

ENTROPY VERSUS ORBIT EQUIVALENCE FOR MINIMAL HOMEOMORPHISMS

MIKE BOYLE AND DAVID HANDELMAN

Every minimal homeomorphism of a Cantor set is strongly orbit equivalent to a homeomorphism of zero entropy. The dyadic adding machine is strongly orbit equivalent to homeomorphisms of all entropies.

1. Introduction and definitions. We construct, using the lexicographic or “adic” maps of Vershik (as modified by Herman, Putnam, and Skau [P], [HPS], [Sk]), examples of minimal homeomorphisms of all possible (topological) entropies (including infinite) which are “strongly” orbit equivalent to the dyadic adding machine. In particular, this answers in a decisively negative way the question as to whether orbit equivalent homeomorphisms have the same entropy. We also show that any minimal homeomorphism of the Cantor set is strongly orbit equivalent to one of zero entropy.

We recall some definitions. A source for these is the survey article of Skau [Sk]. A Bratteli diagram B is a directed graph whose vertex set decomposes into finite subsets, “levels” or “rows”, B_k ($k = 0, 1, \dots$), together with edges from vertices in B_k to vertices in B_{k+1} ; additionally, every vertex of B_k is joined to a vertex of B_{k+1} . The Bratteli diagram is *simple* if for all k , there exists $k' > k$ such that for every vertex v in B_k and every vertex v' in $B_{k'}$, there is a path from v to v' . The Bratteli diagram is *pointed* if $|B_0| = 1$, that is, there is a distinguished top vertex. The set of infinite paths, usually denoted X , is called the *Bratteli compactum*, and is a compact zero dimensional metric space, i.e., a compact Boolean space.

One can describe a Bratteli diagram using rectangular matrices with nonnegative integer entries. If $|B_k| = w(k)$ (with $w(0) = 1$), define a matrix M_k in $\mathbf{Z}^{w(k+1) \times w(k)}$ as follows. Index the vertices in B_k by an initial segment of the positive integers, and set the i, j entry to be the number of edges from point j of B_k to point i of B_{k+1} . One may then form the limit *dimension group*, the direct limit, as ordered abelian groups,

$$(1.1) \quad G_B = \lim \mathbf{Z}^{w(0)} \xrightarrow{M_0} \mathbf{Z}^{w(1)} \xrightarrow{M_1} \mathbf{Z}^{w(2)} \xrightarrow{M_2} \dots$$

Dimension groups may be defined in this fashion; it is clear how to recover the original Bratteli diagram. Note moreover, that there is a distinguished positive element, the image of the generator of the zeroth level, and this must be an order unit (a positive element of the dimension group such that every element of the dimension group is dominated by some multiple). A *telescoping* (also called “contraction” or “blocking”) of a Bratteli diagram is the diagram obtained by telescoping the maps; replace M_0 by $M'_0 = M_{n(0)}M_{n(0)-1} \cdots M_1M_0 : \mathbf{Z}^{w(0)} \rightarrow \mathbf{Z}^{w(n(0)+1)}$, M_1 by $M'_1 = M_{n(1)}M_{n(1)-1} \cdots M_{n(0)+s}M_{n(0)+1} : \mathbf{Z}^{w(n(0)+1)} \rightarrow \mathbf{Z}^{w(n(1)+1)}$, etc., where $\{n(0), n(1), \dots\}$ is a strictly increasing sequence of positive integers.

We consider certain partial orderings on the path space. An *ordered Bratteli diagram* is a Bratteli diagram with a partial ordering in which (infinite) cofinal paths are linearly ordered, but two paths which are not cofinal are not comparable. One fairly general construction is the reverse lexicographic ordering, given in [HPS; Definition 2.3]. Suppose at each level there is a partial ordering on the edges, in which all the edges to the same vertex are linearly ordered, but edges with differing target vertices are not comparable. Then the finite paths, say from level k to level l (with $l > k$), can be ordered lexicographically. Then the ordering extends to cofinal paths (finite or infinite), as follows: if x and y have the same initial vertex, then $x > y$ if there exists n such that $x_n > y_n$ and $x_i = y_i$ for all i exceeding n . In particular, if two paths are comparable in this ordering, then they must be cofinal. Maximal and minimal paths are defined in the obvious manner. Telescoping of ordered Bratteli diagrams preserve the ordering on the infinite paths.

