

Leibniz complexity of Nash functions on differentiations

By Goo ISHIKAWA and Tatsuya YAMASHITA

(Received Dec. 14, 2016)

(Revised Jan. 18, 2018)

Abstract. The derivatives of Nash functions are Nash functions which are derived algebraically from their minimal polynomial equations. In this paper we show that, for any non-Nash analytic function, it is impossible to derive its derivatives algebraically, i.e., by using linearity and Leibniz rule finite times. In fact we prove the impossibility of such kind of algebraic computations, algebraically by using Kähler differentials. Then the notion of Leibniz complexity of a Nash function is introduced in this paper, as a computational complexity on its derivative, by the minimal number of usages of Leibniz rules to compute the total differential algebraically. We provide general observations and upper estimates on Leibniz complexity of Nash functions, related to the binary expansions, the addition chain complexity, the non-scalar complexity and the complexity of Nash functions in the sense of Ramanakorasina.

1. Introduction.

Let $f = f(x_1, \dots, x_n)$ be a C^∞ function on an open subset $U \subset \mathbf{R}^n$. Then f is called a *Nash function* on U if f is analytic-algebraic on U , i.e. if f is analytic on U and there exists a non-zero polynomial $P(x, y) \in \mathbf{R}[x, y]$, $x = (x_1, \dots, x_n)$, such that $P(x, f(x)) = 0$ for any $x \in U$ ([15], [19], [3]). If U is semi-algebraic, then, f is a Nash function if and only if f is analytic and the graph of f in $U \times \mathbf{R} \subset \mathbf{R}^{n+1}$ is a semi-algebraic set ([3]). For further significant progress on global study of Nash functions, see [8].

An analytic function f on U is called *transcendental* if it is not a Nash function. Then in this paper we show that, for any transcendental function, it is impossible to algebraically derive its derivatives by using linearity and Leibniz rule (product rule) finite times, even by using any C^∞ function. In fact an analytic function f is a Nash function if and only if its derivatives $\partial f/\partial x_1, \dots, \partial f/\partial x_n$ are computable algebraically (Theorem 2.1). For example, for the transcendental function $f(x) = e^x$, the formula

$$\frac{d}{dx}e^x = e^x$$

is never proved algebraically but is proved only by a “transcendental” method. The statement above is formulated in terms of Kähler differential exactly.

We begin with the simple example of a Nash function $f(x) = \sqrt{x^2 + 1}$ of one variable. Then $f^2 - (x^2 + 1) = 0$. By differentiating both sides of the relation, we have $2f'f - 2x = 0$

2010 *Mathematics Subject Classification.* Primary 14P20; Secondary 58A07, 32C07.

Key Words and Phrases. Kähler differentials, complexity, extension theorem of Efroymsom.

The first author was supported by JSPS KAKENHI Grant No.JP15H03615 and No.JP15K13431.

where $f' = df/dx$. Here we have used Leibniz rule three times to get $(f^2)' = 2f'f$, $(x^2)' = 2x$ and $1' = 0$ by setting $dx/dx = 1$. Then we have $f'(x) = x/f(x) = x/\sqrt{x^2 + 1}$. If we suppose $c' = 0$ for a constant function c , then the usage of Leibniz rule is counted to be twice.

In general, let f be a Nash function on $U \subset \mathbf{R}^n$. Then there is a non-zero polynomial $P(x, y) \in \mathbf{R}[x, y]$, $x = (x_1, \dots, x_n)$ such that $P(x, f(x)) = 0$ for any $x \in U$. We impose the condition that $(\partial P/\partial y)(x, f(x))$ is not identically zero on U . The condition is achieved by choosing P which has the minimal total degree or the minimal degree on y , among polynomials P satisfying $P(x, f(x)) = 0$ on U . Then, by using Leibniz rule several times, we have

$$\frac{\partial P}{\partial x_i}(x, f(x)) + \frac{\partial P}{\partial y}(x, f(x)) \frac{\partial f}{\partial x_i}(x) = 0 \quad (1 \leq i \leq n).$$

Therefore we have the formula

$$\frac{\partial f}{\partial x_i}(x) = - \frac{\partial P}{\partial x_i}(x, f(x)) \Big/ \frac{\partial P}{\partial y}(x, f(x)) \quad (1 \leq i \leq n).$$

By our assumption that f is a Nash function and the assumption on P , $(\partial P/\partial y)(x, f(x))$ is a Nash function which is not identically zero. Note that the above formula needs not give the value of $(\partial f/\partial x_i)(x)$ for any $x \in U$, but almost all $x \in U$, because $(\partial P/\partial y)(x, f(x))$ may have a zero point in U .

The problem on differentiations reminds us of the problem on integrations. Note that the partial derivatives of Nash functions are Nash functions, while the integrals of Nash functions need not be Nash functions. This fact was one of the reasons to introduce the class of elementary functions in classical calculus. For related results, say, Liouville's theorem on integrals of elementary functions, etc., refer to [18] for instance. There the theory of differential fields plays a significant role likewise in the present paper (Proofs of Lemma 2.3 and Theorem 2.1).

Then *Leibniz complexity* $LC(f)$ of f is defined as the minimal number of usages of Leibniz rules to compute the total differential df algebraically. The Leibniz complexity $LC(f)$ of a Nash function f is a kind of computational complexity. Assume any algorithm to compute the differentials of Nash functions using C^∞ functions possibly. Then $LC(f)$ gives a lower bound of usage count of Leibniz rule in any such algorithm. Actually we will define three variants of Leibniz complexities \widetilde{LC} , LC and lc in Section 3. In particular, Nash functions are characterized by the finiteness of Leibniz complexity LC (Theorem 2.1).

We remark that our complexity is closely related to the addition chain complexity [11] and to other several known computational complexities [1], [12]. We also remark that our complexity of Nash functions is of different kind from the complexity for the description or encoding of a Nash function defined in [7].

In general it is a difficult problem to determine the exact value of the Leibniz complexity for a given Nash function. In Section 3, we provide general observations and estimates on Leibniz complexity of Nash functions using the binary expansions (Proposition 3.13) and discuss their relations with known notions of complexity of Nash functions

([16]).

In Section 4, we generalize Theorem 2.1 to Nash functions on an affine Nash manifold (Theorem 4.1), by using the global results on Nash functions ([6], [9], [8]).

The authors thank anonymous referees for their valuable comments and suggestions. In particular the relations of Leibniz complexity with the addition chain complexity ([11]) and the non-scalar complexity ([1], [12]), and moreover, the results, Lemma 3.5, Remark 3.6, Lemma 3.10 and Remark 3.11 are suggestions to the authors by one of the referees.

The authors dedicate this paper to the memory of Professor Masahiro Shiota, who passed away in January 2018.

2. Algebraic computability of differentials.

Let $\mathcal{C}^\infty(U)$ (resp. $\mathcal{C}^\omega(U)$, $\mathcal{N}^\omega(U)$) denote the set of all C^∞ functions (resp. analytic functions, Nash functions) on an open subset $U \subset \mathbf{R}^n$. The notation $\mathcal{N}^\omega(U)$ is used in [19].

Regarding $A = \mathcal{C}^\infty(U)$ (resp. $\mathcal{C}^\omega(U)$, $\mathcal{N}^\omega(U)$) as an \mathbf{R} -algebra, we take the space Ω_A of Kähler differentials of A and the universal derivation $\mathbf{d} : A \rightarrow \Omega_A$.

In fact, for any \mathbf{R} -algebra A , Ω_A can be constructed as follows: First consider the free A -module \mathfrak{F}_A generated by elements $\mathbf{d}f$, for any $f \in A$, regarded as just symbols.

Second consider the sub- A -module $\mathfrak{R}_A \subset \mathfrak{F}_A$ generated by the set R of all relations of algebraic derivations:

$$\mathbf{d}(h + k) - \mathbf{d}h - \mathbf{d}k, \mathbf{d}(\lambda\ell) - \lambda\mathbf{d}\ell, \mathbf{d}(1), \mathbf{d}(pq) - p\mathbf{d}q - q\mathbf{d}p,$$

$h, k, \ell, p, q \in A, \lambda \in \mathbf{R}$. Note that an element of \mathfrak{R}_A is a finite sum $\sum h_i r_i$ where $h_i \in A, r_i \in R$. Each $h_i r_i$ is called a *term* of the element. The first two kinds of generators of \mathfrak{R}_A in R correspond to the linearity, $\mathbf{d}(1)$ corresponds to the annihilation of $\mathbf{R} \subset A$, and the last kind of generators correspond to the Leibniz rule. We will count just the number of terms involving the last kind of generators. Here we add $\mathbf{d}(1)$, which is generated from $\mathbf{d}(1 \cdot 1) - 1\mathbf{d}(1) - 1\mathbf{d}(1)$, as a generator of \mathfrak{R}_A because we want to use the annihilation of $\mathbf{R} \subset A$ freely.

Third we set $\Omega_A = \mathfrak{F}_A/\mathfrak{R}_A$ and define $\mathbf{d} : A \rightarrow \Omega_A$ by mapping each $f \in A$ to the class of $\mathbf{d}f$ in $\mathfrak{F}_A/\mathfrak{R}_A$. Thus, if an element $\alpha \in \mathfrak{F}_A$ reduces to zero in Ω_A , then there exists an element $\sum h_i r_i \in \mathfrak{R}_A$, which is called an *expression* of α , such that $\alpha = \sum h_i r_i$ in \mathfrak{F}_A .

