

## INTEGERS REPRESENTABLE AS THE SUM OF POWERS OF THEIR PRIME FACTORS

JEAN-MARIE DE KONINCK<sup>1</sup> & FLORIAN LUCA<sup>2</sup>

**Abstract:** Given an integer  $\alpha \geq 2$ , let  $S_\alpha$  be the set of those positive integers  $n$ , with at least two distinct prime factors, which can be written as  $n = \sum_{p|n} p^\alpha$ . We obtain general results

concerning the nature of the sets  $S_\alpha$  and we also identify all those  $n \in S_3$  which have exactly three prime factors. We then consider the set  $T$  (resp.  $T_0$ ) of those positive integers  $n$ , with at least two distinct prime factors, which can be written as  $n = \sum_{p|n} p^{\alpha_p}$ , where the exponents

$\alpha_p \geq 1$  (resp.  $\alpha_p \geq 0$ ) are allowed to vary with each prime factor  $p$ . We examine the size of  $T(x)$  (resp.  $T_0(x)$ ), the number of positive integers  $n \leq x$  belonging to  $T$  (resp.  $T_0$ ).

**Keywords:** Prime factorization

### 1. Introduction

Identifying all those positive integers  $n$  such that

$$n = \sum_{p|n} p^\alpha \tag{1}$$

for some integer  $\alpha \geq 2$  is certainly a difficult problem. Since prime powers  $p^\alpha$  (with  $\alpha \geq 2$ ) trivially satisfy (1), we shall examine the set  $S_\alpha$ , namely the set of those positive integers  $n$  satisfying (1) but which have at least two distinct prime factors.

We first obtain general results concerning the nature of the sets  $S_\alpha$ . We then identify all those  $n \in S_3$  which have exactly 3 prime factors. We further consider the more general equation

$$n = \sum_{p|n} p^{\alpha_p}, \tag{2}$$

---

**2001 Mathematics Subject Classification:** 11A41, 11A25

<sup>1</sup> Research supported in part by a grant from NSERC.

<sup>2</sup> Research supported in part by projects SEP-CONACYT 37259-E and 37260-E.

where the exponents  $\alpha_p$  are allowed to vary with each prime factor  $p$ . Clearly all prime powers have such a representation (2). So let us define  $T$  (resp.  $T_0$ ) as the set of all positive integers  $n$  having a representation (2) with each  $\alpha_p \geq 1$  (resp.  $\alpha_p \geq 0$ ) but with at least two distinct prime divisors. We obtain a non trivial upper bound for the number  $T_0(x)$  of positive integers  $n \leq x$  belonging to  $T_0$ .

Finally, we give a heuristic argument yielding lower and upper estimates for  $T(x)$ , the number of positive integers  $n \leq x$  belonging to  $T$ .

## 2. General observations

For each integer  $n \geq 2$ , let  $\omega(n)$  stand for the number of distinct prime factors of  $n$  and let  $P(n)$  stand for the largest prime factor of  $n$ . We first make the following observations. Given  $\alpha \geq 2$  and  $n \in S_\alpha$ , we have:

- (i)  $P(n) < n^{1/\alpha}$ .
- (ii) Letting  $r = \omega(n)$ , then  $r \geq 3$  and  $r$  is odd; this is easily established by considering separately the cases “ $n$  odd” and “ $n$  even”.
- (iii) If  $\alpha$  is even, then  $\omega(n)$  cannot be a multiple of 3; one can see this by considering separately the cases “ $3|n$ ” and “ $3 \nmid n$ ”.
- (iv) If  $\omega(n) = \alpha$ , then  $n$  cannot be squarefree, since otherwise, comparing the arithmetic mean with the geometric mean of the prime factors of  $n$ , we get

$$n = q_1 q_2 \dots q_\alpha = q_1^\alpha + q_2^\alpha + \dots + q_\alpha^\alpha \geq \alpha q_1 q_2 \dots q_\alpha = \alpha n,$$

a contradiction, since  $\alpha \geq 2$ .

- (v) If  $n \in S_2$ , then, in view of (ii) and (iii),  $r := \omega(n)$  is odd,  $r \geq 5$ ; moreover:
  - \* if  $r = 5$ , then  $n \equiv 5$  or  $8 \pmod{24}$ ,
  - \* if  $r = 7$ , then  $n \equiv 7, 10, 15$  or  $18 \pmod{24}$ ,
  - \* otherwise  $r \geq 11$ .
- (vi) A computer search shows that  $S_3$  contains at least 6 elements, namely:

$$378 = 2 \cdot 3^3 \cdot 7 = 2^3 + 3^3 + 7^3,$$

$$2548 = 2^2 \cdot 7^2 \cdot 13 = 2^3 + 7^3 + 13^3,$$

$$2836295 = 5 \cdot 7 \cdot 11 \cdot 53 \cdot 139 = 5^3 + 7^3 + 11^3 + 53^3 + 139^3,$$

$$4473671462 = 2 \cdot 13 \cdot 179 \cdot 593 \cdot 1621 = 2^3 + 13^3 + 179^3 + 593^3 + 1621^3,$$

$$23040925705 = 5 \cdot 7 \cdot 167 \cdot 1453 \cdot 2713 = 5^3 + 7^3 + 167^3 + 1453^3 + 2713^3,$$

$$21467102506955 = 5 \cdot 7^3 \cdot 313 \cdot 1439 \cdot 27791 = 5^3 + 7^3 + 313^3 + 1439^3 + 27791^3.$$

