SOME NUMERIC RESULTS ON ROOT SYSTEMS

Jian-yi Shi

Let Φ be an irreducible root system (sometimes we denote Φ by $\Phi(X)$ to indicate its type X). Choose a simple root system Π in Φ . Let Φ^+ (resp. Φ^-) be the corresponding positive (resp. negative) root system of Φ . By a subsystem Φ' of Φ (resp. of Φ^+), we mean that Φ' is a subset of Φ (resp. of Φ^+) which itself forms a root system (resp. a positive root system). We refer the readers to Bourbaki's book for the detailed information about root systems. Among all subsystems of Φ , the subsystems of Φ of rank 2 and of type $\neq A_1 \times A_1$ are of particular importance in the theory of Weyl groups and affine Weyl groups (see the papers by Jian-yi Shi). In the present paper, we shall compute the number of such subsystems of Φ for an irreducible root system Φ of any type. Some interesting properties of Φ are also obtained.

1. The number $h(\alpha)$. Let \langle , \rangle be an inner product of the euclidean space *E* spanned by Φ . For any $\alpha \in \Phi$, we denote by $|\alpha|$ the length of α , by α^{\vee} the dual root $2\alpha/\langle \alpha, \alpha \rangle$ of α and by s_{α} the reflection in *E* which sends any vector $v \in E$ to $s_{\alpha}(v) = v - \langle v, \alpha^{\vee} \rangle \alpha$. For $\alpha, \beta \in \Phi$, we write $\alpha < \beta$ if $\beta - \alpha$ is a sum of some positive roots.

For $\alpha \in \Phi$, we define the sets $D(\alpha) = \{\beta \in \Phi | \alpha + \beta \in \Phi\}$, $D^+(\alpha) = D(\alpha) \cap \Phi^+$ and $D^-(\alpha) = D(\alpha) \cap \Phi^-$. Let $d(\alpha)$ be the cardinality of the set $D^+(\alpha)$. Also, we denote by $ht(\alpha)$ the height of α , i.e. $ht(\alpha) = \sum_{\beta \in \Pi} a_\beta$ if $\alpha = \sum_{\beta \in \Pi} a_\beta \beta$ with $a_\beta \in \mathbb{Z}$.

For any $\alpha \in \Phi^+$, there exists a sequence ξ of roots $\alpha_1 = \alpha, \alpha_2, \ldots, \alpha_r$ in Φ^+ such that $\alpha_r \in \Pi$ and for every $i, 1 < i \leq r$, we have $\alpha_{i-1} > \alpha_i = s_{\delta_i}(\alpha_{i-1})$ for some $\delta_i \in \Pi$. Such a sequence ξ is called a root path from α to Π . We denote by $h(\alpha, \xi)$ the length r of ξ . We shall deduce a formula for the number $h(\alpha, \xi)$, from which we shall see that $h(\alpha, \xi)$ is actually independent on the choice of a root path ξ from α to Π but only dependent on the root α .

Note that if the root system Φ contains roots of two different lengths and if $\alpha = \sum_{\beta \in \Pi} a_{\beta}\beta$ is a long root of Φ with $a_{\beta} \in \mathbb{Z}$ then each coefficient a_{β} with β short is divisible by $|\alpha|^2/|\beta|^2$.

LEMMA 1.1. Let $\alpha = \sum_{\beta \in \Pi} a_{\beta}\beta$, $a_{\beta} \in \mathbb{Z}$, be a root of Φ^+ and let ξ be a root path from α to Π . Then

JIAN-YI SHI

(i) If either all the roots of Φ have the same length or α is a short root of Φ with Φ containing roots of two different lengths, then $h(\alpha, \xi) = ht(\alpha)$;