In [V], Vershik developed these ideas in order to obtain a map on the path space, defined except at the maximal paths, that simply sends each path to the smallest path exceeding it. Herman, Putnam and Skau considered the special case that there are unique maximal and unique minimal paths. They observe that in that case the map could be extended to a homeomorphism of the entire path space by sending the maximal path to the minimal path. In this case, we shall call it the *lexicographic map* (obtained from the ordering). They showed that every minimal homeomorphism of a zero dimensional compact metrizable space arises from this construction [HPS].

The orbits of a path under the original (Vershik) map consists of the paths cofinal to it; under the extended map, the orbit of the maximal path consists of everything cofinal with either it or the minimal path.

Two spaces with homeomorphisms, (X_i, ϕ_i) ($i = 1, 2$) are *orbit equivalent* if there exists a homeomorphism $F: X_1 \rightarrow X_2$ that sends orbits to orbits. When there is an orbit equivalence, for each point x in X_1 , there exists an integer $n(x)$ such that $F\phi_1 x = \phi_2^{n(x)} Fx$. This defines a function $n: X_1 \rightarrow \mathbf{Z}$; similarly there is a function $m: X_2 \rightarrow \mathbf{Z}$ such that $F\phi_1^{m(x)} x = \phi_2 Fx$. Suppose that the spaces are compact and metrizable, and the homeomorphisms are topologically transitive (that is, have dense orbits). If n is continuous, ϕ_1 is topologically conjugate to ϕ_2 or $(\phi_2)^{-1}$ [B]; if there are no periodic points, n is continuous if it is bounded. The two homeomorphisms (X_1, ϕ_1) and (X_2, ϕ_2) are said to be *strongly orbit equivalent* if they are orbit equivalent, and the homeomorphism F has associated cocycles m and n each with at most one point of discontinuity.

Giordano, Putnam, and Skau have shown that there is a strong orbit equivalence between the lexicographic maps if and only if the limit dimension groups with distinguished order unit (1.1) are order isomorphic. More recently, they also showed that the lexicographic maps are orbit equivalent if and only if the dimension groups are order isomorphic after factoring out their infinitesimal elements [GPS].

The dyadic odometer (or “adding machine”) arises from the simple Bratteli diagram consisting of a single vertex and two edges at each level. There is (up to relabelling of the two edges at each level, which does not change the topological conjugacy class) just one ordering, and the lexicographic map is the adding machine. The limiting dimension group is obviously $\mathbf{Z}[\frac{1}{2}]$ with the natural ordering (as a subgroup of the reals), and with 1 in $\mathbf{Z}[\frac{1}{2}]$ as the order unit. We exhibit examples of simple Bratteli diagrams with orderings (with unique maximal and minimal paths) whose lexicographic maps have arbitrary entropy, and for which the limit dimension groups are all $\mathbf{Z}[\frac{1}{2}]$. In particular, they are all (strongly) orbit equivalent to the adding machine. As a particular consequence, the natural question as to whether entropy is an invariant for orbit equivalence is answered negatively in drastic fashion.

We focus on the dimension group $\mathbf{Z}[\frac{1}{2}]$ for technical simplicity. We conjecture that any minimal homeomorphism of the Cantor set is strongly orbit equivalent to homeomorphisms of all entropies. In fact we conjecture something much stronger. If T is a minimal homeomorphism of the Cantor set and S is an ergodic measure-preserving isomorphism of nonatomic Lebesgue probability space, then there should exist a homeomorphism T' strongly orbit equivalent to T

and an invariant probability for T' with respect to which T' is measurably isomorphic to S . This would generalize Dye's Theorem [D] (all such S are orbit equivalent in the measurable category). As orbit equivalent homeomorphisms have the same space of invariant Borel probabilities, the conjecture would also generalize Krieger's theorem ([K], see [DGS]) that every such S is measurably isomorphic to some minimal uniquely ergodic homeomorphism of the Cantor set; every measurable system S would be compatible with every space of invariant measures possible for minimal homeomorphisms of the Cantor set. These spaces are precisely the metrizable Choquet simplices, as follows from a theorem of Lazar and Lindenstrauss [LL] (a snappy proof of which can be deduced from [EHS; Theorem 2.2]) and the work of Herman, Putnam and Skau.

2. Finite entropy. We begin this section by showing that any Bratteli diagram can be telescoped so that the resulting diagram admits an ordering such that the corresponding lexicographic map has entropy zero. Then we construct lexicographic maps with arbitrary finite entropy, but with underlying dimension group $\mathbf{Z}[\frac{1}{2}]$; in particular, these lexicographic maps are all strongly orbit equivalent to the 2-adic adding machine, which of course has zero entropy.