If B is any A -module and $D : A \rightarrow B$ is any derivation, i.e. D is an \mathbf{R} -linear map satisfying $D(gh) = gD(h) + hD(g)$ for any $g, h \in A$, then there exists a unique A -homomorphism $\rho : \Omega_A \rightarrow B$ such that $D = \rho \circ \mathbf{d}$.

Suppose U is connected.

Consider the set $S \subset \mathcal{N}^\omega(U)$ of non-zero Nash functions i.e. Nash functions which are not identically zero on U . Then S is closed under the multiplication. For $A = \mathcal{C}^\infty(U)$ (resp. $\mathcal{C}^\omega(U)$, $\mathcal{N}^\omega(U)$), let $\tilde{A} = \tilde{\mathcal{C}}^\infty(U)$ (resp. $\tilde{\mathcal{C}}^\omega(U)$, $\tilde{\mathcal{N}}^\omega(U)$) denote the localization A_S of A by S . Note that any element $k \in \tilde{A}$ is expressed as $k = (1/g)h$ for a $g \in \mathcal{N}^\omega(U)$, $g \neq 0$, and $h \in A$ and, in general, k needs not belong to A if g has a zero point in U . In particular $\tilde{\mathcal{N}}^\omega(U) = \mathcal{N}^\omega(U)_S$ is the quotient field $Q(\mathcal{N}^\omega(U))$.

Then we consider the space $\Omega_{\tilde{A}}$ of Kähler differentials of the \mathbf{R} -algebra \tilde{A} for $A = \mathcal{C}^\infty(U), \mathcal{C}^\omega(U), \mathcal{N}^\omega(U)$,

Then we have:

THEOREM 2.1. *Let U be a semi-algebraic connected open subset of \mathbf{R}^n . Let $A = \mathcal{C}^\infty(U)$ (resp. $\mathcal{C}^\omega(U), \mathcal{N}^\omega(U)$). Then the following 10 conditions on an analytic function $f \in \mathcal{C}^\omega(U)$ are equivalent to each other:*

- (1) f is a Nash function on U .
- (2)_A There exists a non-zero Nash function $g \in \mathcal{N}^\omega(U)$ such that

$$g \left(\mathbf{d}f - \sum_{i=1}^n \frac{\partial f}{\partial x_i} \mathbf{d}x_i \right) = 0,$$

in the space Ω_A of Kähler differentials of A .

- (3)_A $\mathbf{d}f = \sum_{i=1}^n (\partial f / \partial x_i) \mathbf{d}x_i$, in the space $\Omega_{\tilde{A}}$ of Kähler differentials of \tilde{A} .
- (4)_A There exist $f_1, \dots, f_n \in \tilde{A}$ such that $\mathbf{d}f = \sum_{i=1}^n f_i \mathbf{d}x_i$, in the space $\Omega_{\tilde{A}}$ of Kähler differentials of \tilde{A} .

We will show the implications

$$\begin{array}{ccccccc} (1) \Rightarrow (2)_{\mathcal{N}^\omega(U)} & \Rightarrow & (2)_{\mathcal{C}^\omega(U)} & \Rightarrow & (2)_{\mathcal{C}^\infty(U)} & & \\ & & \Downarrow & & \Downarrow & & \\ (3)_{\mathcal{N}^\omega(U)} & \Rightarrow & (3)_{\mathcal{C}^\omega(U)} & \Rightarrow & (3)_{\mathcal{C}^\infty(U)} & & \\ & & \Downarrow & & \Downarrow & & \\ (4)_{\mathcal{N}^\omega(U)} & \Rightarrow & (4)_{\mathcal{C}^\omega(U)} & \Rightarrow & (4)_{\mathcal{C}^\infty(U)} & \Rightarrow & (1) \end{array}$$

to have the equivalence of the 10 conditions.

To show Theorem 2.1, we first recall the following known basic result on Nash functions, which is formulated in more general setting than we are going to use.

LEMMA 2.2. *Let $U \subset \mathbf{R}^n$ be a semi-algebraic open subset and $f \in \mathcal{C}^\omega(U)$ be an analytic function on U . Then the following conditions are equivalent to each other:*

- (i) f is a Nash function on U , i.e. there exists a non-zero polynomial $P(x, y)$ such that $P(x, f(x)) = 0$ for any $x \in U$.
- (ii) The graph of f in $U \times \mathbf{R} \subset \mathbf{R}^{n+1}$ is a semi-algebraic set.
- (iii) For any $a \in U$, the Taylor series $j^\infty f(a)$ of f at a is algebraic in the formal power series algebra $\mathbf{R}[[x - a]]$ over the polynomial algebra $\mathbf{R}[x - a] = \mathbf{R}[x]$, in other words, there exists a non-zero polynomial $P(x, y)$ such that $j^\infty P(x, f)(a) = 0$.
- (iv) For any connected component U' of U , there exists a point $a \in U'$ such that the Taylor series $j^\infty f(a)$ of f at a is algebraic in the formal power series algebra $\mathbf{R}[[x - a]]$ over the polynomial algebra $\mathbf{R}[x - a]$.

PROOF. The equivalence between (i) and (ii) is well-known (see for instance [3]). The implications (i) \Rightarrow (iii) \Rightarrow (iv) are clear. To show the implication (iv) \Rightarrow (i), suppose

(iv). Note that the number of connected components of U is finite. Let U_1, \dots, U_r be the connected components of U . Let $1 \leq i \leq r$. Then there exists $a_i \in U_i$ such that f is expressed by the Taylor series at a_i in a neighborhood $W \subset U_i$ of a_i and there exists a non-zero polynomial $P_i(x, y)$ such that $P_i(x, f(x)) = 0$ for any $x \in W$. Since the function $P_i(x, f(x))$ is analytic on U_i and U_i is connected, $P_i(x, f(x)) = 0$ for any $x \in U_i$. Then it suffices to take $P = \prod_{i=1}^r P_i$ to get (i). \square

Also we need the general algebraic lemma to show the implication $(4)_{\mathcal{C}^\infty(U)} \Rightarrow (1)$ of Theorem 2.1.

LEMMA 2.3. *Let $K \subset L$ be a field extension. Assume that $\mathbf{R} \subset K$. Let $f \in L$ be a transcendental element over K . Then, for any derivation $D_0 : K \rightarrow K$ and for any $u \in L$, there exists a unique derivation $D_u : K(f) \rightarrow L$ satisfying*

$$D_u|_K = D_0, \quad D_u(f) = u.$$

Moreover if L is finitely generated over K , then the derivation D_u extends to a derivation $D : L \rightarrow L$.

PROOF. Since f is transcendental over K , we can define a derivation $D_u : K(f) \rightarrow L$ on the extension field $K(f)$ over K by f , by $D_u|_K = D_0$ and $D_u(f) = u$. Suppose L is finitely generated over K and $L = K(f, h_1, \dots, h_m)$ for some $h_1, \dots, h_m \in L$. Then we define a derivation $D_{u1} : K(f, h_1) \rightarrow L$, $D_{u1}|_{K(f)} = D_u$ as follows: If h_1 is transcendental over $K(f)$, then we set $D_{u1}(h_1) = 0$. If h_1 is algebraic over $K(f)$, then we set $D_{u1}(h_1)$ as the element in $K(f, h_1)$ which is determined by the algebraic relation of h_1 over $K(f)$ and D_u . In fact, if $\sum_{k=0}^m a_k h_1^{m-k} = 0$, $a_k \in K(f)$, is a minimal algebraic relation of h_1 over $K(f)$, then we would have

$$\sum_{k=0}^m D_u(a_k) h_1^{m-k} + \left(\sum_{k=0}^{m-1} (m-k) a_k h_1^{m-k-1} \right) D_{u1}(h_1) = 0.$$

Since $\sum_{k=0}^{m-1} (m-k) a_k h_1^{m-k-1} \neq 0$ by the minimality assumption, $D_{u1}(h_1)$ is uniquely determined by

$$D_{u1}(h_1) = - \left(\sum_{k=0}^m D_u(a_k) h_1^{m-k} \right) / \left(\sum_{k=0}^{m-1} (m-k) a_k h_1^{m-k-1} \right).$$

Thus we extend D_u into a derivation $D = D_{um} : L \rightarrow L$ by a finite number of steps. Note that we need not to use Zorn's lemma to show the existence of extension of derivation. \square

PROOF OF THEOREM 2.1. $(1) \Rightarrow (2)_{\mathcal{N}^\omega(U)}$: Let $f \in \mathcal{C}^\omega(U)$ be a Nash function and $P(x, y)$ be a non-zero polynomial satisfying $P(x, f) = 0$ and $(\partial P / \partial y)(x, f) \neq 0$. Then, by taking Kähler differential on both sides of the polynomial equality $P(x, f) = 0$, we have in $\Omega_{\mathcal{N}^\omega(U)}$,

$$\begin{aligned} 0 &= \mathbf{d}(P(x, f)) = \sum_{i=1}^n \frac{\partial P}{\partial x_i}(x, f) \mathbf{d}x_i + \frac{\partial P}{\partial y}(x, f) \mathbf{d}f \\ &= \sum_{i=1}^n \left(-\frac{\partial P}{\partial y}(x, f) \frac{\partial f}{\partial x_i} \right) \mathbf{d}x_i + \frac{\partial P}{\partial y}(x, f) \mathbf{d}f \\ &= \frac{\partial P}{\partial y}(x, f) \left(\mathbf{d}f - \sum_{i=1}^n \frac{\partial f}{\partial x_i} \mathbf{d}x_i \right), \end{aligned}$$

and that $(\partial P/\partial y)(x, f)$ is a non-zero Nash function on U .