(ii) If α is a long root of Φ with Φ containing roots of two different lengths, then

$$h(x, \xi) = \sum_{\beta \in \Pi} \frac{|\beta|^2}{|\alpha|^2} a_{\beta}.$$

Proof. Let $\alpha_1 = \alpha$, α_2 , ..., α_r be a root path from α to Π . Then in case (i), we have $ht(\alpha_i) = ht(\alpha_{i+1}) + 1$ for any $i, 1 \le i < r$, by the fact that $\langle \alpha_i, \delta_i^{\vee} \rangle = 1$, where $\delta_i \in \Pi$ satisfies the relation $\delta_i(\alpha_{i-1}) = \alpha_i$. So assertion (i) follows immediately by applying induction on $ht(\alpha) \ge 1$. Next assume that we are in case (ii). Again apply induction on $ht(\alpha) \ge 1$. If $ht(\alpha) = 1$, then $\alpha \in \Pi$ and the result is obviously true. Now assume $ht(\alpha) > 1$. Let $\xi : \alpha_1 = \alpha, \alpha_2, \ldots, \alpha_r$ be a root path from α to Π . Then $\xi' : \alpha_2, \alpha_3, \ldots, \alpha_r$ is a root path from α_2 to Π with $ht(\alpha_2) < ht(\alpha)$ and $\alpha_2 = s_{\delta}(\alpha)$ for some $\delta \in \Pi$. Note that α_2 is a long root of Φ . Write

$$lpha_2 = \sum_{eta \in \Pi} a'_eta eta \ , \qquad a'_eta \in \mathbb{Z}.$$

Then by inductive hypothesis, we have

$$h(\alpha_2, \xi') = \sum_{\beta \in \Pi} \frac{|\beta|^2}{|\alpha_2|^2} a'_{\beta}.$$

Since $\langle \alpha, \delta^{\vee} \rangle = |\alpha|^2 / |\delta|^2$ by the assumption $s_{\delta}(\alpha) < \alpha$, we have

$$\alpha = \alpha_2 + \frac{|\alpha|^2}{|\delta|^2} \delta = \sum_{\substack{\beta \in \Pi \\ \beta \neq \delta}} a'_{\beta} \beta + \left(a'_{\delta} + \frac{|\alpha|^2}{|\delta|^2}\right) \delta.$$

This implies that

$$\begin{split} h(\alpha, \xi) &= h(\alpha_2, \xi') + 1 = \sum_{\substack{\beta \in \Pi}} \frac{|\beta|^2}{|\alpha_2|^2} a'_{\beta} + 1 \\ &= \sum_{\substack{\beta \in \Pi \\ \beta \neq \delta}} \frac{|\beta|^2}{|\alpha_2|^2} a'_{\beta} + \frac{|\delta|^2}{|\alpha_2|^2} \left(a'_{\delta} + \frac{|\alpha_2|^2}{|\delta|^2} \right) = \sum_{\substack{\beta \in \Pi \\ |\alpha|^2}} \frac{|\beta|^2}{|\alpha|^2} a_{\beta} \end{split}$$

by noting $|\alpha| = |\alpha_2|$.

 \Box

We see from Lemma 1.1 that, for any $\alpha \in \Phi^+$, the length of a root path ξ from α to Π is only dependent on α but not on the choice of the path ξ . So we can denote $h(\alpha, \xi)$ simply by $h(\alpha)$.

Let Φ^{\vee} be the dual root system of Φ , i.e. $\Phi^{\vee} = \{\alpha^{\vee} | \alpha \in \Phi\}$. Then $\Pi^{\vee} = \{\alpha^{\vee} | \alpha \in \Pi\}$ and $(\Phi^{\vee})^+ = \{\alpha^{\vee} | \alpha \in \Phi^+\}$ are a simple root system and the corresponding positive root system of Φ^{\vee} , respectively. We can define the number $h^{\vee}(\alpha^{\vee})$ for any $\alpha^{\vee} \in (\Phi^{\vee})^+$ in the same way as that for a root of Φ . That is, $h^{\vee}(\alpha^{\vee})$ is the length of a root path from α^{\vee} to Π^{\vee} in $(\Phi^{\vee})^+$.

LEMMA 1.2. For any $\alpha \in \Phi^+$, we have $h(\alpha) = h^{\vee}(\alpha^{\vee})$.

