# ON A THEOREM OF RADEMACHER-TURÁN 

Dedicated to Hans Rademacher on the occasion of his seventieth birthday

BY<br>P. Erdös

A set of points some of which are connected by an edge will be called a graph $G$. Two vertices are connected by at most one edge, and loops (i.e., edges whose endpoints coincide) will be excluded. Vertices will be denoted by $\alpha, \beta, \cdots$, edges will be denoted by $e_{1}, e_{2}, \cdots$ or by $(\alpha, \beta)$ where the edge ( $\alpha, \beta$ ) connects the vertices $\alpha$ and $\beta$.
$G-e_{1}-\cdots-e_{k}$ will denote the graph from which the edges $e_{1}, \cdots, e_{k}$ have been omitted, and $G-\alpha_{1}-\cdots-\alpha_{k}$ denotes the graph from which the vertices $\alpha_{1}, \cdots, \alpha_{k}$ and all the edges emanating from them have been omitted; similarly $G+e_{1}+\cdots+e_{k}$ will denote the graph to which the edges $e_{1}, \cdots, e_{k}$ have been added (without generating a new vertex).

The valency $v(\alpha)$ of a vertex will denote the number of edges emanating from it. $\quad G_{u}^{(v)}$ will denote a graph having $v$ vertices and $u$ edges. The graph $G_{\binom{(2)}{2}}^{(k . e ., ~ t h e ~ g r a p h ~ o f ~} k$ vertices any two of which are connected by an edge) will be called the complete $k$-gon.

A graph is called even if every circuit of it has an even number of edges.
Turán ${ }^{1}$ proved that every

$$
G_{V+1}^{(n)}, \quad V=\frac{k-2}{2(k-1)}\left(n^{2}-r^{2}\right)+\binom{r}{2}
$$

for $n=(k-1) t+r, 0 \leqq r<k-1$, contains a complete $k$-gon, and he determined the structure of the $G_{V}^{(n)}$,s which do not contain a complete $k$-gon. Thus if we put $f(2 m)=m^{2}, f(2 m+1)=m(m+1)$, a special case of Turán's theorem states that every $G_{f(n)+1}^{(n)}$ contains a triangle.

In 1941 Rademacher proved that for even $n$ every $G_{f(n)+1}^{(n)}$ contains at least [ $n / 2$ ] triangles and that [ $n / 2$ ] is best possible. Rademacher's proof was not published. Later on ${ }^{2}$ I simplified Rademacher's proof and proved more generally that for $t \leqq 3, n>2 t$, every $G_{f(n)+t}^{(n)}$ contains at least $t[n / 2]$ triangles. Further I conjectured that for $t<[n / 2]$ every $G_{f(n)+t}^{(n)}$ contains at least $t[n / 2]$ triangles. It is easy to see that for $n=2 m, 2 m>4$, the conjecture is false for $t=n / 2$. To see this, consider a graph $G_{m}^{(2 m)}$ whose vertices are

[^0]$\alpha_{1}, \cdots, \alpha_{2 m}$ and whose edges are
$$
\left(\alpha_{i}, \alpha_{j}\right), \quad 1 \leqq i \leqq m+1<j \leqq 2 m
$$
and the $m+1$ further edges
$$
\left(\alpha_{i}, \alpha_{i+1}\right), \quad 1 \leqq i \leqq m, \quad \text { and } \quad\left(\alpha_{1}, \alpha_{m+1}\right)
$$

It is easy to see that this graph contains $m^{2}-1$ triangles (for $2 m=4$ an unwanted triangle ( $\alpha_{1}, \alpha_{2}, \alpha_{3}$ ) enters and ruins the counting, and in fact it is easy to see that for $2 m=4$ the conjecture holds for $t=m=2$ ). For odd $n=2 m+1$ perhaps every $G_{f(2 m+1)+t}^{(2 m+1)}, t \leqq 2 m-2$, contains at least tm triangles. But here is a $G_{f(2 m+1)+2 m-1}^{(2 m+1)}, 2 m+1 \geqq 9$, which contains fewer than $m(2 m-1)$ triangles. The vertices of our graph are $\alpha_{1}, \cdots, \alpha_{2 m+1}$, the edges are

$$
\left(\alpha_{i}, \alpha_{j}\right), \quad 1 \leqq i \leqq m+2<j \leqq 2 m+1
$$

