A New Variant of the Gödel-Malcev Theorem for the Classical Propositional Calculus and Correction to My Paper—The Connective of Necessity of Modal Logic S₅ is Metalogical

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The Gödel-Malcev theorem says that every consistent system has a model. The theorem can be formulated as follows:

(1) If X is a consistent system then there exists such a 0-1 valuation h that $hX \subseteq \{1\}$.

It has the following syntactical version:

(2) If X is a consistent system then there exists such a substitution s that $sX \subseteq CPC$.

(CPC = the set of all classical theses.)

Our variant is a strengthening of (2) and has the following form:

(3) If X is a consistent system then there exists such a substitution s that $X = s^{-1}CPC$.

This theorem can be treated as a theorem on representation of consistent systems by substitutions.

Let $Ln = (Fr, \neg, \land, \lor, \rightarrow, \equiv)$ be an algebra of formulas formed by means of propositional variables p_0, p_1, \ldots and the connectives: \neg (negation), \land (conjunction), \lor (disjunction), \rightarrow (implication), and \equiv (equivalence). By the symbol $Fr^{(k)}$ ($k = 1, 2, \ldots$) we denote the set of all formulas formed by means of variables p_0, \ldots, p_{k-1} (and connectives), and by the symbol $Ln^{(k)}$ we denote the subalgebra $(Fr^{(k)}, \neg, \land, \lor, \rightarrow, \equiv)$ of algebra Ln determined by variables p_0, \ldots, p_{k-1} . By the symbol Cn we denote an operation of consequence (consequence for short) determined by the classical theses, and the rule of modus ponens defined on all subsets of the set Fr (i.e., $Cn(X) = \cap \{Y \subseteq Fr: X \cup CPC \subseteq Y \text{ and the set } Y \text{ is closed under the rule of modus ponens} \}$ ($X \subseteq Fr$).

A set of formulas X is a system if X = Cn(X). A system X is consistent if $X \neq Fr$. The symbol M denotes the classical 0-1 matrix. The expressions Hom(Ln, Ln), Hom(Ln, M), $Hom(Ln^{(k)}, Ln^{(k)})$, $Hom(Ln^{(k)}, M)$ denote sets of homomorphisms between the indicated algebras (in the second and fourth cases it would be more accurate to speak on the algebra of the matrix M). The first and the second set of homomorphisms will be called substitutions and (0-1) valuations, respectively. Without further mention we will use the following well-known equivalence: $a \in Cn(X)$ iff for every $h \in Hom(Ln, M)$ if $hX \subseteq \{1\}$ then ha = 1 ($a \in Fr$, $X \subseteq Fr$).

Theorem If X is a consistent system there exists such a substitution s that $X = s^{-1}CPC$.

Proof: If A, B are any subsets of any space and f is any function defined in this space, then the following equality holds:

$$A = f^{-1}B$$
 iff $fA \subseteq B$ and $f(-A) \subseteq -B$.

Hence the following conjunction is an equivalent form of the equality from the theorem:

(*)
$$sX \subseteq CPC$$
 and $s(Fr - X) \cap CPC = \emptyset$.

Let k be any fixed natural number ≥ 1 . Let $x_0, \ldots, x_{2^{k-1}}, y_0, \ldots, y_{k-1} \in \{0,1\}$ and let

$$F_{(x_0,\ldots,x_{2^k-1})}^k \colon \{0,1\}^k \to \{0,1\}$$

be such a function that the following equality is satisfied:

$$(1) F_{(x_0,\ldots,x_{2^k-1})}^k(y_0,\ldots,y_{k-1}) = x_{|y_0,\ldots,y_{k-1}|}$$

where
$$|y_0, \dots, y_{k-1}| = y_0 2^0 + \dots + y_{k-1} 2^{k-1}$$
.

Since the classical matrix is functionally complete, the above function is definable in it. To the functions of this type correspond k-ary connectives denoted by the same symbols, which are definable by the primitive connectives of the language Ln.

