

# DECISION METHODS IN THE THEORY OF ORDINALS<sup>1</sup>

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For an ordinal  $\alpha$ , let  $RS(\alpha)$ , the restricted second order theory of  $[\alpha, <]$ , be the interpreted formalism containing the first order theory of  $[\alpha, <]$  and quantification on monadic predicate variables, ranging over all subsets of  $\alpha$ . For a cardinal  $\gamma$ ,  $RS(\alpha, \gamma)$  is like  $RS(\alpha)$ , except that the predicate variables are now restricted to range over subsets of  $\alpha$  of cardinality less than  $\gamma$ .  $\omega = \omega_0$  and  $\omega_1$  denote the first two infinite cardinals. In this note I will outline results concerning  $RS(\alpha, \omega_0)$ , which were obtained in the Spring of 1964 (detailed proofs will appear in [8]), and the corresponding stronger results about  $RS(\alpha, \omega_1)$ , which were obtained in the Fall of 1964.

The binary expansion of natural numbers can be extended to ordinals. If  $x < 2^\alpha$ , let  $\phi x$  be the finite subset  $\{u_1, \dots, u_n\}$  of  $\alpha$ , given by  $x = 2^{u_1} + \dots + 2^{u_n}$ ,  $u_n < \dots < u_1$ .  $\phi$  is a one-to-one map of  $2^\alpha$  onto all finite subsets of  $\alpha$ . Let  $Exy$  stand for  $(\exists u)[x = 2^u \wedge u \in \phi y]$ , and note that the algorithm  $i + j = s$ , for addition in binary notation can be expressed in  $RS(\alpha, \omega_0)$ . It now is easy to see that the first order theory  $FT[2^\alpha, +, E]$  is equivalent to  $RS(\alpha, \omega_0)$ , in the strong sense that the two theories merely differ in the choice of primitive notions; the binary expansion  $\phi$  yields the translation. Similarly,  $RS(\alpha, \gamma)$  can be reinterpreted as a first order theory. We will state our results in one of the two forms, and leave it to the reader to translate.

**THEOREM 1.** *For any  $\alpha$ , there is a decision method for truth of sentences in  $RS(\alpha, \omega_0)$ . The same sentences are true in  $RS(\alpha, \omega_0)$  and  $RS(\beta, \omega_0)$ , if and only if,  $\alpha = \beta < \omega^\omega$  or else  $\alpha, \beta \geq \omega^\omega$  and have the same  $\omega$ -tail.*

If  $\alpha = z + \omega^y + \omega^n c_n + \dots + \omega^0 c_0$ ,  $y \geq \omega$ , then  $z + \omega^y$  is called the  $\omega$ -head of  $\alpha$ , and  $\omega^n c_n + \dots + \omega^0 c_0$  is called the  $\omega$ -tail of  $\alpha$ .

**THEOREM 2.** *For any ordinals  $\beta > \alpha > \omega^\omega$ ,  $[2^\beta, +, E]$  is an elementary extension of  $[2^\alpha, +, E]$ , if and only if,  $\alpha$  and  $\beta$  have the same  $\omega$ -tail. The elementary embedding is then given by  $h(2^{\alpha_0} x + y) = 2^{\beta_0} x + y$ , whereby  $x < 2^\tau$ ,  $y < 2^{\alpha_0}$ ,  $\tau$  is the common  $\omega$ -tail of  $\alpha$  and  $\beta$ ,  $\alpha_0$  and  $\beta_0$  are respectively the  $\omega$ -heads of  $\alpha$  and  $\beta$ .*

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