and the following $2 m+1$ edges:

$$
\left(\alpha_{1}, \alpha_{k}\right), \quad\left(\alpha_{2}, \alpha_{k}\right), \quad\left(\alpha_{3}, \alpha_{4}\right), \quad\left(\alpha_{3}, \alpha_{5}\right), \quad\left(\alpha_{3}, \alpha_{6}\right)
$$

$$
3 \leqq k \leqq m+2
$$

It is easy to see that this graph contains $2 m^{2}-m-1<m(2 m-1)$ triangles. For $2 m+1=5$ we must have $t \leqq 4$, and it is easy to see that the conjecture holds for all these $t$. For $2 m+1=7, t \leqq 9$, and by a little longer argument one can easily convince oneself that the conjecture holds for all these $t$.

In the present paper we are going to prove the following
Theorem. There exists a constant $c_{1}>0$ so that for $t<c_{1} n / 2$ every $G_{f(n)+t}^{(n)}$ contains at least $t[n / 2]$ triangles.

First we need three lemmas.
Lemma 1. Every $G_{f(n-1)+2}^{(n)}$ which is not even contains a triangle.
Lemma 1 was found jointly by Gallai and myself. (The lemma was also found by Mr. Andrásfai independently.)

Let $G$ be a graph with $n$ vertices which is not even and contains no triangle. Let $\alpha_{1}, \cdots, \alpha_{2 k+1}$ be the vertices of the odd circuit of our graph having the least number of vertices. We can assume $3<2 k+1 \leqq n$. The subgraph of $G$ spanned by $\alpha_{1}, \cdots, \alpha_{2 k+1}$ can have no other edges; otherwise our graph would contain an odd circuit having fewer than $2 k+1$ edges. Let $\beta_{1}, \cdots$, $\beta_{n-2 k-1}$ be the other vertices of $G$. Any of the $\beta$ 's can be connected with at most two of the $\alpha$ 's, for otherwise $G$ contains an odd circuit of fewer than $2 k+1$ edges. Finally by Turán's theorem the subgraph of $G$ spanned by $\beta_{1}, \cdots, \beta_{n-2 k-1}$ can have at most $f(n-2 k-1)$ edges. Thus the number of edges of $G$ is at most

$$
2 k+1+2(n-2 k-1)+f(n-2 k-1) \leqq f(n-1)+1
$$

by a simple calculation (equality only for $2 k+1=5$ ). This completes the proof of our lemma.

Our proof in fact gives that a graph $G$ of $n$ vertices whose smallest odd circuit has $2 k+1$ vertices, $k>1$, has at most $2 n-2 k-1+f(n-2 k-1)$ edges, and the following simple example shows that this result is best possible. Let the vertices of $G$ be

$$
\begin{aligned}
\alpha_{1}, \cdots, \alpha_{v}, & \beta_{1}, \cdots, \beta_{u}, \quad \gamma_{1}, \cdots, \gamma_{2 k+1} \\
v & =\left[\frac{n-2 k-1}{2}\right], \quad u=n-\left[\frac{n-2 k-1}{2}\right] .
\end{aligned}
$$

The edges of $G$ are $\left(\alpha_{i}, \beta_{j}\right),\left(\gamma_{1}, \alpha_{i}\right),\left(\gamma_{3}, \alpha_{i}\right), 1 \leqq i \leqq v, \quad\left(\gamma_{2}, \beta_{i}\right),\left(\gamma_{4}, \beta_{i}\right)$, $1 \leqq i \leqq u$, further the edges $\left(\gamma_{i}, \gamma_{i+1}\right), 1 \leqq i \leqq 2 k$, $\left(\gamma_{1}, \gamma_{2 k+1}\right)$.

Lemma 2. There exists a constant $c_{2}>0$ so that every $G_{f(n)+1}^{(n)}$ contains at least $\left[c_{2} n\right]$ triangles having a common edge ( $\alpha_{1}, \alpha_{2}$ ) (i.e., all the edges ( $\alpha_{1}, \beta_{i}$ ), $\left(\alpha_{2}, \beta_{i}\right),\left(\alpha_{1}, \alpha_{2}\right), 1 \leqq i \leqq\left[c_{2} n\right]$, are in $\left.G_{f(n)+1}^{(n)}\right)$.

