (13)  $T \vdash \xi_{\delta}(k+1) \rightarrow \eta'_{\delta}(k)$ ,
- (14) if  $T \vdash \delta(u) \rightarrow u > k$ , then  $T \vdash \eta'_{\delta}(k) \leftrightarrow \eta_{\delta}(k)$ ,
- (15) if  $T \vdash \delta(u) \rightarrow u > k$  and  $T \vdash \eta_{\delta}(k)$ , then  $T \vdash \xi_{\delta}(k+1)$ .
- (11) follows from (7). (12) follows from (7) by the same argument as in the proof of Lemma 5.8. (13) follows by predicate logic from (3) and (8). (14) is obvious.

To prove (15), assume T $\vdash \delta(u) \to u > k$  and T $\vdash \eta_{\delta}(k)$ . Then, by (14), T $\vdash \eta'_{\delta}(k)$ . Also, by (1), T $\vdash \neg \exists z \gamma(\eta_{\delta}(k), z)$ . By (8), it follows that

$$T \vdash \forall uv(\delta(u) \rightarrow \mu(\delta, u, k+1, u - (k+1), v))$$

and so, by (3),  $T \vdash \xi_{\delta}(k+1)$ . This proves (15).

It can now be shown that

if  $\exists u \delta(u)$  is true, then  $T \not\vdash \xi_{\delta}(0)$ .

The proof of this from (4), (12), (15) is the same as that of (6) in the proof of Lemma 5.8.

As in the proof of Lemma 5.8 we can now find a PR formula  $\delta'(x)$  such that  $\exists u \delta'(u)$  is false and  $T \not\vdash \xi_{\delta'}(0)$ . Let  $\xi(x) := \xi_{\delta'}(x)$ ,  $\eta(x) := \eta_{\delta'}(x)$ ,  $\chi_k := \chi_{\delta'(k)}$ .

The verification of (i) – (iv) is now straightforward or much the same as in the proof of Lemma 5.8; this is where (13) is needed.

To prove (v), we first note that  $\{\xi(k): k \in \mathbb{N}\} \leq \{\chi_k: k \in \mathbb{N}\}$  follows from (10), (14), (11). Next suppose  $T + \{\chi_k: k \in \mathbb{N}\} \vdash \pi$ . There is then an m such that  $T + \chi_0 \vdash \chi_1 \land \dots \land \chi_m \to \pi$ . By (i) and (ii),  $T \nvdash \eta(0)$ . Hence, by (2),  $\exists z \gamma(\eta(0), z)$  is  $\Pi_1$ —conservative over  $T + \eta(0)$ . But then, by (10) and (14),  $T + \eta(0) \vdash \chi_1 \land \dots \land \chi_m \to \pi$ . By (10), (14), (ii), (iii),  $T + \chi_1 \vdash \eta(0)$ . It follows that  $T + \chi_1 \land \dots \land \chi_m \vdash \pi$ . Continuing in this way we eventually get  $T + \eta(m) \vdash \pi$  and so, by (iii),  $T + \xi(m+1) \vdash \pi$ . This shows that  $\{\chi_k: k \in \mathbb{N}\} \leq \{\xi(k): k \in \mathbb{N}\}$  and so (v) is proved.  $\blacksquare$ 

**Proof of Theorem 16.** Let  $\xi(x)$  and  $\eta(x)$  be as in Lemma 22. Let  $a = d(\{\xi(k): k \in N\})$ . Then, by Lemma 22 (i),  $a > 0_T$ . Also, by Lemma 22 (v) and Lemma 13, a is  $\Sigma_1$ . By Lemma 22 (ii),  $d(\xi(k)) \le d(\eta(k))$  for every k. That  $d(\xi(k))$  doesn't cup to  $d(\eta(k))$  now follows from Lemma 22 (iv).

