

48. On the Highest Degree of Absolute Polynomials of Alternating Links

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§ 1. Statement of results. Let L be a tame link in S^3 . For a regular projection \tilde{L} of L , $c(\tilde{L})$ denotes the number of crossings of \tilde{L} and $c(L)$ denotes the minimum number of crossings among all regular projections of L . \tilde{L} divides S^2 into finitely many domains, which can be colored by two colors (black and white) like a chess-board such that domains of the same color meet only at crossing points. Let $g(\tilde{L})$ and $g^*(\tilde{L})$ be the graphs of \tilde{L} such that vertices of $g(\tilde{L})$ and $g^*(\tilde{L})$ correspond to the white and the black domains, respectively, and each edge of $g(\tilde{L})$ and $g^*(\tilde{L})$ corresponds to a crossing of \tilde{L} . A vertex v of a graph is called a *stump* if the valency of v is equal to one, and a *twig* if the valency of v is equal to two.

R. D. Brandt, W. B. R. Lickorish and K. C. Millett defined in [1] an unoriented link invariant $Q(L)$ called the *absolute polynomial*. We refer the reader to [1] for the details.

In this paper we prove the following

Theorem. *Let L be a tame alternating link and \tilde{L} a regular alternating projection of L . Then the following conditions are equivalent:*

- (1) *The graphs $g(\tilde{L})$ and $g^*(\tilde{L})$ are connected without stumps, loops and cut-vertices.*
- (2) *The highest degree of the absolute polynomial $\tilde{Q}(L)$ of L is equal to $c(\tilde{L}) - 1$, and the coefficient of the term of the highest degree of $Q(L)$ is positive.*

K. Murasugi proved in [5] that if \tilde{L} is a regular connected proper alternating projection of an alternating link L , then the reduced degree of Jones polynomial [2] of L is equal to $c(\tilde{L})$ and $c(\tilde{L}) = c(L)$. If L is a prime alternating link, then L has a regular projection \tilde{L} which satisfies the condition (1) of Theorem. Therefore we have

Corollary 1. *If L is a prime alternating link then the highest degree of $Q(L)$ is equal to $c(L) - 1$.*

The following is a part of Theorem 1 of W. Menasco [4], for which we will give an alternative proof.

Corollary 2. *If L_1 and L_2 are alternating links and $L = L_1 \# L_2$ is also an alternating link, then, for any connected regular proper alternating projection \tilde{L} of L , $g(\tilde{L})$ and $g^*(\tilde{L})$ have cut-vertices.*

In [3], M. E. Kidwell independently obtained the similar results to our theorem. His method of the proof is different from ours.

§ 2. Proofs. We work in the PL category and all projections of links we consider are assumed to be regular.

The absolute polynomial $Q(L)$ has the following properties :

Theorem A ([1, Properties 1 and 8]). *Let L, L_1 and L_2 be tame links in S^3 then the following formulae hold.*

- (a) $Q(L_1 \# L_2) = Q(L_1)Q(L_2)$, where $L_1 \# L_2$ denotes any connected sum of L_1 and L_2 .
- (b) $Q(L_1 \circ L_2) = \mu Q(L_1)Q(L_2)$, where $\mu = 2x^{-1} - 1$ and $L_1 \circ L_2$ denotes split union.
- (c) $h\text{-deg } Q(L) \leq c(L) - 1$, where $h\text{-deg } Q(L)$ denotes the highest degree of $Q(L)$.

To prove Theorem, we prove the following lemmas.

Lemma 1. *We suppose that \tilde{L} is an alternating projection of a link L such that $c(\tilde{L}) \geq 3$ and that $g(\tilde{L})$ is a connected graph without stumps, loops and cut-vertices. Let \tilde{L}_0 and \tilde{L}_∞ be the projections in [1, Theorem]. Then \tilde{L}_0 and \tilde{L}_∞ are again alternating projections and either $g(\tilde{L}_0)$ or $g(\tilde{L}_\infty)$ is a connected graph without stumps, loops and cut-vertices.*

Proof. Fig. 1 shows that \tilde{L}_0 and \tilde{L}_∞ are alternating projections.

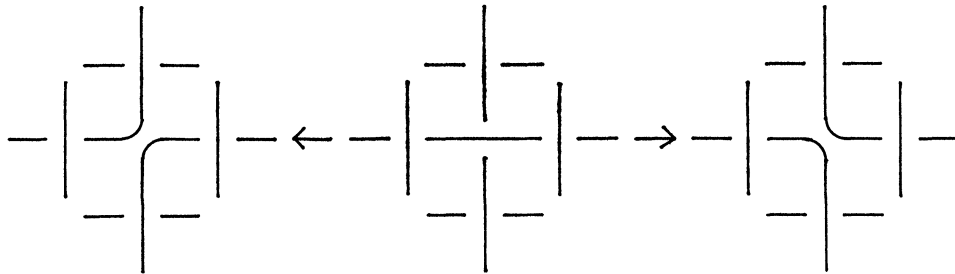


Fig. 1

Without loss of generality, we may assume that \tilde{L}_0 and \tilde{L}_∞ are as shown in Fig. 2.