Since $\mathcal{N}^\omega(U) \subset \mathcal{C}^\omega(U) \subset \mathcal{C}^\infty(U)$, the implications (j) $_{\mathcal{N}^\omega(U)} \Rightarrow$ (j) $_{\mathcal{C}^\omega(U)} \Rightarrow$ (j) $_{\mathcal{C}^\infty(U)}$ are clear, for $j = 2, 3, 4$.

(2) $_A \Rightarrow$ (3) $_A$, $A = \mathcal{C}^\infty(U), \mathcal{C}^\omega(U), \mathcal{N}^\omega(U)$: Since $1/g$ belongs to the localization \tilde{A} , we have that, if $g(\mathbf{d}f - \sum_{i=1}^n (\partial f/\partial x_i) \mathbf{d}x_i) = 0$ in Ω_A , then $\mathbf{d}f - \sum_{i=1}^n (\partial f/\partial x_i) \mathbf{d}x_i = 0$ in $\Omega_{\tilde{A}}$.

The implications (3) $_A \Rightarrow$ (4) $_A$, $A = \mathcal{C}^\infty(U), \mathcal{C}^\omega(U), \mathcal{N}^\omega(U)$, are clear.

(4) $_{\mathcal{C}^\infty(U)} \Rightarrow$ (1) : Suppose f is not a Nash function on U and $\mathbf{d}f - \sum_{i=1}^n f_i \mathbf{d}x_i = 0$ in $\Omega_{\tilde{\mathcal{C}}^\infty(U)}$. Since f is not a Nash function, by Lemma 2.2, there exists a point $a \in U$ such that $f \in \mathbf{R}[[x-a]] \subset Q(\mathbf{R}[[x-a]])$ is not algebraic. Here $\mathbf{R}[[x-a]] = \mathcal{C}_{\mathbf{R}^n, a}^\infty/\mathfrak{m}_{\mathbf{R}^n, a}^\infty$ is the \mathbf{R} -algebra of formal series, $M = Q(\mathbf{R}[[x-a]])$ is its quotient field and the Taylor series of f at a is written also by the same symbol f . Moreover, we have $\mathbf{d}f - \sum_{i=1}^n f_i \mathbf{d}x_i = 0$ in the Kähler differentials Ω_M of M , via the homomorphism $\tilde{\mathcal{C}}^\infty(U) \rightarrow M$ defined by taking the Taylor series. Then, in the free M -module \mathfrak{F}_M generated by elements $\{\mathbf{d}h \mid h \in M\}$, $\mathbf{d}f - \sum_{i=1}^n f_i \mathbf{d}x_i$ is a finite sum of elements of types

$$a(\mathbf{d}(h+k) - \mathbf{d}h - \mathbf{d}k), \quad b(\mathbf{d}(\lambda\ell) - \lambda\mathbf{d}\ell), \quad c(\mathbf{d}(pq) - p\mathbf{d}q - q\mathbf{d}p).$$

Here $a, h, k, b, \ell, c, p, q \in M, \lambda \in \mathbf{R}$. Now we take the subfield $L \subset M$ generated over the rational function field $K = \mathbf{R}(x)$ by $f, f_i (1 \leq i \leq n)$ and those $a, h, k, b, \ell, c, p, q$ which appear in the above expression of $\mathbf{d}f - \sum_{i=1}^n f_i \mathbf{d}x_i$: $L = K(f, h_1, \dots, h_m)$, which is a finitely generated field over K by f and for some $h_1, \dots, h_m \in M$. Then we have $\mathbf{d}f - \sum_{i=1}^n f_i \mathbf{d}x_i = 0$ also in Ω_L .

Take any non-zero element $u \in L$ and fix it. Set $D_0 = 0$. Then, by Lemma 2.3, we have a derivation $D : L \rightarrow L$ with $D(f) = u$. Then by the universality of the Kähler differentials, there exists an L -linear map $\rho : \Omega_L \rightarrow L$ such that $\rho \circ \mathbf{d} = D : L \rightarrow L$. Here $\mathbf{d} : L \rightarrow \Omega_L$ is the universal derivation. Then we have

$$0 = \rho \left(\mathbf{d}f - \sum_{i=1}^n f_i \mathbf{d}x_i \right) = D(f) = u.$$

This leads to a contradiction with the assumption $u \neq 0$. Thus we have that f is a Nash function. □

REMARK 2.4. If Zorn's lemma is used, then the fact that a transcendental basis of $M = Q(\mathbf{R}[[x-a]])$ forms a basis of Ω_M as an M -vector space (Theorem 26.5 of [13]) will give a shorter proof of the part (4) $_{\mathcal{C}^\infty(U)} \Rightarrow$ (1) of proof of Theorem 2.1. In fact if $f \in M$

is transcendental, then there exists a transcendental basis containing f, x_1, \dots, x_n and therefore we have that $\mathbf{d}f, \mathbf{d}x_1, \dots, \mathbf{d}x_n$ are linearly independent over M , which leads to a contradiction. (The remark is based on an anonymous reviewer’s comment informed to the authors.) The same remark is applied also to the proof of our Theorem 4.1.

REMARK 2.5. If U is not connected, then Theorem 2.1 does not hold. In fact, let $U = \mathbf{R} \setminus \{0\}$ and set $f(x) = e^x$ if $x > 0$ and $f(x) = 1$ if $x < 0$. Then $f \in \mathcal{C}^\omega(U)$ and $f \notin \mathcal{N}^\omega(U)$. However the condition (2) is satisfied if we take as g the non-zero Nash function on U defined by $g(x) = 0(x > 0), g(x) = 1(x < 0)$.

3. Estimates on Leibniz complexity.

Let $U \subset \mathbf{R}^n$ be a semi-algebraic connected open subset. Let $f \in \mathcal{N}^\omega(U)$ be a Nash function on U .

Then by the equivalence of (1) and $(2)_{\mathcal{N}^\omega(U)}$ in Theorem 2.1, there exists a non-zero Nash function $g \in \mathcal{N}^\omega(U)$ such that $g(\mathbf{d}f - \sum_{i=1}^n (\partial f / \partial x_i) \mathbf{d}x_i) \in \mathfrak{R}_{\mathcal{N}^\omega(U)} \subset \mathfrak{F}_{\mathcal{N}^\omega(U)}$. Then define $\text{LC}_g(f)$ as the minimal number of terms corresponding to Leibniz rule among all expressions of $g(\mathbf{d}f - \sum_{i=1}^n (\partial f / \partial x_i) \mathbf{d}x_i) \in \mathfrak{R}_{\mathcal{N}^\omega(U)}$. We define the *Leibniz complexity* $\text{LC}(f)$ of f by the minimum of $\text{LC}_g(f)$ for all such non-zero $g \in \mathcal{N}^\omega(U)$.

Note that we do not care about the number of terms corresponding to linearity of the differential. Moreover we do not count the term generated by the relation $\mathbf{d}(1 \cdot 1) - 1\mathbf{d}(1) - 1\mathbf{d}(1)$. Therefore we use the relation $\mathbf{d}(c) = 0$ for $c \in \mathbf{R}$ freely.

Similarly we define $\widetilde{\text{LC}}(f)$, related to Theorem 2.1 $(3)_{\mathcal{N}^\omega(U)}$, as the minimal number of terms corresponding to Leibniz rule among all expressions of $\mathbf{d}f - \sum_{i=1}^n (\partial f / \partial x_i) \mathbf{d}x_i \in \mathfrak{R}_{\mathcal{N}^\omega(U)}$.

Moreover, if $\mathbf{d}f - \sum_{i=1}^n (\partial f / \partial x_i) \mathbf{d}x_i \in \mathfrak{R}_{\mathcal{N}^\omega(U)}$, then we define $\text{lc}(f) = \text{LC}_1(f)$, simply as the minimal number of terms corresponding to Leibniz rule among all its expressions in $\mathfrak{R}_{\mathcal{N}^\omega(U)}$. Note that, if f is a polynomial function, then $\text{lc}(f) < \infty$. However in general $\mathbf{d}f - \sum_{i=1}^n (\partial f / \partial x_i) \mathbf{d}x_i$ may not belong to $\mathfrak{R}_{\mathcal{N}^\omega(U)}$. Then we set $\text{lc}(f) = \infty$.

Hereafter, for $f \in \mathcal{N}^\omega(U)$, we set

$$\zeta(f) := \mathbf{d}f - \sum_{i=1}^n \frac{\partial f}{\partial x_i} \mathbf{d}x_i,$$

regarded as an element in \mathfrak{F}_A for $A = \mathcal{N}^\omega(U)$ or for its localization $A = \widetilde{\mathcal{N}}^\omega(U) = \mathcal{N}^\omega(U)_S$ where $S = \mathcal{N}^\omega(U) \setminus \{0\}$.

First we show general basic inequalities:

LEMMA 3.1. For any $f \in \mathcal{N}^\omega(U)$, we have $\widetilde{\text{LC}}(f) \leq \text{LC}(f) \leq \text{lc}(f)$.