We first recall subshifts and their entropies. Suppose that A is a finite set. Let $A^{\mathbf{Z}}$ be the set of bisequences $x = \dots x_{-1}x_0x_1\dots$ (with each x_i in A), equipped with the product topology. Then $A^{\mathbf{Z}}$ is a compact metrizable totally disconnected space, and the shift map $\sigma : A^{\mathbf{Z}} \rightarrow A^{\mathbf{Z}}$ given by $(\sigma x)_i = x_{i+1}$ is a homeomorphism. The restriction of σ to a closed shift invariant subset Y of $A^{\mathbf{Z}}$ is a *subshift*. If S is the restriction of σ to Y , then the topological entropy of S , $h(S)$, is the growth rate of the number of words of length n occurring in points of Y :

$$h(S) = \limsup_n \frac{\log \#\mathscr{W}_n(S)}{n},$$

where $\mathscr{W}_n(S) = \{y_1y_2\cdots y_n \mid \text{there exists } y = (y_i)_{i \in \mathbf{Z}} \text{ in } Y\}$. We suggest [Wa] or [DGS] as an introduction to topological entropy. This definition is given in [Wa; Theorem 7.13] or [DGS; Proposition 16.11].

If \mathscr{W} is a set of words, let $S(\mathscr{W})$ be the subshift whose domain Y is the set of all bisequences formed by concatenation of words in \mathscr{W} . For the purposes of this paper, it will suffice to assume that all words in \mathscr{W} have the same length, say l , except in Proposition 2.2, below. Thus a bisequence y is in Y if and only if there is a bisequence

... $W(i_{-1})W(i_0)W(i_1)\dots$ and an integer k with $0 \leq k < l$ such that for all j

$$y_{k+jl} \cdots y_{k+(j+1)l-1} = W(i_j).$$

LEMMA 2.1. *Suppose \mathscr{W} is a set consisting of m distinct words of length l ; form the subshift $S(\mathscr{W})$. Then*

$$h(S) = \frac{1}{l} \log m.$$

Proof. For positive integers k and n , suppose $kl \leq n < (k+1)l$. Then $m^k \leq \#\mathscr{W}_n(S) \leq lm^{(k+2)}$ (the last inequality holds because any word of length n must be a subword of a concatenation of $k+2$ words from \mathscr{W} , beginning at one of the l positions in the first word). Thus $\lim \log(\#\mathscr{W}_n(S))/n$ exists and equals $\frac{1}{l} \log m$. \square

We establish some notation and parameters for this and the following section. Let S be a lexicographic map on an ordered Bratteli compactum X . Let $\{1, \dots, m_k\}$ represent and order \mathscr{V}_k , the set of vertices of level k . Let A_k be the set of paths from \mathscr{V}_0 to \mathscr{V}_k . We use π_k to denote any obvious truncation map, obtained by restriction of paths to their initial segments (that is, starting at \mathscr{V}_0 and truncating at \mathscr{V}_k). Given k , to each point x in X , associate a bisequence $\{\pi_k(S^n x)\}$ in $A_k^{\mathbb{Z}}$, and let S_k be the subshift with domain $\{\{\pi_k(S^n x)\} \mid x \in X\}$. (We remark that S is the inverse limit of $\{S_k\}$.) Then

$$h(S) = \lim_k h(S_k)$$

(for example, [Wa; Theorem 7.6]).

An ordering on the path space of a Bratteli diagram is a *consecutive ordering* if whenever edges u , v , and w have the same terminal vertex, u and w have the same initial vertex, and $u < v < w$, then u and v have the same initial vertex.

PROPOSITION 2.2. *Let S be a lexicographic map of a Bratteli compactum with a consecutive ordering. Let m_k be the number of vertices at level k of the Bratteli diagram and let n_k be the minimum number of arcs from a vertex at level $k-1$ to a vertex at level k . Suppose that*

$$\lim_{k \rightarrow \infty} \frac{\log(n_{k+1} \cdot m_k)}{n_{k+1}} = 0.$$

Then S has zero entropy.