- (vii) If  $n \in S_4$ , then  $\omega(n) = 7$  or  $\omega(n) \geq 11$ . To show this, first let  $r = \omega(n)$ . We know from (ii) that  $r \geq 3$  and odd; but from (iii), it follows that  $r \neq 3$ ; hence,  $r \geq 5$ . But  $r \neq 5$ ; indeed, if  $r = 5$ , then first assume that  $5|n$ ; in this case, since  $p^4 \equiv 1 \pmod{5}$  for all primes  $p \neq 5$ ,

$$n = 625 + q_2^4 + q_3^4 + q_4^4 + q_5^4 \equiv 0 + 4 = 4 \pmod{5},$$

which contradicts  $5|n$ ; on the other hand, if  $n$  is not a multiple of 5, then  $n \equiv 5 \pmod{5}$ , again a contradiction. Hence,  $r \geq 7$ . Finally, in view of (iii),  $r \neq 9$ . Hence, we may conclude that  $r = 7$  or  $r \geq 11$ .

- (vii) It is not known if  $T$  is an infinite set. However, if there exist infinitely many primes  $p$  of the form  $p = \frac{2^k + 3^\ell}{5}$ , then  $\#T = +\infty$ , the reason being that in this case, we have  $2 \cdot 3 \cdot p = 2^k + 3^\ell + p$ .
- (viii) Using a parity argument, it is clear that any number  $n \in T$  has an odd number of distinct prime divisors. One can check that the smallest element of  $T$  is 30; in fact, 30 has two representations of type (2), namely

$$30 = 2 \cdot 3 \cdot 5 = 2 + 3 + 5^2 = 2^4 + 3^2 + 5.$$

Letting  $T(x) := \#\{n \leq x : n \in T\}$ , a computer search shows that  $T(100) = 6$ ,  $T(10^3) = 42$ ,  $T(10^4) = 109$ ,  $T(10^5) = 321$  and  $T(10^6) = 973$ . On the other hand, the smallest odd element of  $T$  is 915, in which case we have

$$915 = 3 \cdot 5 \cdot 61 = 3^6 + 5^3 + 61.$$

### 3. Identifying those $n \in S_3$ with $\omega(n) = 3$

**Theorem 1.** *If  $n \in S_3$  and  $\omega(n) = 3$ , then  $n = 2 \cdot 3^3 \cdot 7$  or  $n = 2^2 \cdot 7^2 \cdot 13$ .*

**Proof.** We prove this in 9 steps.

1. Write  $x < y < z$  for the three distinct prime factors of  $n$ . Note that the given relation forces  $z|y^3 + x^3$ , so that  $z|y + x$  or  $z|y^2 - yx + x^2$ , and similarly  $y|z + x$ , or  $y|z^2 - zx + x^2$ , and  $x|z + y$ , or  $x|z^2 - zy + y^2$ .

2. Assume  $z|y + x$ . Since  $y + x < 2y < 2z$ , this is possible only when  $z = y + x$ . If  $x > 2$ , then  $y + x$  is even, and so it cannot be an odd prime. Thus,  $x = 2$ ,  $z = y + 2$ , but then

$$x^3 + y^3 + z^3 = 8 + y^3 + (y + 2)^3 \equiv 16 \pmod{y},$$

which is impossible. Thus,  $z \nmid y + x$ , and  $z|y^2 - yx + x^2$ . Since  $z > 3$ , we also conclude that  $z \equiv 1 \pmod{3}$ , because the relation  $y^2 - yx + x^2 \equiv 0 \pmod{z}$  implies that  $(2y - x)^2 \equiv -3x^2 \pmod{z}$ , which means that  $\left(\frac{-3}{z}\right) = 1$ , which is equivalent to the fact that  $z \equiv 1 \pmod{3}$ . Here, and in what follows, for an odd prime  $p$  and an integer  $a$  we use  $\left(\frac{a}{p}\right)$  for the Legendre symbol of  $a$  in respect to  $p$ .

3. Assume that  $z^2|n$ . In this case, we then get  $z^2|y^3 + x^3$ , and by the previous arguments, it follows that  $z^2|y^2 - yx + x^2$ . This is impossible because  $y^2 - yx + x^2 = y^2 - x(y - x) < y^2 < z^2$ . Thus,  $z \nmid n$ .

4. Assume that  $y|z+x$ . Write  $z := \lambda y - x$ , with some positive integer  $\lambda$ . Clearly  $\lambda \geq 2$ . We then get  $x \equiv \lambda y \pmod{z}$ . Since we also have  $y^2 - yx + x^2 \equiv 0 \pmod{z}$ , we get  $y^2 - y(\lambda y) + (\lambda y)^2 \equiv 0 \pmod{z}$ . Thus,  $z|y^2(1 - \lambda + \lambda^2)$ , and therefore  $z|1 - \lambda + \lambda^2$ . If  $\lambda = 2$ , we get  $z|1 - 2 + 2^2 = 3$ , which is impossible. If  $\lambda = 3$ , we get  $z|1 - 3 + 3^2 = 7$ . Thus,  $z = 7$ , and therefore  $7 = 3y - x$ . Since  $y$  is odd, we get  $x = 2$  and therefore  $y = 3$ , which does give the solution

$$2^3 + 3^3 + 7^3 = 2 \cdot 3^3 \cdot 7$$

mentioned in the statement of our theorem.