Proof. For any $\delta \in \Pi$, we have the following equivalence.

$$(1) \qquad s_{\delta}(\alpha) < \alpha \Leftrightarrow \langle \alpha , \delta^{\vee} \rangle > 0 \Leftrightarrow \langle \alpha^{\vee} , \delta \rangle > 0 \Leftrightarrow s_{\delta^{\vee}}(\alpha^{\vee}) < \alpha^{\vee}.$$

Apply induction on $h(\alpha) \ge 1$. When $h(\alpha) = 1$, we have $\alpha \in \Pi$ and hence $\alpha^{\vee} \in \Pi^{\vee}$. So $h^{\vee}(\alpha^{\vee}) = 1$, and the result is true in this case. Now assume $h(\alpha) > 1$. Then there exists some $\delta \in \Pi$ with $\langle \alpha, \delta^{\vee} \rangle > 0$. So $h(s_{\delta}(\alpha)) = h(\alpha) - 1$. By inductive hypothesis, we have

(2)
$$h(s_{\delta}(\alpha)) = h^{\vee}((s_{\delta}(\alpha))^{\vee}) = h^{\vee}(s_{\delta^{\vee}}(\alpha^{\vee})).$$

But by (1), we have

$$h^{\vee}(s_{\delta^{\vee}}(\alpha^{\vee})) = h^{\vee}(\alpha^{\vee}) - 1$$

Thus we get $h(\alpha) = h^{\vee}(\alpha^{\vee})$.

2. The number $d(\alpha)$. We shall deduce a formula for the number $d(\alpha)$ for any $\alpha \in \Phi^+$.

For $\alpha, \beta \in \Phi$, we call all roots of the form $\alpha + i\beta$ $(i \in \mathbb{Z})$ the β -string through α . Let $\alpha \in \Phi^+$ and $\delta \in \Pi$ satisfy the inequality $\langle \alpha, \delta^{\vee} \rangle > 0$. Then it is easily seen that $\alpha, \alpha - \delta, \ldots, \alpha - \langle \alpha, \delta^{\vee} \rangle \delta$ is the δ -string through α except for the case when α is the highest short root of the root system of type G_2 .

LEMMA 2.1. Given $\alpha \in \Phi^+$ and $\delta \in \Pi$ with $\langle \alpha, \delta^{\vee} \rangle > 0$. Let $\alpha' = s_{\delta}(\alpha)$. Then (i) $D(\alpha') = s_{\delta}(D(\alpha))$.

(ii) $s_{\delta}(D^+(\alpha')) = D^+(\alpha) \cup \{-\delta\}$, provided that α is not the highest short root of the root system of type G_2 ;

(iii) $d(\alpha') = d(\alpha) + 1$ under the same assumption as that in (ii).

Proof. (i) $\beta \in D(\alpha') \Leftrightarrow \beta + \alpha' \in \Phi \Leftrightarrow s_{\delta}(s_{\delta}(\beta) + \alpha) \in \Phi \Leftrightarrow s_{\delta}(\beta) + \alpha \in \Phi \Leftrightarrow s_{\delta}(\beta) \in D(\alpha) \Leftrightarrow \beta \in s_{\delta}(D(\alpha))$.

(ii) First we shall show $s_{\delta}(D^+(\alpha)) \subset D^+(\alpha')$. Let $\beta \in s_{\delta}(D^+(\alpha))$. Then $\beta \in D(\alpha')$ by (i). If $\beta \in D^-(\alpha') \subseteq \Phi^-$, then by the fact $s_{\delta}(\beta) \in D^+(\alpha) \subseteq \Phi^+$, we have $\beta = -\delta$. Since $\alpha, \alpha - \delta, \ldots, \alpha - \langle \alpha, \delta^{\vee} \rangle \delta$ is the δ -string through α by the above remark, we see that $\alpha + s_{\delta}(\beta) = \alpha + \delta \notin \Phi$ which contradicts the condition $s_{\delta}(\beta) \in D^+(\alpha)$. Thus we have $\beta \in D^+(\alpha')$ and so $s_{\delta}(D^+(\alpha)) \subset D^+(\alpha')$, i.e. $D^+(\alpha) \subset s_{\delta}(D^+(\alpha'))$.