PROOF. Suppose $\text{lc}(f) < \infty$ and there exists an expression of $\zeta(f)$ in $\mathfrak{R}_{\mathcal{N}^\omega(U)}$ such that the number of terms involving Leibniz rule is equal to $\text{lc}(f)$. Then setting $g = 1$, $g\zeta(f)$ has the same expression in $\mathfrak{R}_{\mathcal{N}^\omega(U)}$, and therefore we have $\text{LC}(f) \leq \text{lc}(f)$. Next, by the definition of $\text{LC}(f)$, there exist a $g \in \mathcal{N}^\omega(U) \setminus \{0\}$ and an expression of $g\zeta(f)$ in $\mathfrak{R}_{\mathcal{N}^\omega(U)}$ such that the number of terms involving Leibniz rule is equal to $\text{LC}(f)$. Then,

dividing by g , we have an expression of $\zeta(f)$ in $\mathfrak{R}_{\widetilde{\mathcal{N}}^\omega(U)}$ such that the number of terms involving Leibniz rule is equal to $\text{LC}(f)$. Therefore, by the definition of $\widetilde{\text{LC}}(f)$, we have $\widetilde{\text{LC}}(f) \leq \text{LC}(f)$. \square

LEMMA 3.2. For $f, g \in \mathcal{N}^\omega(U)$, we have

$$(1) \text{LC}(f + g) \leq \text{LC}(f) + \text{LC}(g). \quad (2) \text{LC}(fg) \leq \text{LC}(f) + \text{LC}(g) + 1.$$

The same inequalities hold for $\widetilde{\text{LC}}$ and lc .

PROOF. Let $h\zeta(f) \in \mathfrak{R}_{\mathcal{N}^\omega(U)}$ (resp. $k\zeta(g) \in \mathfrak{R}_{\mathcal{N}^\omega(U)}$) be expressed using the terms of Leibniz rule minimally i.e. $\text{LC}(f)$ -times (resp. $\text{LC}(g)$ -times), for a non-zero $h \in \mathcal{N}^\omega(U)$ (resp. a non-zero $k \in \mathcal{N}^\omega(U)$). Then $hk\zeta(f + g) = k(h\zeta(f)) + h(k\zeta(g)) \in \mathfrak{R}_{\mathcal{N}^\omega(U)}$ is expressed using Leibniz rule at most $\text{LC}(f) + \text{LC}(g)$ times. Therefore we have (1). Moreover, by using Leibniz rule once, we have

$$hk\mathbf{d}(fg) = hk(\mathbf{g}\mathbf{d}f + f\mathbf{d}g) = kg(\mathbf{h}\mathbf{d}f) + hf(\mathbf{k}\mathbf{d}g)$$

in $\Omega_{\mathcal{N}^\omega(U)}$. Then, using Leibniz rule $\text{LC}(f) + \text{LC}(g)$ times, we compute $\mathbf{h}\mathbf{d}f$ and $\mathbf{k}\mathbf{d}g$, and thus $hk\mathbf{d}(fg)$. Therefore we have (2).

For $\widetilde{\text{LC}}$ and lc , the inequalities are proved similarly or more easily. \square

By the definition of Leibniz complexity, we have the affine invariance:

LEMMA 3.3. Let $f \in \mathcal{N}^\omega(U)$ and $\varphi : \mathbf{R}^n \rightarrow \mathbf{R}^n$ be an affine isomorphism. Then $f \circ \varphi \in \mathcal{N}^\omega(\varphi^{-1}(U))$ satisfies $\text{LC}(f \circ \varphi) = \text{LC}(f)$, $\widetilde{\text{LC}}(f \circ \varphi) = \widetilde{\text{LC}}(f)$ and $\text{lc}(f \circ \varphi) = \text{lc}(f)$.

PROOF. By the definition of Leibniz complexity $h(\mathbf{d}f - \sum_{i=1}^n (\partial f / \partial x_i) \mathbf{d}(x_i))$ is zero in $\Omega_{\mathcal{N}^\omega(U)}$ by using Leibniz rule $\text{LC}(f)$ -times, for a non-zero $h \in \mathcal{N}^\omega(U)$. Let $x' = (x'_1, \dots, x'_n)$ be new affine coordinate system on \mathbf{R}^n defined by $x' = \varphi^{-1}(x)$. Then $(h \circ \varphi)(\mathbf{d}(f \circ \varphi) - \sum_{i=1}^n (\partial f / \partial x_i) \circ \varphi \mathbf{d}(\varphi_i))$ is zero in $\Omega_{\mathcal{N}^\omega(U)}$ by using Leibniz rule $\text{LC}(f)$ -times. Since we do not count the usage of Leibniz rule for $\mathbf{d}(c) = 0, c \in \mathbf{R}$, we have that $(h \circ \varphi)(\mathbf{d}(f \circ \varphi) - \sum_{i=1}^n (\partial(f \circ \varphi) / \partial x'_i) \mathbf{d}(x'_i))$ is zero in $\Omega_{\mathcal{N}^\omega(U)}$ by using Leibniz rule the same $\text{LC}(f)$ -times. Note that $h \circ \varphi \in \mathcal{N}^\omega(U)$ is non-zero. Therefore we have $\text{LC}(f \circ \varphi) \leq \text{LC}(f)$. Similarly, we have $\text{LC}(f) = \text{LC}((f \circ \varphi) \circ \varphi^{-1}) \leq \text{LC}(f \circ \varphi)$. Thus we have the required equality. The equality for $\widetilde{\text{LC}}$ (resp. $\text{lc}(f)$) is proved similarly or more easily. \square

In general it is a difficult problem to determine the exact value of the Leibniz complexity even for an polynomial function.

EXAMPLE 3.4. Let $n = 1$ and write $x = x_1$. Then we have $\widetilde{\text{LC}}(x + c) = \text{LC}(x + c) = \text{lc}(x + c) = 0$. $\widetilde{\text{LC}}(x^2 + bx + c) = \text{LC}(x^2 + bx + c) = \text{lc}(x^2 + bx + c) = 1$. $\widetilde{\text{LC}}(\sqrt{x^2 + 1}) = \text{LC}(\sqrt{x^2 + 1}) = \text{lc}(\sqrt{x^2 + 1}) = 2$.

Let $n = 2$. For $\lambda \in \mathbf{R}$, we have

$$\text{LC}(x_1^2 + x_2^2 + \lambda x_1 x_2) = \begin{cases} 1 & \text{if } |\lambda| \geq 2 \\ 2 & \text{if } |\lambda| < 2. \end{cases}$$

In fact, $x_1^2 + x_2^2 + \lambda x_1 x_2 = (x_1 + (\lambda/2)x_2)^2 + (1 - \lambda^2/4)x_2^2$. Moreover $x_1^2 + x_2^2 + \lambda x_1 x_2 = (\alpha x_1 + \beta x_2)(x_1 + \beta x_2)$ for some $\alpha, \beta \in \mathbf{R}$ if and only if $|\lambda| \geq 2$. The same results hold for $\widetilde{\text{LC}}$ and lc .

Let $n = 1$ and write $x = x_1$. We consider Leibniz complexity of a monomial x^k . For example, $\text{lc}(x^0) = \text{lc}(1) = 0, \text{lc}(x) = 0, \text{lc}(x^2) = 1, \text{lc}(x^3) = 2, \text{lc}(x^4) = 2$. Also for LC and $\widetilde{\text{LC}}$ we have the same results. For example we calculate $\mathbf{d}(x^4) = 2x^2\mathbf{d}(x^2) = 4x^3\mathbf{d}(x)$ by using Leibniz rule twice, and we can check that it is impossible to calculate $\mathbf{d}(x^4)$ by using Leibniz rule just once.

To observe the essence of the problem to estimate the Leibniz complexity, let us digress to consider “the problem of strips”. Let k be a positive integer. Suppose we have a sheet of paper having width k and, using a pair of scissors, we make k -strips of width 1. We may cut several sheets of the same width at once by piling them. Then the problem is to minimize the total number of cuts. Clearly it is at most $k - 1$.

The exact answer to the above problem is given by the *addition chain complexity* $\ell(k)$ (see [11]). An *addition chain* of k is a sequence of integers

$$1 = a_0, a_1, a_2, \dots, a_r = k$$

satisfying that, for any $i = 1, 2, \dots, r$, there exist j and m with $0 \leq j \leq m < i$, such that $a_i = a_j + a_m$. Then $\ell(k)$ is defined as the minimum of the length r for all addition chain of k .

A process of making k -strips as above corresponds to an addition chain bijectively. Therefore the minimum of the total number of cuts is given by $\ell(k)$.

LEMMA 3.5. *For a positive integer k , we have*

$$\widetilde{\text{LC}}(x^k) \leq \text{LC}(x^k) \leq \text{lc}(x^k) \leq \ell(k).$$

PROOF. Let $1 = a_0, a_1, a_2, \dots, a_r = k$ be an addition chain of k . Since $k = a_r = a_j + a_m$ for some $0 \leq j \leq m < k$, we have one relation

$$\mathbf{d}(x^k) - x^{a_j}\mathbf{d}(x^{a_m}) - x^{a_m}\mathbf{d}(x^{a_j})$$

in $\mathfrak{A}_{\mathcal{N}^\omega(U)}$. Thus we have

$$\mathbf{d}(x^k) = x^{a_j}\mathbf{d}(x^{a_m}) + x^{a_m}\mathbf{d}(x^{a_j})$$

in $\Omega_{\mathcal{N}^\omega(U)}$ using Leibniz rule once. If $j < m$, we apply this procedure to $\mathbf{d}(x^{a_m})$. Then, using a relation

$$\mathbf{d}(x^k) - x^{a_j}\mathbf{d}(x^{a_m}) - x^{a_m}\mathbf{d}(x^{a_j}) + x^{a_j}(\mathbf{d}(x^{a_m}) - x^{a_{j'}}\mathbf{d}(x^{a_{m'}}) - x^{a_{m'}}\mathbf{d}(x^{a_{j'}}))$$

with two terms, in the sense of Section 2, in $\mathfrak{A}_{\mathcal{N}^\omega(U)}$ for some $0 \leq j' \leq m' < m$, we have

$$\mathbf{d}(x^k) = x^{a_j+a_{j'}}\mathbf{d}(x^{a_{m'}}) + x^{a_j+a_{m'}}\mathbf{d}(x^{a_{j'}})$$

in $\Omega_{\mathcal{N}^\omega(U)}$ using Leibniz rule twice. If $j = m$, then $x^{aj} \mathbf{d}(x^{am}) + x^{am} \mathbf{d}(x^{aj}) = 2x^{am} \mathbf{d}(x^{am})$, and then a similar procedure is applied to $\mathbf{d}(x^{am})$. Thus we see that, by using a relation with s -terms involving Leibniz rule, $\mathbf{d}(x^k)$ is reduced to a functional linear combination of $\mathbf{d}(x^{a_0}), \mathbf{d}(x^{a_1}), \dots, \mathbf{d}(x^{a_{r-s}})$ in $\Omega_{\mathcal{N}^\omega(U)}$, $s = 1, 2, \dots, r$. Therefore we have $\text{lc}(x^k)(= \text{LC}_1(x^k)) \leq r$, for any addition chain of k . Hence we have $\text{lc}(x^k) \leq \ell(k)$. Other inequalities follow from Lemma 3.1. \square