Assume now that  $\lambda \geq 4$ . Then,

$$z = \lambda y - x > (\lambda - 1)y = \lambda y \cdot \frac{\lambda - 1}{\lambda} \geq \frac{3\lambda y}{4}.$$

Since  $z|1 - \lambda + \lambda^2$ , we also get

$$\lambda^2 > 1 - \lambda + \lambda^2 \geq z > \frac{3\lambda y}{4},$$

and therefore that

$$\lambda > \frac{3y}{4}.$$

Thus,

$$z > \frac{3\lambda y}{4} > \frac{9y^2}{16}.$$

Since we also have  $z|y^2 - yx + x^2$ , we get that

$$\delta = \frac{y^2 - yx + x^2}{z}$$

is a positive integer. However,

$$\delta < \frac{y^2}{z} < \frac{16}{9} < 2,$$

therefore  $\delta = 1$ , and so

$$z = y^2 - yx + x^2.$$

Thus,

$$n = x^3 + y^3 + z^3 = (y+x)(y^2 - yx + x^2) + z^3 = z(y+x) + z^3,$$

therefore

$$\frac{n}{z} = y + x + z^2.$$

Looking at this last relation modulo  $y$ , we get  $x+z^2 \equiv 0 \pmod{y}$ . Since  $y|x+z$ , we also get  $z \equiv -x \pmod{y}$  and therefore  $z^2 \equiv x^2 \pmod{y}$ . Thus,  $x^2+x \equiv 0 \pmod{y}$ ; hence,  $y|x(x+1)$ . This is possible only when  $y = x+1$  and  $x = 2$ . Thus,  $x = 2$ ,  $y = 3$ ,  $z = 3^2 - 2 \cdot 3 + 2^2 = 7$ , so that  $\lambda = 3$ , contradicting the fact that  $\lambda \geq 4$ .

5. From now on, we may assume that  $y \nmid z+x$  and therefore that  $y|z^2 - zx + x^2$ . If  $y = 3$ , then  $x = 2$ , in which case  $z|2^3 + 3^3 = 35$ ; hence,  $z = 7$  (because  $z \equiv 1 \pmod{3}$ ), which is a case already treated. Thus, we may assume that  $y > 3$ , and since  $y|z^2 - zx + x^2$ , an argument similar to the one employed at step 2 tells us that  $y \equiv 1 \pmod{3}$ .

6. Here, we observe that  $x \equiv 2 \pmod{3}$ . Indeed, for if not, we must either have  $x = 3$ , which is impossible because then  $3|n$ , but  $x^3 + y^3 + z^3 \equiv 2 \pmod{3}$ , or  $x \equiv 1 \pmod{3}$ , therefore  $3 \nmid n$ , while  $x^3 + y^3 + z^3 \equiv 0 \pmod{3}$ .

7. Write  $n := x^\alpha y^\beta z$ . Since we already know that  $x \equiv 2 \pmod{3}$  and  $y \equiv z \equiv 1 \pmod{3}$ , we reduce the relation

$$x^3 + y^3 + z^3 = x^\alpha y^\beta z$$

modulo 3 to get  $1 \equiv 2^\alpha \pmod{3}$ . This shows that  $\alpha$  is even.

8.1. Assume that  $x = 2$ . We first show that  $\alpha = 2$ . Indeed, for if not, we would first get  $8 | y^3 + z^3$  and hence that  $8|(z+y)(z^2 - zy + y^2)$ . Since  $z^2 - zy + y^2$  is odd, we get  $8|y+z$ . Thus,  $(y, z) \in \{(1, 7), (7, 1), (3, 5), (5, 3)\} \pmod{8}$ .

We know that  $z|y^3 + 2^3$ , and  $y|z^3 + 2^3$ . In particular,  $-2y \equiv (4/y)^2 \pmod{z}$ , and so

$$\left(\frac{-2y}{z}\right) = 1,$$

and in a similar way one deduces that

$$\left(\frac{-2z}{y}\right) = 1.$$

Hence, we have

$$\begin{aligned} 1 &= \left(\frac{-1}{z}\right) \left(\frac{-1}{y}\right) \left(\frac{2}{y}\right) \left(\frac{2}{z}\right) \left(\frac{y}{z}\right) \left(\frac{z}{y}\right) = (-1)^{\binom{z-1}{2} + \binom{y-1}{2} + \binom{z^2-1}{8} + \binom{y^2-1}{8} + \binom{(y-1)(z-1)}{4}} \\ &= (-1)^{1+0+0} = -1, \end{aligned}$$

a contradiction. Therefore,  $\alpha = 2$ .

8.2. Here, we show that  $\beta \in \{2, 3\}$ . If  $\beta = 1$ , we get

$$4yz = 2^3 + y^3 + z^3 > 3(2 \cdot y \cdot z) = 6yz,$$

which is impossible, the above inequality following from the AGM-inequality. Using now the fact that  $z|y^2 - yx + x^2$  (see step 2), together with the fact that  $y^2 - xy + x^2 = y^2 - x(y-x) < y^2$ , we learn that  $z < y^2$ . Since

$$3z^3 > x^3 + y^3 + z^3 = 4y^\beta z,$$

we get

$$y^\beta < \frac{3z^2}{4} < \frac{3y^4}{4} < y^4,$$

and therefore that  $\beta < 4$ .