It is obvious that $\{-\delta\} \subseteq s_{\delta}(D^+(\alpha'))$. Thus it remains to show the reversing inclusion. Now assume $\beta \in s_{\delta}(D^+(\alpha'))$. Then $s_{\delta}(\beta) \in D^+(\alpha')$. This implies that $s_{\delta}(\beta) + \alpha' \in \Phi$ and $s_{\delta}(\beta) \in \Phi^+$. Hence $\beta + \alpha \in \Phi$ and $s_{\delta}(\beta) \in \Phi^+$. But then we have either $\beta \in D^+(\alpha)$ or $\beta = -\delta$, which implies $s_{\delta}(D^+(\alpha')) \subseteq D^+(\alpha) \cup \{-\delta\}$.

(iii) This is an immediate consequence of (ii).

REMARK. In the case when the type of Φ is G_2 , let $\Pi = \{\gamma, \delta\}$ with δ short. Then $D^+(2\delta + \gamma) = \{\delta, \delta + \gamma\}$, $D^+(\delta + \gamma) = \{\delta, 2\delta + \gamma\}$ and $\delta + \gamma = s_{\delta}(2\delta + \gamma)$. Thus the results (ii), (iii) of Lemma 2.1 do not hold in this case.

In Φ^+ , let α^l be the highest long root and let α^s be the highest short root, where we stipulate $\alpha^s = \alpha^l$ in the case when all the roots of Φ have the same length.

Theorem 2.2. Given $\alpha \in \Phi^+$.

(i) If α is short and if the type of Φ is not G_2 , then

 $h(\alpha) + d(\alpha) = \operatorname{ht}(\alpha^l).$

(ii) If α is long, then

$$h(\alpha) + d(\alpha) = \operatorname{ht}(\alpha^s).$$

Proof. First assume that the result has been shown to be true in the case when $\alpha = \alpha^s$ in (i) and $\alpha = \alpha^l$ in (ii). Apply reversing induction on $h(\alpha) \le h(\alpha^s)$ in (i) and on $h(\alpha) \le h(\alpha^l)$ in (ii). Now assume that α is either short with $h(\alpha) < h(\alpha^s)$ or long with $h(\alpha) < h(\alpha^l)$. Then there must exist some $\delta \in \Pi$ with $\langle \alpha, \delta^{\vee} \rangle < 0$. So $\alpha' = s_{\delta}(\alpha) > \alpha$ with $h(\alpha') = h(\alpha) + 1$. We see $\langle \alpha', \delta^{\vee} \rangle > 0$. By Lemma 2.1(iii), we

have $d(\alpha') = d(\alpha) - 1$. So by inductive hypothesis, we get

$$h(\alpha) + d(\alpha) = (h(\alpha') - 1) + (d(\alpha') + 1)$$
$$= h(\alpha') + d(\alpha')$$
$$= \begin{cases} ht(\alpha^l) & \text{if } \alpha \text{ is short,} \\ ht(\alpha^s) & \text{if } \alpha \text{ is long,} \end{cases}$$

by noting $|\alpha| = |\alpha'|$.

Thus it remains to show that assertion (i) is true for $\alpha = \alpha^s$ and that assertion (ii) is true for $\alpha = \alpha^l$.

In the case when the Dynkin diagram is simply laced, we have $h(\alpha^s) = ht(\alpha^s)$ by Lemma 1.1(i). Clearly, $d(\alpha^s) = 0$. So our result is true in this case. Now assume that Φ contains roots of two different lengths. If Φ has type B_n , then $h(\alpha^s) = n$, $d(\alpha^s) = n - 1$, $ht(\alpha^l) = 2n - 1$, $d(\alpha^l) = 0$ and $h(\alpha^l) = h^{\vee}((\alpha^l)^{\vee}) = ht((\alpha^l)^{\vee}) = ht(\alpha^s) = 2n - 2$ by Lemmas 1.2 and 1.1(i). If Φ has type C_n , then $h(\alpha^s) = 2n - 2$, $d(\alpha^s) = 1$, $ht(\alpha^l) = 2n - 1$ and $d(\alpha^l) = 0$. We also have