REMARK 3.6. We can define, naturally, a kind of Leibniz complexity lc_{poly} by using the Kähler differential Ω_A of polynomial algebra $A = \mathbf{R}[x_1, \dots, x_n]$. Then the proof of Lemma 3.5 gives also the inequalities $\text{lc}(x^k) \leq \text{lc}_{\text{poly}}(x^k) \leq \ell(k)$. The authors conjecture, at least, the equality $\text{lc}_{\text{poly}}(x^k) = \ell(k)$, but they have no proof of that.

Now we show one known strategy to obtain an explicit estimate. Consider the binary expansion of k :

$$k = 2^{\mu_r} + 2^{\mu_{r-1}} + \dots + 2^{\mu_1},$$

for some integers $\mu_r > \mu_{r-1} > \dots > \mu_1 \geq 0$. We set $\mu = \mu_r$. Then *the number of digits* ('1' or '0') is given by $\mu + 1$, while r is *the number of units*, '1', appearing in the binary expansion. Then first we cut the sheet into r sheets of width $2^\mu, 2^{\mu_{r-1}}, \dots, 2^{\mu_1}$ by $(r - 1)$ -cuts. Second, divide the sheet of width 2^μ into sheets of width $2^{\mu_{r-1}}$ by $\mu - \mu_{r-1}$ -cuts. Third, divide the piled sheets of width $2^{\mu_{r-1}}$ into sheets of width $2^{\mu_{r-2}}$ by $\mu_{r-1} - \mu_{r-2}$ -cuts, and so on. Iterating the process, we have sheets of width 2^{μ_1} , which we divide into strips of width 1 by μ_1 -cuts finally. The total number of cuts by this method is given by $\mu + r - 1$.

Thus we have by Lemma 3.5:

COROLLARY 3.7. *For a positive integer k , we have*

$$\widetilde{\text{LC}}(x^k) \leq \text{LC}(x^k) \leq \text{lc}(x^k) \leq \ell(k) \leq \mu + r - 1.$$

REMARK 3.8. The estimate in Corollary 3.7 is, by no means, best possible. For example, let $k = 31$. Then $31 = 2^4 + 2^3 + 2^2 + 2^1 + 2^0$. Therefore $r = 5$ and $\mu = 4$. Therefore $\mu + r - 1 = 8$. Moreover we have the addition chain complexity $\ell(31) = 7$. However $\text{LC}(x^{31}) \leq 6$. In fact, since $32 = 2^5$, we have by Lemma 3.7,

$$x\mathbf{d}(x^{31}) = \mathbf{d}(x^{32}) - x^{31}\mathbf{d}(x) = 32x^{31}\mathbf{d}(x) - x^{31}\mathbf{d}(x) = 31x^{31}\mathbf{d}(x),$$

by using Leibniz rule 6 times. Then we have $\mathbf{d}(x^{31}) = 31x^{30}\mathbf{d}(x)$ in $\Omega_{\tilde{\mathcal{C}}^\infty(U)}$.

Related to Corollary 3.7, we observe

LEMMA 3.9. *For $f \in \mathcal{N}^\omega(U)$ and a natural number $k \geq 1$, we have $\text{LC}(f^k) \leq \text{LC}(f) + \text{LC}(x^k)$.*

PROOF. If f is a constant function, then $\text{LC}(f^k) = 0$, so the inequality holds trivially. We suppose f is not a constant function. By definition, for some non-zero $g \in$

$\mathcal{N}^\omega(\mathbf{R})$, $g\mathbf{d}(x^k)$ is deformed into $g kx^{k-1}\mathbf{d}x$ in $\Omega_{\mathcal{N}^\omega}(\mathbf{R})$ using Leibniz rules $\text{LC}(x^k)$ -times. Using the same procedure, $(g \circ f)\mathbf{d}(f^k)$ is deformed into $(g \circ f)kf^{k-1}\mathbf{d}f$ in $\Omega_{\mathcal{N}^\omega(U)}$ using Leibniz rules $\text{LC}(x^k)$ -times. Note that $g \circ f$ is non-zero in $\mathcal{N}^\omega(U)$. Moreover, using Leibniz rules $\text{LC}(f)$ times, $h(g \circ f)kf^{k-1}\mathbf{d}f$ is deformed into $h(g \circ f)\sum_{i=1}^n kf^{k-1}(\partial f/\partial x_i)\mathbf{d}x_i$ for some non-zero $h \in \mathcal{N}^\omega(U)$. Since $g \circ f$ is non-zero, $h(g \circ f)$ is non-zero. \square

In general we have

LEMMA 3.10. *Let $g_1, \dots, g_m \in \mathcal{N}^\omega(U)$ and $P(y_1, \dots, y_m) \in \mathbf{R}[y_1, \dots, y_m]$ be a polynomial regarded as a function on \mathbf{R}^m . Then, for the Leibniz complexity of $f = P(g_1, \dots, g_m)$, we have*

$$\text{LC}(f) \leq \text{lc}(P) + \sum_{i=1}^m \text{LC}(g_i), \quad \widetilde{\text{LC}}(f) \leq \text{lc}(P) + \sum_{i=1}^m \widetilde{\text{LC}}(g_i), \quad \text{lc}(f) \leq \text{lc}(P) + \sum_{i=1}^m \text{lc}(g_i).$$

PROOF. We give a proof of the first inequality only. The remaining inequalities are proved similarly or more easily.

Using Leibniz rule $\text{lc}(P)$ times, we have

$$\mathbf{d}(f) = \sum \frac{\partial P}{\partial y_i}(g_1, \dots, g_m) \mathbf{d}(g_i),$$

in $\Omega_{\mathcal{N}^\omega(U)}$. For each $i = 1, \dots, m$, there exists non-zero Nash function h_i such that

$$h_i \mathbf{d}(g_i) = h_i \sum_{j=1}^n \frac{\partial g_i}{\partial x_j} \mathbf{d}(x_j)$$

by an $\text{LC}(g_i)$ times usage of Leibniz rule. Therefore, we have

$$h_1 \cdots h_m \mathbf{d}(f) = h_1 \cdots h_m \left(\sum_{j=1}^n \frac{\partial f}{\partial x_j} \mathbf{d}(x_j) \right),$$

in $\Omega_{\mathcal{N}^\omega(U)}$, using Leibniz rule $\text{lc}(P) + \sum_{i=1}^m \text{LC}(g_i)$ times in total. Therefore, we have $\text{LC}(f) \leq \text{lc}(P) + \sum_{i=1}^m \text{LC}(g_i)$. \square

REMARK 3.11. The Leibniz complexity $\text{lc}(P)$ or $\text{lc}_{\text{poly}}(P)$ (see Remark 3.6) for polynomials P is closely related to the *non-scalar complexity* of P ([12], [1]). The non-scalar complexity of a polynomial P is defined roughly as follows. Consider any program to produce polynomials in $\mathbf{R}[x_1, \dots, x_n]$ by scalar multiplications, additions and products, without divisions, starting from the 0-th stage 1, x_1, \dots, x_n (depth 0), and making some pair of linear combinations of polynomials appeared in previous stages of depth $\leq r$ and, as the next stage, making the product of them (depth $r + 1$) and so on. Then the non-scalar complexity $L_{\text{ns}}(P)$ is defined as the minimal depth of the polynomial P in all such programs producing P . Then we have

$$\text{lc}(P) \leq \text{lc}_{\text{poly}}(P) \leq L_{\text{ns}}(P).$$

The proof is similar to that of Lemma 3.5. The authors conjecture also that the equality $lc_{\text{poly}}(P) = L_{\text{ns}}(P)$ holds, but they have no proof of the equality.

In [1], the non-scalar complexity of rational functions for programs allowing divisions is considered. Moreover, for any rational function f , it is given an estimate of the non-scalar complexity of partial derivatives $\partial f / \partial x_i$ by means of that of f . It is interesting to estimate the Leibniz complexity of partial derivatives of higher order by using Baur-Strassen's result [1].