Proof. Recall that S_k denotes the subshift which is the quotient of S obtained by restricting to the first k levels. A symbol in this subshift is a path from the top level to a vertex at level k . Suppose u is a vertex at level k . Let p_1, \dots, p_s be the paths from the top level to u , listed in increasing order. Let $W(u)$ be the word of length s which is the concatenation $p_1 p_2 \cdots p_s$. Now also assume that v is a vertex at level $k+1$, that a_1, \dots, a_t are the edges from u to v , and that y is an infinite path in the Bratteli diagram such that $y_1 \cdots y_{k+1} = p_1 a_1$. If y' is the image of y in S_k , then it follows that $y'_0 \cdots y'_{st-1}$ is the word $W(u)$ concatenated t times. Therefore S_k is contained in the subshift $T = S(\mathscr{W})$, where

$$\mathscr{W} = \{(W(u))^t \mid u \text{ is a vertex at level } k \text{ where } n_{k+1} \leq t < 2n_{k+1}\}.$$

If there is a length-decreasing bijection from \mathscr{W} to a set \mathscr{W}' of distinct words, then $h(S(\mathscr{W}')) \geq h(S(\mathscr{W}))$. Let \mathscr{W}' be a collection of $m_k n_{k+1}$ distinct words of length n_{k+1} . Then

$$h(S_k) \leq h(S(\mathscr{W})) \leq h(S(\mathscr{W}')) = \frac{\log(m_k n_{k+1})}{n_{k+1}},$$

where the equality follows from Lemma 2.1. Therefore $h(S) = \lim h(S_k) = 0$. \square

THEOREM 2.3. *Any minimal homeomorphism of a Cantor set is strongly orbit equivalent to a homeomorphism of zero entropy.*

Proof. By [HPS], we may assume without loss of generality that our homeomorphism is given as a lexicographic map T defined by means of a nontrivial simple ordered pointed Bratteli diagram B . Given integers $t > s$, let $n(s, t)$ be the minimum number of paths from a vertex at level s to a vertex at level t . Then for s fixed, $n(s, t)$ goes to infinity as t does. Therefore we may define recursively a sequence $\{j(k)\}$ such that if B' is the telescoped diagram whose vertex set at level k is the vertex set of B at level $j(k)$, then in the notation of Proposition 2.2, we have $\lim_k \log(m_k n_{k+1}) / (n_{k+1}) = 0$.

The original map T is conjugate to the lexicographic map T' determined from B' with the order induced from B . To define the desired zero entropy map S , we choose any consecutive ordering on B' whose minimal and maximal paths match those of T' . Then $h(S) = 0$ and S is strongly orbit equivalent to T . \square

Our entropy zero result above automatically yields the existence of strongly orbit equivalent minimal homeomorphisms with different entropies, independently of the examples in this section and §3, merely from the existence of minimal homeomorphisms of arbitrary entropy [DGS]. (Of course, it is interesting to construct examples explicitly, and moreover, examples strongly orbit equivalent to the odometer.)

Now we exhibit lexicographic maps of all possible finite entropies strongly orbit equivalent to the dyadic odometer. To this end, we make a special assumption, which has the effect of guaranteeing that the underlying dimension group is of rank one: Suppose that for $k \geq 1$, there are n_k arcs from each vertex in \mathcal{V}_{k-1} to each vertex v in \mathcal{V}_k .

Let l_k be the number of paths from the unique vertex in \mathcal{V}_0 to each vertex in \mathcal{V}_k , so $l_k = n_k m_{k-1} l_{k-1}$. Suppose that for each v in \mathcal{V}_{k+1} we have a linear ordering on the $m_k n_{k+1}$ arcs to v from \mathcal{V}_k . This ordering determines an $m_k n_{k+1}$ -tuple σ_v with entries from $\{1, 2, \dots, m_k\}$ — $\sigma_v(i)$ is the initial vertex in \mathcal{V}_k of the i th arc to v . We say the arcs to vertices of \mathcal{V}_{k+1} have *distinct orderings* if for all v and v' in \mathcal{V}_{k+1} ,

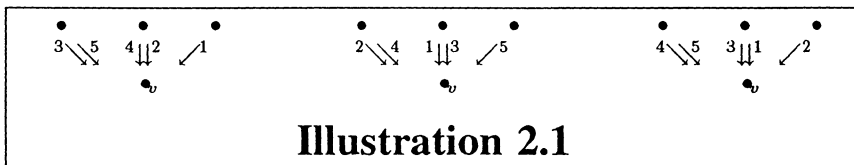
$$\sigma_v = \sigma_{v'} \quad \text{implies} \quad v = v'.$$

For example, in Illustration 2.1 below, for the leftmost configuration, $\sigma_v = (3, 2, 1, 2, 1)$; for the middle one, $\sigma_v = (2, 1, 2, 1, 3)$, and for the rightmost one, $\sigma_v = (2, 3, 2, 1, 1)$.