8.3. Assume that  $\beta = 3$ . Rewrite the equation

$$8 + y^3 + z^3 = 4y^3z$$

as

$$y^3 = \frac{z^3 + 8}{4z - 1}.$$

Let  $D := 4z - 1$ . Thus  $z \equiv 4^{-1} \pmod{D}$ . Since we also have  $z^3 + 8 \equiv 0 \pmod{D}$ , we get  $4^{-3} + 8 \equiv 0 \pmod{D}$  and therefore that  $D | 1 + 8 \cdot 4^3 = 513 = 3^3 \cdot 19$ . Thus,  $D \in \{1, 3, 3^2, 3^3, 19, 3 \cdot 19, 3^2 \cdot 19, 3^3 \cdot 19\}$ . Since  $z$  must be at least the second prime number which is congruent to 1 modulo 3, we have that  $D \geq 4 \cdot 13 - 1 = 51$ , and since we also have that  $D \equiv -1 \pmod{4}$ , it follows that in fact only the instance  $D = 3^2 \cdot 19$  is possible. Therefore  $z = \frac{D+1}{4} = \frac{3^2 \cdot 19 + 1}{4} = 43$ . However, for this value of  $z$ , the number  $\frac{z^3 + 8}{4z - 1} = \frac{43^3 + 8}{4 \cdot 43 - 1} = 465$  is not the cube of a prime number.

8.4. Assume that  $\beta = 2$ . In this case,

$$z^3 < x^3 + y^3 + z^3 = 4y^2z,$$

so that

$$z^2 < 4y^2,$$

which implies that  $z < 2y$ . But we also have that  $y^2 | (x^3 + z^3)$ , and since  $y$  does not divide  $x + z$ , it follows that  $y^2 | z^2 - zx + x^2 = z^2 - 2z + 4$ . Since  $z \equiv 1 \pmod{3}$ , we also have that  $3 | z^2 - 2z + 4$ , and since  $y > 3$ , we have that  $y^2 | (z^2 - 2z + 4)/3$ . Now write

$$y^2 = \frac{z^2 - 2z + 4}{3\delta},$$

where  $\delta$  is a positive integer. We then get

$$\delta = \frac{z^2 - 2z + 4}{3y^2} < \frac{z^2}{3y^2} < \frac{4y^2}{3y^2} = \frac{4}{3} < 2,$$

which means that  $\delta = 1$ . Thus,  $3y^2 = z^2 - 2z + 4$ . The original relation becomes

$$4y^2z = 8 + y^3 + z^3 = y^3 + (z + 2)(z^2 - 2z + 4) = y^3 + 3y^2(z + 2),$$

so that

$$4z = y + 3(z + 2) = 3z + y + 6,$$

which implies that  $z = y + 6$ . Thus,  $y \equiv -6 \pmod{z}$ , and since  $z \mid y^2 - yx + x^2 = y^2 - 2y + 4$ , we get  $z \mid (-6)^2 - 2(-6) + 4 = 52 = 4 \cdot 13$ . Thus,  $z = 13$ ,  $y = z - 6 = 7$ , and we have obtained the solution

$$2^3 + 7^3 + 13^3 = 2^2 \cdot 7^2 \cdot 13$$

mentioned in the statement of our theorem.

9. From now on, we assume that  $x > 2$ . The relation  $x \mid y^3 + z^3$  implies that  $y^3 \equiv -z^3 \pmod{x}$  and therefore  $-yz \equiv \left(\frac{z^2}{y}\right)^2 \pmod{x}$ , and so

$$\left(\frac{-yz}{x}\right) = 1. \tag{3}$$

In a similar way, using the facts that  $y \mid x^3 + z^3$  and  $z \mid x^3 + y^3$ , one gets

$$\left(\frac{-xz}{y}\right) = \left(\frac{-xy}{z}\right) = 1.$$

Thus,

$$1 = \left(\frac{-yz}{x}\right) = \left(\frac{-1}{x}\right) \cdot \left(\frac{y}{x}\right) \cdot \left(\frac{z}{x}\right) = (-1)^{\frac{x-1}{2}} \cdot \left(\frac{y}{x}\right) \cdot \left(\frac{z}{x}\right),$$

and similarly

$$1 = (-1)^{\frac{y-1}{2}} \cdot \left(\frac{x}{y}\right) \cdot \left(\frac{z}{y}\right),$$

and

$$1 = (-1)^{\frac{z-1}{2}} \left(\frac{x}{z}\right) \cdot \left(\frac{y}{z}\right).$$

Write  $a := \frac{x-1}{2}$ ,  $b := \frac{y-1}{2}$ ,  $c := \frac{z-1}{2}$ . Multiplying the three relations above side by side and using quadratic reciprocity we get

$$1 = (-1)^{a+b+c+ab+ac+bc},$$

which means that

$$S := a + b + c + ab + ac + bc$$

must be an even number. Let us notice that it is not possible that all three numbers  $a$ ,  $b$ ,  $c$  are even. Indeed, if this were so, then  $x \equiv y \equiv z \equiv 1 \pmod{4}$ , and reducing the equation

$$x^3 + y^3 + z^3 = n$$

modulo 4, we would get  $3 \equiv 1 \pmod{4}$ , which is impossible. Thus, at least one of the numbers  $a$ ,  $b$ ,  $c$  is odd. This, together with the fact that  $S$  is even implies that all three numbers  $a$ ,  $b$ ,  $c$  are odd, therefore  $x \equiv y \equiv z \equiv 3 \pmod{4}$ . We reduce now the relation

$$x^3 + y^3 + z^3 = x^\alpha y^\beta z$$

modulo 4, and since  $\alpha$  is even (see step 7), we get  $1 \equiv 3^{\beta+1} \pmod{4}$  and therefore that  $\beta$  is odd. Thus, we may write our original equation as

