

## LOCATING VERTICES OF TREES

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Let  $X$  be a nonempty set and let  $R$  and  $S$  be binary relations on  $X$ . Let  $x, y, x_1, x_2, y_1, y_2$  be arbitrary elements of  $X$ , then, where  $R(R)$  is the range of  $R$  and  $\omega$  is the set of nonnegative integers,  $\langle X, R, S \rangle$  is called a (*dyadic ordered*) *tree* if the following hold:

- (1) If  $x_1Ry$  and  $x_2Ry$ , then  $x_1 = x_2$ .
- (2) For each  $x \in X$ ,  $xRy$  for at most two  $y$ .
- (3)  $X - R(R)$  is a unit set,  $\{x_0\}$ .
- (4)  $y_1Sy_2$  iff (a)  $y_1 \neq y_2$ , (b)  $y_2Ry_1$ , and (c) for some  $x \in X$ , both  $xRy_1$  and  $xRy_2$ .
- (5) There exists a function  $l: X \rightarrow \omega$  with the properties: (a)  $l(x_0) = 0$  and (b) if  $xRy$ , then  $l(y) = l(x) + 1$ .

This definition, with minor modifications, is essentially the one given in [1].

If  $\langle X, R, S \rangle$  is a tree, then the elements of  $X$  are called *points* or *vertices*. If  $xRy$  holds for a unique  $y \in X$ ,  $x$  is called a *simple point*; if  $xRy$  holds for two distinct  $y$ ,  $x$  is called a *junction point*. Whenever  $xRy$  then  $y$  is said to be an *immediate successor* of  $x$ . The relation  $S$ , in effect, selects one of the two immediate successors of a junction point. Thus if  $xRy_1, xRy_2$  and  $y_1Sy_2$ , we say that  $y_1$  is the *left successor* and  $y_2$  the *right successor* of  $x$ .

$l(x)$  is called the *level* of  $x$ .  $l_n$  will denote the set of vertices of level  $n$ ,  $n \in \omega$ . Each  $l_n$  has at most  $2^n$  vertices; hence for any tree  $\langle X, R, S \rangle$ ,  $X$  must be countable. Note that  $\langle X, R, S \rangle$  has no junction points iff  $l$  is one-one iff  $S = \emptyset$ .

A *path* of a tree  $\langle X, R, S \rangle$  is a finite sequence  $[a_0, a_1, \dots, a_n]$  or a denumerable sequence  $[a_0, a_1, \dots, a_n, \dots]$  with the properties:

- (1) for each  $a_k$  appearing in the sequence,  $a_k \in X$  and
- (2) if  $a_{k+1}$  also appears in the sequence, then  $a_{k+1}$  is an immediate successor of  $a_k$ .

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The path  $[a_0, a_1, \dots, a_n]$  is called a *path from  $a_0$  to  $a_n$* . For any distinct  $x_i, x_j \in X$ , we say  $x_i$  *precedes*  $x_j$  if there is a path from  $x_i$  to  $x_j$ .

**Theorem 1.** (*Induction Principle for Trees*) Let  $\langle X, R, S \rangle$  be a tree. Let  $A$  be a subset of  $X$  which satisfies: (1)  $l_0 \subseteq A$  and, (2) whenever  $l_n \subseteq A$ , then  $l_{n+1} \subseteq A$ , then  $A = X$ .

**Theorem 2.** Let  $\mathcal{P}$  be the set of positive integers. Let  $\langle X, R, S \rangle$  be a tree. Then there is a unique function  $L : X \rightarrow \mathcal{P}$  such that (1)  $L(x_0) = 1$  and such that (2) whenever  $y \in R(R)$  and  $xRy$ , then (a)  $L(y) = 2L(x)$  if either  $x$  is a simple point or  $y$  is the left successor of  $x$ , or (b)  $L(y) = 2L(x) + 1$  if  $y$  is the right successor of  $x$ .

