#### AN AXIOM-SYSTEM FOR {K;N}-PROPOSITIONAL CALCULUS RELATED TO SIMONS' AXIOMATIZATION OF S3

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In [4] Simons has shown that Lewis' system S3 can be axiomatized with six mutually independent axiom schemata and the rule of detachment for material implication. As I mentioned in [5], p. 52, it is clear that this formalization of Simons can be reformulated in such a way that instead of axiom schemata the analogous proper axioms

- Α1 ΝΜΚ<sub>p</sub>NK<sub>p</sub>p
- A2 NMKKpqNq
  - A3 NMKKKrpNKqrNKpNq
  - A4 NKNMpNNp
  - A5 NMKpNMp
  - A6 NMKNMKpNqNNMKNMqNNMp

are adopted together with the following two rules of procedure

I. The rule of substitution ordinarily used in the propositional calculus, but adjusted to the primitive functors  $K^{n}$ ,  $N^{n}$  and  $M^{n}$ .

II. The rule of detachment adjusted to the primitive functors "K" and "N", viz.:

If the formulas "NK $\alpha$ N $\beta$ " and " $\alpha$ " are theses of the system, then formula " $\beta$ " is also a thesis of this system.

In this note I like to stress a rather interesting fact that the following four theses

- B1 NKpNKpp
- B2 NKKpqNq
- B3 NKKKrpNKqrNKpNq
- B4 NKNKpNqNNKNqNNp

i.e. the formulas which we can obtain by deleting the modal functor M in the axioms A1, A2, A3 and A6 of Simons, taken together with the rules of procedure I and II constitute an axiom-system for the complete classical  $\{K; N\}$ -

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propositional calculus. It should be noted that this axiomatization which possesses a certain peculiar property does not appear on the list assembled by Porte in [2] of the known axiom-systems for  $\{K; N\}$ -propositional calculus.

In order to prove that the discussed axiomatization constitute the complete classical  $\{K; N\}$ -propositional calculus we proceed as follows:

# METARULE OF PROCEDURE SI

# **SI** If $\vdash \alpha$ and $\vdash NK\alpha NNK\beta Ny$ , then $\vdash NKNyNN\beta$

Proof:

α)	$\vdash \alpha$	[The assumption]
b)	– ΝΚαΝΝΚβΝγ	[The assumption]
c)	μ ΝΚβΝγ	[b; a]
b)	$\vdash NKNYNN\beta$	$[B4, p/\beta, q/\gamma; c]$
		Q. E. D.
B5	NKNKpNqNNKKrpNKqr	[B4, p/KKrpNKqr, q/KpNq; B3]
B6	NKNpp	[B5, q/Kpp, r/Np; B1; SI; B2, q/p]
<b>B</b> 7	NKNKprNNKrNNp	[B5, p/NNp, q/p; B6, p/Np; SI]
B8	NKpNNNp	[B7, p/Np, r/p; B6]
B9	ΝΚΝΚρρΝΝρ	[B7, r/NKpp; B1]
B10	NKKrpNKNNpr	[ <i>B5</i> , <i>q/NNp</i> ; <i>B8</i> ]
B11	NKNKNNprNNKrp	[B7, p/Krp, r/NKNNpr; B10]
B12	NKNNpNKNNpp [B5, p/NK	NNpp, q/NKpp, r/NNp; B11, r/p; SI; B9]
B13	NKNpNNp [B5, p/NNp, q/	<pre>(KNNpp, r/Np; B12; SI; B2, p/NNp, q/p]</pre>
B14	NKpNp	[B5, p/Np, q/Np, r/p; B13; SI; B6]
B15	ΝΚΚρqΝΚqp	[B5, p/q, r/p; B14, p/q]
B16	ΝΚΝΚϥϷΝΝΚϷϥ	[B4, p/Kpq, q/Kqp; B15]

#### METARULE OF PROCEDURE SII

**SII** If  $\vdash$  NK $\alpha$ N $\beta$  and  $\vdash$  NK $\beta$ Ny, then  $\vdash$  NK $\alpha$ Ny

Proof:

a) b) c) b)	⊢ ΝΚαΝβ ⊢ ΝΚβΝγ ⊢ ΝΚΝγα ⊢ ΝΚαΝγ	[The assumption] [The assumption] [B5, $p/\alpha$ , $q/\beta$ , $r/Ny$ ; $\alpha$ ; <b>SI</b> ; $b$ ] [B16, $p/\alpha$ , $q/Ny$ ; c] Q. E. D.
B17	NKNKpqNNKpNNq	[B16, p/q, q/p; B7, p/q, r/p; SII]
B18	NKNKpNNqNNKpq	[B16, p/NNq, q/p; B11, p/q, r/p; SII]
B19	NKNKqpNNKpNNq	[B16; B17; SII]
B20	NKNKqpKpNNq	[B18, p/NKqp, q/KpNNq; B19]
B21	NKK pNNqNKqp	[B16, p/KpNNq, q/NKqp; B20]
B22	NKKpqNp	[B15; B2, p/q, q/p; SII]
B23	NKKNKqrNNKrpNKpNq	[B21, p/NKqr, q/Krp; B3; SII]
B24	NKNKpNqNNKNKqrNNKrp	[B4, p/KNKqrNNKrp, q/KpNq; B23]

Thus, we have the theses B1, B22, and B24, and, therefore, we obtained Rosser's axiom-system of  $\{K;N\}$ -propositional calculus, cf. [3], p. 12 and pp. 54-76. Hence, the proof is completed.

The argumentations concerning the independence of the formulas H1, H2 and H3 given by Simons, [4], pp. 314-315, show also that each of the axioms B1, B2 and B3 does not follow from the remaining postulates of our axiom-system. It is evident that B4 is independent of the other axioms, since B1-B3 and the rules of procedure I and II are such that they allow only: a) if  $\alpha$  is one of the mentioned axioms, to deduce  $K\alpha\alpha$  (by B1) and b) if  $K\beta\beta$  is a formula already proved, to deduce either  $\beta$  (by B2) or  $KK\beta\beta K\beta\beta$  (by B1). And, since no thesis being a consequence of B1-B3 can have a form  $K\alpha N\beta$ , axiom B3 cannot be used at all. Hence, our axioms are mutually independent.

As it was mentioned above the discussed axiom-system possesses a certain peculiar property. Namely, instead of B4 we can adopt, obviously, as an axiom, the following simpler thesis

#### B4\* NKNKpqNNKqNNp

since B4 follows from B4\* by a direct substitution  $(B4^*, q/Nq)$ . Thus, we obtain two inferentially equivalent axiom-systems, viz. {B1; B2; B3; B4} and {B1; B2; B3; B4\*}, such that the former implies B4\* in a rather complicated way, while no use of B1-B3 is needed in order to deduce B4 in the latter system. This situation resembles the cases of the, so called, "generalizing deduction" analyzed by Eukasiewicz in [1].

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