As above, we consider “the problem of strips” starting from several number of sheets, say, s , having width k_s, k_{s-1} , and k_1 respectively. Then we have

LEMMA 3.12. *Let $P = P(x) = a_s x^{k_s} + a_{s-1} x^{k_{s-1}} + \dots + a_1 x^{k_1} \in \mathbf{R}[x]$ be a polynomial function of one variable, where $a_j \neq 0$ ($1 \leq j \leq s$) and $k_s > k_{s-1} > \dots > k_1 \geq 0$. Regarding the binary expansion, let μ be (the number of digits of k_s) -1 , and r_j the number of units of k_j , $1 \leq j \leq s$. Then, by using Leibniz rule $(\mu + \sum_{j=1}^s (r_j - 1))$ -times together with linearity, and by supposing $\mathbf{d}(c) = 0, c \in \mathbf{R}$, we have $\mathbf{d}(P) = (dP(x)/dx)\mathbf{d}(x)$ in $\Omega_{\mathcal{N}^\omega(U)}$. In particular we have*

$$\widetilde{\text{LC}}(P) \leq \text{LC}(P) \leq \text{lc}(P) \leq \mu + \sum_{j=1}^s (r_j - 1).$$

PROOF. Let $\mu = \mu_t > \mu_{t-1} > \dots > \mu_1 \geq 0$ be all of the exponents appearing in the binary expansions of k_s, k_{s-1}, \dots, k_1 . First, by using Leibniz rule $\sum_{j=1}^s (r_j - 1)$ -times, we modify $\mathbf{d}(P)$ into a linear combination of $\mathbf{d}(x^\ell), \ell = 2^\mu = 2^{\mu_t}, 2^{\mu_{t-1}}, \dots, 2^{\mu_1}$. Second, by using Leibniz rule $\mu - \mu_{t-1}$ -times, we modify $\mathbf{d}(x^\ell), \ell = 2^\mu$ into $\mathbf{d}(x^{\ell'}), \ell' = 2^{\mu_{t-1}}$. Repeating the procedure, we modify $\mathbf{d}(P)$ into a multiple of $\mathbf{d}(x^\ell), \ell = 2^{\mu_1}$. Finally, by using Leibniz rule μ_1 -times, we modify $\mathbf{d}(P)$ into a multiple of $\mathbf{d}(x)$. \square

We estimate the Leibniz complexity for a polynomial of n -variables. Let $P(x) = P(x_1, \dots, x_n) \in \mathbf{R}[x_1, \dots, x_n]$. We set $P(x) = \sum b_\alpha x^\alpha, b_\alpha \in \mathbf{R}$, by using multi-index $\alpha = (\alpha_1, \dots, \alpha_n)$ of non-negative integers. It is trivial that $\text{lc}(P)$ is at most the total number of multiplications of variables:

$$\sum_{b_\alpha \neq 0} \max\{|\alpha| - 1, 0\}.$$

Instead we consider the number

$$\sigma(P) := \sum_{b_\alpha \neq 0} \max\{\#\{i \mid 1 \leq i \leq n, \alpha_i > 0\} - 1, 0\},$$

which is needed just to separate the variables on differentiation, and we try to save the additional usage of Leibniz rule.

Suppose that, by arranging terms with respect to x_i for each $i, 1 \leq i \leq n$,

$$P(x) = a_{i,s(i)} x_i^{k_{i,s(i)}} + a_{i,s(i)-1} x_i^{k_{i,s(i)}-1} + \dots + a_{i,1} x_i^{k_{i,1}},$$

where $a_{i,j}$ is a non-zero polynomial of x_1, \dots, x_n without x_i , ($1 \leq j \leq s(i)$), and $k_{i,s(i)} > k_{i,s(i)-1} > \dots > k_{i,1} \geq 0$. The maximal exponent $k_{i,s(i)}$ is written as $\text{deg}_{x_i} P$, the degree of P in the variable x_i . For the binary expansion of $\text{deg}_{x_i} P$, let μ_i denote (the number of digits of $\text{deg}_{x_i} P$) -1 . Moreover let $r_{ij}, 1 \leq j \leq s(i)$ denote the number of units of the exponent k_{ij} for the binary expansion. Then we have

LEMMA 3.13. *By using the linearity, $\mathbf{d}(c) = 0, c \in \mathbf{R}$, and Leibniz rule $(\sigma(P) + \sum_{i=1}^n (\mu_i + \sum_{j=1}^{s(i)} (r_{ij} - 1)))$ -times together with linearity, we have $\mathbf{d}(P) = \sum_{i=1}^n (\partial P(x)/\partial x_i) \mathbf{d}(x_i)$ in $\Omega_{\mathcal{N}^\omega(U)}$. In particular we have the estimate*

$$\text{lc}(P) \leq \sigma(P) + \sum_{i=1}^n \left(\mu_i + \sum_{j=1}^{s(i)} (r_{ij} - 1) \right).$$

REMARK 3.14. We have, for any polynomial $P(x) = \sum b_\alpha x^\alpha$,

$$\sigma(P) + \sum_{i=1}^n \left(\mu_i + \sum_{j=1}^{s(i)} (r_{ij} - 1) \right) \leq \sum_{b_\alpha \neq 0} \max\{|\alpha| - 1, 0\}.$$

and in almost cases the inequality is strict.

PROOF OF LEMMA 3.13. By applying Leibniz rule to each term of P , we can express $\mathbf{d}(P)$ as a sum of the products of the forms $a_{i,j} \mathbf{d}(x_i^{k_{i,j}})$, in which we have the differential of one variable x_i and a function $a_{i,j}$ of the other variables. For this process we need to use Leibniz rule $\sigma(P)$ -times. Then $\mathbf{d}(P)$ is the sum of the form

$$a_{i,s(i)} \mathbf{d}(x_i^{k_{i,s(i)}}) + a_{i,s(i)-1} \mathbf{d}(x_i^{k_{i,s(i)-1}}) + \dots + a_{i,1} \mathbf{d}(x_i^{k_{i,1}}),$$

($i = 1, \dots, n$). By Lemma 3.12, for each $i = 1, \dots, n$, the form is deformed into $\sum_{i=1}^n (\partial P/\partial x_i) \mathbf{d}x_i$ by using Leibniz rule $(\mu_i + \sum_{j=1}^{s(i)} (r_{ij} - 1))$ -times. Thus we have the estimate. □

Now we give an upper estimate of Leibniz complexities for Nash functions by those for polynomial functions in terms of its polynomial relation. Let $f \in \mathcal{N}^\omega(U)$ be a Nash function on a connected open subset U of \mathbf{R}^n . Let $P(x, y) = P(x_1, \dots, x_n, y)$ be a polynomial such that $P(x, f(x)) = 0$ on U and $(\partial P/\partial y)(x, f(x))$ is not identically zero. We set $x_0 = y$. Suppose that, by arranging with respect to x_i for each $i, 0 \leq i \leq n$, we have

$$P(x, y) = a_{i,s(i)} x_i^{k_{i,s(i)}} + a_{i,s(i)-1} x_i^{k_{i,s(i)-1}} + \dots + a_{i,1} x_i^{k_{i,1}},$$

where $a_{i,j}$ is a non-zero polynomial of x_0, x_1, \dots, x_n without x_i ($1 \leq j \leq s(i)$), and $k_{i,s(i)} > k_{i,s(i)-1} > \dots > k_{i,1} \geq 0$. For the binary expansion, let μ_i (resp. $r_{ij}, 1 \leq j \leq s(i)$) be (the number of digits of $\text{deg}_{x_i} P$) -1 (resp. the number of units of k_{ij}), $0 \leq i \leq n$, respectively. Write $\text{deg}_{x_i} P$ the degree of P with respect to $x_i, 0 \leq i \leq n$ and use the same notation $\sigma(P)$ as in Lemma 3.13 for the polynomial P of $n + 1$ variables.

PROPOSITION 3.15. *Under the above notations, we have the estimate*

$$LC(f) \leq \sigma(P) + \sum_{i=0}^n \left(\mu_i + \sum_{j=1}^{s(i)} (r_{ij} - 1) \right).$$

In particular we have

$$LC(f) \leq \sigma(P) + \sum_{i=0}^n \{(\deg_{x_i} P + 2)(\log_2(\deg_{x_i} P) - 1)\} + n + 1.$$

EXAMPLE 3.16. Let $n = 1, f = \sqrt{x^2 + 1}$ and $P(x, y) = y^2 - x^2 - 1$. Then $\sigma(P) = 0, \mu_0 = \mu_1 = 1$ and $r_{ij} = 1$. Therefore the first inequality gives us that $LC(f) \leq 2$ as is seen in Introduction.