In order to accommodate the unique maximal and minimal path condition, we additionally require for each v in \mathcal{V}_k that $\sigma_v(1) = 1$ and $\sigma_v(m_k n_k) = m_k$ (note that none of the three examples in the illustration satisfy this!). As a technical convenience, we formulate the following condition:

- (2.4) If $k \geq 1$ and i and j are (not necessarily distinct) vertices in \mathcal{V}_k , there is a vertex v in \mathcal{V}_{k+1} and $t < m_k n_{k+1}$ such that $\sigma_v(t) = i$ and $\sigma_v(t + 1) = j$.

The l_k paths from \mathcal{V}_0 to a vertex v in \mathcal{V}_k are ordered as explained in §1. For v at some level, say k , let $\mathcal{W}(v)$ denote the word of length



l_k on symbols of A_k whose i th symbol is the i th path from \mathcal{V}_0 to v . Let $S_k(t) = S(\mathcal{W})$ for the collection of words $\mathcal{W} = \{\pi_k \mathcal{W}(v) \mid v \in \mathcal{V}_{k+t}\}$. So $S_k(t)$ is the subshift of A_k^Z generated by the truncations to the first k levels of the paths of length $k+t$ starting at \mathcal{V}_0 (and terminating at \mathcal{V}_{k+t}).

LEMMA 2.5. *Let S be a lexicographic map on an ordered Bratteli compactum as above, satisfying (2.4). Then for $k \geq 1$,*

$$h(S_k) = \lim_t h(S_k(t)).$$

Proof. Clearly $S_k(0) \supset S_k(1) \supset \dots$ and $S_k \subset \bigcap_{t \geq 1} S_k(t)$. Conversely, suppose y belongs to $\bigcap_{t \geq 1} S_k(t)$. By compactness, y belongs to S_k if for any positive integer n , the word $y_{-n} \dots y_n$ appears in a point of S_k . It suffices to show $y_{-n} \dots y_n$ occurs in the subshift generated by $\pi_k \mathcal{W}(v)$ for some vertex v .

Choose $i \geq k$ such that $l_i > 2n$. Since y belongs to $S_k(i-k)$, $y_{-n} \dots y_n$ is a subword of $\pi_k \mathcal{W}(v') \pi_k \mathcal{W}(v'')$ for some vertices v' and v'' in \mathcal{V}_i . By condition (2.4), $\pi_k \mathcal{W}(v') \pi_k \mathcal{W}(v'')$ is a subword of $\pi_k \mathcal{W}(v)$ for some v in \mathcal{V}_{i+1} .

Therefore S_k is the nested intersection of the subshifts $S_k(t)$. It is not hard to check that this implies $h(S_k) = \lim_t h(S_k(t))$ [DGS; p. 113, Corollary to Proposition 16.16]. \square

LEMMA 2.6. *Let S be a lexicographic map on an ordered Bratteli compactum as above, satisfying (2.4). Suppose that for all $k \geq 1$, the arcs from vertices of \mathcal{V}_{k+1} have distinct orderings. Then*

$$h(S) = \lim_k \frac{\log m_k}{l_k}.$$

Proof. Suppose $k \geq 1$ and $t \geq 0$. Let A_k^* denote the set of words which are concatenations of words in A_k . The hypothesis of distinct orderings implies that the map $\pi_k: \{\mathcal{W}(v) \mid v \in \mathcal{V}_{k+t}\} \rightarrow A_k^*$ is one to one. Therefore

$$h(S_k(t)) = h(S_{k+t}(0)) = \frac{1}{l_{k+t}} \log m_{k+t},$$

$$h(S) = \lim_k h(S_k) = \lim_k \lim_t h(S_k(t)) = \lim_k \lim_t \frac{1}{l_{k+t}} \log m_{k+t},$$

which finishes the proof. In fact, $h(S) = h(S_k)$ when $k \geq 1$. \square

To use Lemma 2.6, we control the differences, $\frac{1}{l_k} \log m_k - \frac{1}{l_{k+1}} \log m_{k+1}$. We require some notation to frame our estimates. Suppose m and n are integers exceeding 1. Let $M(m, n)$ be the number of mn -tuples σ whose entries are from $\{1, 2, \dots, m\}$ and satisfy the following conditions:

- (i) $\sigma(1) = 1$ and $\sigma(mn) = m$,
- (ii) each element of $\{1, 2, \dots, m\}$ occurs exactly n times in σ .

