## 114. On B-covers and the Notion of Independence in Lattices

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Introduction. In [3], L. M. Kelley has introduced the concept of *B*-covers as metric-between in a normed lattice. We have extended this notion to the case of general lattices in [4] and studied the geometries in lattices by means of *B*-covers and *B*<sup>\*</sup>-covers in [5]. In the first section of this paper we shall show that the relation "relative modularity" or "relative independence" which is derived from Wilcox [1] has a close connection with the *J*-cover or the *CJ*-cover which is a part of the *B*-cover in lattices. In the second section we shall consider the relations between the *B*-covers and independent sets in lattices.

For any two elements a, b of a lattice L, we shall define as follows.

 $J(a,b) = \{x \mid (a \frown x) \smile (b \frown x) = x, x \in L\}, CJ(a,b) = \{x \mid (a \smile x) \frown (b \smile x) = x, x \in L\}.$  J(a,b) is called the *J*-cover of *a* and *b*, and if  $x \in J(a,b)$ , then we shall write J(axb). Similarly we shall define *CJ*-cover and *CJ(axb)*.

 $B(a, b) = J(a, b) \frown CJ(a, b)$  is called the *B*-cover of *a* and *b* and we shall write *axb* when  $x \in B(a, b)$  (cf. [4, 5]).

1. Relative modular pairs and J-covers (CJ-covers). Following L. R. Wilcox [1], (a, b) is called a modular pair when  $x \leq b$  implies  $(x \sim a) \neg b = x \sim (a \neg b)$ , and in this case we write (a, b)M. In [5] we defined a relative modular pair  $(a, b)M^*$  to be a pair (a, b) such that  $a \neg b \leq x \leq b$  implies  $(x \sim a) \neg b = x \sim (a \neg b)$ .

*B*-covers treat "between" in lattices (cf. [4, 5]), while *J*-covers and *CJ*-covers may be considered as describing "semi-between" in lattices.

In the following L is always assumed to be a lattice.

Lemma 1.1. The following statements are equivalent in case  $b' \leq b$ :

- (a) (b' a) b = b' (a b) = b. ((b' a) b = b' (a b) = b').
- (b) J(abb') (CJ(ab'b)).

Proof. If  $(b' \multimap a) \frown b = b' \smile (a \frown b) = b$ , then we have  $(a \frown b) \smile (b \frown b') = (a \frown b) \smile b' = b$ , that is J(abb'). Conversely if J(abb'), then we have  $b = (a \frown b) \smile (b \frown b') \leq b \frown (a \smile b') \leq b$ , and hence we have  $(b' \smile a) \frown b = b = b' \cup (a \frown b)$ . Similarly we can treat the other case.

Theorem 1.1. If J(abb') (resp. CJ(ab'b)) holds for any b' with  $b' \leq b$  then we have (a, b)M.