*Proof.* For each  $x \in X$ , let  $A(x)$  be the set of all  $a \in x$  such that  $a$  is a right successor (of some junction point) and either  $a = x$  or  $a$  precedes  $x$ . Let

$$L(x) = 2^{l(x)} + \sum_{a \in A(x)} 2^{l(x)-l(a)}.$$

We have

$$L(x_0) = 2^{l(x_0)} + \sum_{a \in \emptyset} 2^{l(x_0)-l(a)} = 1.$$

Let  $xRy$  and suppose that either  $x$  is a simple point or else  $y$  is the left successor of  $x$ . Then  $A(y) = A(x)$ , whereas,

$$(1) \quad l(y) = l(x) + 1.$$

Hence,

$$\begin{aligned} L(y) &= 2^{l(y)} + \sum_{a \in A(y)} 2^{l(y)-l(a)} \\ &= 2^{l(x)+1} + \sum_{a \in A(x)} 2^{l(x)+1-l(a)} \\ &= 2 \left( 2^{l(x)} + \sum_{a \in A(x)} 2^{l(x)-l(a)} \right) \\ &= 2L(x). \end{aligned}$$

Now suppose that  $y$  is the right successor of  $x$ . Then  $A(y) = A(x) \cup \{y\}$ ; again, (1) holds.

This time

$$\begin{aligned} L(y) &= 2^{l(y)} + \sum_{a \in A(y)} 2^{l(y)-l(a)} \\ &= 2^{l(x)+1} + \sum_{a \in A(x)} 2^{l(x)+1-l(a)} + 2 \\ &= 2L(x) + 1. \end{aligned}$$

If  $L^* : X \rightarrow \omega$  is any function satisfying conditions (1) and (2) of the present theorem, then by the Principle of Induction for trees it follows that  $L^* = L$ . For let  $B = \{x \in X : L^*(x) = L(x)\}$ . Then  $l_0 \subseteq B$  because by (1),  $L^*(x_0) = L(x_0) = 1$ . Suppose  $l_n \subseteq B$ . If  $l_{n+1} = \emptyset$ , then surely  $l_{n+1} \subseteq B$ . Otherwise, let  $y \in l_{n+1}$  and let  $x$  be the unique vertex satisfying  $xRy$ . We have

$x \in L_n \subseteq B$ ; hence  $L^*(x) = L(x)$ . If  $x$  is a simple point or if  $y$  is the left successor of  $x$ , then

$$L^*(y) = 2L^*(x) = 2L(x) = L(y).$$

In case  $y$  is the right successor of  $x$ , then

$$L^*(y) = 2L^*(x) + 1 = 2L(x) + 1 = L(y).$$

Thus  $l_{n+} \subseteq B$  and by theorem 1,  $B = X$ .

*Corollary.*  $2^{l(x)} \leq L(x) < 2^{l(x)+1}$  for all  $x \in X$ .

By means of  $L$  we can locate elements of the tree; thus we call  $L$  the *location function* of the tree.  $L$  is especially useful in trees whose vertices are (occurrences of) (1) subformulas of a given formula or (2) probability events. In particular, in the case of an analytic tableau for a formula  $P$  (see [1]), various subformulas of  $P$  are repeated again and again. It might be convenient to index the subformulas of the tableau by their locations. Thus if  $Q$  is a subformula of  $P$  which occurs in the tableau, we replace each occurrence of  $Q$  by  $Q_n$ —so that  $L(Q_n) = n$ . Moreover,  $L$  can be used to specify the relations  $R$  and  $S$  (in the definition of “tree”) in the following sense.

**Theorem 3.** *Let  $X$  be a countable set. Let  $L : X \rightarrow \mathcal{P}$  be any one-one function satisfying (a)  $1 \in R(L)$  and (b) for any  $n \geq 1$ , whenever  $2n + 1 \in R(L)$ , then  $2n \in R(L)$  and whenever  $2n \in R(L)$ , then  $n \in R(L)$ . Then there is a unique tree  $\langle X, R, S \rangle$  for which*

- (i)  $xRy$  iff either  $L(y) = 2L(x)$  or  $L(y) = 2L(x) + 1$ , and
- (ii)  $y_1Sy_2$  iff  $L(y_1)$  is even and  $L(y_2) = L(y_1) + 1$ .