Let $M'(m, n)$ be the largest power of two less than or equal $M(m, n)$. Obviously $M'(m, n) \leq m^{mn}$.

LEMMA 2.7.

$$\lim_n \left(\frac{M'(m, n)}{m^{mn}} \right)^{1/mn} = 1.$$

Proof. We first note that $M'(m, n) \geq \frac{1}{2}M(m, n)$, so it suffices to show the limit is 1 when $M(m, n)$ replaces $M'(m, n)$. The number of mn -tuples on $\{1, 2, \dots, m\}$ satisfying (ii) above is $(mn)!/(n!)^m$. The fraction of such in which 1 is the first entry and m is the last is given by

$$c(m, n) = \frac{(mn-2)!}{(mn)!} \left(\frac{(n-1)!}{n!} \right)^2 = \frac{1}{(mn)(mn-1)n^2}.$$

Clearly $c(m, n)^{1/mn} \rightarrow 1$ as $n \rightarrow \infty$. So we are reduced to showing

$$\lim_n \frac{((mn)!/(n!)^m)^{1/mn}}{m} = 1.$$

The crude inequalities (valid for N at least 2),

$$(N/e)^N < N! < ((N+1)/e)^{N+1}$$

yield,

$$\begin{aligned} \frac{((mn)!/(n!)^m)^{1/mn}}{m} &> \frac{1}{m} \frac{(mn/e)}{((n+1)/e)((n+1)/e)^{1/n}} \\ &= \frac{n}{n+1} \left(\frac{e}{n+1} \right)^{1/n}, \end{aligned}$$

which obviously goes to 1 as n goes to infinity. The other inequality is similarly dealt with. \square

THEOREM 2.8. *Suppose $0 < \log \alpha < \infty$. There exists a homeomorphism S strongly orbit equivalent to the dyadic adding machine such that $h(S) = \log \alpha$.*

Proof. By the comments in §1, we need only construct a lexicographic map whose underlying dimension group is $\mathbf{Z}[\frac{1}{2}]$. To this end, choose m_1 a power of 2 such that $\log m_1 > \log \alpha$. Begin the construction of a Bratteli diagram by setting $\#\mathcal{V}_0 = 1$, $\mathcal{V}_1 = \{1, 2, \dots, m_1\}$, and adjoining one edge to each vertex in \mathcal{V}_1 from \mathcal{V}_0 . Now we describe the recursive step of the construction.

Begin with the ordered Bratteli diagram from \mathcal{V}_0 to \mathcal{V}_k with $\frac{1}{l_k} \log m_k - \log \alpha = \varepsilon > 0$. Pick a power of 2, n_{k+1} , such that the following hold:

$$(2.9) \quad (M'(m_k, n_{k+1})^{1/m_k n_{k+1}})^{1/l_k} > \alpha,$$

$$(2.10) \quad \frac{\log 2}{n_{k+1}} < \frac{\varepsilon}{2},$$

$$(2.11) \quad \alpha^{l_{k+1}} > m_k^2,$$

where $l_{k+1} = n_{k+1} m_k l_k$; we can achieve (2.9) as a consequence of Lemma 2.7.

Next, choose a power of 2, m_{k+1} , such that $m_{k+1} \leq M'(m_k, n_{k+1})$ and

$$0 < \frac{\log m_{k+1}}{l_{k+1}} - \log \alpha < \frac{\varepsilon}{2}.$$

This is possible since

$$\frac{\log M'(m_k, n_{k+1})}{l_{k+1}} = \log M'(m_k, n_{k+1})^{1/n_{k+1} m_k l_k} > \log \alpha \quad (\text{by 2.9})$$

and

$$\frac{\log 2^{i+1}}{l_{k+1}} - \frac{\log 2^i}{l_{k+1}} = \frac{\log 2}{n_{k+1} m_k l_k} < \frac{\varepsilon}{2} \quad (\text{by 2.10}).$$

Now assign to each vertex v of \mathcal{V}_{k+1} an $m_k n_{k+1}$ -tuple σ_v , determining the ordering of its incoming arcs. As $m_{k+1} \leq M'(m_k, n_{k+1})$, we can insist that $\sigma_v(1) = 1$ and $\sigma_v(m_k n_{k+1}) = m$, and the map $v \mapsto \sigma_v$ be injective. By (2.11) and since $m_{k+1} > \alpha^{l_{k+1}}$, we have enough vertices in \mathcal{V}_{k+1} to assign orderings of arcs which satisfy the technical condition (2.4).