Suppose b is  $\Pi_1$  and  $b \le a$ . Then, by Lemma 22 (ii) and (iii),  $b \le d(\xi(k))$ , for some k, and  $d(\eta(k)) \le a$ . Since  $d(\xi(k))$  doesn't cup to  $d(\eta(k))$ , it follows that b doesn't cup to a.

Note that if a is as in Theorem 16, then a does not cup to any  $\Pi_1$  degree. Indeed, let b be  $\Pi_1$  and  $\geq$  a. If a cups to b, there is a  $\Pi_1$  degree  $c \leq a$  which cups to b. But then c cups to a, contrary to assumption.

Finally, we prove Theorem 5 (and a bit more). We have already observed that  $d(\neg \pi)$  is the p.c. of  $d(\pi)$ . Thus, every  $\Pi_1$  degree has a p.c. It follows that, in terms of our classification of degrees, the following result is the best we can do.

**Theorem 17.** There is a  $\Sigma_1$  degree which has no p.c.

This is a consequence of the following strengthening of Lemma 20.

**Lemma 23.** There is a sentence  $\sigma$  such that  $\{b \ge d(\neg \sigma): b \text{ is } \Sigma_1\}$  has no g.l.b.

To prove this, we need another:

**Lemma 24.** Suppose  $\{\pi_k : k \in \mathbb{N}\}$  is r.e. and let  $G = \{d(\pi_k) : k \in \mathbb{N}\}$ . Suppose there is no finite subset H of G such that  $\bigcap$ H is a lower bound of G. Then G has no g.l.b.

**Proof.** Let  $X = \{\pi: T + \pi_k \vdash \pi \text{ for every } k\}$ . X is not r.e. This can be seen as follows. Let R(k,m) be a primitive recursive relation such that  $Y = \{k: \forall mR(k,m)\}$  is not r.e. and let  $\rho(x,y)$  be a PR binumeration of R(k,m). We may assume that  $Z = \{\pi_k: k \in N\}$  is primitive recursive; let  $\zeta(x)$  be a PR binumeration of Z. Finally, let  $\eta(x) := \{x \in N\}$ 

$$\forall z (\neg \rho(x,z) \to \exists u \le z (\zeta(u) \land Tr_{\Pi_1}(u)).$$

It is sufficient to show that

(1)  $Y = \{k: \eta(k) \in X\}.$ 

If  $k \in Y$ , then, clearly,  $\eta(k) \in X$ . Suppose  $k \notin Y$ . Let m be such that not R(k,m). Then  $T + \eta(k) \vdash \bigvee Z \mid m$ . By assumption, there is an n such that  $T + \pi_n \not\vdash \bigvee Z \mid m$  and so  $\eta(k) \notin X$ . Thus, (1) holds and so X is not r.e.

Suppose  $d(A) \le d(\pi_k)$  for every k. Then  $Th(A) \cap \Pi_1 \subseteq X$ . Since X is not r.e., it follows that there is a  $\pi \in X$  such that  $A \not\vdash \pi$ . But then  $\pi \le \pi_k$  for every k and  $d(\pi) \not\subseteq d(A)$ . Thus, d(A) is not the g.l.b. of G.

**Proof of Lemma 23.** From the proof of Theorem 11 it is clear that there are (primitive) recursive functions f(n) and g(n) such that if  $\pi$  is any  $\Pi_1$  sentence, then  $f(\pi)$  is a  $\Pi_1$  sentence,  $g(\pi)$  is a  $\Sigma_1$  sentence, and if  $T \not\vdash \pi$ , then  $T < T + f(\pi) \equiv T + g(\pi) \leq T + \pi$ .

We now define  $\pi_k$  and  $\sigma_k$  as follows. Let  $\pi_0$  be any  $\Pi_1$  sentence not provable in T. Next suppose  $\pi_k$  has been defined and  $T \not\vdash \pi_k$ . Let  $\psi$  be a  $\Pi_1$  sentence undecidable in  $T + \neg \pi_k$ . Then  $T < T + \pi_k \lor \psi < T + \pi_k$ . Let  $\sigma_k := g(\pi_k \lor \psi)$  and  $\pi_{k+1} := f(\pi_k \lor \psi)$ . Then  $T \not\vdash \pi_{k+1}$ .

