## A DECISION PROCEDURE FOR A CLASS OF FORMULAS OF FIRST ORDER PREDICATE CALCULUS

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1. Introduction. In [4] (c.f. also case V' page 256 of [1]) J. Herbrand provides a decision procedure which is equivalent to a decision procedure for determining for a fixed contradiction C and for any first order prenex formula  $\Gamma$  whose matrix is a conjunction of signed atomic formulas, whether  $\Gamma \rightarrow C$  is valid. In this paper we define a class a of first order formulas and then provide a decision procedure for determining for any first order prenex formula  $\Gamma$  whose matrix is a conjunction of signed atomic formulas and and member  $\varDelta$  of the class, whether  $\Gamma \rightarrow \varDelta$  is valid. Although the class of formulas  $\varDelta$  that we consider is rather large, it is clear that some restriction is necessary since a decision procedure for the class itself is obtained by using for  $\Gamma$  a single propositional parameter that does not occur in  $\varDelta$ .

The formulas we consider are those of any system of pure first order predicate calculus without equality and without function symbols. We use  $\lor$ ,  $\land$ ,  $\neg$ , and  $\rightarrow$  for the propositional connectives disjunction, conjunction, denial, and the conditional, respectively. The symbols  $\Gamma$ ,  $\Delta$ ,  $\Gamma_0$ ,  $\Delta_0$ ,  $\Gamma_1$ ,  $\Delta_1$ ,  $\cdots$  shall range over arbitrary formulas, P, Q,  $P_1$ ,  $Q_1$ ,  $\cdots$ over prefixes, and  $M, N, M_1, N_1, \cdots$  over matrices. A propositional parameter or predicate parameter together with its attached individual variables or individual parameters will be called an atomic formula. An occurrence of an atomic formula in a formula  $\Gamma$  is called an *atomic* part of  $\Gamma$ . Two prenex formulas are similar if their matrices differ only in the symbols occupying individual variable places of the atomic formulas. Two prenex formulas are *congruent* if they differ only by equivalent replacements of bound variables. We indicate that  $\varDelta$  is a logical consequence of  $\Gamma$  by writing  $\Gamma \models \Delta$ . If  $\Gamma \models \Delta$  then there exists a symmetric L-deduction of  $\Delta$  from  $\Gamma$  as described in [2]. For any formulas  $\Gamma$  and  $\Delta$  an L-deduction of  $\Delta$  from  $\Gamma$  is an ordered (n + 1)tuple  $\langle \Gamma_0, \cdots, \Gamma_n \rangle$  where  $\Gamma_0 = \Gamma$  and  $\Gamma_n = \Delta$ , together with a specification of how, for any m < n,  $\Gamma_{m+1}$  results from  $\Gamma_m$  by an application of an L-rule. The reader is referred to pages 252 and 253 of [2] for the definitions of the eleven L-rules. An L-deduction is symmetric if and only if the order in which the different kinds of L-rules are applied satisfies conditions (iii) through (vi) on page 257 of [2]. In addition, for convenience, we require that a symmetric L-deduction have exactly one application of the operation matrix change.

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