$$x^3 + y^3 + z^3 = m^2yz, \quad (4)$$

where  $m := x^{\alpha/2}y^{(\beta-1)/2}$  is an integer. Write  $x + y = 2\ell$ . Notice that since  $x \equiv y \equiv 3 \pmod{4}$ , we have that  $\ell$  is an odd number. Let  $p$  be an arbitrary prime divisor of  $\ell$ . Reducing the above equation mod  $p$ , we get  $z^3 \equiv m^2yz \pmod{p}$ , therefore  $y \equiv \left(\frac{z}{m}\right)^2 \pmod{p}$ . Thus,

$$\left(\frac{y}{p}\right) = 1.$$

Since  $y \equiv -x \pmod{p}$ , we get that

$$1 = \left(\frac{y}{p}\right) = \left(\frac{-x}{p}\right) = \left(\frac{-1}{p}\right) \cdot \left(\frac{x}{p}\right) = (-1)^{\frac{p-1}{2}} \cdot (-1)^{\frac{x-1}{2} \cdot \frac{p-1}{2}} \cdot \left(\frac{p}{x}\right) = \left(\frac{p}{x}\right),$$

where in the above computation we used the quadratic reciprocity law together with the fact that  $x \equiv 3 \pmod{4}$ . Since the above formula holds for all prime divisors  $p$  of  $\ell$ , we get, by multiplying all these relations, that

$$1 = \left(\frac{\ell}{x}\right) = \left(\frac{(y+x)/2}{x}\right) = \left(\frac{4}{x}\right) \cdot \left(\frac{(y+x)/2}{x}\right) = \left(\frac{2y+2x}{x}\right) = \left(\frac{2y}{x}\right).$$

In the above argument, we used only equation (4) (which is symmetric in  $y$  and  $z$ ), together with the fact that  $x \equiv y \equiv z \equiv 3 \pmod{4}$  (which is also symmetric in  $y$  and  $z$ ), but we did not use size arguments (i.e. the fact that  $y < z$ ). Thus, an identical argument can be carried through to show that

$$\left(\frac{2z}{x}\right) = 1.$$

Multiplying these last two relations we get

$$1 = \left(\frac{2y}{x}\right) \cdot \left(\frac{2z}{x}\right) = \left(\frac{4}{x}\right) \cdot \left(\frac{yz}{x}\right) = \left(\frac{yz}{x}\right),$$

which together with the fact that

$$\left(\frac{-yz}{x}\right) = 1$$

(see equation (3)), implies that

$$\left(\frac{-1}{x}\right) = 1,$$

contradicting the fact that  $x \equiv 3 \pmod{4}$ .

This completes the proof of Theorem 1.

### 4. An upper bound for $T_0(x)$

**Theorem 2.** As  $x \rightarrow \infty$ , we have

$$T_0(x) \leq x \exp \left\{ -(1 + o(1)) \sqrt{\frac{1}{6} \log x \log \log x} \right\}.$$

**Proof.** First recall the estimate

$$\Psi(x, y) := \#\{n \leq x : P(n) \leq y\} \ll x \exp\{-(1 + o(1))u \log u\}, \tag{5}$$

where  $u = \log x / \log y$  (see for instance Tenenbaum [4]). Now let

$$y = \exp \left\{ \sqrt{\frac{3}{2} \log x \log \log x} \right\} \tag{6}$$

and set

$$u = \frac{\log x}{\log y} = \sqrt{\frac{2}{3} \frac{\log x}{\log \log x}} \quad \text{so that} \quad u \log u = (1 + o(1)) \sqrt{\frac{1}{6} \log x \log \log x}. \tag{7}$$

It follows from (5), (6) and (7) that

$$\begin{aligned} \#\{n \leq x : n \in T_0, P(n) \leq y\} &\ll x \exp\{-(1 + o(1))u \log u\} \\ &\ll x \exp \left\{ -(1 + o(1)) \sqrt{\frac{1}{6} \log x \log \log x} \right\}. \end{aligned} \tag{8}$$

We shall therefore assume from now on that  $P(n) > y$ .

Let  $x$  be a large number with the corresponding  $y$  and  $u$  defined by (6) and (7). Then, using Stirling's formula, as well as the fact that

$$\sum_{p \leq y} \frac{1}{p} = \log \log y + O(1)$$

holds as  $y$  tends to infinity, we get

$$\begin{aligned} \#\{n \leq x : \omega(n) \geq u\} &\leq \sum_{p_1 \dots p_{[u]} \leq x} \frac{x}{p_1 \dots p_{[u]}} \leq \frac{x}{[u]!} \left( \sum_{p \leq x} \frac{1}{p} \right)^{[u]} \\ &\leq x \left( \frac{e \log \log x + O(1)}{[u]} \right)^{[u]} \\ &\leq x \exp \{-(1 + o(1))u \log u\} \\ &\ll x \exp \left\{ -(1 + o(1)) \sqrt{\frac{1}{6} \log x \log \log x} \right\}. \end{aligned} \tag{9}$$

Hence, from here on we may assume that  $\omega(n) < u$ .

