

## RECURSIVE SEQUENCES AND FAITHFULLY FLAT EXTENSIONS

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**ABSTRACT.** For any faithfully flat morphism  $A \rightarrow B$  of Noetherian normal domains, a power series with coefficients in  $A$  which is rational over  $B$ , is already rational over  $A$ . The proof uses the fact that a sequence is recursive whenever it is recursive over some faithfully flat extension.

**0. Introduction.** It is well known that a necessary condition for a ring morphism  $A \rightarrow B$  to be faithfully flat is that *any linear system of equations with coefficients from  $A$  which has a solution over  $B$ , must already have a solution over  $A$* . In fact, if we strengthen this condition to *any solution over  $B$  comes from solutions over  $A$  by base change*, then this also becomes a sufficient condition for being faithfully flat. We could paraphrase the necessary condition as follows: *any linear system of equations over  $A$  which is solvable over a faithfully flat extension  $B$  of  $A$ , is already solvable over  $A$* .

In this paper I present another necessary condition of the same flavor. The key definition is that of a (linear) recursive sequence  $(x_n)_n$  over a ring  $A$ , as a sequence satisfying some fixed linear relation over  $A$  among  $t$  consecutive terms. I show that if  $A \rightarrow B$  is faithfully flat and  $(x_n)_n$  is a sequence of elements in  $A$  satisfying a linear recursion relation with coefficients in  $B$ , then it already satisfies such a recursion relation (of the same length) with coefficients in  $A$ . As there is a strong connection between recursive sequences and rational power series, I obtain the following corollary. Assume, moreover, that  $A$  and  $B$  are normal domains; then any power series over  $A$  which is rational (meaning that it can be written as a quotient of two polynomials) over  $B$ , is already rational over  $A$ . Any direct attempt, however, to prove this corollary just using faithfully flatness seems to fail, as far as I can tell.

**1. Definition.** Let  $A$  be a Noetherian ring and let  $\mathbf{x} = (x_n)_{n < \omega}$  be a (countable) sequence of elements of  $A$ . We say that  $\mathbf{x}$  is *recursive over  $A$*

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(of length  $t$ ), if there exists a linear form  $L(X) = r_0X_0 + \cdots + r_{t-1}X_{t-1} \in A[X]$ , where  $t \in \mathbf{N}$  and  $X = (X_i)_{i < t}$ , such that for all  $n \gg 0$  we have that

$$L(x_n, \dots, x_{n+t-1}) = x_{n+t}.$$

Here we have used the notation  $n \gg 0$  as an abbreviation for *there exists  $n_0$ , such that for all  $n \geq n_0$* .

**2. Theorem.** *Let  $A \rightarrow B$  be a morphism of Noetherian rings and let  $\mathbf{x}$  be a sequence in  $A$ . If  $B$  is faithfully flat over  $A$  and  $\mathbf{x}$  is recursive over  $B$ , then  $\mathbf{x}$  is already recursive over  $A$  (of the same length).*

*Proof.* Let  $\mathcal{L}_t(\mathbf{x})$  be the collection of all linear forms  $L(X) \in A[X]$ , where  $X = (X_i)_{i \leq t}$  and  $t \in \mathbf{N}$ , such that for all  $n \gg 0$ , we have that  $L(x_n, \dots, x_{n+t}) = 0$ . Evidently,  $\mathcal{L}_t(\mathbf{x})$  carries the structure of an  $A$ -module. Let  $e$  be the  $(t+1)$ -tuple  $(0, \dots, 0, -1)$ , and let

$$\mathfrak{a}_t^A(\mathbf{x}) = \{L(e) \mid L \in \mathcal{L}_t(\mathbf{x})\}.$$

As the latter is the image of  $\mathcal{L}_t(\mathbf{x})$  under the morphism  $A[X] \rightarrow A$  defined by substituting  $e$  for  $X$ , it is an ideal of  $A$ .

We claim that  $\mathbf{x}$  is recursive over  $A$  of length  $t$ , if and only if,  $\mathfrak{a}_t^A(\mathbf{x}) = A$ . Indeed, if  $\mathbf{x}$  is recursive over  $A$  of length  $t$ , then there is a linear form  $L(X) \in A[X]$ , with  $X = (X_i)_{i < t}$  such that for all  $n \gg 0$  we have that  $L(x_n, \dots, x_{n+t-1}) = x_{n+t}$ . Let  $L'(X, X_t) = L(X) - X_t$  so that  $L' \in \mathcal{L}_t(\mathbf{x})$ . Since  $1 = L'(e)$ , we proved one direction and the converse follows along the same lines.

Now, by assumption,  $\mathbf{x}$  is recursive over  $B$  of length  $t$  so that by the criterion we just established,  $\mathfrak{a}_t^B(\mathbf{x}) = B$ . Hence let  $L(X) = b_0X_0 + \cdots + b_tX_t \in B[X]$  be a witness to this, where  $X = (X_i)_{i \leq t}$ , i.e., such that for some  $n_0 \in \mathbf{N}$ , we have, for all  $n \geq n_0$ , that  $L(x_n, \dots, x_{t+n}) = 0$  and  $L(e) = 1$ . Therefore  $b_t = -1$ . Put  $b = (b_i)_{i < t} \in B^t$ . For every  $n \geq n_0$ , let  $M_n(Y) \in A[Y]$ , where  $Y = (Y_i)_{i \leq t}$ , be defined as

$$M_n(Y) = x_nY_0 + \cdots + x_{n+t}Y_t.$$

By Noetherianity, there exists some  $n_1 \geq n_0$  such that each  $M_n$ , for  $n \geq n_0$ , lies in the ideal of  $A[Y]$  generated by  $M_{n_0}, \dots, M_{n_1}$ . In other