For every k,

 $(1) \qquad \pi_{k+1} \le \sigma_k < \pi_k.$ 

By Theorem 5.4 (a), there is a sentence  $\sigma$  such that

(2)  $T + \sigma$  is a  $\Pi_1$ -conservative extension of  $T + \{ \neg \pi_k : k \in \mathbb{N} \}$ .

By (1) and (2),

(3)  $\neg \sigma \leq \sigma_k$ .

Moreover

(4) if b is  $\Sigma_1$  and  $b \ge d(\neg \sigma)$ , there is a k such that  $b \ge d(\pi_k)$ .

For suppose  $b = d(\chi) \ge d(\neg \sigma)$ , where  $\chi$  is  $\Sigma_1$ . Then  $T + \chi \vdash \neg \sigma$ , whence  $T + \sigma \vdash \neg \chi$ . But then, by (2), there is a k such that  $T + \neg \pi_k \vdash \neg \chi$ , whence  $T + \chi \vdash \pi_k$  and so  $b \ge d(\pi_k)$ .

Let  $G = \{d(\pi_k): k \in \mathbb{N}\}$ . If  $\{b \ge d(\neg \sigma): b \text{ is } \Sigma_1\}$  has a g.l.b. c, then, by (1), (3), (4), c is the g.l.b. of G. But from (1) it follows that no  $d(\pi_k)$  is a lower bound of G. Hence, by Lemma 24, G has no g.l.b. and so  $\{b \ge d(\neg \sigma): b \text{ is } \Sigma_1\}$  has no g.l.b.

**Proof of Theorem 17.** Let  $\sigma$  be as in Lemma 23. By Lemma 6, for all B,  $(T + \sigma) \downarrow B \le T$  iff  $B \le T + \chi$  for all  $\Sigma_1$  sentences  $\chi$  such that  $T + \chi \vdash \neg \sigma$ . But then the p.c. of  $d(\sigma)$ , if it had one, would also be the g.l.b. of  $\{b \ge d(\neg \sigma): b \text{ is } \Sigma_1\}$ . Thus, by Lemma 23,  $d(\sigma)$  has no p.c.

Every  $\Sigma_1$  degree is the p.c. of some degree. It is an open problem if the converse of this is true. If it is, the  $\Sigma_1$  degrees can be characterized in a purely algebraic way as those degrees that are p.c.s.

#### Exercises for Chapter 7.

In the following exercises we assume that PA¬T and that A, B, etc. are extensions of T.

1. Suppose  $G \subseteq D_T$ . G is independent if for any disjoint finite subsets  $G_0$  and  $G_1$  of G,  $\bigcap G_0 \not\leq \bigcup G_1$ . ( $\bigcap \emptyset = 1_T$ ,  $\bigcup \emptyset = 0_T$ .) (Thus, for example,  $\emptyset$  is independent and  $\{a\}$  is independent iff  $0_T < a < 1_T$ .) Show that for every finite independent set G, there are degrees  $b_0$ ,  $b_1$  such that  $G \cup \{b_i\}$  is independent, i = 0, 1, and  $b_0 \cap b_1 = 0_T$ . Conclude that every finite independent set is included in  $2^{\aleph_0}$  many maximal independent sets.