We now neglect those integers  $n \leq x$ ,  $n \in T_0$  with  $P(n) > y$  and such that  $P(n)^2 | n$ , since the number of such integers is

$$\begin{aligned} &\ll \#\{n \leq x : P(n) > y, P(n)^2 | n\} \leq \sum_{p > y} \frac{x}{p^2} \\ &\ll \frac{x}{y} = x \exp \left\{ -\sqrt{\frac{3}{2} \log x \log \log x} \right\}. \end{aligned} \quad (10)$$

From here on, we shall therefore assume that  $Q := P(n) \nmid n$  and write  $n = mQ$ . Now, writing (2) as

$$n = mQ = p_1^{b_1} + \dots + p_k^{b_k}, \quad (11)$$

where  $p_1 < \dots < p_k = Q$  are the prime factors of  $n$  and each  $b_i$  is non negative, we get from (11) that

$$p_1^{b_1} + \dots + p_{k-1}^{b_{k-1}} + \delta \equiv 0 \pmod{Q}, \quad (12)$$

where  $\delta$  is 0 or 1, depending if  $b_k > 0$  or  $b_k = 0$ . The number appearing on the left hand side of (12) depends only on the prime factors of  $m$  and does not depend on  $Q$ , and moreover, each one of these numbers has at most  $\log x$  factors. Thus, we may fix  $m \leq x/y$  and count how many candidates there may be for a given prime number  $Q$ . Since  $n$  is not a prime power, we have  $k \geq 2$ , and therefore the left hand side of congruence (12) is a positive integer. Since  $p_i^{b_i} < n \leq x$ , it follows that  $b_i \ll \log x + 1$ . In fact,  $b_i < \log x + 1$  always holds except when  $i = 1$  and  $p_1 = 2$ , in which case  $b_1 \leq \frac{\log x + 1}{\log 2}$ . Thus, the total number of integers which can appear on the left hand side of (12) is  $\ll (\log x + 1)^{\omega(n)} \ll (\log x + O(1))^u \ll \exp\{(1 + o(1))u \log \log x\}$ , which means that

$$\begin{aligned} &\#\{n \leq x : n \in T_0, P(n) > y, P(n) \nmid n, \omega(n) < u\} \\ &\ll \frac{x \log x}{y} \exp\{(1 + o(1))u \log \log x\} \\ &\ll x \exp \left\{ -(1 + o(1)) \sqrt{\frac{1}{6} \log x \log \log x} \right\}. \end{aligned} \quad (13)$$

Theorem 2 then follows from (8), (9), (10) and (13).

## 5. Empirical lower and upper bounds for $T(x)$

Although we cannot prove that  $T$  is an infinite set, a heuristic argument shows that

$$\exp \left( \frac{2}{e} (1 + o(1)) \frac{\log x}{(\log \log x)^2} \right) \leq T(x) \leq x^{1/2 + o(1)}. \quad (14)$$

Our argument goes as follows. First, we will show that, heuristically,

$$T(x) = \frac{1}{2} \sum_{n \leq x} f(n), \quad \text{where } f(n) := \frac{1}{n} \prod_{p|n} \left\lfloor \frac{\log n}{\log p} \right\rfloor, \quad (15)$$

from which we will show that (14) follows.

Indeed, given a positive integer  $n$  such that  $\omega(n)$  is odd and writing  $n = q_1^{\alpha_1} \dots q_r^{\alpha_r}$ , then in order to have  $n \in T$ , we must find a representation of the form

$$n = q_1^{\alpha_1} + \dots + q_r^{\alpha_r}. \quad (16)$$

Now, for each exponent  $\alpha_i$ , there are  $\lfloor \log n / \log q_i \rfloor$  possible choices. Hence, if a representation of the form (16) is possible, then the exponents  $\alpha_i$  have been chosen in the interval  $[1, \lfloor \log n / \log q_i \rfloor]$ . Therefore, since there are  $\prod_{i=1}^r \lfloor \log n / \log q_i \rfloor$  possible choices for the right hand side of (16), we should ‘expect’ that a representation of the form (16) will be possible with a ‘probability’ equal to  $\frac{1}{n} \prod_{p|n} \left\lfloor \frac{\log n}{\log p} \right\rfloor$ ,

thus establishing (15); note that the factor  $\frac{1}{2}$  comes from the fact that a randomly chosen number has an odd “ $\omega(n)$ ” with a probability  $\frac{1}{2}$ .

It remains to prove that (14) follows from (15).

First we prove the lower bound. Let  $x$  be a large positive real number and let  $k \geq 1$  be an integer.

Let  $p_1 < \dots < p_k$  be the first  $k$  primes. We shall consider only the contribution to  $T(x)$  of those positive integers  $n = p_1 \dots p_k p \leq x$ , where  $p > p_k$  is a prime number. We first get rid of the integer parts. Clearly, if  $i \in \{1, \dots, k\}$ , then

$$\left\lfloor \frac{\log n}{\log p_i} \right\rfloor = \frac{\log n}{\log p_i} - \left\{ \frac{\log n}{\log p_i} \right\} > \frac{\log n}{\log p_i} \left( 1 - \frac{\log p_i}{\log n} \right) > \frac{\log n}{\log p_i} \exp \left( -2 \frac{\log p_i}{\log n} \right),$$

where in the above inequalities we used the fact that  $\log p_i / \log n \leq 1/2$  and that the inequality  $1 - t > \exp(-2t)$  holds for  $t \in (0, 1/2)$ . Together with the fact that  $\lfloor \log n / \log p \rfloor \geq 1$ , we get

$$f(n) \geq \left( \prod_{i=1}^k \frac{\log n}{\log p_i} \right) \exp \left( -2 \sum_{i=1}^k \frac{\log p_i}{\log n} \right) > \exp(-2) \prod_{i=1}^k \frac{\log n}{\log p_i} \gg \frac{(\log p)^k}{\log p_1 \dots \log p_k}.$$