words, there exists  $p_{n,k}(Y) \in A[Y]$  such that, for all  $n \geq n_0$  and all  $k$  with  $n_0 \leq k \leq n_1$ , we have that

$$(1) \quad M_n(Y) = \sum_{k=n_0}^{n_1} p_{n,k}(Y)M_k(Y).$$

By construction,  $M_k(b, -1) = 0$  for  $n_0 \leq k \leq n_1$ . By flatness, we can find finitely many  $a^{(j)} \in A^t$ ,  $e^{(j)} \in A$  and  $\beta^{(j)} \in B$  such that

$$(2) \quad M_k(a^{(j)}, e^{(j)}) = 0,$$

for all  $j < s$  and  $n_0 \leq k \leq n_1$ , and

$$(3) \quad b = \sum_{j < s} \beta^{(j)} a^{(j)} \quad -1 = \sum_{j < s} \beta^{(j)} e^{(j)}.$$

However, from (1) and (2), it then follows that  $M_n(a^{(j)}, e^{(j)}) = 0$  for all  $j < s$  and all  $n \geq n_0$ . This means that  $L^{(j)}(X) = a_0^{(j)}X_0 + \dots + a_{t-1}^{(j)} + e^{(j)}X_t$  lies in  $\mathcal{L}_t(\mathbf{x})$ , for each  $j < s$ , where  $a^{(j)} = (a_i^{(j)})_{i < t}$ . Hence  $-e^{(j)} \in \mathfrak{a}_t^A(\mathbf{x})$ . Together with (3), we therefore conclude that  $\mathfrak{a}_t^A(\mathbf{x})B = B$ . But faithfully flatness then implies that  $\mathfrak{a}_t^A(\mathbf{x}) = A$ , which by the above criterion means that  $\mathbf{x}$  is recursive over  $A$  of length  $t$ .  $\square$

**3. Proposition.** *Let  $A$  be a Noetherian ring and  $\mathbf{x}$  a sequence in  $A$ . Let  $\xi_{\mathbf{x}}(T) \in A[[T]]$  be the generating series of  $\mathbf{x}$ , i.e.,*

$$\xi_{\mathbf{x}}(T) = \sum_{n=0}^{\infty} x_n T^n.$$

*Then  $\mathbf{x}$  is recursive over  $A$ , if and only if,  $\xi_{\mathbf{x}}(T)$  lies in  $A(T)$ , where the latter ring is the localization of  $A[[T]]$  with respect to the multiplicative set  $1 + (T)A[[T]]$ .*

*Proof.* Suppose that  $\mathbf{x}$  is recursive and let  $\xi(T) = \xi_{\mathbf{x}}(T)$ . There is some  $n_0$  and some  $a_k \in A$  for  $k < t$  such that

$$(4) \quad x_n = a_0 x_{n-t} + \dots + a_{t-1} x_{n-1}$$

for all  $n \geq n_0$ . We want to show that

$$(5) \quad \xi(T) \equiv Q(T)\xi(T) \pmod{A[T]}$$

for some polynomial  $Q(T) \in (T)A[T]$ . Indeed, if (5) holds, then

$$\xi(T) = Q(T)\xi(T) + P(T)$$

for some  $P(T) \in A[T]$  and hence  $\xi(T) = P(T)/(1 - Q(T))$  as required.

We work in the  $A$ -module  $A[[T]]/A[T]$ . Clearly

$$\xi(T) \equiv T^{n_0} \cdot \xi_{\mathbf{x}'}(T) \pmod{A[T]},$$

where  $\mathbf{x}'$  is the sequence obtained from  $\mathbf{x}$  by deleting the first  $n_0$  elements. Hence in order to prove (5) we may work with this new recursive sequence and hence assume from the start that  $n_0 = 0$ . We have, using (4), that

$$\begin{aligned} \xi(T) &= \sum_{n < t} x_n T^n + \sum_{n \geq t} x_n T^n \\ &\equiv \sum_{n \geq t} (a_0 x_{n-t} + \cdots + a_{t-1} x_{n-1}) T^n \pmod{A[T]} \\ &\equiv a_0 T^t \xi(T) + \cdots + a_{t-1} T \xi(T) \pmod{A[T]}. \end{aligned}$$

This proves (5).

The converse is an easy exercise and is left to the reader.  $\square$

**4. Corollary.** *Let  $\varphi : A \rightarrow B$  be a morphism of Noetherian rings. If  $\varphi$  is faithfully flat, then*

$$(6) \quad A(T) = A[[T]] \cap B(T),$$

for  $T$  a single variable.

*Proof.* The  $\subset$ -inclusion is immediate, hence take  $F$  in the righthand side of (6). Taking the coefficients of  $F$  as the members of a sequence  $\mathbf{x}$  over  $A$ , we have that  $F(T) = \xi_{\mathbf{x}}(T)$ . Since  $F(T) \in B(T)$ , the sequence  $\mathbf{x}$  is recursive over  $B$  by Proposition 3. Therefore, by faithfully flatness

and Theorem 2,  $\mathbf{x}$  is also recursive over  $A$  which by Proposition 3 again means that  $F \in A(T)$ .  $\square$

**5. Corollary.** *Let  $A \rightarrow B$  be a faithfully flat morphism of Noetherian normal domains. Then any power series with coefficients in  $A$  which is rational over  $B$ , is already rational over  $A$ .*

*Proof.* A Noetherian normal domain  $A$  has the Fatou property by Chabert [1, Section 3], meaning that  $A(T)$  is equal to the intersection of the fraction field of  $A[T]$  with  $A[[T]]$ , i.e.,  $A(T)$  is the ring of rational power series and we can apply Corollary 4.

#### REFERENCES

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