For all i, j $(0 \le i < k, 0 \le j < 2^k)$ we pick $x_j^i \in \{0, 1\}$ in such a way that the following equality is satisfied:

(2)
$$|x_i^0, \dots, x_i^{k-1}| = j$$

and let $f_0, \ldots, f_{2^k-1} \in Hom(Ln^{(k)}, M)$ be picked in such a way that the following equality is satisfied:

$$(3) f_j p_i = x_j^i.$$

Note that the homomorphisms f_0, \ldots, f_{2^k-1} are all different 0-1 valuations for the variables p_0, \ldots, p_{k-1} . Since X is a consistent system, then by virtue of the Gödel-Malcev theorem it follows that there exists such a 0-1 valuation g that $gX \subseteq \{1\}$. Let G be the set of all valuations of this kind and let G_k be the set of all homomorphisms of the type $Ln^{(k)} \to M$ having an extension in G. We pick a sequence $g_0, \ldots, g_{2^k-1} \in G_k$ in such a way that

(4)
$$G_k = \{g_0, \ldots, g_{2^k-1}\}.$$

Let $t \in Hom(Ln^{(k)}, Ln^{(k)})$ satisfy the following condition:

(5)
$$tp_i = F_{(g_0p_1,\ldots,g_{2^k-1}p_i)}^k(p_0,\ldots,p_{k-1}).$$

Now we show that

$$g_j p_i = f_j t p_i.$$

The proof of this equality has the following form

$$f_j t p_i = F_{(g_0 p_i, \dots, g_{2^k-1} p_i)}^k (f_j p_0, \dots, f_j p_{k-1})$$

$$= F_{(g_0 p_i, \dots, g_{2^k-1} p_i)}^k (x_0^0, \dots, x_j^{k-1}) = g_j p_i.$$

Since f_j , g_j , t are homomorphisms, for every $a \in Fr^{(k)}$ the following equality holds:

(6) $g_i a = f_i t a$.

Now we prove the following equivalence:

- (7) $ta \in CPC \text{ iff } g_0 a = \ldots = g_{2^k 1} a = 1.$
- (→) Let $g_j a \neq 1$ for some $0 \leq j < 2^k$. By (6), we obtain $f_j t a \neq 1$. Hence $t a \notin CPC$. (←) Let $g_0 a = \ldots = g_{2^k-1} a = 1$. By (6), we have $f_0 t a = \ldots = f_{2^k-1} t a = 1$. As has been mentioned, f_0, \ldots, f_{2^k-1} are all the 0-1 valuations for the variables p_0, \ldots, p_{k-1} . Since $t a \in Fr^{(k)}$, then $t a \in CPC$. In this way the proof of (7) has been completed.

From the definition of G and G_k it follows that

$$g_0(X \cap Fr^{(k)}) = \ldots = g_{2^k-1}(X \cap Fr^{(k)}) \subseteq \{1\}.$$

Hence, by virtue of (7), we have

$$t(X \cap Fr^{(k)}) \subseteq CPC.$$

Assume, now, that $a \in (Fr - X) \cap Fr^{(k)}$. Since $a \notin Cn(X)$ then there exists such a valuation g that $gX \subseteq \{1\}$ and $ga \ne 1$. Hence by the definition of G we have $g \in G$. Let g' be the restriction of g to the variables p_0, \ldots, p_{k-1} . Obviously $g' \in G_k$. Since $g'a \ne 1$ then by virtue of (4) and (7) we have $ta \notin CPC$. So $t((Fr - X) \cap Fr^{(k)}) \cap CPC = \emptyset$. Then

(8)
$$t(X \cap Fr^{(k)}) \subseteq CPC$$
 and $t((Fr - X) \cap Fr^{(k)}) \cap CPC = \emptyset$.

Next we prove that the substitution t can be extended into a substitution t^+ defined on p_k in such a way that (8) is satisfied for k+1. The sets G_k and G_{k+1} are picked in such a way that every homomorphism belonging to G_k has an extension onto the variable p_k which belongs to G_{k+1} and G_{k+1} includes only such homomorphisms that are extensions of some homomorphisms from G_k . If now $h \in G_k$, then by the symbols h', h'' we denote all extensions of h onto the variable p_k that belong to G_{k+1} . (If h has only one extension then h' = h''.) Hence, and by virtue of (4), we obtain

$$(4') \{g'_0,\ldots,g'_{2^k-1},g''_0,\ldots,g''_{2^k-1}\}=G_{k+1}.$$

Let, now, $u \in Hom(Ln^{(k+1)}, Ln^{(k+1)})$ satisfy the following condition

$$(5') \ up_i = F_{(g'_0p_i,\ldots,g'_{2^k-1}p_i,g'_0p_i,\ldots,g''_{2^k-1}p_i)}^{k+1}(p_0,\ldots,p_k)(0 \le i \le k).$$

As by (4) and (5), we have proved (8) (for t and k). In an analogous way, by (4') and (5') we prove that

(8')
$$u(X \cap Fr^{(k+1)}) \subseteq CPC$$
 and $u((Fr - X) \cap Fr^{(k+1)}) \cap CPC = \emptyset$.