# 2. Suppose a < b.

- (a) c cups to b above a if there is a d such that  $a \le d < b$  and  $c \cup d = b$ . Show that there is a  $c \in (a,b]$  which doesn't cup to b above a.
- (b) c caps to a below b if there is a d such that  $a < d \le b$  and  $c \cap d = a$ . Show that there is a  $c \in [a,b)$  which doesn't cap to a below b.
- 3. Suppose a < b and  $b < 1_T$  if T is  $\Sigma_1$ -sound. For  $c \in [a,b]$ , let  $c^*$  be the *complement of* c *in* [a,b] if it exists, i.e.  $c \cap c^* = a$  and  $c \cup c^* = b$ . (Complements are unique.) Let  $Cpl_{a,b}$  be the set of degrees in [a,b] having complements in [a,b].
  - (a) Show that  $Cpl_{a,b}$  is closed under  $\cap$ ,  $\cup$ , and \*.
  - Let  $Cpl_{a,b} = (Cpl_{a,b}, \cap, \cup, *, a, b)$ . Then  $Cpl_{a,b}$  is a Boolean algebra.
- (b) Show that if c,  $d \in Cpl_{a,b}$  and c < d, there is an  $e \in Cpl_{a,b}$  such that c < e < d. (It follows that the Boolean algebras  $Cpl_{a,b}$  are (denumerable and) atomless and therefore isomorphic.)
- (c) Show that if  $a \le c < d \le b$ , there is an  $e \in [c,d)$  such that  $Cpl_{a,b} \cap [e,d) = \emptyset$ . [Hint:  $Cpl_{a,b} \cap [c,d] \subseteq Cpl_{c,d}$ .]
- (d) Show that if  $a \le c < e < d \le b$  and  $e \notin Cpl_{a,b}$ , there are c', d' such that  $c \le c' < e < d' \le d$  and  $Cpl_{a,b} \cap [c',d'] = \emptyset$ .

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- 4. Suppose a is  $\Sigma_1$ .
  - (a) Show that if  $a < b < 1_T$ , then a caps to  $0_T$  below b.
- (b) Show that if a < b and b is high, then a << b. Conclude that if  $b_i > a$ , i = 0, 1, and  $b_0 \cap b_1 = a$ , then  $b_0$  and  $b_1$  are low. ((a) and (b) are true of every a which is the p.c. of some degree,)
- 5. Show that for every low degree a, there is a low  $\Pi_1$  degree  $\geq$  a. [Hint: Let B = T +  $\sigma$  and  $\sigma$  :=  $\exists x \delta(x)$ , where  $\delta(x)$  is PR, be such that a  $\leq$  d(B) < 1<sub>T</sub>. We may assume that B $\not\vdash$   $\neg$ Con<sub>R</sub>. Let

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\begin{split} \theta &:= \forall y ( \operatorname{Prf}_B(\bot, y) \to \exists x \leq y \delta(x)), \\ \chi &:= \exists x (\delta(x) \land \forall y \leq x \neg \operatorname{Prf}_B(\bot, y)). \end{split} Then \sigma \leq \theta \leq \chi and T + \chi is consistent.]
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6. Referring to the proof of Theorem 4, show that there is a primitive recursive function g such that  $\psi$  can be replaced by the sentence