Let $\left(\alpha_{i}, \beta_{i}, \gamma_{i}\right), 1 \leqq i \leqq r$, be a maximal system of disjoint triangles of our $\operatorname{graph} G_{f(n)+1}^{(n)}$. In other words if we omit the vertices $\alpha_{i}, \beta_{i}, \gamma_{i}, 1 \leqq i \leqq r$, the subgraph of $G_{f(n)+1}^{(n)}$ spanned by the remaining $n-3 r$ vertices contains no triangle and has therefore at most $f(n-3 r)$ edges (by Turán's theorem).

Denote by $G(n, i)$ the graph $G_{f(n)+1}^{(n)}-\sum_{j=1}^{i-1}\left(\alpha_{j}+\beta_{j}+\gamma_{j}\right)$, and let $v^{(i)}\left(\alpha_{i}\right), v^{(i)}\left(\beta_{i}\right), v^{(i)}\left(\gamma_{i}\right)$ be the valencies of $\alpha_{i}, \beta_{i}, \gamma_{i}$ in $G(n, i)$. Now we show that for some $i, 1 \leqq i \leqq r$, we must have

$$
\begin{equation*}
v^{(i)}\left(\alpha_{i}\right)+v^{(i)}\left(\beta_{i}\right)+v^{(i)}\left(\gamma_{i}\right)>n\left(1+9 c_{2}\right)-3 i \tag{1}
\end{equation*}
$$

For if (1) would not hold for any $i, 1 \leqq i \leqq r$, then the number of edges of $G_{f(n)+1}^{(n)}$ would be not greater than

$$
\begin{equation*}
\sum_{i=1}^{r}\left(n\left(1+9 c_{2}\right)-3 i\right)+f(n-3 r)<f(n) \tag{2}
\end{equation*}
$$

by a simple calculation for sufficiently small $c_{2}$. But (2) is an evident contradiction since $G_{f(n)+1}^{(n)}$ has by definition $f(n)+1$ edges.

Thus (1) holds for say $i=i_{0}$. Then a simple computation shows that there are at least $3\left[c_{2} n\right]$ vertices of $G\left(n, i_{0}\right)$ which are connected with more than one of the vertices $\alpha_{i_{0}}, \beta_{i_{0}}, \gamma_{i_{0}}$. Therefore there are at least [ $c_{2} n$ ] of them which are connected with the same pair, i.e., $G\left(n, i_{0}\right)$, and therefore $G_{f(n)+1}^{(n)}$, contains the configuration required by Lemma 2, which completes the proof of the lemma.

By more careful considerations we can prove that every $G_{f(n)+1}^{(n)}$ contains $n / 6+O(1)$ triangles $\left(\alpha_{1}, \alpha_{2}, \beta_{i}\right), 1 \leqq i \leqq n / 6+O(1)$, and that this result is best possible.

Lemma 3. Let $\delta>0$ be a fixed number. Consider any graph

$$
G_{u}^{(n)}, \quad u>f(n)-(n / 2)(1-\delta), \quad n>n_{0}(\delta)
$$

which contains a triangle. Then $G_{u}^{(n)}$ contains an edge $\left(\alpha_{1}, \alpha_{2}\right)$ and $\left[c_{3} n\right]+1, c_{3}=$ $c_{3}(\delta)$, vertices $\beta_{1}, \cdots, \beta_{r}, r=\left[c_{3} n\right]+1$, so that all the triangles $\left(\alpha_{1}, \alpha_{2}, \beta_{i}\right)$, $1 \leqq i \leqq r$, are in $G_{u}^{(n)}$.

By assumption $G_{u}^{(n)}$ contains a triangle $\left(\alpha_{1}, \alpha_{2}, \alpha_{3}\right)$. Assume first that

$$
\begin{equation*}
v\left(\alpha_{1}\right)+v\left(\alpha_{2}\right)+v\left(\alpha_{3}\right)>n\left(1+9 c_{3}\right)+9 \tag{3}
\end{equation*}
$$

Then as in the proof of Lemma 2 we can show that $G_{u}^{(n)}$ contains the required configuration.