*Proof.* Suppose  $R$  and  $S$  are binary relations on  $X$  defined by (i) and (ii) of theorem 3. In order to show that  $\langle X, R, S \rangle$  is a tree we must show that clauses (1)-(5) of the definition of “tree” hold.

(1) Suppose  $x_1Ry$  and  $x_2Ry$ . Then for  $i = 1, 2$ ,  $L(y) = 2L(x_i)$  or  $L(y) = 2L(x_i) + 1$ . Parity considerations assure that either  $L(y) = 2L(x_1) = 2L(x_2)$  or else  $L(y) = 2L(x_1) + 1 = 2L(x_2) + 1$ . Thus  $x_1 = x_2$  because  $L$  is one-one.

(2) follows from (i).

(3) According to (a),  $1 \in R(L)$ ; since  $L$  is one-one, there is a unique element  $x_0$  in  $X$  such that  $L(x_0) = 1$ .  $1 \notin R(R)$ , by (i). Suppose  $x \in X - \{x_0\}$ . Then  $L(x) > 1$ ; hence for some  $n \geq 1$ ,  $L(x) = 2n$  or  $L(x) = 2n + 1$ . Thus  $\{n, L(x)\} \subseteq R(L)$  by (b), and by (i),  $L^{-1}(n)Rx$ . Consequently,  $X - R(R) = \{x_0\}$ .

(4) Suppose  $y_1Sy_2$ . Then  $L(y_2) = L(y_1) + 1$ . Since  $L$  is a function,  $y_1 \neq y_2$ . Were  $y_2Sy_1$  as well as  $y_1Sy_2$  to hold, then  $L(y_1) = L(y_2) + 1 = L(y_1) + 2$ . Contradiction! Finally, let  $L(y_1) = 2n$ . Then  $n \in R(L)$  and  $L^{-1}(n)$  is unique.  $L(y_1) = 2L(L^{-1}(n))$  and  $L(y_2) = 2L(L^{-1}(n)) + 1$ ; hence  $L^{-1}(n)Ry_1$  and  $L^{-1}(n)Ry_2$ .

Now suppose that (a), (b), and (c) of (4) hold, and let  $x$  be as in (c). Then for  $i = 1, 2$ ,  $xRy_i$ ; hence by (i),  $L(y_i) = 2L(x)$  or  $L(y_i) = 2L(x) + 1$ . Since  $L$  is one-one and  $y_1 \neq y_2$ , we have either  $L(y_1) = 2L(x)$  and  $L(y_2) =$

$2L(x) + 1$ , or else  $L(y_2) = 2L(x)$  and  $L(y_1) = 2L(x) + 1$ . In the latter case, by (ii), we would have  $y_2Sy_1$ , contradicting (b) of (4). Thus the former case holds and  $y_1Sy_2$ .

(5) Define  $l: X \rightarrow \omega$  by  $l(x) = m$  if  $2^m \leq L(x) < 2^{m+1}$ . Then  $l(x_0) = 0$  because  $L(x_0) = 2$ . Let  $xRy$ . First suppose  $L(y) = 2L(x)$ . Then  $2^m \leq L(x) < 2^{m+1}$  iff  $2^{m+1} \leq L(y) < 2^{m+2}$ . Now suppose  $L(y) = 2L(x) + 1$ . Then  $2^m \leq L(x) < 2^{m+1}$  iff  $2^{m+1} \leq L(y) - 1 < 2^{m+2}$  iff

$$(2) \quad 2^{m+1} < L(y) < 2^{m+2} + 1.$$

Since  $L(y)$  is odd, (2) holds iff

$$2^{m+1} \leq L(y) < 2^{m+2}.$$

Thus in either case for  $L(y)$ , we have  $L(y) = l(x) + 1$ .

The uniqueness of the tree  $\langle X, R, S \rangle$  follows from the fact that the relations  $R$  and  $S$  are completely determined in terms of the given function  $L$ .

#### REFERENCE

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