PROOF OF PROPOSITION 3.15. We write the right hand side by ψ of the first inequality. By Lemma 3.13, we have, by using Leibniz rule ψ -times,

$$d(P(x, y)) = \sum_{i=1}^n \frac{\partial P}{\partial x_i}(x, y) dx_i + \frac{\partial P}{\partial y}(x, y) dy,$$

modulo several linearity relations and $dc, c \in \mathbf{R}$ in $\Omega_{\mathcal{N}^\omega(U \times \mathbf{R})}$. Then, substituting y by f , we have that

$$0 = d(P(x, f)) = \sum_{i=1}^n \frac{\partial P}{\partial x_i}(x, f) dx_i + \frac{\partial P}{\partial y}(x, f) df,$$

in $\Omega_{\mathcal{N}^\omega(U)}$, therefore that

$$\frac{\partial P}{\partial y}(x, f) \left(df - \sum_{i=1}^n \frac{\partial f}{\partial x_i} dx_i \right) = 0,$$

in $\Omega_{\mathcal{N}^\omega(U)}$, by using Leibniz rule at most ψ -times. Thus we have the first inequality. The second equality is obtained from the first equality combined with the inequalities derived by the definitions:

$$2^{\mu_i} \leq \deg_{x_i} P < 2^{\mu_i+1}, \quad s(i) \leq \deg_{x_i} P + 1, \quad \text{and} \quad r_{ij} \leq \mu_i,$$

$$(1 \leq j \leq s(i), 0 \leq i \leq n). \quad \square$$

In [16], the complexity $C(f)$ of a Nash function f is defined as the minimum the total degree $\deg P$ of non-zero polynomials $P(x, y)$ with $P(x, f) = 0$. Moreover we define

$$S(f) := \min\{\sigma(P \circ \psi) \mid P(x, f) = 0, \deg P = C(f), \psi \text{ is an affine isomorphism on } \mathbf{R}^{n+1}\},$$

i.e. the minimum of the number σ for any defining polynomial P of f with minimal total degree under any choice of affine coordinates. We can regard $S(f)$ a complexity for the separation of variables in differentiation of f . Then we have the following result:

COROLLARY 3.17. *Let $f \in \mathcal{N}^\omega(U)$ be a Nash function on a connected open set $U \subset \mathbf{R}^n$. Then we have an estimate on the Leibniz complexity $\text{LC}(f)$ by the Ramanakora's complexity $\text{C}(f)$ and another complexity $\text{S}(f)$,*

$$\text{LC}(f) \leq \text{S}(f) + (n + 1)(\text{C}(f) + 2)(\log_2 \text{C}(f) - 1) + n + 1.$$

PROOF. Since $\deg_{x_i} P \leq \text{C}(f)$ ($0 \leq i \leq n$) we have the above estimate by Proposition 3.15 and Lemma 3.3. □

Naturally we would like to pose a problem to obtain any lower estimate of Leibniz complexity.

4. Algebraic differentiation on Nash manifolds.

Let U be a connected semi-algebraic open subset of \mathbf{R}^n and $M \subset U$ a Nash submanifold ([3], [19]). Suppose M is a closed connected subset in U . We consider the quotient \mathbf{R} -algebra $\mathcal{N}^\omega(U)/I$ by the ideal I of $\mathcal{N}^\omega(U)$ consisting of Nash functions on U which vanish on M .

Since $\mathcal{N}^\omega(U)$ is Noetherian ([17], [14]), I is generated by a finite number of Nash functions $g_1, \dots, g_\ell \in \mathcal{N}^\omega(U)$ over $\mathcal{N}^\omega(U)$.

An element $[f] \in \mathcal{C}^\omega(U)/IC^\omega(U)$ is called *Nash* if there exists a polynomial $P(x, y) = a_m(x)y^m + a_{m-1}(x)y^{m-1} + \dots + a_1(x)y + a_0(x) \in \mathbf{R}[x, y]$ satisfying that at least one of $a_m([x]), a_{m-1}([x]), \dots, a_1([x]), a_0([x])$ is not zero in $\mathcal{N}^\omega(U)/I$ and that $P([x], [f]) = 0$ in $\mathcal{C}^\omega(U)/IC^\omega(U)$. The condition is equivalent to that $[f]$ is algebraic over $\mathbf{R}(x)$ via the composition $\mathbf{R}(x) \hookrightarrow \mathcal{C}^\omega(U) \rightarrow \mathcal{C}^\omega(U)/IC^\omega(U)$ of natural homomorphisms. Also the condition is equivalent to that $[f]$ is algebraic over $\mathcal{N}^\omega(U)/I$ via the natural homomorphism $\mathcal{N}^\omega(U)/I \rightarrow \mathcal{C}^\omega(U)/IC^\omega(U)$. Then there exist a non-zero polynomial $P(x, y)$ and $h_j \in \mathcal{C}^\omega(U), 1 \leq j \leq \ell$ such that

$$P(x, f(x)) = \sum_{j=1}^{\ell} h_j(x)g_j(x),$$

for any $x \in U$ and that $(\partial P/\partial y)(x, f) \notin IC^\omega(U)$. By differentiating both sides of the relation by x_i , we have that

$$\frac{\partial P}{\partial x_i}(x, f(x)) + \frac{\partial P}{\partial y}(x, f(x)) \frac{\partial f}{\partial x_i} = \sum_{j=1}^{\ell} g_j(x) \frac{\partial h_j}{\partial x_i}(x) + \sum_{j=1}^{\ell} h_j(x) \frac{\partial g_j}{\partial x_i}(x),$$

so that

$$\frac{\partial P}{\partial y}([x], [f]) \left[\frac{\partial f}{\partial x_i} \right] = - \frac{\partial P}{\partial x_i}([x], [f]),$$

in $\mathcal{C}^\infty(U)/(I + \langle \partial g_1/\partial x_i, \dots, \partial g_\ell/\partial x_i \rangle_{\mathcal{C}^\infty(U)})$, for $1 \leq i \leq n$. Note that $(\partial P/\partial y)([x], [f])$ is non-null in $\mathcal{C}^\omega(U)/IC^\omega(U)$ and algebraic over $\mathcal{N}^\omega(U)/I$.

We consider the space Ω_A of Kähler differentials of $A = \mathcal{C}^\infty(U)$ (resp. $\mathcal{C}^\omega(U), \mathcal{N}^\omega(U)$). Note that $\Omega_{A/IA} \cong \Omega_A / (\text{Ad}I + I\Omega_A)$, as an A/IA -module. For the set S of non-zero Nash elements in $\mathcal{C}^\omega(U)/IC^\omega(U)$, $\widetilde{A/IA} = (A/IA)_S$ denote the localization of $A/IA = \mathcal{C}^\infty(U)/IC^\infty(U)$ (resp. $\mathcal{C}^\omega(U)/IC^\omega(U), \mathcal{N}^\omega(U)/I$) by S .

An ideal I of $\mathcal{N}^\omega(U)$ is called *locally formally prime* if, for each $a \in U$, the ideal I_a in the formal algebra $\mathbf{R}[[x - a]]$ generated by $\{j^\infty h(a) \mid h \in I\}$ is prime.

Then we have:

THEOREM 4.1. *Let U be a connected semi-algebraic open subset of \mathbf{R}^n and I a locally formally prime ideal in $\mathcal{N}^\omega(U)$. Let $A = \mathcal{C}^\infty(U)/IC^\infty(U)$, $\mathcal{C}^\omega(U)/IC^\omega(U)$ or $\mathcal{N}^\omega(U)/I$. Then the following 10 conditions on $[f] \in \mathcal{C}^\omega(U)/IC^\omega(U)$ are equivalent to each other:*

- (1) $[f]$ is Nash.
- (2)_A There exists a non-zero Nash element $[g] \in \mathcal{C}^\omega(U)/IC^\omega(U)$ such that

$$[g] \left(\mathbf{d}[f] - \sum_{i=1}^n \left[\frac{\partial f}{\partial x_i} \right] \mathbf{d}[x_i] \right) = 0,$$

in the space Ω_A of Kähler differentials of A .

(3)_A $\mathbf{d}[f] = \sum_{i=1}^n [\partial f / \partial x_i] \mathbf{d}[x_i]$, in the space $\Omega_{\widetilde{A}}$ of Kähler differentials of the localization \widetilde{A} of A by the set of non-zero Nash elements.

(4)_A There exist $\alpha_1, \dots, \alpha_n \in \widetilde{A}$ such that $\mathbf{d}[f] = \sum_{i=1}^n \alpha_i \mathbf{d}[x_i]$, in the space $\Omega_{\widetilde{A}}$.

REMARK 4.2. If I is the ideal of Nash functions vanishing on a connected closed Nash submanifold $M \subset U$, then I is locally formally prime and $IC^\omega(U)$ is prime in $\mathcal{C}^\omega(U)$.

To show Theorem 4.1, we need the following characterization of Nash function. It is proved using the extension theorem due to Efroymsen or its generalization [9]:

LEMMA 4.3. *Let $U \subset \mathbf{R}^n$ be a connected semi-algebraic open subset and $I \subset \mathcal{N}^\omega(U)$ be an ideal. For any $f \in \mathcal{C}^\omega(U)$ the following conditions are equivalent to each other:*

- (i) $[f] \in \mathcal{C}^\omega(U)/IC^\omega(U)$ is Nash.
- (ii) For any $a \in U$, the Taylor series $j^\infty f(a)$ of f at a is algebraic in $\mathbf{R}[[x - a]]/I_a$, in other words, there exists a polynomial $P(x, y) \in \mathbf{R}[x, y]$, $\deg_y P > 0$, which possibly depends on a , such that $j^\infty P(x, f)(a) \in I_a$, where I_a is the ideal in $\mathbf{R}[[x - a]]$ generated by $\{j^\infty h(a) \mid h \in I\}$.
- (iii) There exists a Nash function $g \in \mathcal{N}^\omega(U)$ such that $[g] = [f] \in \mathcal{C}^\omega(U)/IC^\omega(U)$.

PROOF. The implication (i) \Rightarrow (ii) is clear.

(ii) \Rightarrow (iii): Let \mathcal{I} be the finite ideal sheaf generated by I in the sheaf \mathcal{N}^ω_U of Nash functions. Then f defines a section of the quotient sheaf $\mathcal{N}^\omega_U / \mathcal{I}$. By the extension theorem ([6], [9]) in non-compact case, there exists $g \in \mathcal{N}^\omega(U)$ which defines the same section of $\mathcal{N}^\omega_U / \mathcal{I}$ with that defined by f . Therefore $f - g \in \mathcal{C}^\omega(U)$ defines a section of \mathcal{IC}^ω_U , the ideal sheaf generated by \mathcal{I} in the sheaf \mathcal{C}^ω_U of analytic functions. Then $f - g \in IC^\omega(U)$, by Cartan's theorem A for real analytic functions ([5]). Thus we have (iii).