If (3) is not satisfied, then $G_{u}^{(n)}-\alpha_{1}-\alpha_{2}-\alpha_{3}$ has $n-3$ vertices and at least $u-n\left(1+9 c_{3}\right)-9$ edges. But if $c_{3}<\delta / 18$, then for $n>n_{0}$

$$
\begin{aligned}
& u-n\left(1+9 c_{3}\right)-9 \\
& \quad>f(n)-(n / 2)(1-\alpha)-n\left(1+9 c_{3}\right)-9>f(n-3)
\end{aligned}
$$

Thus by Lemma $2, G_{u}^{(n)}-\alpha_{1}-\alpha_{2}-\alpha_{3}$, and therefore $G_{u}^{(n)}$, contains the configuration required by Lemma 3, which completes the proof of Lemma 3.

Now we can prove our Theorem. Let there be given a $G_{f(n)+t}^{(n)}, t<c_{1} n / 2$. Assume first that after the omission of any $r=\left[c_{1} n / 2 c_{3}\right], c_{3}=c_{3}\left(\frac{1}{4}\right)\left(\delta=\frac{1}{4}\right.$ in Lemma 3), edges the graph will still contain a triangle. For sufficiently small $c_{1}, c_{1} / 2 c_{3}<\frac{1}{4}$; thus it will be permissible to apply Lemma 3 during the omission of these edges.

By Lemma 3 (or Lemma 2) there exists an edge $e_{1}$ which is contained in at least $\left[c_{3} n\right]+1$ triangles of $G_{f(n)+t}^{(n)}$; again by Lemma 3 in $G_{f(n)+t}^{(n)}-e_{1}$ there exists an edge $e_{2}$ which is contained in $\left[c_{3} n\right]+1$ triangles of $G_{f(n)+t}^{(n)}-e_{1}$. Suppose we have already chosen the edges $e_{1}, \cdots, e_{r}$ each of which is contained in at least $\left[c_{3} n\right]+1$ triangles. By our assumption $G_{f(n)+t}^{(n)}-e_{1}-\cdots-e_{r}$ contains at least one triangle; thus by Lemma 3 there is an edge $e_{r+1}$ in $G_{f(n)+t}^{(n)}-e_{1}-\cdots-e_{r}$ which is contained in at least $\left[c_{3} n\right]+1$ triangles in this graph. These triangles incident on the edges $e_{1}, \cdots, e_{r+1}$ are evidently distinct; thus $G_{f(n)+t}^{(n)}$ contains at least $(r+1)\left(\left[c_{3} n\right]+1\right)>c_{1} n^{2} / 2>t n / 2$ triangles, which proves our Theorem in this case.

Therefore we can assume that there are $l \leqq r<n / 4$ edges $e_{1}, \cdots, e_{l}$ so that $G=G_{f(n)+t}^{(n)}-e_{1}-\cdots-e_{l}$ contains no triangle, and we can assume that $l$ is the smallest integer with this property. $\mathrm{By} l \leqq r<n / 4, G$ has

$$
f(n)+t-l>f(n)-n / 4>f(n-1)+1
$$

edges. Thus by Lemma $1, G$ is even.
By Turán's theorem, $l \geqq t$. Assume first $l=t$ (it is not necessary to treat the cases $l=t$ and $l>t$ separately, but perhaps it will be easier for the reader to do so). Then $G$ has $f(n)$ edges, and by Turán's theorem $G$ is of the following form: The vertices of $G$ are $\alpha_{1}, \cdots, \alpha_{[n / 2]}, \beta_{1}, \cdots, \beta_{n-[n / 2]}$, and the edges are $\left(\alpha_{i}, \beta_{j}\right), 1 \leqq i \leqq[n / 2], 1 \leqq j \leqq n-[n / 2]$. A simple argument shows that the addition of every further edge introduces at least
[ $n / 2$ ] triangles and that these triangles are distinct. Thus $G_{f(n)+t}^{(n)}$ contains at least $t[n / 2]$ triangles, and our Theorem is proved in this case too.

Assume finally $l=t+w, 0<w<n / 4$ (since $l<n / 4$ ). It will be more convenient to assume first that $n$ is even. Put $n=2 m$. Since $G$ is even, it is contained in a graph $G(E, u)$ whose vertices are $\alpha_{1}, \cdots, \alpha_{m-u}, \beta_{1}, \cdots, \beta_{m+u}$ and whose edges are $\left(\alpha_{i}, \beta_{j}\right), 1 \leqq i \leqq m-u, 1 \leqq j \leqq m+u$ (since $G$ has more than $f(2 m)-m / 2=m^{2}-m / 2$ edges, we have $\left.0 \leqq u<(m / 2)^{1 / 2}\right)$.

