A NOTE ON THE AXIOMATIZATION OF RUBIN'S SYSTEM (5)

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Rubin, in [3], suggests that the axiomatic for system (5) may be simplified. It is here shown that

R5 If
$$\vdash \alpha$$
 in (S) then $\vdash \Box_2 \alpha$ in (S)

and

$$R*5$$
 If $\vdash \alpha$ in (S) then $\vdash \Box$, α in (S)

are derivable from the other axioms and rules of (5).

This paper presupposes [3] and adopts the same primitive basis and definitions for (5). Thus the axioms and rules of (5) are

- A1 $(\alpha \wedge \beta) \Longrightarrow (\beta \wedge \alpha)$
- $A2 \quad (\alpha \land \beta) \Longrightarrow \alpha$
- $A3 \quad \alpha \Longrightarrow (\alpha \land \alpha)$
- $A4 \quad ((\alpha \land \beta) \land \gamma) \Longrightarrow (\alpha \land (\beta \land \gamma))$
- $A5 \quad \alpha \Longrightarrow \sim \sim \alpha$
- $A6 \quad ((\alpha \Rightarrow \beta) \land (\beta \Rightarrow \gamma)) \Rightarrow (\alpha \Rightarrow \gamma)$
- $A7 \quad (\alpha \land (\alpha \Longrightarrow \beta)) \Longrightarrow \beta$
- $A8 \quad \square_2 \alpha \Longrightarrow \square_2 \square_2 \alpha$
- $A12 \square_2 \alpha \Longrightarrow \square_1 \alpha$
- **R1** If $\vdash \alpha$ and $\vdash (\alpha \implies \beta)$ then $\vdash \beta$.
- **R2** If $\vdash \alpha$ and $\vdash \beta$ then $\vdash (\alpha \land \beta)$.
- **R3** If $\vdash (\alpha \iff \beta)$ and $\vdash \gamma$ and δ results from γ by replacing α for β (or β for α) then $\vdash \delta$.

together with A*1-A*8, R*1-R*3. (T* is the wff obtained from T by replacing all the " \diamondsuit_2 's" by " \diamondsuit_1 's".)

The following theorems, S1-S7, follow from A1-A8 and R1-R3, and their proofs can be found in [1].

$$SI \quad \alpha \Rightarrow \alpha$$

$$S2 \quad ((\alpha \supset \beta) \land (\beta \supset \gamma)) \Rightarrow (\alpha \supset \gamma)$$

$$S3 \quad \square_2 \alpha \Rightarrow \alpha$$

$$S4 \quad (\alpha \Rightarrow \beta) \Leftrightarrow \square_2(\alpha \supset \beta)$$

$$[18.42]$$

$$[18.7]$$

$$S5 \qquad (\square_2 \alpha \wedge \square_2 \beta) \iff \square_2 (\alpha \wedge \beta)$$

$$S6 \qquad (\square_2 \alpha \wedge \square_2 \beta) \implies (\alpha \iff \beta)$$

$$S7 \qquad \square_2 \square_2 \alpha \iff \square_2 \alpha$$

$$[19.81]$$

$$[19.84]$$

$$[C10]$$

Some additional theorems of (5) follow

S8
$$(\alpha \Rightarrow \Box_2 \beta) \Rightarrow (\alpha \Rightarrow \beta)$$

S9 $(\alpha \Rightarrow \beta) \Rightarrow (\Box_2 \alpha \Rightarrow \Box_2 \beta)$

S8 was proved by Parry [2] p. 139, and S9 is given by Rubin [3] p. 308. Any further theorems of (5) that are needed will be given without proof—though their proofs are indicated by the bracketed information.

The following lemma will be needed for the proof of R5.