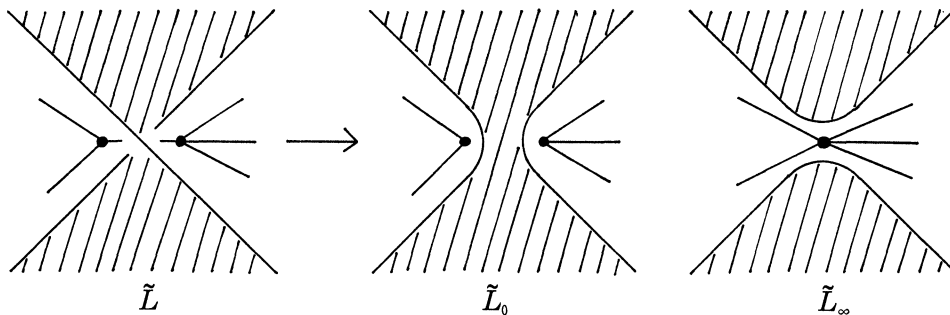


Fig. 2

By the condition of $g(\tilde{L})$, both $g(\tilde{L}_0)$ and $g(\tilde{L}_\infty)$ are connected, and $g(\tilde{L}_0)$ has no loops and $g(\tilde{L}_\infty)$ has no stumps. There are two cases to be considered.

Case 1. The graph $g(\tilde{L}_0)$ has a stump: In this case, $g(\tilde{L})$ has a twig. Since $c(\tilde{L}) \geq 3$ and $g(\tilde{L})$ has no cut-vertices, $g(\tilde{L})$ is as in Fig. 3.

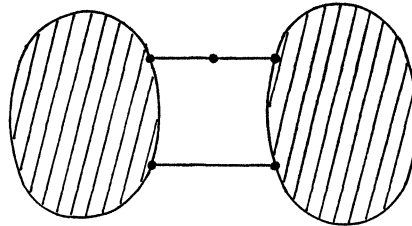


Fig. 3

Then $g(\tilde{L}_\infty)$ has none of stumps, loops and cut-vertices.

Case 2. $g(\tilde{L}_0)$ has a cut-vertex: In this case $g(\tilde{L}_\infty)$ has none of stumps, loops and cut-vertices. See Fig. 4. □

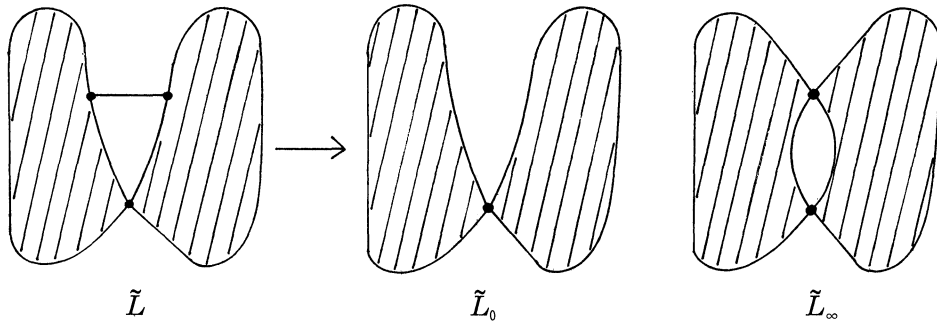


Fig. 4

Lemma 2. If \tilde{L} is a non-alternating projection of a link L , then $h\text{-deg } Q(L) \leq c(\tilde{L}) - 2$.

Proof. If \tilde{L} is disconnected, then by Theorem A the assertion holds. Therefore it is sufficient to consider the case \tilde{L} is connected. Since \tilde{L} is a non-alternating projection, we may assume that \tilde{L} has two successive under crossings p_1 and p_2 .

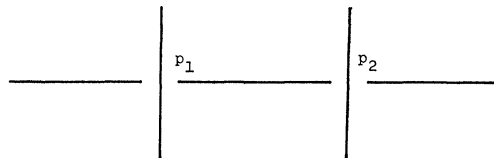


Fig. 5

We can obtain a sequence of link projections $\tilde{L} = \tilde{L}_0 \rightarrow \tilde{L}_1 \rightarrow \tilde{L}_2 \rightarrow \dots \rightarrow \tilde{L}_m$ such that \tilde{L}_{i+1} is obtained from \tilde{L}_i by switching one of the crossings except p_1 and p_2 and that \tilde{L}_m is an ascending projection of \tilde{L} . Since we do not switch the crossings p_1 and p_2 in all \tilde{L}_i 's, all \tilde{L}_i 's are non-alternating projections.