This finishes the construction. Clearly $\lim_k \frac{1}{l_k} \log m_k = \log \alpha$ and by Lemma 2.6, this equals $h(S)$. Finally, since the number of vertices (m_k) at each level, and the multiplicity of arcs from each vertex to each vertex at the next level (n_k) are all powers of 2, the limiting dimension group is $\mathbf{Z}[\frac{1}{2}]$.

3. Infinite entropy. In this section, we construct an infinite entropy lexicographic map S strongly orbit equivalent to the dyadic adding machine. This is an elaboration of the finite entropy construction. As

in §2, \bar{m}_k is the number of vertices at level k , n_k is the number of arcs to each vertex in \mathcal{V}_k from each vertex in \mathcal{V}_{k-1} , and l_k is the number of paths from \mathcal{V}_0 to each vertex in \mathcal{V}_k . In addition, \mathcal{V}_k will be the disjoint union of sets $\mathcal{V}_k(t)$, $1 \leq t \leq j_k$ each of cardinality \bar{m}_k . The construction is recursive.

We begin with $m_1 = n_1 = 2$ and $j_1 = 1$. Given the ordered graph from \mathcal{V}_0 through \mathcal{V}_k , we construct \mathcal{V}_{k+1} and the ordered arcs from \mathcal{V}_k to \mathcal{V}_{k+1} ; we choose n_{k+1} , let $\bar{m}_{k+1} = M'(m_k, n_{k+1})$, and to each of the \bar{m}_{k+1} vertices v in $\mathcal{V}_{k+1}(1)$ assign a distinct ordering of the $m_k n_{k+1}$ arcs from \mathcal{V}_k to v , with each ordering beginning at the first vertex of \mathcal{V}_k (for some arbitrary linear ordering on \mathcal{V}_k) and ending at the last vertex of \mathcal{V}_k . This determines $l_{k+1} = n_{k+1} l_k m_k$. Then choose j_{k+1} with $m_{k+1} = j_{k+1} \bar{m}_{k+1}$ such that $\lim_k \frac{1}{l_{k+1}} \log j_{k+1} \bar{m}_{k+1} = \infty$.

For each t , $2 \leq t \leq j_{k+1}$, we assign the same set of arc orderings to $\mathcal{V}_{k+1}(t)$ as were assigned to $\mathcal{V}_{k+1}(1)$. We require n_{k+1} and j_{k+1} to be powers of 2, so S will be strongly orbit equivalent to the dyadic adding machine.

The basic idea is that if the numbers n_k are large enough, then $h(S_k)$ is not much less than $\frac{1}{k} \log m_k$, so $h(S) = \lim_k h(S_k) = \infty$. To carry this out, we first specify the parameters at the recursive step. Define positive numbers ε_k by $\prod_{i=1}^k (1 - \varepsilon_i) = \frac{1}{2} + 2^{-k-1}$, so

$$(3.1) \quad \prod_{i=1}^{\infty} (1 - \varepsilon_i) = \frac{1}{2}.$$

Suppose we are given the ordered Bratteli diagram from \mathcal{V}_0 through \mathcal{V}_k . By Lemma 2.7, we may choose a power of 2, n_{k+1} such that

$$(M'(m_k, n_{k+1}))^{1/m_k n_{k+1}} > (1 - \varepsilon_{k+1}) m_k.$$

Now $\bar{m}_{k+1} = M'(m_k, n_{k+1})$, $l_{k+1} = n_{k+1} l_k m_k$, and

$$(3.2) \quad (\bar{m}_{k+1})^{1/l_{k+1}} > (1 - \varepsilon_{k+1})^{1/l_k} (m_k)^{1/l_k}.$$

Then choose a power of 2, j_{k+1} , such that $(j_{k+1})^{1/l_{k+1}} > 2$. Now $m_{k+1} = \bar{m}_{k+1} j_{k+1}$. This finishes the construction.