$$h(\alpha^l) = h^{\vee}((\alpha^l)^{\vee}) = \operatorname{ht}((\alpha^l)^{\vee}) = \operatorname{ht}(\alpha^s) = n$$

by Lemmas 1.2 and 1.1(i). If Φ has type F_4 , then $h(\alpha^s) = 8$, $d(\alpha^s) = 3$, $ht(\alpha^l) = 11$ and $d(\alpha^l) = 0$. By the same reason as above, we have

$$h(\alpha^l) = h^{\vee}((\alpha^l)^{\vee}) = \operatorname{ht}((\alpha^l)^{\vee}) = \operatorname{ht}(\alpha^s) = 8.$$

If Φ has type G_2 , then $d(\alpha^l) = 0$ and $h(\alpha^l) = ht(\alpha^s) = 3$. Thus in all the cases, our result is true.

COROLLARY 2.3. Assume that the type of Φ is not G_2 . Then for any short root α of Φ^+ , we have the equation

$$ht(\alpha) + d(\alpha) = h - 1,$$

where h is the Coxeter number of Φ .

Proof. We have $h(\alpha) = ht(\alpha)$ by Lemma 1.1(i). Since $ht(\alpha^l) = h - 1$, our result follows immediately from Theorem 2.2(i).

3. The number of certain rank 2 subsystems in Φ . Let $g(\Phi)$ be the number of subsystems of Φ of rank 2 and of type other than $A_1 \times A_1$. Then $g(\Phi)$ is also equal to the number of positive subsystems of Φ^+ of rank 2 and of type $\neq A_1 \times A_1$. In this section, we shall compute the number $g(\Phi)$ for Φ of any type. JIAN-YI SHI

LEMMA 3.1. If the Dynkin diagram of Φ is simply laced, then

(3)
$$g(\Phi) = \frac{1}{2} \sum_{\alpha \in \Phi^+} d(\alpha).$$

Proof. Under our assumption, the only possible type for a subsystem of Φ^+ of rank 2 and of type $\neq A_1 \times A_1$ is A_2 . Each of such subsystems could be obtained by first taking a root $\alpha \in \Phi^+$ and then taking any root β in the set $D^+(\alpha)$ to form a subsystem $\{\alpha, \beta, \alpha + \beta\}$. Since such a subsystem is obtained twice in the above way, this implies the required formula (3) for the number $g(\Phi)$.

Define

$$H(\Phi) = \sum_{\alpha \in \Phi^+} \operatorname{ht}(\alpha), \quad H^s(\Phi) = \sum_{\substack{\alpha \in \Phi^+ \\ \text{short}}} \operatorname{ht}(\alpha) \text{ and}$$
$$H^l(\Phi) = \sum_{\substack{\alpha \in \Phi^+ \\ \text{long}}} \operatorname{ht}(\alpha).$$

These numbers could be computed for any irreducible root system Φ . Define $\binom{m}{n} = \frac{m!}{n!(m-n)!}$ for any integers $m, n, 0 \le n \le m$.

Lemma 3.2.

Type of Φ	$H(\mathbf{\Phi})$	$H^{s}(\mathbf{\Phi})$	$H^{l}(\mathbf{\Phi})$
$A_n \ (n \ge 1)$	$\binom{n+2}{3}$		
$B_n \ (n \ge 2)$	$\frac{n(n+1)(4n-1)}{6}$	$\binom{n+1}{2}$	$4\binom{n+1}{3}$
$C_n \ (n \ge 2)$	$\frac{n(n+1)(4n-1)}{6}$	$\frac{n(n-1)(4n+1)}{6}$	<i>n</i> ²
$D_n \ (n \ge 4)$	$\frac{n(n-1)(2n-1)}{3}$		
<i>E</i> ₆	156		
E_7	399		
<i>E</i> ₈	1240		
<i>F</i> ₄	110	46	64
<i>G</i> ₂	16	6	10

Now we can compute the numbers $g(\Phi)$ for Φ of types A_n , $n \ge 1$, D_m , $m \ge 4$, and E_i , i = 6, 7, 8 as follows.