Let h be any 0-1 valuation. By (5) and (5') we have

$$htp_i = F_{(g_0p_i,\ldots,g_{2^k-1}p_i)}^k(hp_0,\ldots,hp_{k-1})$$

$$hup_i = F_{(g'_0p_1, \dots, g'_{2^k-1}p_i, g''_{p_i}, \dots, g''_{2^k-1}p_i)}^{k+1}(hp_0, \dots, hp_k)(0 \le i < k).$$

Note that if the number $|hp_0, \ldots, hp_{k-1}|$ is denoted by the letter j, then

$$htp_i = g_j p_i$$

$$hup_i = \begin{cases} g'_j p_i \text{ if } hp_k = 0 & (0 \le i < k) \\ g''_j p_i \text{ if } hp_k = 1 \end{cases}$$

For i < k we have $g_j p_i = g'_j p_i = g''_j p_i$. Then for i < k we have $htp_i = hup_i$. Since h is any 0 - 1 valuation and h, t, u are homomorphisms, then

(9)
$$\lceil ta \equiv ua \rceil \in CPC$$
 for every $a \in Fr^{(k)}$

Let, now, $t^+ \in Hom(Ln^{(k+1)}, Ln^{(k+1)})$ satisfy the following condition

$$t^+p_i = \begin{cases} tp_i \text{ if } i < k \\ up_k \text{ if } i = k \end{cases} (0 \le i \le k).$$

Hence by virtue of (8') and (9), it is easy to see that t^+ is an extension of t onto the variable p_k such that

$$(8'') \ t^+(X \cap F^{(k+1)}) \subseteq CPC \text{ and } t^+((Fr-X) \cap Fr^{(k+1)}) \cap CPC = \emptyset.$$

Now we come to the final part of our proof. By (8) and (8"), there exists a sequence of homomorphisms $t_1 \in Hom(Ln^{(1)}, Ln^{(1)}), t_2 \in Hom(Ln^{(2)}, Ln^{(2)}), \ldots$ such that each next one is an extension of the preceding one onto succeeding variables satisfying (8) for $k = 1, 2, \ldots$ Let s be a substitution (i.e., $s \in Hom(Ln, Ln)$) obtained by joining those homomorphisms. We prove that s is the substitution we search for, i.e., (*) holds.

Suppose that for some $a \in X$ it is false that $sa \in CPC$. Let n be such that $a \in Fr^{(n)}$. By (8) (for k = n and $t = t_n$), we have $sa \in t_n(X \cap Fr^{(n)}) \subseteq CPC - a$ contradiction. Similarly we prove the second part of (*). QED

Finally, we point out two errors in [1]. The proof of Lemma 3 was done incorrectly (see the second application of the rule r_2). By virtue of (*) from the proof of the theorem we obtain the following implication:

If
$$\{a_1, \ldots, a_k\} \cap Cn(\{a\}) = \emptyset$$
, then there exists such a substitution $s \in Hom(Ln, Ln)$ that $s\{a_1, \ldots, a_k\} \cap CPC = \emptyset$ and $sa \in CPC$ $(a, a_1, \ldots, a_k \in Fr)$.

Since the deduction theorem holds for Cn, the above implication can be expressed as follows:

If $a \to a_1 \in CPC$, ..., $a \to a_k \in CPC$, then there exists such a substitution $s \in Hom(Ln, Ln)$ that $sa \in CPC$ and it is not true $sa_1 \in CPC$ or ... or $sa_k \in CPC$.

From this implication, Lemma 3 follows trivially.

The second error is connected with the proof of Theorem 3. Details aside, the point is that in order to prove that the effective equivalence of Stone's theorem and the theorem that the rule r^* holds, in the schema of this rule we have to use a nondenumerable rather than denumerable number of formulas. After this correction, the schema of r^* is artificial and uninteresting. Therefore I abstain from correcting this proof.

NOTE

1. For the proof of the implication (1) \rightarrow (2), it is enough for the substitution s to be defined as follows

$$sp_i = \begin{cases} p_0 \to p_0 & \text{if } hp_i = 1 \\ \neg (p_0 \to p_0) & \text{otherwise} \end{cases}$$

 $(i = 0, 1, ...; p_i)$ is a propositional variable). And for the proof of the implication $(2) \rightarrow (1)$ it is enough for h to be defined as follows: $h = h' \circ s$ where h' is any 0 - 1 valuation.

REFERENCES

[1] Dywan, Z., "The connective of necessity of modal logic S₅ is metalogical," *Notre Dame Journal of Formal Logic*, vol. 24 (1983), pp. 410-414.

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