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\chi:=\forall u(Prf_B(\bot,u)\to\exists z{<}g(u)Prf_T(\bot,z)), similar to \theta. [Hint: Define g in such a way that PA\vdash \neg\phi\to\chi.]
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- 7. (a) Show that there is an r.p. degree a which is not  $\Pi_1$  (compare Lemma 11). [Hint: Let  $\kappa(x)$  be as in Exercise 2.11 and let  $a = d({\kappa(k):k \in \mathbb{N}})$ ].
- (b) Improve (a) by showing that there is a non- $\Pi_1$   $\Sigma_1$  degree a which is r.p. (compare Exercise 16 (c)). [Hint: Define  $\pi_k$  and  $\sigma_k$  so that  $T \not\vdash \pi_k$ ,  $T \vdash \kappa(k) \to \pi_k$ , where  $\kappa(x)$  is as in (a), and  $\pi_0 \land ... \land \pi_{k-1} \land \sigma_k \equiv \pi_0 \land ... \land \pi_k$ . Let  $a = d(\{\sigma_k : k \in N\})$ .]
- 8. Suppose A $\dashv$  B. Show that there is a  $\Delta_2$  sentence  $\varphi$  such that A +  $\varphi \simeq$  B (compare Corollary 6.10 and Theorem 8).
- 9. Show that there is a  $\Sigma_1$  sentence  $\sigma$  such that  $0_T < d(\sigma) < 1_T$  and for every  $\Sigma_1$  sentence  $\chi$ , if  $\sigma \le \chi$ , then  $T + \chi \vdash \sigma$ . [Hint: Let  $\sigma$  be such that  $d(\neg \sigma)$  is  $\Sigma_1$ .]
- 10. Let  $\langle \sigma_k \rangle_{k < \omega}$  and  $\sigma$  be as in Lemma 14. Show that every  $\Pi_1$  degree  $\leq$  d( $\sigma$ ) caps to  $0_T$  below d( $\sigma$ ) (compare Exercise 26 ( $\sigma$ )).
- 11. (a) Show that for every  $\Pi_1$  sentence  $\pi$ ,  $d(\pi) \cup d(\neg \pi)$  is high; in fact, if b is the p.c. of a, then  $a \cup b$  is high.
- (b) Let  $a = d(\sigma) \cup d(\neg \sigma)$ . a is high. Let  $\langle \sigma_k \rangle_{k < \omega}$  and  $\sigma$  be as in Lemma 14. Show that if b is  $\Pi_1$  and b  $\leq$  a, then b is low. [Hint: Use Exercise 4 (a).]
- 12. Show that there is a degree of the form  $d(\sigma \lor \pi)$  which is neither  $\Sigma_1$  nor  $\Pi_1$ .
- 13. Show that there is a  $\Pi_1$  degree a such that for every  $\Pi_1$  degree b, if  $a \cap b = 0_T$ , then  $a \cup b$  is low (compare Exercise 18 (c)).

- 14. Suppose  $a < b < 1_T$ . Show that
  - (a) there is a degree  $c < 1_T$  such that for every d, if  $b \cap d = a$ , then  $d \le c$ ,
  - (b) there is a degree  $c > 0_T$  such that for every d, if  $a \cup d = b$ , then  $d \ge c$ .
- 15. (a) Verify that in any distributive lattice, for any a, b, the intervals  $[a \cap b,a]$  and  $[b,a \cup b]$  are isomorphic.
- (b) Show that there are degrees a, b, c, d such that a  $\ll$  b, c < d, not c  $\ll$  d, and [a,b] and [c,d] are isomorphic. [Hint: Use Exercises 4 (b) and 11 (a).]
- 16. (a) Verify that in any distributive lattice, if a < b < c and [a,c] satisfies the reduction principle, so does [b,c].
- (b) Show that for each degree  $a < 1_T$ , there is a b such that  $a < b < 1_T$  and [a,b] does not satisfy the reduction principle.
- (c) The non–r.p. degree a defined in the proof of Lemma 12 is high (cf. Exercise 11 (a)). Show that there is a  $\Sigma_1$  degree which is not r.p. Conclude from Exercise 7 (b) that there are non– $\Pi_1$   $\Sigma_1$  degrees such that  $[0_T$ ,a] and  $[0_T$ ,b] are not isomorphic. [Hint: Use Theorem 14' (a).]
- 17. (a) Suppose  $\varphi$  and X are as in Lemma 16. Show that if  $\varphi \leq X$ , then  $\varphi \ll X$ .
- (b) Suppose a < b. Show that there are c, d such that a  $\leq$  c < d  $\leq$  b and [c,d] contains no B<sub>1</sub> degree.
- 18. (a) Show, by combining the proofs of Theorem 4 and Lemma 15, that there are cupping degrees  $a_0$  and  $a_1$  which are  $\Sigma_1$  and  $\Pi_1$  and such that  $a_0 \cap a_1 = 0_T$ . Conclude that there are low cupping degrees. (This also follows from Theorem 14′ (b).)
- (b) Show that there is a high  $(\Pi_1)$  degree a which is not cupping. [Hint: Suppose  $d(Con_T) < 1_T$ . Let  $a = d(\pi)$  where  $\pi$  is  $\Sigma_1$ -conservative over  $T + \neg Con_T$  and  $\neg \pi$  is  $\Pi_1$ -conservative over  $T + \neg Con_T$ .]
- (c) Show that there is a low  $(\Pi_1)$  degree a such that for every degree b, if  $a \cap b = 0_T$ , then  $a \cup b$  is not cupping (compare Exercise 13). [Hint: Let  $d(\pi)$  be as in (b). Define a sentence  $\sigma$  such that  $d(\sigma) > 0_T$  and  $d(\sigma) \cup d(\neg \sigma) \le d(\pi)$ ; use Theorem 11. Let  $a = d(\neg \sigma)$ .]
- 19. Show that there are degrees a, b such that a is  $\Sigma_1$ , b is both  $\Sigma_1$  and  $\Pi_1$ , and a  $\cup$  b is not  $B_1$ .
- 20. Prove Lemma 15 by letting  $\theta_0$  be a  $\Pi_1$  Rosser sentence for T and  $\theta_1 := \forall u(\Prf_T(\neg \theta_0, u) \to \exists z \leq u\Prf_T(\theta_0, z))$ . Conclude that  $d(\theta_0)$  is  $\Sigma_1$  (compare Exercise 6.9).
- 21. (a) Suppose  $a \in E_T$  and  $a > 0_T$ . Show that there is a degree b < a such that  $[b,a] \subseteq$