THEOREM 3.3.

Type of Φ	$g(\Phi)$	
$A_n \ (n \ge 1)$	$\binom{n+1}{3}$	
$D_n \ (n \ge 4)$	$4\binom{n}{3}$	
E_6	120	
<i>E</i> ₇	336	
E_8	1120	

Proof. By Corollary 2.3 and Lemma 3.1, we have

$$g(\Phi) = \frac{1}{2} \sum_{\alpha \in \Phi^+} d(\alpha) = \frac{1}{2} \sum_{\alpha \in \Phi^+} (h - 1 - ht(\alpha))$$

= $\frac{1}{2} ((h - 1)|\Phi^+| - H(\Phi)).$

Thus we have $g(\Phi(A_n)) = \frac{1}{2}(n\binom{n+1}{2} - \binom{n+2}{3}) = \binom{n+1}{3}$ for $n \ge 1$. For $n \ge 4$, we have

$$g(\Phi(D_n)) = \frac{1}{2} \left((2n-3)n(n-1) - \frac{n(n-1)(2n-1)}{3} \right) = 4 \binom{n}{3}.$$

Also, we have $g(\Phi(E_6)) = \frac{1}{2}(11 \cdot 36 - 156) = 120$,

$$g(\Phi(E_7)) = \frac{1}{2}(17 \cdot 63 - 399) = 336,$$

- $\frac{1}{2}(29 \cdot 120 - 1240) = 1120$

and $g(\Phi(E_8)) = \frac{1}{2}(29 \cdot 120 - 1240) = 1120$.

Now assume that Φ contains roots of two different lengths and that the type of Φ is not G_2 . Then the possible types for a subsystem Φ' of Φ of rank 2 and of type $\neq A_1 \times A_1$ are A_2 and B_2 . Let $u(\Phi)$ be the cardinality of the set

 $\{\{\alpha, \beta\} \mid \alpha, \beta \in \Phi^+ \text{ have different lengths with } \alpha + \beta \in \Phi^+\}.$ Then it is easily seen that the following formula for $g(\Phi)$ holds.

(4)
$$g(\Phi) = \frac{1}{2} \sum_{\alpha \in \Phi^+} d(\alpha) - u(\Phi).$$

JIAN-YI SHI

First let us consider the case when Φ has type C_n , $n \ge 2$. We see that a subsystem Φ' of Φ has type A_2 only if all the roots in Φ' are short. This implies that for each long root $\beta \in \Phi^+$, the set $D^+(\beta)$ contains no long root and hence $u(\Phi) = \sum_{\beta \in \Phi^+ \text{ long }} d(\beta)$. So by (4), we get

$$g(\Phi) = \frac{1}{2} \sum_{\alpha \in \Phi^+} d(\alpha) - \sum_{\substack{\beta \in \Phi^+ \\ \text{long}}} d(\beta) = \frac{1}{2} \left(\sum_{\substack{\alpha \in \Phi^+ \\ \text{short}}} d(\alpha) - \sum_{\substack{\beta \in \Phi^+ \\ \text{long}}} d(\beta) \right)$$
$$= \frac{1}{2} \left(\sum_{\substack{\alpha \in \Phi^+ \\ \text{short}}} (h - 1 - \text{ht}(\alpha)) - \sum_{i=1}^n (i - 1) \right)$$

by Theorem 2.2, Corollary 2.3 and Lemma 1.2. Then by Lemma 3.2, we have

$$g(\Phi) = \frac{1}{2} \left((2n-1)n(n-1) - \frac{n(n-1)(4n+1)}{6} - \frac{n(n-1)}{2} \right)$$
$$= \frac{n(n-1)(4n-5)}{6}.$$

Since the root system of type B_n is the dual of the one of type C_n , there exists a bijection from the set of subsystems of the root system of type C_n to that of type B_n by sending Φ' to ${\Phi'}^{\vee}$. Such a bijective map preserves the ranks of subsystems and also preserves the types of them whenever their ranks are not greater than 2. This implies that we also have $g(\Phi) = \frac{n(n-1)(4n-5)}{6}$ when Φ has type B_n .

