

TWO NOTES ON ACKERMANN'S SET THEORY

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We give solutions to two problems which concern Ackermann's set theory, **A**.* This theory was introduced in [1] and it is now formulated in the first-order predicate calculus with identity, using ϵ for membership and an individual constant, \forall , for the class of all sets. We use the letters ϕ , Ψ , θ , and χ to stand for formulae which do not contain \forall and capital Greek letters to stand for any formulae. Then the axioms of **A** are the universal closures of

$$A1 \quad \forall t (t \in x \leftrightarrow t \in y) \rightarrow x = y,$$

$$A2 \quad \exists z \forall t (t \in z \leftrightarrow t \in \forall \wedge \Theta),$$

$$A3 \quad x \in \forall \wedge (t \in x \vee \forall u (u \in t \rightarrow u \in x)) \rightarrow t \in \forall,$$

$$A4 \quad x, y \in \forall \wedge \forall t (\Psi(x, y, t) \rightarrow t \in \forall) \rightarrow \exists z \epsilon \forall \forall t (t \in z \leftrightarrow \Psi(x, y, t)),$$

where all free variables are shown in *A4* and z does not occur in the Θ of *A2*. **A*** is **A** augmented by the axiom

$$A5 \quad x \in \forall \wedge \exists u u \in x \rightarrow \exists u \epsilon x \forall t \epsilon u t \notin x.$$

Firstly, we shall solve a problem from [3], by extending some of the work on permutation models (see [2], for instance) to models of **A**.

Definition 1 A functional formula $y = F(x)$ is said to be a permutation if it represents a bijection of the universe onto itself. If $y = F(x)$ is a permutation then we write $x \epsilon_F y$ for $F(x) \epsilon y$ and Ψ_F for the formula Ψ with all instances of ϵ replaced by ϵ_F .

Theorem 2 *If $y = F(x)$ is a functional ϵ -formula such that*

- (i) *F is a permutation,*
- (ii) *$x \in \forall$ iff $F(x) \in \forall$,*

*then we can interpret **A** in **A** using ϵ_F for the membership relation and \forall for \forall .*

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Proof: Firstly, note that $x \in_F V \leftrightarrow F(x) \in V \leftrightarrow x \in V$, and we often use this in showing that the interpretations of the axioms hold.

(A1)_F Extensionality holds as F is a bijection.

(A2)_F We show that $\exists z \forall t (t \in_F z \leftrightarrow t \in V \wedge \Psi_F)$. By A2 and (ii) $\exists z \forall p (p \in z \leftrightarrow \exists t \in V (p = F(t) \wedge \Psi_F))$. Then $t \in_F z \leftrightarrow F(t) \in z \leftrightarrow t \in V \wedge \Psi_F$, as required.

(A3)_F We firstly show that $x \in_F y \in V \rightarrow x \in V$. $x \in_F y \rightarrow F(x) \in y \rightarrow F(x) \in V \rightarrow x \in V$ by A3 and (ii). Now it remains to show that

$$\forall t (t \in_F x \rightarrow t \in_F y) \wedge y \in V \rightarrow x \in V.$$

$\forall t (t \in_F x \rightarrow t \in_F y) \rightarrow \forall t (F(t) \in x \rightarrow F(t) \in y)$ by definition,
 $\rightarrow x \subseteq y$ by (ii),
 $\rightarrow x \in V$ by A3.

(A4)_F Suppose that $x, y \in V$ and $\forall t (\phi_F(x, y, t) \rightarrow t \in V)$, and we show that $\exists z \in V \forall t (t \in_F z \leftrightarrow \phi_F(x, y, t))$. Let $\Psi(x, y, z)$ be the formula $\exists t (p = F(t) \wedge \phi_F(x, y, t))$ and then from (ii) we get $\forall p (\Psi(x, y, p) \rightarrow p \in V)$. Then, by A4, $\exists z \in V \forall p (p \in z \leftrightarrow \Psi(x, y, p))$. Now, $t \in_F z \leftrightarrow F(t) \in z \leftrightarrow \phi_F(x, y, t)$, so that (A4)_F holds. Q.E.D.