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E<sub>T</sub>. [Hint: We may assume that a  $\nleq$  d(Con<sub>T</sub>). Let b = a  $\cap$  d(Con<sub>T</sub>) and use Theorem 14' (b). (By Lemma 17, no member of [b,a) is  $\Sigma_1$ ).]

- (b) Suppose there is a  $\Sigma_1$  degree which cups to a. Show that there is a b < a such that for every  $c \in [b,a]$ , there is a  $\Sigma_1$  degree which cups to c.
- 22. (a) Let  $E_T'$  be the set of degrees obtained from  $0_T$  by taking l.u.b.s, g.l.b.s, and  $\Sigma_1$ -extensions. Show that if  $a \in E_T'$  there is a least  $\Sigma_1$  degree  $\geq a$ . Conclude that there is a  $\Pi_1$  degree not in  $E_T'$  (This improves Theorem 12.)
- (b) Let  $F_T'$  be the set of degrees obtained from  $E_T'$  and the  $\Pi_1$  degrees by taking l.u.b.s and  $\Sigma_1$ —extensions. Show that the degree defined in the proof of Theorem 13 is not in  $F_T'$  Conclude that there is a degree which is not the l.u.b. of a finite set of degrees of the form  $d(\pi \wedge \sigma)$ . (This improves Theorem 13.)
- 23. Show that for any a, if there is a member of  $G_T$  which cups to a, then there is a  $\Sigma_1$  degree which cups to a. (This improves Theorem 15.)
- 24. (a) Show that not all non– $\Pi_1$   $\Sigma_1$  degrees are as stated in Theorem 16.
- (b) Improve Theorem 16 by showing that for every degree  $b>0_T$ , there is a  $\Sigma_1$  degree a such that  $0_T< a< b$  and no  $\Pi_1$  degree cups to a. [Hint: By Theorem 11, there are sentences  $\pi$  and  $\sigma$  such that  $0_T< d(\pi)=d(\sigma)< b$ . Let  $C=T+\neg\pi$ . By the proof of Lemma 22, with T replaced by C, there are  $\Pi_1$  formulas  $\xi(x)$ ,  $\eta(x)$  and  $\Sigma_1$  sentences  $\chi_k$  such that (i) (iv) hold with T replaced by C and  $C+\{\xi(k): k\in N\}\equiv C+\{\chi_k: k\in N\}$ . Let  $a=d(\{\xi(k)\lor\pi: k\in N\})$ .]
- 25. Show that in contrast to Lemma 24 we have the following: There is a set  $G = \{d(\sigma_k): k \in N\}$  of  $\Sigma_1$  degrees, where  $\{\sigma_k: k \in N\}$  is (primitive) recursive, such that  $\bigcap H > 0_T$  for every finite subset H of G and  $\bigcap G = 0_T$ . [Hint: Let a be high and such that there is no high  $\Pi_1$  degree  $\leq a$  (cf. Exercise 11 (b)). Let  $A \in a$  and let  $\sigma_k := \neg Con_{A \mid k}$ .]
- 26. (a) Show that there is a PR formula  $\delta(u)$  such that if  $\theta$  is defined as in the proof of Theorem 11, then  $d(\neg \theta)$  isn't  $\Pi_1$ .
- (b) Let  $\theta$  be as in (a). Show that  $d(\neg \theta)$  has a p.c. Conclude that there is a non- $\Pi_1$   $\Sigma_1$  degree which has a p.c.
- (c) Let  $\theta$  be as in (a). Show that there is a  $\Pi_1$  degree < d( $\neg \theta$ ) which does not cap to  $0_T$  below d( $\neg \theta$ ) (compare Exercise 10).