Clearly every one of the edges $e_{1}, \cdots, e_{l}$ connect two $\alpha$ 's or two $\beta$ 's. For if say $e_{i}$ would connect an $\alpha$ with a $\beta$, then

$$
G_{f(n)+t}^{(n)}-e_{1}-\cdots-e_{1-1}-e_{i+1}-\cdots-e_{l}
$$

would be even, and hence would contain no triangle, which contradicts the minimum property of $l$.

By our assumption $G$ is a subgraph of $G(E, u)$. Assume that $G$ is obtained from $G(E, u)$ by the omission of $x$ edges. Then we evidently have

$$
l=x+u^{2}+t \quad\left(\text { or } w=x+u^{2}\right)
$$

and $G_{f(n)+t}^{(n)}$ is obtained from $G$ by adding $l$ edges $e_{1}, \cdots, e_{l}$ which are all of the form $\left(\alpha_{i_{1}}, \alpha_{i_{2}}\right)$ or $\left(\beta_{i_{1}}, \beta_{i_{2}}\right)$. Put $e_{i}=\left(\beta_{i_{1}}, \beta_{i_{2}}\right)$, and let us estimate the number of triangles $\left(\beta_{i_{1}}, \beta_{i_{2}}, \alpha_{j}\right)$ in $G(E, u)+e_{i}$. Clearly at most $x$ of the edges $\left(\beta_{i_{1}}, \alpha_{j}\right),\left(\beta_{i_{2}}, \alpha_{j}\right)$ are not in $G(E, u)$; thus $G(E, u)+e_{i}$ contains at least

$$
m-u-x
$$

triangles (if $e_{i}$ connects two $\alpha$ 's, then $G(E, u)+e_{i}$ contains at least $m+u-x$ triangles). For different $e_{i}$ 's these triangles are clearly different; thus

$$
G_{f(n)+t}^{(n)}=G+e_{1}+\cdots+e_{l}
$$

contains at least

$$
\begin{equation*}
(m-u-x) l=(m-u-x)\left(x+u^{2}+t\right) \geqq t m=t(n / 2) \tag{4}
\end{equation*}
$$

triangles. (4) follows by simple computation from $l=u^{2}+x+t<m / 2$. (4) completes the proof of our Theorem for $n=2 m$. For $n=2 m+1$ the proof is almost identical and can be omitted. Thus the proof of our Theorem is complete.

It seems possible that a slight improvement of this proof will give the conjecture that every $G_{f(n)+t}^{(n)}, t<[n / 2]$ contains at least $t[n / 2]$ triangles, but I have not been successful in doing this.

I have not succeeded in formulating a reasonable conjecture about the minimum number of trianglesa $G_{f(n)+t}^{(n)}$ must contain if $[n / 2] \leqq t \leqq\binom{ n}{2}-f(n)$. It is easy to see that if $t$ is close to $\binom{n}{2}-f(n)$, then $G_{f(n)+t}^{(n)}$ contains more than $t[n / 2]$ triangles, and it would be easy to obtain a best possible result in this case. But I have not investigated the range of $t$ for which this is possible. I just remark that $G_{\left(\begin{array}{l}(2)-l\end{array}\right.}^{(n)}, l \leqq 2$, contains at least $\binom{n}{3}-l(n-2)$ triangles
and that $G_{\binom{n}{2}-3}^{(n)}$ contains at least $\binom{n}{3}-3(n-2)+1$ triangles, and that these results are best possible. The simple proofs are left to the reader.

Turán's theorem implies that every $G_{3 n^{2}+1}^{(3 n)}$ contains a complete 4 -gon. As an analogue of the theorem of Rademacher I can prove by very much more complicated arguments that it contains at least $n^{2}$ complete 4 -gons; this result is easily seen to be best possible.

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[^0]:    Received March 20, 1961.
    ${ }^{1}$ P. Turán, Matematikai és Fizikai Lapok, vol. 48 (1941), pp. 436-452 (in Hungarian); see also On the theory of graphs, Colloq. Math., vol. 3 (1954), pp. 19-30.
    ${ }^{2}$ P. Erdös, Some theorems on graphs, Riveon Lematematika, vol. 9 (1955), pp. 13-17 (in Hebrew with English summary).