This implies that

$$\begin{aligned} T(x) &= \frac{1}{2} \sum_{n \leq x} f(n) \gg \sum_{\substack{p_1 \dots p_k p \leq x \\ p > p_k}} \frac{1}{p_1 \log p_1 \dots p_k \log p_k} \frac{(\log p)^k}{p} \\ &= \frac{1}{p_1 \log p_1 \dots p_k \log p_k} \sum_{p_k < p \leq x/p_1 \dots p_k} \frac{(\log p)^k}{p} \end{aligned} \quad (17)$$

$$\begin{aligned}
&\gg \frac{1}{p_1 \log p_1 \dots p_k \log p_k} \int_{p_k}^{x/p_1 \dots p_k} \frac{(\log t)^{k-1}}{t} dt \\
&\gg \frac{1}{p_1 \log p_1 \dots p_k \log p_k} \frac{(\log(x/p_1 \dots p_k))^k}{k} \\
&= \exp \left( k \log \log x - \sum_{i=1}^k (\log p_i + \log \log p_i) + O \left( \frac{k(\log p_1 + \dots + \log p_k)}{\log x} \right) \right).
\end{aligned}$$

The above chain of inequalities holds when  $k$  is such that

$$\log(x/p_1 \dots p_k) - \log p_k \gg \log(x/p_1 \dots p_k),$$

which in turn is true when

$$\log p_k + \frac{\log p_1 + \dots + \log p_k}{\log x} = o(\log x),$$

which holds when

$$\log p_k + \frac{k \log p_k}{\log x} = o(\log x). \quad (18)$$

We now use the fact that, as  $k$  tends to infinity,

$$p_k \leq k \log k + k \log \log k - k + o(k)$$

(see Théorème A (v) in [1]), together with the well known estimate

$$\sum_{p \leq y} \log p = \sum_{n \leq y} \Lambda(n) + O(y^{1/2}) = y + O \left( \frac{y}{\exp(c\sqrt{\log y})} \right) = y + O \left( \frac{y}{(\log y)^2} \right),$$

where  $c$  is some positive constant and  $\Lambda$  denotes the von Mangoldt function, to conclude that

$$\sum_{i=1}^k \log p_i = p_k + O \left( \frac{p_k}{(\log k)^2} \right) \leq k \log k + k \log \log k - k + o(k). \quad (19)$$

Since  $p_k < 2k \log k$  holds for all sufficiently large  $k$ , we also have that

$$\begin{aligned}
\sum_{i=1}^k \log \log p_i &\leq k \log \log p_k \leq k \log (\log k + \log(2 \log k)) \\
&\leq k \log \log k + O \left( \frac{k \log \log k}{\log k} \right) \\
&= k \log \log k + o(k).
\end{aligned} \quad (20)$$

Introducing inequalities (19) and (20) into (17), we get

$$T(x) \geq \exp \left( k \log \log x - k \log k - 2k \log \log k + k + o(k) + O \left( \frac{k^2 \log k}{\log x} \right) \right) \quad (21)$$

$$= \exp \left( k \log \left( \frac{\log x}{k(\log k)^2} \right) + k + o(k) + O \left( \frac{k^2 \log k}{\log x} \right) \right).$$

In order to maximize the main term of the above inequality, we should choose  $k$  versus  $x$  in such a way that the expression  $k \log \left( \frac{\log x}{k(\log k)^2} \right)$  should be as large as possible. Thus, we choose  $k := \left\lfloor \frac{1}{e} \frac{\log x}{(\log \log x)^2} \right\rfloor$ . We note that  $k$  is in the acceptable range; i.e.,  $p_1 \dots p_k < x$ , that condition (18) is fulfilled, that with this choice of  $k$  we have

$$k \log \left( \frac{\log x}{k(\log k)^2} \right) = (1 + o(1))k,$$

and that the error term is

$$\frac{k^2 \log k}{\log x} = \frac{k}{\log k} \frac{k(\log k)^2}{\log x} = O \left( \frac{k}{\log k} \right) = o(k).$$

Hence, we may replace (21) by

$$T(x) \geq \exp(2(1 + o(1))k) = \exp \left( \frac{2}{e}(1 + o(1)) \frac{\log x}{(\log \log x)^2} \right),$$

which proves the left hand side of inequality (14).

We now prove the upper bound.

Fix a large number  $k$  and write

$$T(x) < \sum_{\substack{n \leq x \\ \omega(n) < k}} \frac{1}{n} \prod_{p|n} \frac{\log n}{\log p} + \sum_{\substack{n \leq x \\ \omega(n) \geq k}} \frac{1}{n} \prod_{p|n} \frac{\log n}{\log p} = T_1(x) + T_2(x), \quad (22)$$

say. We have

$$T_1(x) \leq \sum_{\substack{n \leq x \\ \omega(n) < k}} \frac{1}{n} (\log n)^{\omega(n)} \leq \sum_{n \leq x} \frac{(\log n)^k}{n} \ll \frac{(\log x)^{k+1}}{k+1}. \quad (23)$$

In particular,

$$T_1(x) < (\log x)^{k+1} \quad (24)$$

holds if  $k$  is sufficiently large.