Question 4.24(c) of [3] asks "If we add the following schema of downward reflection to \mathbf{A} , then do we get \mathbf{A}^* ?"

DR *If ϕ has exactly two free variables, then $y \in V \wedge \phi(V, y) \rightarrow \exists x \in V \phi(x, y)$."*

We shall answer this question negatively, provided that \mathbf{ZF} is consistent. From [3], the consistency of \mathbf{ZF} implies the consistency of \mathbf{A}^* and $\mathbf{A}^* \vdash \mathbf{DR}$ so that it suffices to give an interpretation of $\mathbf{A} + \mathbf{DR}$ in $\mathbf{A} + \mathbf{DR}$ in such a way that the interpretation of A5 fails. We do this as follows:

Let $y = F(x)$ be a functional ϵ -formula which says that

$$\begin{aligned} F(72) &= \{72\}, \\ F(\{72\}) &= 72, \end{aligned}$$

and

$$F(x) = x, \text{ otherwise.}$$

F obviously satisfies the hypothesis of Theorem 2 so that result shows that we can interpret \mathbf{A} in $\mathbf{A} + \mathbf{DR}$, using ϵ_F for membership. An instance of (DR)_F becomes

$$y \in V \wedge \phi_F(V, y) \rightarrow \exists x \in V \phi_F(x, y).$$

This is just another instance of **DR** so that we can interpret $\mathbf{A} + \mathbf{DR}$ in $\mathbf{A} + \mathbf{DR}$, using ϵ_F for membership. The interpretation of A5 does not hold as $x \in_F \{72\} \leftrightarrow F(x) \in \{72\} \leftrightarrow F(x) = 72 \leftrightarrow x = \{72\}$, as required.

Now we shall consider an extension of \mathbf{A} which was suggested by Wang. On page 428 of [4] he suggests that it might be possible to allow any formula to occur in A4 if we modify the axiom to

$$A4_w \quad x, y \in V \wedge \forall t(\Phi(x, y, t) \rightarrow t \in V) \wedge \exists t \in V \neg \Phi(x, y, t) \rightarrow \\ \exists z \in V \forall t(t \in z \leftrightarrow \Phi(z, u, t)),$$

where all free variables are shown. He also mentions that it might be necessary to add the existence of the empty set as an axiom, but it is straightforward to check that there is a model of such axioms in which $V = \{\phi\}$. However, Theorem 3 shows that if we add an axiom asserting the existence of at least two sets, then the theory becomes inconsistent.

Theorem 3 *The theory with axioms A1, A2, A3, A4, and $\exists x \in V \exists y \in V x \neq y$ is inconsistent.*

Proof: In this theory we firstly prove

$$\forall x \in V \quad x \notin x. \quad (*)$$

Suppose that $\exists x \in V \quad x \in x$ and let $\Phi(x)$ be the formula $x \in V \wedge x \notin x$. Then, by $A4_w$, $z = \{t \mid t \in V \wedge t \notin t\} \in V$, but $z \in z \leftrightarrow z \notin z$, so that (*) holds. From (*), we know that

$$V \notin V. \quad (**)$$

Let $x \in V$ and we can suppose that $x \neq \phi$. Using $t \neq t$ in $A4_w$, we see that $\phi \in V$. Let $\Psi_1(t)$ be the formula $t \in V \wedge t \neq \phi$, so that by $A4_w$ $z' = \{x \mid x \in V \wedge x \neq \phi\} \in V$. Let $\Psi_2(t)$ be $t \in z' \vee t = \phi$. Then $\forall t(\Psi_2(t) \rightarrow t \in V)$, $z' \notin z'$ by (*) and $z' \neq \phi$ by our assumption that there are at least two sets. Hence we can use Ψ_2 in $A4_w$ and this gives $\{x \mid x \in z' \vee x = \phi\} \in V$, i.e., $V \in V$, and this contradicts (**). Q.E.D.

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