The implication (iii) \Rightarrow (i) is clear. □

PROOF OF THEOREM 4.3. (1) \Rightarrow (2) $_{\mathcal{N}^\omega(U)/I}$: Suppose (1). We take a representative f which belongs to $\mathcal{N}^\omega(U)$ by Lemma 4.3. Then we have

$$\begin{aligned} 0 &= \mathbf{d}(P([x], [f])) = \sum_{i=1}^n \frac{\partial P}{\partial x_i}([x], [f])\mathbf{d}[x_i] + \frac{\partial P}{\partial y}([x], [f])\mathbf{d}[f] \\ &= \sum_{i=1}^n \left(-\frac{\partial P}{\partial y}([x], [f]) \left[\frac{\partial f}{\partial x_i} \right] \right) \mathbf{d}[x_i] + \frac{\partial P}{\partial y}([x], [f])\mathbf{d}[f] \\ &= \frac{\partial P}{\partial y}([x], [f]) \left(\mathbf{d}[f] - \sum_{i=1}^n \frac{\partial P}{\partial x_i}([x], [f])\mathbf{d}[x_i] \right), \end{aligned}$$

in $\Omega_{\mathcal{N}^\omega(U)/I}$, and $(\partial P/\partial y)([x], [f]) \in \mathcal{N}^\omega(U)/I$ is non-zero and algebraic over $\mathcal{N}^\omega(U)/I$.

The implications (j) $_{\mathcal{N}^\omega(U)/I} \Rightarrow$ (j) $_{\mathcal{C}^\omega(U)/IC^\omega(U)} \Rightarrow$ (j) $_{\mathcal{C}^\infty(U)/IC^\infty(U)}$ are clear, for $j = 2, 3, 4$.

The implications (2) $_A \Rightarrow$ (3) $_A$, for $A = \mathcal{N}^\omega(U)/I, \mathcal{C}^\omega(U)/IC^\omega(U), \mathcal{C}^\infty(U)/IC^\infty(U)$, are clear, since $[g] \in S$.

The implications (3) $_A \Rightarrow$ (4) $_A$ for $A = \mathcal{N}^\omega(U)/I, \mathcal{C}^\omega(U)/IC^\omega(U), \mathcal{C}^\infty(U)/IC^\infty(U)$ are clear.

(4) $_{\mathcal{C}^\infty(U)/IC^\infty(U)} \Rightarrow$ (1): Suppose (4) $_{\mathcal{C}^\infty(U)/IC^\infty(U)}$ and $[f]$ is not Nash. Then, by Lemma 4.3, there exists a point $a \in U$ such that $[f]$ is transcendental in $\mathbf{R}[[x-a]]/I_a$ via the \mathbf{R} -algebra homomorphism $\varphi_a : \mathcal{N}^\omega(U)/I \rightarrow \mathbf{R}[[x-a]]/I_a$, where I_a is the ideal in the formal power series ring $\mathbf{R}[[x-a]]$ generated by g_1, \dots, g_ℓ . Let $K = Q(\varphi_a(\mathcal{N}^\omega(U)/I))$ be the quotient field of the image of $\mathcal{N}^\omega(U)/I$ by φ_a . Moreover let $L = K([f], [h_1], \dots, [h_m])$ be the extended field of K which is generated by all elements which appear in the relation $\mathbf{d}[f] - \sum_{i=1}^n \alpha_i \mathbf{d}[x_i] = 0$ in $\Omega_{\mathbf{R}[[x-a]]/I_a}$. Then the relation holds also in Ω_L .

Let u be any non-zero element of L . We extend the zero derivation $D_0 = 0 : K \rightarrow L$ to $D_u : K([f]) \rightarrow L$ by setting $D_u([f]) = u$, for the given non-zero element $u \in L$. Moreover we extend D_u to a derivation $D : L \rightarrow L$. Then for an L -homomorphism $\rho : \Omega_L \rightarrow L$ we have $D = \rho \circ \mathbf{d} : L \rightarrow L$. Then we have

$$0 = \rho \left(\mathbf{d}[f] - \sum_{i=1}^n \alpha_i \mathbf{d}[x_i] \right) = D([f]) = u.$$

This leads to a contradiction. Thus we have (1). □

For a Nash element $[f] \in \mathcal{C}^\omega(U)/IC^\omega(U)$, we define the Leibniz complexity of $[f]$ by the minimal number of terms corresponding to Leibniz rule for $[g] (\mathbf{d}[f] - \sum_{i=1}^n [\partial f/\partial x_i] \mathbf{d}[x_i])$ in the free $\mathcal{N}^\omega(U)/I$ -module $\mathfrak{F}_{\mathcal{N}^\omega(U)/I}$ among all expressions for all non-zero Nash element $[g] \in \mathcal{C}^\omega(U)/IC^\omega(U)$. The definition is based on the statement (2) $_{\mathcal{N}^\omega(U)/I}$ of Theorem 4.1. We do not care about the number of terms corresponding to linearity of the differential. Moreover we will do not count the term generated by the relation $\mathbf{d}([1 \cdot 1]) - [1]\mathbf{d}([1]) - [1]\mathbf{d}([1])$. Therefore we use the relation $\mathbf{d}([c]) = 0$ for $c \in \mathbf{R}$ freely.

Let $\text{LC}([f])$ denote the Leibniz complexity of $[f]$. Similarly to Proposition 3.15 we have an upper estimate:

PROPOSITION 4.4. *Let U be a connected semi-algebraic open subset of \mathbf{R}^n , I a locally formally prime ideal in $\mathcal{N}^\omega(U)$ and $[f] \in \mathcal{C}^\omega(U)/IC^\omega(U)$ Nash. Let $P(x, y)$ be a polynomial such that $P(x, f) \in IC^\omega(U)$ and $(\partial P/\partial y)(x, f) \notin IC^\omega(U)$. Then we have*

$$\text{LC}([f]) \leq \sigma(P) + \sum_{i=0}^n \left(\mu_i + \sum_{j=1}^{s(i)} (r_{ij} - 1) \right).$$

References

- [1] W. Baur and V. Strassen, The complexity of partial derivatives, [Theoretical Computer Science](#), **22** (1983), 317–330.
- [2] R. Benedetti and J.-J. Risler, Real algebraic and semi-algebraic sets, Hermann, 1990.
- [3] J. Bochnak, M. Coste and M.-F. Roy, Géométrie algébrique réelle, [Ergebnisse Der Mathematik Und Ihrer Grenzgebiete 3 Folge](#), Springer-Verlag, 1987.
- [4] J. Bochnak and G. Efroyimson, Real algebraic geometry and the 17th Hilbert problem, [Math. Ann.](#), **251** (1980), 213–241.
- [5] H. Cartan, Variétés analytiques réelles et variétés analytiques complexes, [Bull. Soc. Math. France](#), **85** (1957), 77–99.
- [6] M. Coste, J. M. Ruiz and M. Shiota, Separation, factorization and finite sheaves on Nash manifolds, [Compositio Math.](#), **103** (1996), 31–62.
- [7] M. Coste, J. M. Ruiz and M. Shiota, Uniform bounds on complexity and transfer of global properties of Nash functions, [J. Reine Angew. Math.](#), **536** (2001), 209–235.
- [8] M. Coste, J. M. Ruiz and M. Shiota, Global problems on Nash functions, [Rev. Mat. Complut.](#), **17** (2004), 83–115.
- [9] M. Coste and M. Shiota, Nash functions on noncompact Nash manifolds, [Ann. Sci. École Norm. Sup.](#), **33** (2000), 139–149.
- [10] D. Joyce, Algebraic geometry over C^∞ -rings, arXiv: 1001.0023, (2016).
- [11] D. E. Knuth, The art of computer programming, **2**, Seminumerical Algorithms, 3rd edition, Addison-Wesley, 1998.
- [12] T. Krick and L. M. Pardo, A computational method in Diophantine approximation, In: [Progress in Math.](#), **143**, Birkhäuser, 1996, 193–253.
- [13] H. Matsumura, Commutative ring theory, Cambridge University Press, 1989.
- [14] T. Mostowski, Some properties of the ring of Nash functions, [Annali Scuola Norm. Super. Pisa](#), **3** (1976), 245–266.
- [15] J. Nash, Real algebraic manifolds, [Annals of Math.](#), **56** (1952), 405–421.
- [16] R. Ramanakorasina, Complexité des fonctions de Nash, [Comm. Algebra](#), **17** (1989), 1395–1406.
- [17] J.-J. Risler, Sur l’anneau des fonctions de Nash globales, [Ann. Sci. École Norm. Sup. \(4\)](#), **8** (1975), 365–378.
- [18] M. Rosenlicht, Integration in finite terms, [Amer. Math. Monthly](#), **79** (1972), 963–972.
- [19] M. Shiota, Nash manifolds, [Lecture Notes in Math.](#), **1269**, Springer-Verlag, Berlin, 1987.
- [20] T. Yamashita, Vector fields on differentiable schemes and derivations on differentiable rings, [RIMS Kôkyûroku](#), Singularity theory of differential maps and its applications, **1948** (2015), 57–64.

Goo ISHIKAWA

Department of Mathematics

Hokkaido University

Sapporo 060-0810, Japan

E-mail: ishikawa@math.sci.hokudai.ac.jp

Tatsuya YAMASHITA

Department of Mathematics

Hokkaido University

Sapporo 060-0810, Japan

E-mail: tatsuya-y@math.sci.hokudai.ac.jp