#### Notes for Chapter 7.

The lattice  $\mathbf{D}_{\mathrm{T}}$  was introduced by Lindström (1979), (1984b); a related lattice  $\mathbf{V}_{\mathrm{T}}$  (degrees of finite extensions of T) has been defined by Švejdar (1978) (see also Jeroslow (1971a)). (By Theorem 6.11 (a),  $\mathbf{V}_{\mathrm{T}}$  and  $\mathbf{D}_{\mathrm{T}}$  are isomorphic.) Theorem 1 is due to Lindström (1979), (1984b) and (for  $\mathbf{V}_{\mathrm{T}}$ ) to Švejdar (1978). Corollary 1 is,

modulo Theorem 6.6, a restatement of the equivalence of Exercise 2.22 (i) and (ii). The proof of Theorem 4 was suggested by the proof of a related result in Hájková II (1971). Theorem 7 is new; the term "reduction principle" is borrowed from descriptive set theory and recursion theory (cf. Soare (1987)). (The only way of showing that intervals are isomorphic known so far is given in Exercise 15 (a) and works in all distributive lattices.) The remaining results of § 1 are due to Lindström (1979), (1984b). In connection with the proof of Theorem 4, see Exercise 6. Lemmas 11 and 12 lead to the question if there is a non– $\Pi_1$  r.p. degree; this question is answered in Exercise 7.

Theorem 8 (with a slightly different proof; see Exercise 6.12 (a)) is due to Montagna (cf. Lindström (1993)). Theorem 9 is due to Lindström (1979), (1984b), (1993); (a) and (c) were also proved by Švejdar (1978); for a different proof of Theorem 9 (d), see Exercise 12.

Theorem 10 is due to Lindström (1979), (1984b); (a) and the first half of (b) were also proved by Švejdar (1978). Theorems 14 and 16 are new, they were announced in Lindström (1993), where a weaker form of Theorem 16 is proved; Theorem 16 leads to the question if there is a  $\Sigma_1$  degree a such that no  $\Pi_1$  degree caps to a; this is answered negatively in Exercise 5; in connection with Theorem 16, see also Exercise 24. The remaining results of § 3 are due to Lindström (1984b), (1993). The definition of the sentences  $\phi_n$  and  $\psi_n$  in the proof of Theorem 14 (a) and the observations concerning these sentences, except (8), were first used by Misercque (1982) in a different context. For improvements of Theorems 12, 13, 15, and 16, see Exercises 22 (a), 22 (b), 23, 24 (b). Theorem 17 leads to the question if no non– $\Pi_1$   $\Sigma_1$  degree has a p.c.; this question is answered in Exercise 26 (b).

For a proof of Exercise 26 (a), see Lindström (1993).