Hence we can obtain a resolution R of \tilde{L} such that all projections of R are non-alternating projections. Since a non-alternating projection whose number of crossings is two is a trivial link, we can obtain trivial links by smoothing at most $c(\tilde{L})-2$ crossings. Hence $h\text{-deg } Q(L) \leq c(\tilde{L})-2$. \square

Proof of Theorem. First, we prove necessity. If $g(\tilde{L})$ has a stump or a loop, then there exists the projection \tilde{L}' of L such that $c(\tilde{L}')=c(\tilde{L})-1$. Therefore $h\text{-deg } Q(L) \leq c(\tilde{L}')-1=c(\tilde{L})-2$. If $g(\tilde{L})$ is disconnected then $\tilde{L}=\tilde{L}_1 \cup \tilde{L}_2 \cup \dots \cup \tilde{L}_n$, where \tilde{L}_i is a connected component and $n \geq 2$. By Theorem A we have $h\text{-deg } Q(L) = \sum_{i=1}^n h\text{-deg } Q(L_i) \leq \sum_{i=1}^n (c(\tilde{L}_i) - 1) = \sum_{i=1}^n c(\tilde{L}_i) - n = c(\tilde{L}) - n < c(\tilde{L}) - 1$. If $g(\tilde{L})$ has a cut-vertex then $g(\tilde{L})$ is one point union of subgraphs Γ_1 and Γ_2 . Let \tilde{l}_i denote the projection corresponding to Γ_i ($i=1, 2$). Since L is a split or connected sum of links l_1 and l_2 , we have $h\text{-deg } Q(L) = \sum_{i=1}^2 h\text{-deg } Q(l_i) \leq \sum_{i=1}^2 (c(\tilde{l}_i) - 1) = \sum_{i=1}^2 c(\tilde{l}_i) - 2 < c(\tilde{L}) - 1$.

Now we prove sufficiency of Theorem by induction on $c(\tilde{L})$. In case $c(\tilde{L})=2$, L is the Hopf link and $Q(L) = -2x^{-1} + 1 + 2x$. In case $c(\tilde{L}) \geq 3$, by Lemma 1, at least one of \tilde{L}_0 and \tilde{L}_∞ has a connected graph without stumps, loops, and cut-vertices. If $g(\tilde{L}_0)$ and $g(\tilde{L}_\infty)$ have stumps, loops, or cut-vertices, then, by Theorem A, $h\text{-deg } Q(\tilde{L}_0) < c(\tilde{L}_0) - 1$ and $h\text{-deg } Q(\tilde{L}_\infty) < c(\tilde{L}_\infty) - 1$, respectively. If both $g(\tilde{L}_0)$ and $g(\tilde{L}_\infty)$ are connected graphs without stumps, loops and cut-vertices, by the hypothesis of induction the coefficient of the highest degree term is positive. Therefore $h\text{-deg } (x(Q(\tilde{L}_0) + Q(\tilde{L}_\infty))) = c(\tilde{L}) - 1$. Since \tilde{L}_- is obtained from \tilde{L} by switching a crossing, $h\text{-deg } Q(\tilde{L}_-) \leq c(\tilde{L}) - 2$ by Lemma 2. Since $Q(\tilde{L}_+) = -Q(\tilde{L}_-) + x(Q(\tilde{L}_0) + Q(\tilde{L}_\infty))$ and the coefficient of the highest degree term is positive, we have that $h\text{-deg } Q(\tilde{L}_+) = c(\tilde{L}_+) - 1$ and the coefficient of the highest degree term is positive. For the graph $g^*(\tilde{L})$, we can prove similarly and the proof is complete. \square

Proof of Corollary 2. If there exists \tilde{L} such that $g(\tilde{L})$ and $g^*(\tilde{L})$ have no cut-vertices, then $h\text{-deg } Q(L) = c(\tilde{L}) - 1 = c(L) - 1$ by Theorem.

By Corollary 4 of [5] and Theorem A we have $c(L) - 1 = h\text{-deg } Q(L) = h\text{-deg } Q(L_1) + h\text{-deg } Q(L_2) \leq c(L_1) - 1 + c(L_2) - 1 = c(L_1) + c(L_2) - 2 = c(L) - 2$. This is a contradiction. \square

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