Lemma I If $\vdash \Box_2 \alpha$ and $\vdash (\alpha \rightarrow \beta)$ then $\vdash \Box_2 \beta$

Proof. (1)	$\square_2 \alpha$	[Hypothesis]
(2)	$\alpha \rightarrow \beta$	[Hypothesis]
(3)	$\square_1 \alpha$	[(1), A12, R1]
(4)	$\square_1(\alpha\supset\beta)$	[(2), S*4, R*3]
(5)	$\square_1 \alpha \wedge \square_1(\alpha \supset \beta)$	$[(3), (4), \mathbf{R}^*2]$
(6)	$\alpha \longleftrightarrow (\alpha \supset \beta)$	[(5), <i>S*6</i> , R*1]
(7)	$\square_2 (\alpha \supset \beta)$	[(1), (6), R*3]
(8)	$\alpha \Rightarrow \beta$	[(7), <i>S4</i> , R3]
(9)	$\square_2 \alpha \implies \square_2 \beta$	[(8), S9, R1]
(10)	$\square_2 \beta$	[(1), (9), R1]

It is now possible to prove

R5 If $\vdash \alpha$ then $\vdash \square_2 \alpha$

Proof. **R5** can be established by induction on the length of the proof of α . If α is one of the axioms, AI-A8, A12, then $\square_2\alpha$ follows by S12, **R1** and AI-A8, A12 respectively. If α is one of the axioms A*I-A*5 then $\square_2\alpha$ follows by S11, **R1**, and A1-A5 respectively. Thus it remains to show $\square_2\alpha$ when α is one of the axioms A*6-A*8. But the following are theorems of (5).

$$\begin{array}{lll} S18 & \square_2 \left(\left((\alpha \supset \beta) \land (\beta \supset \gamma) \right) \to (\alpha \supset \gamma) \right) & [S2, S11, \, \mathbb{R}1] \\ S19 & \square_2 \left(\square_1 \left((\alpha \supset \beta) \land (\beta \supset \gamma) \right) \to \square_1 \left(\alpha \supset \gamma \right) \right) & [S18, \, S*9, \, \text{Lemma I}] \\ S20 & \square_2 \left(\left((\alpha \to \beta) \land (\beta \to \gamma) \right) \to (\alpha \to \gamma) \right) & [S19, \, S*4, \, S*5, \, \mathbb{R}^*3] \\ S21 & \square_2 \left(\square_1 \alpha \to \square_1 \alpha \right) & [S13, \, S*9, \, \text{Lemma I}] \\ S22 & \square_2 \left(\square_1 \alpha \to \alpha \right) & [S21, \, S*8, \, \text{Lemma I}] \end{array}$$

$$\begin{array}{lll} S23 & \square_2 \left(\left(\alpha \wedge \left(\alpha \to \beta \right) \right) \to \beta \right) & [S22, S*17, \text{ Lemma I}] \\ S24 & \square_2 \left(\square_1 \alpha \to \square_1 \square_1 \alpha \right) & [S21, S*7, R*3] \end{array}$$

Thus S20, S23, and S24 complete the basis of the induction for the proof of R5.

If α follows from previous theorems of (S) by R1-R3 or R*1-R*3, $\square_2 \alpha$ may be obtained from the induction hypothesis using S3, S5, S9 (in case α follows by R1-R3, or R*2); by S3 and Lemma I (in case α follows by R*1); and by S3 (in case α follows by R*3).

Finally, on the basis of R5, it is possible to prove

R*5 If $\vdash \alpha$ then $\vdash \Box_1 \alpha$.

Proof.(1) α [Hypothesis](2) $\square_2 \alpha$ [(1), R5](3) $\square_1 \alpha$ [(2), A12, R1]

NOTE

1) The numbers in brackets following each of the theorems, S1-S7, refer to the corresponding theorems in Lewis and Langford [1].

BIBLIOGRAPHY

- [1] C. I. Lewis and C. H. Langford, Symbolic Logic, Dover, New York (1932), 506 pp.
- [2] W. T. Parry, Modalities in the "Survey System" of strict implication, The Journal of Symbolic Logic, v. 4 (1939), 137-154.
- [3] J. E. Rubin, Bi-modal logic, double-closure algebras, and Hilbert space. Zeitschrift für Mathematische Logik und Grundlagen der Mathematik, v. 8 (1962), 305-322.

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