Now we observe that

$$\begin{aligned} h(S_{k+1}(0)) &= \frac{1}{l_{k+1}} \log m_{k+1} = \frac{1}{l_{k+1}} \log \bar{m}_{k+1} j_{k+1} \\ &= \log \bar{m}_{k+1}^{1/l_{k+1}} + \log j_{k+1}^{1/l_{k+1}} \\ &> \log[(1 - \varepsilon_{k+1}) m_k^{1/l_k}] + \log j_{k+1}^{1/l_{k+1}} \\ &> h(S_k(0)) + \log 2 + \log(1 - \varepsilon_{k+1}). \end{aligned}$$

Therefore $\lim_k h(S_k(0)) = \infty$.

Next, given $k \geq 1$ and $t \geq 1$, let

$$J = \left(\prod_{i=1}^t j_{k+i}^{1/l_{k+i}} \right)^{l_{k+t}}.$$

We claim

$$(3.3) \quad \#\{\pi_k W(v) \mid v \in \mathcal{V}_{k+t}\} \geq \frac{1}{J} m_{k+t}.$$

To prove (3.3), it suffices to show that the map

$$\begin{aligned} \{W(v) \mid v \in \mathcal{V}_{k+t}\} &\rightarrow \{W(v) \mid v \in \mathcal{V}_k\}^* \\ W &\mapsto \pi_k(W) \end{aligned}$$

is at most J to one. To this end, we remark that for each v in \mathcal{V}_{k+t} , for each i between 1 and t , there are vertices v_r ($1 \leq r \leq a(i)$) such that $\pi_{k+i}(W(v))$ is a concatenation $W(v_1) \cdots W(v_{a(i)})$, where v_r belongs to $\mathcal{V}_{k+i}(t_r)$ and $a(i) = l_{k+t}/l_{k+i}$. Given i , the number of possible sequences $t_1 \cdots t_{a(i)}$ is $j_{k+i}^{a(i)}$. The number of ways to fix these sequences for all i with $1 \leq i \leq t$ is at most $\prod_{i=1}^t j_{k+i}^{a(i)} = J$.

However, given one of these J choices, and given $\pi_k(W(v))$, we can recover $W(v)$, since $\pi_k(W(v))$ and the sequence t_r for $i = 1$ determine $\pi_{k+1}(W(v))$, by the distinctness of the arc orderings chosen for $\mathcal{V}_{k+1}(1)$; similarly, we recover $\pi_{k+2}(W(v)), \dots, \pi_{k+t}(W(v)) = W(v)$. This proves (3.3).

Finally, suppose $k, t \geq 1$. Then

$$\begin{aligned} h(S_k(t)) &= \frac{1}{l_{k+t}} \log(\#\{\pi_k(W(v)) \mid v \in \mathcal{V}_{k+t}\}) \\ &\geq \frac{1}{l_{k+t}} \log \frac{m_{k+t}}{J} \quad (\text{by 3.3}) \end{aligned}$$

and therefore

$$(3.4) \quad h(S_k(t)) \geq \log \left(\frac{m_{k+t}}{J} \right)^{1/l_{k+t}}.$$

By (3.2), we have

$$\begin{aligned} m_{k+t}^{1/l_{k+t}} &= j_{k+t}^{1/l_{k+t}} \bar{m}_{k+t}^{1/l_{k+t}} \\ &> j_{k+t}^{1/l_{k+t}} (1 - \varepsilon_{k+t})^{1/l_{k+t-1}} m_{k+t-1}^{1/l_{k+t-1}} \\ &> \cdots > \left(\prod_{i=1}^t j_{k+i}^{1/l_{k+i}} \right) \left(\prod_{i=1}^t (1 - \varepsilon_{k+i})^{1/l_{k+i-1}} \right) m_k^{1/l_k}. \end{aligned}$$

Therefore

$$m_{k+t}^{1/l_{k+t}} > J^{1/l_{k+t}} \cdot \frac{1}{2} \cdot m_k^{1/l_k}.$$

Substituting this into (3.4) yields

$$h(S_k(t)) > h(S_k(0)) - \log 2;$$

thus $h(S) = \lim_k \lim_t h(S_k(t)) = \infty$.

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UNIVERSITY OF MARYLAND
COLLEGE PARK, MD 20742
E-mail address: mmb@math.umd.edu

AND

UNIVERSITY OF OTTAWA
OTTAWA, ONTARIO, K1N 6N5
CANADA
E-mail address: dehsg@ACADVMI1.uottawa.ca