Next assume that Φ has type F_4 . By Theorem 2.2, Lemma 3.2 and Lemmas 1.1, 1.2, we get

$$\frac{1}{2}\sum_{\alpha\in\Phi^+} d(\alpha) = \frac{1}{2} \left(\sum_{\substack{\alpha\in\Phi^+\\\text{short}}} (\operatorname{ht}(\alpha^l) - \operatorname{ht}(\alpha)) + \sum_{\substack{\beta\in\Phi^+\\\text{long}}} (\operatorname{ht}(\alpha^s) - \operatorname{ht}(\beta^{\vee})) \right)$$
$$= \frac{1}{2} \left(\frac{1}{2} |\Phi^+| (\operatorname{ht}(\alpha^l) + \operatorname{ht}(\alpha^s)) - 2H^s(\Phi) \right)$$
$$= \frac{1}{2} \left(\frac{1}{2} \cdot 24 \cdot (11 + 8) - 92 \right)$$
$$= 68.$$

Also, by a direct computation, we get $u(\Phi) = 18$. So by (4), we have $g(\Phi) = 68 - 18 = 50.$

Finally, it is easily seen that $g(\Phi) = 3$ when Φ has type G_2 . Summing up, we get the following table.

THEOREM 3.4.

Type of Φ	$g(\Phi)$	
B_n or C_n $(n \ge 2)$	$\frac{n(n-1)(4n-5)}{6}$	
F_4	50	
<i>G</i> ₂	3	

From the above discussion, we can deduce even more precise conclusion. We note that in any irreducible root system Φ , there exist at most two different types of subsystems which have rank 2 and types $\neq A_1 \times A_1$. Let $g'(\Phi)$ be the number of subsystems of Φ of type A_2 and let $g''(\Phi)$ be the number of subsystems of Φ of type B_2 or G_2 . Then by Theorem 3.3, we have

$$g'(\Phi(B_n)) = g'(\Phi(C_n)) = g(\Phi(D_n)) = 4\binom{n}{3} \quad \text{for } n \ge 4$$

by noting that all the long (resp. short) roots of $\Phi(B_n)$ (resp. $\Phi(C_n)$) form a root system of type D_n . Hence we also have

$$g''(\Phi(B_n)) = g''(\Phi(C_n)) = g(\Phi(B_n)) - g'(\Phi(B_n))$$

= $\frac{n(n-1)(4n-5)}{6} - 4\binom{n}{3}$
= $\binom{n}{2}$.

On the other hand, we have

$$g''(\Phi(F_4)) = u(\Phi(F_4)) = 18$$
 and
 $g'(\Phi(F_4)) = g(\Phi(F_4)) - g''(\Phi(F_4)) = 50 - 18 = 32.$

Finally, it is obvious that $g'(\Phi(G_2)) = 2$ and $g''(\Phi(G_2)) = 1$. Summing up, we have the following table.

THEOREM 3.5.

Type of Φ	$g'(\mathbf{\Phi})$	$g''(\Phi)$
$B_n, C_n \ (n \geq 2)$	$4\binom{n}{3}$	$\binom{n}{2}$
F_4	32	18
<i>G</i> ₂	2	1

Proof. By the above discussion, it remains to show the result for Φ being of types B_m or C_m , m = 2, 3. But this could be checked directly.

References

- [1] N. Bourbaki, Groupes et Algèbres de Lie, Ch. 4-6, Hermann, Paris, 1968.
- [2] Jian-yi Shi, Alcoves corresponding to an affine Weyl group, J. London Math. Soc., 35 (1987), 42-55.
- [3] ____, Sign types corresponding to an affine Weyl group, J. London Math. Soc., 35 (1987), 56-74.

Received November 15, 1991. Supported by the National Science Foundation of China and by the Science Foundation of the University Doctoral Program of CNEC.

East China Normal University 3663 Zhongshan Road (Northern) Shanghai 200062, China