

A NOTE ON A RESULT OF KAMAE*

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In [3] Kamae proved that for each constant c , there exists a finite string y such that

$$K(y) - K(y|x) > c$$

for all but finitely many finite strings x , where $K(\cdot)$ and $K(\cdot | \cdot)$ are the unconditional and conditional minimal-program complexity measures respectively of Kolmogorov [4]. By considering infinite sequences we are able to obtain a slightly stronger statement of this result.

Let X^∞ denote the set of all infinite binary strings. For $x \in X^\infty$ let x^n denote the initial segment of x of length n , i.e., $x^n = x(1) \dots x(n)$. To simplify matters we will associate with each finite binary string y the integer n whose binary representation is $1y$. By this means we will consider complexity expressions of the form $K(x^n)$, $K(x^n|m)$ and $K(m)$ for $x \in X^\infty$ and integers n and m . Then $K(x^n|n)$ is Kolmogorov's restricted conditional complexity. By a recursively enumerable sequence we mean the characteristic sequence of a recursively enumerable set. By " $\exists n$ " and " $\forall n$ " we mean respectively "There exist infinitely many integers n " and "For all but finitely many integers n ."

Theorem. *There exists a recursively enumerable sequence x , such that*

$$\forall c \exists n \forall m. K(x^n|n) - K(x^n|m) > c.$$

Proof. We need the following two lemmas.

Lemma. *For every recursively enumerable sequence x ,*

$$\exists c_1 \forall n \forall m. K(x^n|m) \leq K(n) + c_1.$$

Proof. Let h be a total recursive function which enumerates the 1's of x , i.e., $x(i) = 1 \Leftrightarrow \exists j. h(j) = i$. Define the total recursive function $f(n, m)$ as follows:
 Step 1: Enumerate via h the first m 1's of the sequence x , i.e. compute $h(1), \dots, h(m)$.
 Step 2: Output the following finite string y of length n — $y(i) = 1 \Leftrightarrow i$ appears on

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