In the sequel, we shall be using the fact that, if  $k$  is sufficiently large, then

$$\prod_{i=1}^k \log p_i > (\log k)^k. \quad (25)$$

Indeed, since  $p_i \geq i \log i$  holds for all  $i \geq 2$  (see [3]), one gets

$$\log p_i \geq \log i + \log \log i \quad (i \geq 3). \quad (26)$$

The inequality  $\log(1+t) > t/2$  holds for all  $t \in (0, 1/2)$ . The function  $t \mapsto \log \log t / \log t$  is decreasing for  $t > e^e$  and its value at  $e^e$  is  $1/e < 1/2$ . Hence,

$$\begin{aligned} \log(\log i + \log \log i) &= \log \log i + \log \left( 1 + \frac{\log \log i}{\log i} \right) > \log \log i + \frac{1}{2} \frac{\log \log i}{\log i} \\ (i > e^e \approx 15.2). \end{aligned}$$

We thus get

$$\begin{aligned} \sum_{i=1}^k \log \log p_i &\geq \sum_{i=16}^k \log(\log i + \log \log i) + O(1) \\ &= \sum_{i=16}^k \log \log i + \sum_{i=16}^k \log \left( 1 + \frac{\log \log i}{\log i} \right) + O(1) \\ &\geq \sum_{i=16}^k \log \log i + \frac{1}{2} \sum_{i=16}^k \frac{\log \log i}{\log i} + O(1) \\ &\geq \int_{16}^k \log \log t dt + \frac{1}{2} \int_{16}^t \frac{\log \log t}{\log t} dt + O(1) \\ &= t \log \log t \Big|_{t=16}^{t=k} - \int_{16}^t \frac{1}{\log t} dt + \frac{1}{2} \int_{16}^t \frac{\log \log t}{\log t} dt + O(1) \\ &> k \log \log k, \end{aligned}$$

where the last inequality follows for large enough  $k$  due to the fact that the function  $\int_{16}^k \left( \frac{1}{2} \frac{\log \log t}{\log t} - \frac{1}{\log t} \right) dt$  tends to infinity with  $k$ , thus establishing (25).

Using (25), we have

$$T_2(x) \leq \sum_{\substack{n \leq x \\ \omega(n) \geq k}} \frac{1}{n} \frac{(\log n)^k}{\prod_{i=1}^k \log p_i} \leq \sum_{\substack{n \leq x \\ \omega(n) \geq k}} \frac{1}{n} \left( \frac{\log n}{\log \omega(n)} \right)^{\omega(n)}. \quad (27)$$

Using the fact that

$$\omega(n) \leq \frac{\log n}{\log \log n} + (1 + o(1)) \frac{\log n}{(\log \log n)^2}$$

(see Pomerance [2]), together with the fact that the function  $t \mapsto \left(\frac{\log n}{\log t}\right)^t$  is increasing for  $t \leq \log n$ , it follows, from (27), that

$$T_2(x) \leq \sum_{\substack{n \leq x \\ \omega(n) \geq k}} \frac{1}{n} \cdot n \cdot e^{O\left(\frac{\log n \cdot \log \log \log n}{\log \log n}\right)} \tag{28}$$

$$\ll \mathcal{N}_k(x) \exp \left\{ O \left( \frac{\log x \log \log \log x}{\log \log x} \right) \right\},$$

where

$$\mathcal{N}_k(x) = \#\{n \leq x \mid \omega(n) \geq k\}.$$

It is easy to see, using Stirling's formula, that

$$\mathcal{N}_k(x) \leq x \sum_{\substack{q_1 < \dots < q_k \\ q_1 \dots q_k \leq x}} \frac{1}{q_1 \dots q_k} \leq \frac{x}{k!} \left( \sum_{q \leq x} \frac{1}{q} \right)^k \ll \frac{x}{\sqrt{k}} \left( \frac{e \log \log x + O(1)}{k} \right)^k. \tag{29}$$

In particular, combining (28) and (29), for large  $x$  and  $k$ , we have that

$$T_2(x) < x \cdot \left( \frac{(\log \log x)^{3/2}}{k} \right)^k \exp \left\{ O \left( \frac{\log x \log \log \log x}{\log \log x} \right) \right\}. \tag{30}$$

We now choose  $k$  such that  $k := \left\lfloor \frac{1}{2} \frac{\log x}{\log \log x} \right\rfloor$ . It is clear that  $k$  is in the acceptable range; i.e.,  $k = \omega(n)$  for some  $n \leq x$ . Furthermore, inequality (24) shows that

$$T_1(x) < x^{1/2+o(1)}, \tag{31}$$

while inequality (30) shows that

$$T_2(x) < x \exp \left( \frac{3}{2} k \log \log \log x - k \log k - O \left( \frac{\log x \log \log \log x}{\log \log x} \right) \right) \tag{32}$$

$$= x \exp \left( -\frac{\log x}{2} + O \left( \frac{\log x \log \log \log x}{\log \log x} \right) \right) = x^{1/2+o(1)}.$$

Using (31) and (32) in (22), we obtain the upper bound in (14).

**Acknowledgement.** The authors would like to thank the referee for some very helpful suggestions.

## References

- [1] J. Massias and G. Robin, 'Bornes effectives pour certaines fonctions concernant les nombres premiers', *J. Théorie Nombres Bordeaux* **8** (1996), 215–242.
- [2] C. Pomerance, 'On the distribution of round numbers', Ootacamund, India 1984, K. Alladi, ed., *Lecture Notes in Math.* **1122** (1985), 173–200.
- [3] J.B. Rosser and L. Schoenfeld, 'Approximate formulas for some functions of prime numbers', *Illinois J. Math.* **6** (1962), 64–94.
- [4] G. Tenenbaum, *Introduction to Analytic and Probabilistic Number Theory*, Cambridge University Press, 1995.

**Addresses:** Jean-Marie De Koninck, Département de mathématiques, Université Laval, Québec G1K 7P4, Canada  
Florian Luca, Mathematical Institute, UNAM, Ap. Postal 61-3 (Xangari), CP 58 089, Morelia, Michoacán, Mexico

**E-mail:** [jmdk@mat.ulaval.ca](mailto:jmdk@mat.ulaval.ca); [fluca@matmor.unam.mx](mailto:fluca@matmor.unam.mx)

**Received:** 30 July 2004; **revised:** 20 December 2004